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1 DataStruct

1.1 基础-单调队列

```
1 <u>int</u> main() {//luogu1440
    rep(i,1,n){
      if(que.empty()) ans[i]=0;
3
      else ans[i]=a[que.front()];
5
      if(!que.empty()&&i-que.front()>=m) que.pop_front();
      while(!que.empty()&&a[i] < a[que.back()]) que.pop_back();</pre>
      que.push_back(i);
    }//min(ai-m,...,ai-1)
9
10 <u>int</u> main(){//luogu2852
11
    rep(i,1,n){
12
      while(!que.empty()&&h[i]<h[que.back()]) que.pop_back();</pre>
13
      que.push_back(i);
      while(!que.empty()&&i-que.front()>=m) que.pop_front();
14
      if(que.empty())ans[i]=0;
15
16
       else ans[i]=h[que.front()];
17
    }//min(ai-m+1,...,ai)
18|}
```

1.2 基础-单调栈

```
1 int main() {//hdu1506
    \underline{cin} >> n;
    rep(i,1,n) <u>cin</u>>>a[i];
4
    ll ans=a[1];
5
    while(!stk.empty()) stk.pop();
6
    rep(i,1,n){
      while(!stk.empty()&&a[stk.top()]>=a[i]) stk.pop();
7
8
       <u>if</u>(!stk.empty()) l[i]=stk.top()+1;
9
      else l[i]=1;
10
       stk.push(i);
11
|12|
    while(!stk.empty()) stk.pop();
13
    per(i,1,n){
14
      while(!stk.empty()&&a[stk.top()]>=a[i]) stk.pop();
15
       if(!stk.empty()) r[i]=stk.top()-1;
16
      else r[i]=n;
17
      stk.push(i);
18
19
    rep(i,1,n) ans=max(ans,a[i]*(r[i]-l[i]+1));
20
     cout<<ans<<endl;</pre>
21|
```

L.3 基础-拉链法哈希表

```
1 const int maxsz=3e5+7;//maxsz 素数表
2|//1e7+19,2e7+3,3e7+23,4e5+9 maxsz 最好为素数
3 //1e6+3,2e6+3,3e6+7,4e6+9,1e5+3,2e5+3,3e5+7
4|//因为是 vector 不需要限制操作次数了
5| //count 操作不增加新节点
6 template typename key, typename val>
  class hash_map{public:
    struct node{key u;val v;int next;};
    vector<node> e;
    int head[maxsz], nume, numk, id[maxsz];
10
    int geths(key &u){return u;}
11
  //geths 是把 key 映射到 [0,maxsz-1] 的函数
12
    bool count(key &u){
13
14
      <u>int</u> hs=geths(u);
15
      for(int i=head[hs];i;i=e[i].next)
       if(e[i].u==u) return 1;
16
      return 0;
17
18
19|
    val& <u>operator</u>[](key &u){//视情况加引用,可能引起 CE
20
      int hs=geths(u);
      for(int i=head[hs];i;i=e[i].next)if(e[i].u==u) return e[i].v;
21
22
      if(!head[hs])id[++numk]=hs;
23
      if(++nume>=e.size())e.resize(nume<<1);</pre>
      return e[nume]=(node){u,0,head[hs]},head[hs]=nume,e[nume].v;
24
25
    void clear(){
```

```
rep(i,0,numk)head[id[i]]=0;
28
      numk=nume=0;
29
    }
30 };
```

基础-可删堆

```
1 // 可删堆
2| //保证 remove 元素被包含在 x 中, 不能多删
  template<typename T>class re_heap{public:
    priority_queue<T> x,y;
5
    void push(T a){x.push(a);}
6
    void remove(T a){y.push(a);}
7
    T top(){
8
      while(y.size()&&x.top()==y.top())
9
      x.pop(),y.pop();return x.top();
10
11
    int size(){return x.size()-y.size();}
12
    void pop(){
13
      while(y.size()&&x.top()==y.top())
|14|
      x.pop(),y.pop();
15
      x.pop();
16
17
    T sectop(){
18
      T a=top();pop();
19
      T b=top();push(a);
20
      return b;
21
22|};
23 re_heap<pii> heap;
```

并查集-非递归 + 按秩合并

```
1 int pre[maxn], rk[maxn];
 2 int find(int a){
    if(pre[a] == -1) return a;
    int t,rt=a;
    while(pre[rt]!=-1) rt=pre[rt];
    while(a!=rt)t=pre[a],pre[a]=rt,a=t;
    return rt;
8
9 #define same(a,b) (find(a)==find(b))
10 void unite(int a, int b){
11
    a=find(a),b=find(b);
12
    if(a==b) return;
13
    <u>if</u>(rk[a]>rk[b])pre[b]=a;
|14|
    else{
15
      pre[a]=b;
16
      <u>if</u>(rk[a]==rk[b])rk[b]++;
17
18
19
  class ufs{public:
    void init(int n){rep(i,0,n) pre[i]=i;}
20
21
    int find(int a){return pre[a] == a?a:pre[a] = find(pre[a]);}
22
    bool same(int a,int b) {return find(a)==find(b);}
23
    void unite(int a,int b){
24
      a=find(a),b=find(b);
25
      if(a==b) return;
26
      pre[b]=a;
27
28 }dsu;
```

并查集-可删除

```
1 #define same(a,b) (find(a)==find(b))
 2 int pre[maxn],id[maxn];
 3
   int numid;
 4 \mid \underline{int} \mid fd(\underline{int} \mid a)  {
 5
       return pre[a] == a?a:pre[a] = find(pre[a]);
 6 }
   int find(int a){return fd(id[a])};
   void unite(int a,int b){
 9
       a=find(id[a]),b=find(id[b]);
       pre[b]=a;
10
11|}
```

Page 3

```
12 <u>void</u> remove(<u>int</u> now){
13
        int pos=find(id[now]);
        id[now] = ++numid;
|14|
15
       pre[numid] = numid;
16|}
```

并查集-可撤销 1.7

```
class ufs{public:
    int fa[maxn],sz[maxn];
    stack<pii> stk;
    void init(int tot){
5
      rep(i,1,tot) fa[i]=i,sz[i]=1;
6
7
    int find(int now){
8
      if(fa[now] == now) return now;
9
      return find(fa[now]);
10
11
    bool unite(int a,int b){
      a=find(a),b=find(b);
12
13
      if(a==b)return 0;
14
      <u>if</u>(sz[a]<sz[b]) fa[a]=b,sz[b]+=sz[a],stk.push(make_pair(a,b));
15
      else fa[b]=a,sz[a]+=sz[b],stk.push(make_pair(b,a));
16
17
18
    void undo(){
19
      if(!stk.empty()){
        auto now=stk.top();
20
21
        stk.pop();
|22|
        fa[now.fi]=now.fi;
23
        sz[now.se] -=sz[now.fi];
24
25
26 }dsu;
```

区间信息-数列分块

```
const int maxn=1e5+10, maxm=sqrt(maxn+1)+10;
  class sqblock{public:
    int id[maxn],sz,cnt;
    struct block{int 1,r;11 sum,tag;}node[maxm];
5
    void init(int n){
6
      sz=sqrt(n+0.5)
|7|
      rep(i,1,n) id[i]=(i-1)/sz+1;cnt=id[n];
8
      rep(i,1,cnt) node[i].l=(i-1)*sz+1;
9
      rep(i,1,cnt-1) node[i].r=i*sz;node[cnt].r=n;
10
      rep(i,1,n) node[id[i]].sum+=arr[i];
11
12
    void update(int s,int t,ll x){
13
      <u>int</u> p1=id[s],p2=id[t];
14
      <u>if</u>(p1==p2)rep(i,s,t)node[p1].sum+=x,arr[i]+=x;
15
        rep(i,s,node[p1].r) arr[i]+=x,node[p1].sum+=x;
16
17
        rep(i,node[p2].1,t) arr[i]+=x,node[p2].sum+=x;
18
        rep(i,p1+1,p2-1) node[i].tag+=x;
19
20
21
    11 query(int s,int t){
22
      <u>int</u> p1=id[s],p2=id[t];ll ans=0;
23
      <u>if</u>(p1==p2) rep(i,s,t) ans+=arr[i]+node[p1].tag;
24
      else {
25
        rep(i,s,node[p1].r) ans+=arr[i]+node[p1].tag;
26
        rep(i,node[p2].1,t) ans+=arr[i]+node[p2].tag;
27
        rep(i,p1+1,p2-1) ans+=node[i].sum+node[i].tag*sz;
28
29
      return ans;
    }
30
31 }square;
```

区间信息-ST 表

```
1 const int maxp=21;
2 class stable [public:
   int logn[maxn],dp[maxp][maxn],pos[maxp][maxn];
```

```
void init(int n=maxn-1){
6
       logn[2]=1;
       rep(i,3,n) logn[i]=logn[i>>1]+1;
8
9
     \underline{\text{void}} \ \text{cal}(\underline{\text{int}} \ *\_a, \underline{\text{int}} \ n) \{//\text{init}(n)
10
       rep(i,1,n) dp[0][i]=a[i],pos[0][i]=i;
|11|
12
       <u>for(int</u> j=1;(1<<j)<=n;j++) <u>for(int</u> i=1;i+(1<<j)-1<=n;++i)
13
         dp[j][i]=min(dp[j-1][i],dp[j-1][i+(1<<(j-1))]);
14
       <u>for(int</u> j=1;(1<<j)<=n;j++) <u>for(int</u> i=1;i+(1<<j)-1<=n;++i)
         <u>if</u>(a[pos[j-1][i]]<=a[pos[j-1][i+(1<<(j-1))]]) pos[j][i]=pos[j-1][i];
15
16
         <u>else</u> pos[j][i]=pos[j-1][i+(1<<(j-1))];
17
18
     inline int query(int 1,int r){
19
       int lg=logn[r-l+1];
20
       return min(dp[lg][l],dp[lg][r-(1<<lg)+1]);</pre>
21
22
     inline int query_pos(int 1,int r){
       int lg=logn[r-l+1];
23
24
       <u>if</u>(a[pos[lg][1]]<=a[pos[lg][r-(1<<lg)+1]])<u>return</u> pos[lg][1];
25
       26
27 }st;
```

1.10 区间信息-树状数组

```
1 //树状数组区间求和,单点修改
  template<typename T> class bit{public:
    T val[maxn];
     \underline{int} lb(\underline{int} x) {\underline{return} x&(-x);}
5
     int len;
     void init(int _n){
      len=_n;
      fill_n(val+1,len+1,INF);
9
10
    void update(int pos,T x){
       for(;pos>0&&pos<=n;pos+=lb(pos)) val[pos]+=x;</pre>
11
12
13
    T psum(<u>int</u> pos){
14
      T sum=0;
15
       for(;pos>0;pos-=lb(pos)) sum+=val[pos];
16
       return sum;
17
18
     T query(int 1,int r){return psum(r)-psum(l-1);}
19
  }tree;
20 //树状数组前缀最小值,支持单点的取 min 覆盖操作
21
  template<typename T> class bit{public:
22
     T val[2*maxn];
23
     \underline{int} lb(\underline{int} x) {\underline{return} x&-x;}
|24|
     int len;
25
     void init(int _n){
26
       len=_n;
27
      fill_n(val+1,len+1,INF);
28
29
     void update(int pos,T x){
       for(;pos<=len;pos+=lb(pos)) val[pos]=min(val[pos],x);</pre>
30
31
32
    T query(int pos){
33
      T res=INF;
       for(;pos;pos-=lb(pos)) res=min(val[pos],res);
34
35
       return res;
36
    }
37 }:
38 bit<<u>int</u>> tree;
```

1.11 区间信息-扫描线算法

```
double dis[maxn];
struct node{
    double x,y1,y2;int tag;
    bool operator <(node &other) const{return x<other.x;}
} seg[maxn];
class segtree{public:
    #define nd node[now]</pre>
```

Page 5

```
#define ndl node[now<<1]</pre>
   #define ndr node[now<<1|1]</pre>
9
10
     struct segnode {
       int l,r,tag;double dis;
11
       inline int mid(){return (r+1)>>1;}
12
13
       inline int len(){return r-l+1;}
     node[maxn << 2|3];
|14|
15
     inline void update(int now){
       if(nd.tag) nd.dis=dis[nd.r+1]-dis[nd.l];
16
17
       else if(nd.len()==1) nd.dis=0;
18
       else nd.dis=ndl.dis+ndr.dis;
19
20
     void maketree(int s,int t,int now=1){
21
       nd={s,t,0,0};
22
       if(s==t) return ;
23
       maketree(s,nd.mid(),now<<1);</pre>
24
       maketree(nd.mid()+1,t,now<<1|1);
25
26
    void update(int s,int t,int x,int now=1){
27
       \underline{if}(s \le nd.l\&t \ge nd.r) {
28
       nd.tag+=x;update(now);
29
       <u>return</u>;
30
31
       <u>if</u>(s<=ndl.r) update(s,t,x,now<<1);</pre>
32
       if(t>ndl.r) update(s,t,x,now<<1|1);</pre>
33
       update(now);
34
35
   }tree;
   int main() {I0; cout << fixed << setprecision(2);</pre>
37
    \underline{\text{while}}((\underline{\text{cin}} >> n) \&\& n) \{
38
    m=0;
39
    rep(i,1,n){
40
       <u>double</u> a,b,c,d;<u>cin</u>>>a>>b>>c>>d;
41
       dis[++m]=b; seg[m]={a,b,d,1};
42
       dis[++m]=d; seg[m]=\{c,b,d,-1\};
43
     sort(dis+1,dis+m+1);
44
45
     sort(seg+1,seg+m+1);
46
     int cnt=unique(dis+1,dis+1+m)-dis-1;
47
     tree.maketree(1,cnt);
48
     double ans=0;
49
     rep(i,1,m-1){
50
       int l=lower_bound(dis+1,dis+1+cnt,seg[i].y1)-dis;
51
       int r=lower_bound(dis+1,dis+1+cnt,seg[i].y2)-dis;
52
53
       if(l<=r) tree.update(l,r,seg[i].tag);</pre>
54
       ans+=tree.node[1].dis*(seg[i+1].x-seg[i].x);
55
     cout<<"Test case #"<<++casn<<endl;</pre>
56
57
     cout<<"Total explored area: "<<ans<<endl<<endl;</pre>
58
59 }
```

1.12 区间信息-动态开点线段树

```
1 int root;
  const 11 rangel=0,ranger=1e9+10;
3 class dsegtree{public:
 4 #define nd node[now]
5 #define ndl node[node[now].son[0]]
 6 #define ndr node[node[now].son[1]]
    struct dsegnode{
      int son[2],mx,tag;
      void update(int x){mx+=x,tag+=x;}
    node[maxn*50];
11
     int cnt;
    void pushup(int now){nd.mx=max(ndl.mx,ndr.mx);}
12
13
    void pushdown(int now){
14
      \underline{if}(nd.tag){
15
        \underline{if}(!nd.son[0]) nd.son[0]=++cnt;
16
        <u>if</u>(!nd.son[1]) nd.son[1]=++cnt;
17
        ndl.update(nd.tag),ndr.update(nd.tag);
18
        nd.tag=0;
19
```

```
20
21
     void update(ll s,ll t,int x,ll l=rangel,ll r=ranger,int &now=root){
       \underline{if}(!now) \{now=++cnt;nd=\{0,0,0,0\};\}
22
23
       if(s<=1\&\&t>=r){
24
         nd.update(x);
25
         return ;
26
27
      pushdown(now);
28
       if(s <= ((1+r) >> 1)) update(s,t,x,1,(1+r) >> 1,nd.son[0]);
29
       <u>if</u>(t>((l+r)>>1)) update(s,t,x,1+((l+r)>>1),r,nd.son[1]);
30
      pushup(now);
31
32| }tree;
```

1.13 区间信息-线段树优化建图

```
1 int cnt;
  class segtree{public:
  #define nd node[now]
  #define ndl node[now<<1]</pre>
  #define ndr node[now<<1|1]</pre>
    int flag;//1==intree,0==outtree
    struct segnode {
8
       int l,r,id;
9
       inline int mid(){return (r+1)>>1;}
10
      inline int len(){return r-l+1;}
    };
11
|12|
    segnode node[maxn<<2|3];
13
    vector<int> v;
|14|
    void init(int n,int flag){
       this->flag=flag;
15
16
      maketree(1,n);
17
18
    void pushup(int now){
19
      if(!flag){
20
        g.add(nd.id,ndl.id,0);
21
        g.add(nd.id,ndr.id,0);
22
       }else {
23
        g.add(ndl.id,nd.id,0);
24
        g.add(ndr.id,nd.id,0);
25
26
27
    void maketree(int s,int t,int now=1){
28
      nd={s,t,++cnt};
29
      \underline{if}(s==t){}
30
        if(flag) g.add(s,nd.id,0);
        else g.add(nd.id,s,0);
31
32
        return ;
33
34
      maketree(s,nd.mid(),now<<1);</pre>
35
      maketree(nd.mid()+1,t,now << 1|1);
36
      pushup(now);
37
38
    vector<int> query(int s,int t){v.clear();find(s,t);return v;}
39
    void find(int s,int t,int now=1){
40
      if(s<=nd.l&&t>=nd.r) {
41
        v.emplace_back(nd.id);
42
        return ;
43
44
       <u>if</u>(s<=ndl.r) find(s,t,now<<1);
45
       \underline{if}(t>ndl.r) find(s,t,now<<1|1);
46
47 }intree, outree;
```

1.14 区间信息-线段树合并

```
int a[maxn],cnt,root[maxn];
vector<int> g[maxn];
ll ans[maxn];
class dsegtree{public:
    #define nd node[now]
    #define ndr node[now].son[0]]
    #define ndr node[node[now].son[1]]
    struct dsegnode {
```

```
int son[2],cnt,id,ans;
10
     node[maxn*50];
11
     void pushup(int now){
12
       if(ndl.cnt>ndr.cnt){
13
        nd.cnt=ndl.cnt;
14
        nd.id=ndl.id;
15
        nd.ans=ndl.ans;
16
       }else if(ndr.cnt>ndl.cnt){
17
        nd.cnt=ndr.cnt;
18
        nd.id=ndr.id;
19
        nd.ans=ndr.ans;
20
       }<u>else</u> {
21
        nd.cnt=ndr.cnt;
|22|
        nd.id=ndr.id;
23
        nd.ans=ndr.ans+ndl.ans;
24
25
     }
26
     void update(int 1,int r,int pos,int &now){
27
       if(!now) now=++cnt;
28
      <u>if</u>(l==r){
29
        nd.id=nd.ans=1;
30
        nd.cnt+=1;
31
        <u>return</u>;
32
       <u>int</u> mid=(l+r)>>1;
33
34
       if(pos<=mid) update(1,mid,pos,nd.son[0]);</pre>
35
       else update(mid+1,r,pos,nd.son[1]);
36
      pushup(now);
37
38
     int merge(int now,int b,int l,int r){
39
      if(!now||!b) return now^b;
       \underline{if}(l==r){}
40
        nd.id=nd.ans=1;
41
42
        nd.cnt+=node[b].cnt;
43
        return now;
44
45
      nd.son[0] = merge(nd.son[0], node[b].son[0], 1, (1+r)/2);
      nd.son[1] = merge(nd.son[1], node[b].son[1], (1+r)/2+1,r);
46
47
      pushup(now);
48
       return now;
49
50| }tree;
51
  void dfs(int now,int fa){
52
     for(int to:g[now]){
53
       if(to==fa) continue;
54
      dfs(to,now);
      tree.merge(root[now],root[to],1,1e5);
55
56
57
     tree.update(1,1e5,a[now],root[now]);
58
     ans[now] = tree.node[root[now]].ans;
59|}
  int main() {
60
61
     cin >> n;
62|
     rep(i,1,n){
63
        <u>cin</u>>>a[i];
64
       root[i]=i;
    }
65
66
     cnt=n;
67
    rep(i,2,n){
68
       <u>int</u> a,b;<u>cin</u>>>a>>b;
69
       g[a].push_back(b);
70
      g[b].push_back(a);
71
    dfs(1,0);
72
73
     rep(i,1,n) <u>cout</u><<ans[i]<<' ';
```

1.15 区间信息-二维线段树

```
class sstree{public:
    #define nd node[nowx][nowy]

struct segnode {int val;};
    int n,m,x1,x2,y1,y2,x,nowx;int segnode node[maxn][maxn];
```

```
void init(int nn,int mm) {
|7|
          n=nn, m=mm;
8
          memset(node ,0,sizeof node);
9
10
      void update(int xx1,int xx2,int yy1,int yy2){
11
          x1=xx1, y1=yy1, x2=xx2, y2=yy2;
12
          updatex(1,n);
13
14
      void updatey(int 1,int r,int nowy=1){
15
          if(y1>r||y2<1) return;
          <u>if</u>(y1<=1&&y2>=r){nd.val^=1;<u>return</u> ;}
16
          updatey(1,(1+r)>>1,nowy<<1); updatey(((1+r)>>1)+1,r,nowy<<1|1);
17
18
19
      void updatex(int 1,int r,int now=1){
20
          if(x1>r||x2<1) return;
21
          <u>if</u>(x1<=1&&x2>=r) {nowx=now;updatey(1,m);<u>return</u> ;}
22
          updatex(1,(1+r)>>1,now<<1); updatex(((1+r)>>1)+1,r,now<<1|1);
23
      int query(int xx,int yy){
24
25
          x1=xx,y1=yy;ans=0;
26
          queryx(1,n);
27
          return ans;
28
29
      void queryy(int l,int r,int nowy=1){
30
          if(y1>r||y1<1) return ;</pre>
          ans^=nd.val;
31
32
          if(l==r) return ;
33
          queryy(1,(1+r)>>1,nowy<<1);queryy(((1+r)>>1)+1,r,nowy<<1|1);
34
35
      void queryx(int 1,int r,int now=1){
36
          \underline{if}(x1>r||x1<1) \underline{return};
37
      nowx=now;queryy(1,m);
38
          if(l==r) return ;
39
          queryx(1,(1+r)>>1,now<<1); queryx(((1+r)>>1)+1,r,now<<1|1);
40
41| }tree;
```

1.16 可持久化数据结构-01 字典树

```
int rt[maxn];
  class ptrie{public:
     int node[maxn*40][2],top
    void init(){rt[0]=node[0][1]=node[0][0]=top=0;}
5
    int add(int pre,int val,int bit=31){
      int now=++top;
 6
      if(bit<0) return now;</pre>
      int t=val>>bit&1;
9
      node[now][t]=add(node[pre][t],val,bit-1);
10
      node[now][t^1]=node[pre][t^1];
11
      return now;
12
13
    int query(int now,int pre,int val,int bit=31,int ans=0){
14
      if(bit<0) return ans;</pre>
15
      int t=val>>bit&1;
      if (node [now] [t^1]-node [pre] [t^1]>0) return query (node [now] [t^1], node [pre] [t^1], val, bit-1, ans | (1<<bit
16
      return query(node[now][t],node[pre][t],val,bit-1,ans);
17
19
20 }tree;
```

1.17 可持久化数据结构-主席树

```
l int rt[maxn];//树根
2 class ptree{public:
3 #define nd node[now]
4 #define mid (s+t)/2
6 int cnt;
7 struct segnode{int l,r,sum;}node[maxn*30];
8 void maketree(int s,int t,int &now=rt[0]){
9 now=++cnt;nd={s,t,0};
10 if(s==t) return;
11 maketree(s,mid,nd.1); maketree(mid+1,t,nd.r);
```

```
12
13
    void update(int pos,int val,int s,int t,int &now,int pre){
14
      now=++cnt;nd=ndp;nd.sum+=val;
15
       <u>if</u>(s==t) <u>return</u> ;
16
       if(pos<=mid) update(pos,val,s,mid,nd.1,ndp.1);</pre>
17
      else update(pos,val,mid+1,t,nd.r,ndp.r);
18
19
    11 query(int 1, int r, int s, int t, int now, int pre){
20
       if(l<=s&&r>=t)return nd.sum-ndp.sum;
21
      11 sum=0;
22
       if(l<=mid) sum+=query(l,r,s,mid,nd.l,ndp.l);</pre>
23
       if(r>mid) sum+=query(l,r,mid+1,t,nd.r,ndp.r);
24
      return sum;
25
    }
26
  #undef mid
27| }tree;
```

1.18 可持久化数据结构-主席树区间第 k 大

```
#include<bits/stdc++.h>
   #define rep(ii,a,b) for(int ii=a;ii<=b;++ii)</pre>
3 #define all(x) x.begin(),x.end()
4 using namespace std;//head
5 <u>const</u> <u>int</u> maxn=2e5+10, maxm=2e6+10;
6 int casn,n,m,k;
7 <u>int</u> rt[maxn];//树根
8 class ptree{public:
9 #define nd node[now]
10 #define ndp node[pre]
11 \frac{\text{#define}}{\text{mid (s+t)}/2}
12
     int cnt;
13
     struct segnode{int 1,r,sum;}node[maxn*20];
|14|
    void init(){cnt=0;}
15
     void update(int pre,int &now,int pos,int s=1,int t=k){
16
       now=++cnt;nd=ndp;nd.sum++;
17
       <u>if</u>(s==t) <u>return</u> ;
18
       if(pos<=mid) update(ndp.1,nd.1,pos,s,mid);</pre>
19
       else update(ndp.r,nd.r,pos,mid+1,t);
20
21
     int query(int pre,int now,int pos,int s=1,int t=k){
22
       if(s==t) return s;
23
       int sum=node[nd.1].sum-node[ndp.1].sum;
24
       if(pos<=sum) return query(ndp.1,nd.1,pos,s,mid);</pre>
25
       else return query(ndp.r,nd.r,pos-sum,mid+1,t);
26
27 #undef mid
28 }tree;
29 int a [maxn];
30 vector<int>pos;
31 int main() {IO;
32
    \underline{cin} >> n >> m;
33
    rep(i,1,n) <u>cin</u>>>a[i];
34
    rep(i,1,n) pos.push_back(a[i]);
35
     sort(all(pos));
36
     pos.erase(unique(all(pos)),pos.end());
37
     k=pos.size();
38
    rep(i,1,n){
39
       int id=lower_bound(all(pos),a[i])-pos.begin();
40
       tree.update(rt[i-1],rt[i],id+1);
41
42
     \underline{\text{while}}(m--) {
43
       <u>int</u> a,b,c;<u>cin</u>>>a>>b>>c;
       <u>cout</u><<pos[tree.query(rt[a-1],rt[b],c)-1]<<'\n';
44
45
    return 0;
46
47|}
```

1.19 可持久化数据结构-bit 套主席树

```
1 namespace dsegtree{
2 #define nd node[now]
3 #define ndl node[node[now].son[0]]
4 #define ndr node[node[now].son[1]]
5 struct dsegnode{int son[2],sum;}node[maxn*200];
```

```
int cnt;
7
     int pos,s,t,x;
     void update(int 1,int r,int &now){
8
9
       if(!now) now=++cnt;
10
       nd.sum+=x;
11
       if(l==r) return
       <u>if</u>(pos<=((l+r)>>1)) update(l,(l+r)>>1,nd.son[0]);
12
13
       else update(1+((1+r)>>1),r,nd.son[1]);
14
15
     void update_1(int _pos,int _x,int &root){
16
       pos=_pos,x=_x;
17
       update(1,n,root);
18
19
     int query(int 1,int r,int now){
20
       if(!now||nd.sum==0)return 0;
21
       if(s<=l&&t>=r) return nd.sum;
22
       int mid=(l+r)>>1;
23
       \underline{\mathbf{if}}(\mathrm{nd.son}[0]\&\&s \le \mathrm{mid})\mathrm{sum} + \mathrm{query}(1,\mathrm{mid},\mathrm{nd.son}[0]);
|24|
25
       if(nd.son[1]&&t>mid) sum+=query(mid+1,r,nd.son[1]);
26
       return sum;
27
28
     int query_1(int _s,int _t,int root){
29
       s=_s, t=_t;
       return query(1,n,root);
30
31
32| \}
33
  namespace bit{
34
     int node[maxn];
35
     inline void update(int pos,int y,int val){
36
       for(;pos<=n;pos+=pos&(-pos))</pre>
37
         dsegtree::update_1(y,val,node[pos]);
38
39
     <u>inline</u> <u>int</u> ask(<u>int</u> pos, <u>int</u> x, <u>int</u> y){
40
       int sum=0;
41
       for(;pos;pos==pos&(-pos))
42
         sum+=dsegtree::query_1(x,y,node[pos]);
43
       return sum;
44
45
     inline int query(int 1,int r,int x,int y){
46
       <u>return</u> ask(r,x,y)-ask(l-1,x,y);
47
48 }
49 int aa[maxn];
50 int main() {
51
     \underline{cin} >> n >> m;
52
     rep(i,1,n) {
53
       cin>>aa[i];
54
       <u>if</u>(aa[i]!=aa[i-1]) bit::update(i,aa[i],1);
55
56
     register int x,y,b,c,d,e,a;
     while (m--){
57
58
       cin>>a;
59
       <u>if</u>(a==1){
60
         cin>>x>>y;
         if(aa[x]==y) continue;
61
62
         if(aa[x]!=aa[x-1]) bit::update(x,aa[x],-1);
63
         if(x+1<n\&\&aa[x+1]!=aa[x]) bit::update(x+1,aa[x+1],-1);
64
         aa[x]=y;
65
         if(aa[x]!=aa[x-1]) bit::update(x,aa[x],1);
         \underline{if}(x+1<n\&\&aa[x+1]!=aa[x]) bit::update(x+1,aa[x+1],1);
66
67
68
         <u>cin</u>>>b>>c>>d>>e;
69
         int ans=bit::query(b,c,d,e);
70
         \underline{if}(aa[b] == aa[b-1] \&\&aa[b] >= d\&\&aa[b] <= e) ans++;
71
         cout << ans << end1
72
73
```

1.20 可持久化数据结构-数组

```
2| <u>int</u> a0 [maxn];//初始数组
3 int rt[maxn];//树根
  template<typename T>class parray{public:
  #define nd node[now]
6
  #define ndp node[pre]
  #define mid (s+t)/2
8
    int cnt;
    <u>static</u> <u>int</u> 10,r0;//数组区间
9
10
    struct segnode{int l,r;T val;}node[maxn*20];
    void init(int s,int t,T a[]){
11
12
       10=s,r0=t,cnt=0;
13
      makearray(a);
14
15
    void makearray(T a[], int &now=rt[0], int s=10, int t=r0){
16
      now=++cnt:
       if(s==t){nd.val=a[s];return;}
17
18
      makearray(a,nd.l,s,mid);makearray(a,nd.r,mid+1,t);
19
20
    void update(int pre,int &now,int pos,T val,int s=10,int t=r0){
21
      now=++cnt;nd=ndp;
22
      if(s==t){nd.val=val;return;}
23
       if(pos<=mid) update(ndp.1,nd.1,pos,val,s,mid);</pre>
24
       else update(ndp.r,nd.r,pos,val,mid+1,t);
25
    T query(<u>int</u> now,<u>int</u> pos,<u>int</u> s=10,<u>int</u> t=r0){
26
27
       if(s==t) return nd.val;
       if(pos<=mid) return query(nd.1,pos,s,mid);</pre>
28
29
       else return query(nd.r,pos,mid+1,t);
30
31 <u>#undef</u> mid
32|};
33 parray<<u>int</u>>arr;
```

1.21 可持久化数据结构-并查集

```
1|//可持久化并查集,主席树实现
2 int rt[maxn];
3 <u>int</u> 10,r0;
4 <u>class</u> pdsu{<u>public</u>:
5 #define nd node[now]
6 #define ndp node[pre]
7 \mid \frac{\text{#define}}{\text{mid}} \mid \text{mid} \mid (\text{s+t})/2
Q.
    struct segnode{int 1,r,fa,dep;}node[maxn*30];
10
    void init(int n){
11
      10=1, r0=n, cnt=0;
|12|
      makearray();
13
    void makearray(int &now=rt[0],int s=10,int t=r0){
14
15
      now=++cnt;nd=\{0,0,s,0\};
16
       if(s==t)return;
17
      makearray(nd.1,s,mid);makearray(nd.r,mid+1,t);
18
    void merge(int pre,int &now,int pos,int fa,int s=10,int t=r0){
19
20
      now=++cnt;nd=ndp;
21
       if(s==t){nd.fa=fa; return;}
22
       if(pos<=mid) merge(ndp.1,nd.1,pos,fa,s,mid);</pre>
23
       else merge(ndp.r,nd.r,pos,fa,mid+1,t);
24
    void update(int now,int pos,int s=10,int t=r0){
25
26
       if(s==t) {nd.dep++;return;}
27
       if(pos<=mid) update(nd.1,pos,s,mid);</pre>
28
      else update(nd.r,pos,mid+1,t);
29
30
    int query(int now,int pos,int s=10,int t=r0){
31
      if(s==t) return now;
32
       if(pos<=mid) return query(nd.1,pos,s,mid);</pre>
33
       else return query(nd.r,pos,mid+1,t);
34
  //找到第 ver 个版本的集合根:
35
36
    int find(int ver,int pos){
37
       int now=query(ver,pos);
38
      if(nd.fa==pos) return now;
```

```
return find(ver,nd.fa);
|40|
  //在 ver1 的基础上,合并 a,b 集合,得到 ver2:
41
    void unite(int ver1,int ver2,int a,int b){
42
       rt[ver2]=rt[ver1];
43
       int fa=find(rt[ver2],a),fb=find(rt[ver2],b);
44
       if(node[fa].fa==node[fb].fa) return;
45
       if(node[fa].dep>node[fb].dep) swap(fa,fb);
46
       merge(rt[ver1],rt[ver2],node[fa].fa,node[fb].fa);
47
48
       if (node[fa].dep==node[fb].dep)
49
         update(rt[ver2],node[fb].fa);
50
  //复制 ver1 的状态到 ver2, 并查询 a,b 是否为同集合
51
    bool same(int ver1,int ver2,int a,int b){
52
      rt[ver2]=rt[ver1];
53
54
      int fa=find(rt[ver2],a),fb=find(rt[ver2],b);
      return node[fa].fa==node[fb].fa;
55
56
  //令 ver2 的状态回到 ver1:
57
    void popback(int ver1,int ver2){
58
59
      rt[ver2]=rt[ver1];
60
61 #undef mid
62| }dsu;
```

1.22 平衡树-老司机树

```
class odt{public:
    struct segnode{
      bool operator<(const segnode &b)const {return l<b.1;}</pre>
5
6
    set < segnode > nd;
    void init(int n=maxn-5){nd.insert(1,n,0);}
    #define iter set<segnode>::iterator
9
    auto split(int pos){
10
      auto it=nd.lower_bound({pos,pos,0});
      if(it!=nd.end()&&it->l==pos) return it;
11
12
13
      int l=it->l,r=it->r,val=it->val;
14
      nd.erase(it);nd.insert({1,pos-1,val});
15
      return nd.insert({pos,r,val}).fi;
16
17
    void update(int 1,int r,int val){
18
      auto itr=split(r+1);auto itl=split(l);
19
      nd.erase(itl,itr);nd.insert({1,r,val});
20
21
    int query(int pos){
22
         auto it=nd.lower_bound({pos,pos,0});
23
          if(it!=nd.end()&&it->l==pos) return it->val;
24
          it--; return it->val;
25
26
  }tree;
```

平衡树-Treap

```
class splaytree{public:
   #define nd node[now]
   #define ndl node[node[now].son[0]]
   #define ndr node[node[now].son[1]]
5
       struct splaynode{
           int son[2],fa,val,size;
           splaynode(){size=1,fa=son[0]=son[1]=0;}
 7
       };
9
       int cnt,root;
10
       vector<splaynode> node;
       <u>inline</u> <u>void</u> pushup(<u>int</u> now){nd.size=ndl.size+ndr.size+1;}
11
12
       <u>inline</u> <u>void</u> pushdown(<u>int</u> now){}
13
       <u>inline</u> <u>int</u> wh(<u>int</u> now){<u>return</u> node[nd.fa].son[1]==now;}
14
       void rotate(int now){
15
           int fa=nd.fa,gf=node[fa].fa,c=wh(now);
16
           pushdown(fa);pushdown(now);
17
           if(gf) node[gf].son[wh(fa)]=now;
18
           nd.fa=gf;
```

```
node[fa].son[c]=nd.son[c^1];
          node [node [fa] . son [c]] .fa=fa;nd.son [c^1]=fa;node [fa] .fa=now;
20
21
          pushup(fa);pushup(now);
22
23
      void splay(int now,int dst=0){
          for(;nd.fa!=dst;rotate(now))
24
25
              if(node[nd.fa].fa!=dst)rotate(wh(now)==wh(nd.fa)?nd.fa:now);
26
              if(!dst) root=now;
27
28
      void insert(int pos){
29
          int now=root,fa=0,val=node[pos].val;
30
          while(now) fa=now,now=val<nd.val?nd.son[0]:nd.son[1];</pre>
31
32
          node[fa].son[val>node[fa].val]=now;
33
          nd.fa=fa;
34
          splay(now);
35
36
      void order(int now){
37
          int l=nd.son[0],r=nd.son[1];
38
          nd.son[0]=nd.son[1]=nd.fa=0;
39
          nd.size=1;
40
          <u>if</u>(1) order(1);
41
          insert(now);
42
          if(r) order(r);
43
44
      void merge(int a,int b){
          if(a==b) return ;
45
46
          splay(a);splay(b);
          if(node[a].size>node[b].size) swap(a,b);
47
48
          pre[a]=b;root=b;
49
          order(a);
50
51
      <u>int</u> kth(<u>int</u> now, <u>int</u> k){
          splay(now); int lsize=0;
52
53
          while(now){
54
              int lsum=lsize+ndl.size;
55
              if(k<=lsum) now=nd.son[0];</pre>
              else if(k==lsum+1) return now;
56
57
              else lsize=lsum+1,now=nd.son[1];
          }
58
59
          return -1;
60
61
      splaytree(int n){
62
          node.resize(n+7,splaynode());
63
          rep(i,1,n) node[i].val=val[i];
64
          node[0].size=0;
65
          root=0,cnt=0;
66
67| };
```

1.24 平衡树-Splay

```
1| <u>int</u> root;//树根
  class splaytree{public:
    int fa[maxn], son[maxn][2], sz[maxn], val[maxn], cnt[maxn];
    int tot;//根权值重复次数子树大小
    inline void pushup(int now){
 6
      sz[now] = sz[son[now][0]] + sz[son[now][1]] + cnt[now];
7
    }//更新当前节点信息
    <u>inline</u> <u>bool</u> getson(<u>int</u> now) {<u>return</u> now==son[fa[now]][1];}//真为右儿子
8
9
    inline void clear(int now){
10
      fa[now] = son[now] [0] = son[now] [1] = sz[now] = val[now] = cnt[now] = 0;
11
    }//清空节点
12
    inline void rotate(int now){
      int f=fa[now],gf=fa[fa[now]],flag=getson(now);
13
14
      son[f][flag]=son[now][flag^1];
      fa[son[now][flag^1]]=f;
15
      son[now][flag^1]=f;
16
      fa[f]=now;fa[now]=gf;
17
      if(gf) son[gf][f==son[gf][1]]=now;
18
19
      pushup(now);pushup(f);
20
    }//旋转一层
21
    void splay(int now){
```

```
for(int f=fa[now];f=fa[now],f;rotate(now))
23
        if(fa[f]) rotate(getson(now)==getson(f)?f:now);
24
      root=now;
    }//旋转到根
25
    void insert(int x,int now=root,int f=0){
26
      if (!now) {
27
28
        now=++tot;
29
        val[now] = x, cnt[now] ++;
30
        if(!root)root=now;
        pushup(now);
31
32
        if(f){
33
          fa[now]=f;
34
          son[f][val[f] < x] = now;
35
          pushup(f);splay(now);
36
37
      }else if(val[now]==x){
38
        ++cnt[now];
39
        pushup(now);pushup(f);
40
        splay(now);
      }else insert(x,son[now][val[now]<x],now);</pre>
41
    }//插入新点
42
    int getrank(int x,int now=root,int ans=0){
43
      while(1){
44
        if(x<val[now]) now=son[now][0];</pre>
45
46
        else {
47
          ans+=sz[son[now][0]];
48
          if(x==val[now]) return splay(now),ans+1;
49
          ans+=cnt[now];now=son[now][1];
50
      }
51
    }//多少个元素小于x
52
53
    int get(int k,int now=root){
      \underline{\text{while}}(1){
54
55
        \underline{if}(son[now][0]\&\&k \le sz[son[now][0]]) now=son[now][0];
56
57
          k-=cnt[now]+sz[son[now][0]];
          if(k<=0) return val[now];</pre>
58
59
          now=son[now][1];
60
      }
61
    }//查询元素
62
63
    int pre(){
      int now=son[root][0];
64
      while(son[now][1]) now=son[now][1];
65
66
      return now;
    }//前驱
67
68
    int nxt(){
      int now=son[root][1];
69
70
      while(son[now][0]) now=son[now][0];
71
      return now;
    }//后缀
72
73
    int lower(int x){
74
      insert(x);
75
      int ans=val[pre()];
76
      erase(x); return ans;
    }//查询前驱
77
    int upper(int x){
78
79
      insert(x);
      int ans=val[nxt()];
80
81
      erase(x); return ans;
    }//查询后缀
82
    void erase(int x){
83
84
      getrank(x);
      if (cnt [root] >1) {
85
        --cnt[root];pushup(root);
86
87
      }else if(!son[root][0]&&!son[root][1]){
88
        clear(root);root=0;
89
      }else if(!son[root][0]||!son[root][1]){
90
        int t=root;
91
        if(!son[root][0]) root=son[root][1];
92
        else root=son[root][0];
```

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```
fa[root]=0;clear(t);
94
       }else {
95
         int now=pre(),t=root;
96
         splay(now);
         fa[son[t][1]]=now;
97
98
         son[now][1]=son[t][1];
99
         clear(t);pushup(root);
100|
101
     }//删除元素
102 }tree;
103 | \underline{int}  main() {
104
     cin>>n;
     \underline{\text{while}}(n--){
105
106
       int opt,x;cin>>opt>>x;
107
        if(opt==1) tree.insert(x);
108
        if(opt==2) tree.erase(x);
109
        if(opt==3) cout<<tree.getrank(x)<<endl;</pre>
110
        <u>if</u>(opt==4) <u>cout</u><<tree.get(x)<<<u>endl</u>;
111
        <u>if</u>(opt==5) <u>cout</u><<tree.lower(x)<<<u>endl</u>;
112
        113
114
     return 0;
115|}
```

2 Graph

2.1 基础-链式前向星

```
class graph{public:
    struct node{int to,next,cost;};
    node e[maxm];int head[maxn],nume;
    void init(int n=maxn-5){nume=0;fill_n(head,n+1,0);}
    void add(int a,int b,int c){e[++nume]={b,head[a],c};head[a]=nume;}
}
```

2.2 最短路-Dijkstra 算法

```
1 namespace dij{
     struct road{
3
       int now; ll dis;
      road(<u>int</u> a=0,ll _dis=0):now(a),dis(_dis){}
5
      bool operator<(const road &x)const {return dis>x.dis};
6
     ll dis[maxn];
8
     bool vis[maxn];
     priority_queue<road>que;
void cal(int st,int n=maxn-5){
9
10
      fill_n(dis,n+1,1e18);fill_n(vis,n+1,0);
11
12
       que.emplace(st,0);dis[st]=0;
13
      while(!que.empty()){
         road t=que.top();que.pop();
14
15
         for(auto e:g[t.now]){
16
           11 cost=t.dis+e.cost;
17
           if(cost<dis[e.to]){</pre>
18
             dis[e.to]=cost;
19
             que.emplace(e.to,cost);
20
21
22
23
     }
24
25
   template<typename T> class shortpath{public:
26
     T dis[maxn];
27
     bool vis[maxn];
28
     #define pdi pair<T, int>
     priority_queue<pdi,vector<pdi>,greater<pdi>>que;
void cal(int st,int n,vector<pdi> g[]){
29
30
      fill_n(dis,n+1,INF);fill_n(vis,n+1,0);
31
32
      que.emplace(0,st);dis[st]=0;
33
      while (!que.empty()){
34
         pdi t=que.top();que.pop();
35
         for(auto e:g[t.se]){
36
           T cost=t.fi+e.fi;
37
           if(cost<dis[e.se]){</pre>
            dis[e.se]=cost;
```

```
最短路-Floyd 求最小环
 1 const int inf=500;
2 <u>const</u> <u>int</u> N=105;
3 int n,m,ans,num;
4 <u>int</u> dis[N][N], mapp[N][N], pre[N][N], path[N];
5 void floyd() {
6
     ans=inf;
     for(int k=1; k<=n; k++) {</pre>
       for(int i=1; i<=n; i++) {</pre>
8
9
         <u>for</u>(<u>int</u> j=1; j<=n; j++)
10
           <u>if</u>(i==j||k==i||k==j) {
11
             continue;
12
13
           <u>if</u>(dis[i][j]+mapp[i][k]+mapp[k][j]<ans) {
14
             ans=dis[i][j]+mapp[i][k]+mapp[k][j];
15
             num=0;
16
             path[num++]=k;
17
             int temp=j;
18
             while(temp!=i) {
               path[num++]=temp;
19
20
               temp=pre[i][temp];
21
22
             path[num++]=i;
23
         }
24
25
26
       for(int i=1; i<=n; i++) {</pre>
27
         <u>for</u>(<u>int</u> j=1; j<=n; j++) {
28
           <u>if</u>(i==j||k==i||k==j) {
29
             continue;
30
           <u>if</u>(dis[i][j]>dis[i][k]+dis[k][j]) {
31
32
             dis[i][j]=dis[i][k]+dis[k][j];
33
             pre[i][j]=pre[k][j];
34
35
36
       }
37
     }
38 }
39
   int main() {
40
     cin >> n >> m;
41
     for(int i=1; i<=n; i++) {</pre>
42
       <u>for(int</u> j=1; j<=n; j++) {
         pre[i][j]=i;
43
44
         <u>if</u>(i==j) {
45
           mapp[i][j]=dis[i][j]=0;
46
         } else {
47
           mapp[i][j]=dis[i][j]=inf;
48
       }
49
     }
50
51
     int u,v,w;
52
     for(int i=1; i<=m; i++) {</pre>
53
       cin>>u>>v>>w;
       mapp[u][v]=mapp[v][u]=dis[u][v]=dis[v][u]=min(w,mapp[u][v]);
54
55
56
     floyd();
57
     \underline{if}(ans==inf) {
58
       printf("No solution.\n");
59
     } <u>else</u> {
60
       for(int i=num-1; i>=0; i--) {
61
         printf("%d",path[i]);
         <u>if</u>(i==0) {
62
           printf("\n");
63
         } else {
64
```

```
printf(" ");
66
67
68
69
     return 0;
70|}
```

强连通-求桥

```
1| struct node{int to,nx;}e[maxm],e2[maxm];
  int head[maxn],head2[maxn],nume,nume2;
3 bool bg[maxm]; int dfn[maxn], low[maxn];
  int numc,cnt,vis1[maxn],belong[maxn];
  void add(int a,int b){
    e[++nume]={b,head[a]};head[a]=nume;
7
8
  void add2(int a,int b){
    e2[++nume2]={b,head2[a]};head2[a]=nume2;
9
10|}
11 void tdfs(int now,int in){
12
    dfn[now] = low[now] = ++cnt;
13
    for(int i=head[now];i;i=e[i].nx){
14
      int to=e[i].to;
15
      if(!dfn[to]) {
16
        tdfs(to,i);
        low[now] = min(low[now], low[to]);
17
        if(low[to]>dfn[now]) bg[i]=bg[i^1]=1;
18
19
      }else if(i!=(in^1))low[now]=min(low[now],dfn[to]);
20
21|
22
  void dfs(int now){
23
    belong[now]=numc;
24
    for(int i=head[now];i;i=e[i].nx){
25
      int to=e[i].to;
26
       27
      dfs(to);
28
29 }
30
  pii dfs2(<u>int</u> now,<u>int</u> fa,<u>int</u> d=0){
31
    vis1[now]=1; pii x={d,now};
32
    for(int i=head2[now];i;i=e2[i].nx){
33
      int to=e2[i].to;
34
      if(to==fa) continue;
      pii t=dfs2(to,now,d+1);
35
36
      \underline{if}(x < t) x = t;
37
38
    return x;
39|}
40 int main() {
41
    cin>>casn;
42
    while(casn--){
43
      \underline{cin} >> n >> m;
44
      rep(i,1,n){
45
        belong[i]=head[i]=head2[i]=0;
46
        low[i]=vis1[i]=dfn[i]=0;
47
48
      rep(i,1,m*2+2)bg[i]=0;
49
      cnt=numc=0, nume=nume2=1;
50
      \underline{\text{while}}(m--)
51
        int a,b;cin>>a>>b;
52
        add(a,b);add(b,a);
53
54
      rep(i,1,n) \underline{if}(!dfn[i]) tdfs(i,0);
55
      rep(i,1,n)
56
        if(!belong[i]) {
57
          numc++;dfs(i);
58
59
60
      for(int i=2;i<=nume;i+=2){</pre>
61
        int a=e[i].to,b=e[i^1].to;
62
        if(belong[a]!=belong[b]){
63
          add2(belong[a],belong[b]);
64
```

```
add2(belong[b],belong[a]);
66
      }
67
68
      int c=0;
69
      rep(i,1,numc){
70
        if(vis1[i]) continue;
        int t=dfs2(i,i).se;
71
72
        c=max(c,dfs2(t,t).fi);
73
74
       cout<<k-c<<endl;</pre>
75
76|}
```

2.5 强连通-求割顶

```
1 void init(int n){
     for(int i=0;i<=n;i++){if(i!=0) f[i]=i;}</pre>
    memset(cut, false, n<<2); memset(low, 0, n<<2);</pre>
3
    memset(dfn,0,n<<2); memset(head,0,n<<2);
5
    nume=cnt=0;
6
  void tarjan(int u,int p){
|7|
     low[u]=dfn[u]=++cnt;
8
     int son=0;
10
     for(int i=head[u];i;i=e[i].next){
       int v=e[i].to;
11
       if(v==p)continue;
12
13
       <u>if</u>(!dfn[v]){
|14|
        son++;
15
        tarjan(v,u);
16
        low[u]=min(low[u],low[v]);
17
         \underline{if}(u!=p\&\&low[v]>=dfn[u]){}
           cut[u]=true;
18
19
       } else low[u]=min(low[u],dfn[v]);
20
21
22
     if(u==p&&son>1) cut[u]=true;
23|}
```

2.6 强连通-缩点

```
int casn,n,m,k;
  struct node {int to,next;}e[maxm];int head[maxn],nume;
  inline void add(int a,int b){e[++nume]=(node){b,head[a]};head[a]=nume;}
  namespace tarjan{
    int stk[maxn],top,cnt,dfn[maxn],low[maxn],numc,belong[maxn],vis[maxn];
5
    vector<int>g[maxn];
    void tdfs(int now,int fa){
8
      dfn[now] = low[now] = ++cnt;
9
      stk[top++]=now,vis[now]=1;
10
      for(int i=head[now];i;i=e[i].next){
11
        int to=e[i].to;
|12|
        if(!dfn[to]){tdfs(to,now);low[now]=min(low[now],low[to]);}
13
        else if(vis[to]) low[now]=min(low[now],dfn[to]);
14
        if(to==fa) continue;
15
16
        if(!dfn[to]){tdfs(to,now);low[now]=min(low[now],low[to]);}
17
        else low[now]=min(low[now],dfn[to]);
18
      }
19
20
      <u>if</u>(low[now] ==dfn[now]){
21
        numc++;
22
        int to;
23
        do{to=stk[--top];
24
          belong[to]=numc;
25
          vis[to]=0;
26
        }while(to!=now);
27
28
29
    void makegraph(int n){
30
      for(int i=1;i<=n;i++) if(!dfn[i]) tdfs(i,i);</pre>
31
      rep(i,1,n){
32
        int a=belong[i];
        for(int j=head[i];j;j=e[j].next){
33
```

2.7 强连通-无向图双连通

```
1 int stk[maxn],top,cnt,dfn[maxn],low[maxn],numc,belong[maxn],vis[maxn],sz[maxn];
  struct node {int to,next;}e[maxm];int head[maxn],nume;
  inline void add(int a,int b){e[++nume]=(node){b,head[a]};head[a]=nume;}
  void tdfs(int now,int pre){
    dfn[now] = low[now] = ++cnt;
    stk[top++]=now,vis[now]=1;
    for(int i=head[now];i;i=e[i].next){
8
      int to=e[i].to;
9
      if(to==pre) continue;
      if(!dfn[to]){tdfs(to,now);low[now]=min(low[now],low[to]);}
10
11
      else low[now] = min(low[now], dfn[to]);
12
13
    if(low[now] ==dfn[now]){
14
      numc++;
15
      int to;
16
      do{to=stk[--top];
        belong[to] = numc;
17
18
        sz[numc]++;
19
      }<u>while</u>(to!=now);
20
21|
```

2.8 树-Prim 算法

```
1| //最小生成树 prim 算法
  int head[30],next[200],point[200],val[200],size,dist[30];
  bool vis[30];
  void add (int a,int b, int v){
5
    int i;
6
    for(i=head[a];~i;i=next[i]){
      if(point[i]==b){
8
        if(val[i]>v)val[i]=v;
9
        return;
10
    }
11
|12|
    point[size]=b;
13
    val[size]=v;
    next[size] = head[a];
|14|
15
    head[a]=size++;
16|}
17 struct cmp{
    bool operator()(pii a,pii b){
18
19
      return a.first>b.first;
    }
20
21 };
22
  void prim(int s){
23
    int i,ans=0;
    memset(dist,-1,<u>sizeof</u>(dist));
24
25
    memset(vis,0,sizeof(vis));
26
    priority_queue<pii, vector<pii>, cmp>q;
27
    for (i=head[s];~i;i=next[i]){
      dist[point[i]]=val[i];
28
      q.push(make_pair(dist[point[i]],point[i]));
29
30
31
    dist[s]=0;
32
    vis[s]=1;
33
    while(!q.empty()){
34
      pii u=q.top();
      q.pop();
35
36
      if(vis[u.second])continue;
37
      vis[u.second]=1;
38
      ans+=u.first;
```

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```
48 }
       树-灭绝树算法
 1 int cntin[maxn];
  struct node{int to,next;};
  class graph{public:
    node e[maxn]; int head[maxn], nume;
5
    void init(int n=maxn-5){nume=0;fill_n(head,n+1,0);}
    void add(int a,int b){e[++nume]={b,head[a]};head[a]=nume;}
6
  }inv,nxt,dom;
  class domtree{public://DAG
    int deep[maxn], anc[maxn] [maxp], que[maxn];
10
    vector<int>edge;
    void init(int n=maxn-5){
11
12
      inv.init(n),nxt.init(n),dom.init(n);
13
      edge.clear();
    }
14
15
    void bit(int &x,int h){
16
      for(int i=0;h>0;++i){
17
        if(h&1) x=anc[x][i];
18
        h>>=1;
      }
19
20
21
    void add(int a,int b){nxt.add(a,b);inv.add(b,a);}
22
    int lca(int a,int b){
23
      if(deep[a] < deep[b]) swap(a,b);</pre>
24
      bit(a,deep[a]-deep[b]);
      if(a==b) return a;
25
26
      per(i,0,maxp-1) <u>if</u>(anc[a][i]!=anc[b][i])
27
        a=anc[a][i],b=anc[b][i];
28
      return anc[a][0];
29
    void tpsort(int n){
30
31
      int tp=0,ed=0;
32
      rep(i,1,n) {
33
        if(!cntin[i]) {
34
          que[ed++]=i;
35
          inv.add(0,i);
36
          nxt.add(i,0);
37
          edge.push_back(i);
38
      }
39
40
      while (ed!=tp){
41
        int now=que[tp++];
42
        forn(i,now,inv.head,inv.e){
43
          int to=inv.e[i].to;
44
          cntin[to]--
45
          if(!cntin[to]) que[ed++]=to,edge.push_back(to);
46
47
    }
48
49
    void maketree(int n){
50
      tpsort(n);
51
      for(auto i:edge){
52
        <u>int</u> fa=-1;
53
        forn(j,i,nxt.head,nxt.e){
54
          int to=nxt.e[j].to;
55
          <u>if</u>(fa==-1) fa=to;
56
          else fa=lca(fa,to);
57
        }fa=fa==-1?0:fa;
58
        deep[i]=deep[fa]+1,anc[i][0]=fa;
59
        rep(j,1,maxp-1) anc[i][j]=anc[anc[i][j-1]][j-1];
60
        dom.add(fa,i);
```

```
2.10 树-支配树算法
 1 int root;
  class domtree{public://dom 为最终的支配树,root 为根,cnt 为每个点的支配点的数量
    int dfn[maxn],rev[maxn],anc[maxn];
    int semi[maxn],idom[maxn];
    int fa[maxn],mi[maxn],clo;
 6
    struct node{int to,next;};
    struct graph{
8
      node e[maxn]; int head[maxn], nume;
9
      void init(int n=maxn-5){nume=0;fill_n(head,n+1,0);}
10
      void add(int a,int b){e[++nume]={b,head[a]};head[a]=nume;}
11
    }inv,nxt,dom;
12
    void init(int n=maxn-5){
13
      clo=0;
14
      rep(i,1,n)fa[i]=mi[i]=semi[i]=i,rev[i]=dfn[i]=anc[i]=idom[i]=0;
15
      nxt.init(n),inv.init(n),dom.init(n);
16
    void add(int a,int b){inv.add(b,a),nxt.add(a,b);}
17
    int find(int now){
18
      if(fa[now] == now) return now;
19
20
      int fx=fa[now],y=find(fa[now]);
21
      if(dfn[semi[mi[fx]]]<dfn[semi[mi[now]]])</pre>
22
        mi[now]=mi[fx];
23
      return fa[now] = y;
24
25
    void tdfs(int now){
      dfn[now]=++clo;rev[clo]=now;
26
27
      forn(i,now,nxt.head,nxt.e)<u>if</u>(!dfn[nxt.e[i].to])
28
        anc[nxt.e[i].to]=now,tdfs(nxt.e[i].to);
29
30
    void maketree(int root,int n=maxn-5){
      tdfs(root);
31
32
      per(i,2,n){
        int now=rev[i],tmp=n;
33
34
        forn(i,now,inv.head,inv.e){
35
          int to=inv.e[i].to;if(!dfn[to]) continue;
          if(dfn[to]<dfn[now]) tmp=min(tmp,dfn[to])</pre>
36
37
          else find(to),tmp=min(tmp,dfn[semi[mi[to]]]);
38
39
        semi[now] = rev[tmp]; fa[now] = anc[now];
        dom.add(semi[now],now);
40
        now=rev[i-1];
41
42
        forn(i,now,dom.head,dom.e){
          int to=dom.e[i].to;find(to);
43
          if(semi[mi[to]] == now) idom[to] = now;
44
45
          else idom[to]=mi[to];
46
      }
47
      rep(i,2,n){
48
49
        int to=rev[i];
50
        if(idom[to]!=semi[to]) idom[to]=idom[idom[to]];
51
52
      dom.init(n);
      rep(i,1,n) if(i!=root)dom.add(idom[i],i);
53
54
55| }tree;
```

2.11 树-最小树形图

```
1 //定根最小树形图
2 struct node{int a,b,c;}e[maxm];
3 int in[maxn],pre[maxn],vis[maxn],id[maxn];
4 ll mdst(){
```

```
11 ans=0; int cnt=0,a,b,laz;
6
     while(1){
7
      rep(i,1,n) in[i]=INF,id[i]=vis[i]=0;
8
      rep(i,1,m) \underline{if}(e[i].a^e[i].b\&\&e[i].c < in[e[i].b])
9
        pre[e[i].b]=e[i].a,in[e[i].b]=e[i].c;
10
       in[k]=0;
11
      rep(i,1,n){
12
         if(in[i] == INF) return -1;
13
        ans+=in[i];
         for(a=i;a^k&&vis[a]^i&&!id[a];a=pre[a])vis[a]=i;
14
15
         if(a^k&&!id[a]){
16
          id[a]=++cnt;
17
           for(b=pre[a];a^b;b=pre[b])id[b]=cnt;
18
19
20
       if(!cnt) return ans;
      rep(i,1,n) <u>if</u>(!id[i]) id[i]=++cnt;
21
22
      rep(i,1,m) {
23
        laz=in[e[i].b];
24
        <u>if</u>((e[i].a=id[e[i].a])^(e[i].b=id[e[i].b]))
25
          e[i].c-=laz;
26
27
      n=cnt; k=id[k], cnt=0;
28
29 }
30 <u>int</u> main() {IO;
31
     cin>>n>>m>>k;
32
    rep(i,1,m)<u>cin</u>>>e[i].a>>e[i].b>>e[i].c;
33
     cout<<mdst()<<endl;</pre>
```

2.12 树-RMQ 求 LCA+ 树上链交

```
1| //rmq 求 lca+ 快速求树上链交
 2 const int maxp=18;
3 class graph {public:
     struct node{int to,next;}e[maxn<<1];</pre>
5
     int head[maxn], nume, dfn[maxn], deep[maxn];
6
     int logn[maxn],pos[maxp][maxn],cnt;
     inline void add(int a, int b){e[++nume]={b,head[a]};head[a]=nume;}
8
     void init(int n){rep(i,1,n) head[i]=0;cnt=0,nume=1;}
9
     void cal(int n){
10
       logn[2]=1;
11
       rep(i,3,n) logn[i]=logn[i>>1]+1;
12
       \underline{\text{for}}(\underline{\text{int}} \ j=1;(1<< j)<=n;j++) \ \underline{\text{for}}(\underline{\text{int}} \ i=1;i+(1<< j)-1<=n;++i)\{
13
         int r = i + (1 << (j-1))
14
         <u>if</u>(deep[pos[j-1][i]]<deep[pos[j-1][r]]) pos[j][i]=pos[j-1][i];
15
         <u>else</u> pos[j][i]=pos[j-1][r];
16
     }
17
18
     void dfs(int now=root,int fa=root,int d=1){
19
       dfn[now]=++cnt;deep[now]=d;pos[0][cnt]=now;
20
       forn(i,now){
21
         if(e[i].to==fa) continue;
22
         dfs(e[i].to,now,d+1);pos[0][++cnt]=now;
23
24
25
     inline int lca(int l,int r){
26
       l=dfn[l],r=dfn[r];<u>if</u>(l>r) swap(l,r);
27
       int lg=logn[r-l+1];
28
       <u>if</u>(deep[pos[lg][1]]<deep[pos[lg][r-(1<<lg)+1]])<u>return</u> pos[lg][1];
29
       30
     <u>inline</u> <u>int</u> getdis(<u>int</u> a, <u>int</u> b){<u>return</u> deep[a]+deep[b]-2*deep[lca(a,b)];}
31
32
     void getlca(){dfs();cal(cnt);}
33
     <u>inline</u> <u>bool</u> check(<u>int</u> a,<u>int</u> b){<u>return</u> lca(a,b)==a;}
34
     int getans(int a1,int a2,int b1,int b2){
35
       int ra=lca(a1,a2);
36
       bool f1=check(ra,b1),f2=check(ra,b2);
37
       <u>if</u>(!f1&&!f2) <u>return</u> 0;
38
       <u>if</u>(f1&&f2){
         int rb=lca(b1,b2);
|39|
40
         if(!( check(rb,a1)||check(rb,a2)))return 0;
```

```
<u>int</u> r1=lca(a1,b1),r2=lca(a1,b2);
42
         <u>int</u> r3=lca(a2,b1),r4=lca(a2,b2);
43
         if(r1==r3&&r2==r4) return 1;
         <u>return</u> getdis(r1==ra?r3:r1,r2==ra?r4:r2)+1;
44
45
46
       <u>if</u>(!f1)swap(b1,b2);
47
       int r1=lca(a1,b1),r3=lca(a2,b1);
48
       return getdis(r1==ra?r3:r1,ra)+1;
49
50|}g;
```

2.13 树-点分治

```
1 <u>namespace</u> graph{
    vector<int>g[maxn];
    int all,sz[maxn],root,maxt;
    bool vis[maxn];
5
    int dfs_root(int now,int fa){
      int cnt=1;
      for(auto to:g[now])if(to!=fa&&!vis[to])
8
        cnt+=dfs_root(to,now);
9
      int tmp=max(cnt-1,all-cnt);
10
      if(maxt>tmp) maxt=tmp,root=now;
      return sz[now]=cnt;
11
12
    }//基础部分
13
    int ans[maxn];
14
    void dfs_col(int now,int fa,int c){
15
      ans[now]=c;
16
      for(auto to:g[now])if(to!=fa&&!vis[to])
        dfs_col(to,now,c);
17
18
19
    void dfs_dv(int now,int d=0){
20
      vis[now]=1;dfs_col(now,now,d);
      for(auto to:g[now]){
21
22
        if(vis[to]) continue;
23
        maxt=root=n+1;all=sz[to];
24
        dfs_root(to,now);dfs_dv(root,d+1);
25
26
27
    void solve(int n){
28
      all=maxt=root=n+1;
29
      dfs_root(1,1);
30
      all-=maxt;
      dfs_dv(root);
31
32
33|}
```

2.14 树-轻重链剖分

```
1 class graph{//按边
     struct node{int from,to,cost,next;}e[maxn<<1];</pre>
     int head[maxn], nume, cnt2;
     <u>inline</u> <u>void</u> add(<u>int</u> a, <u>int</u> b, <u>int</u> c){
5
       e[++nume]={a,b,c,head[a]};head[a]=nume;
6
7
    int fa[maxn],sz[maxn],top[maxn],remp[maxn],ans[maxn];
8
     int son[maxn],in[maxn],cnt,deep[maxn];
9
     void dfs1(int now,int pre,int d){
      deep[now] =d; sz[now] =1; fa[now] =pre;
10
11
      for(int i=head[now];i;i=e[i].next){
        if(e[i].to==pre) continue;
12
13
        dfs1(e[i].to,now,d+1);
14
        sz[now] += sz[e[i].to];
         if(sz[son[now]]<sz[e[i].to]) son[now]=e[i].to;</pre>
15
16
17
18
    void dfs2(int now,int pre,int st){
19
      top[now] = st; in[now] = ++cnt; remp[cnt] = now;
20
       if(son[now]) dfs2(son[now],now,st);
21
      for(int i=head[now];i;i=e[i].next)
22
        <u>if</u>(e[i].to!=pre&&e[i].to!=son[now])
23
          dfs2(e[i].to,now,e[i].to);
24
25
    int query(int a,int b){
```

```
int sum=0;
      while(top[a]!=top[b]){
27
        if(deep[top[a]] < deep[top[b]]) swap(a,b);</pre>
28
29
        sum+=tree.query(in[top[a]],in[a]);
30
        a=fa[top[a]];
31
32
      if(a==b)return sum;
33
      if(deep[a]>deep[b]) swap(a,b);
34
      sum+=tree.query(in[son[a]],in[b]);
35
      return sum;
36
37
    void getchain(){dfs1(1,1,0);dfs2(1,1,1);}
38| }g;
39 <u>int</u> root=1;
40 class graph{public://按点
    struct node{int to,next;}e[maxn<<1];</pre>
41
    int head[maxn], nume, mp[maxn];
42
43
    inline void add(int a,int b){
      e[++nume]={b,head[a]};
44
      head[a]=nume;
45
46
47
    int ltop[maxn],fa[maxn],deep[maxn];
    int sz[maxn],remp[maxn];
48
49
    int son[maxn],cnt;
50
    void init(int n){rep(i,1,n) head[i]=0;cnt=0,nume=1;}
51
    void dfs1(int now=root,int pre=root,int d=0){
52
      deep[now]=d,fa[now]=pre,sz[now]=1,son[now]=0;
53
      forn(i,now){
        int to=e[i].to;
54
55
        if(to!=pre) {
56
          dfs1(to,now,d+1);
57
          sz[now] += sz[to];
58
          if(sz[to]>sz[son[now]]) son[now]=to;
59
      }
60
61
62
    void dfs2(int now=root,int pre=root,int sp=root){
63
      ltop[now] = sp; mp[now] = ++cnt; remp[cnt] = now;
64
        if(son[now]) dfs2(son[now],now,sp);
65
        forn(i,now){
66
          int to=e[i].to;
67
          if(to!=son[now]&&to!=pre) dfs2(to,now,to);
68
69
70
    void getchain(){dfs1();dfs2();}
71
    int lca(int x,int y){
      for(;ltop[x]!=ltop[y];deep[ltop[x]]>deep[ltop[y]]?x=fa[ltop[x]]:y=fa[ltop[y]]);
72
73
      return deep[x]<deep[y]?x:y;</pre>
74
    inline int getdis(int a,int b){return deep[a]+deep[b]-2*deep[lca(a,b)];}
75
76
    inline bool check(int a,int b){return dfn[a]<=dfn[b]&&dfn[a]+sz[a]-1>=dfn[b]+sz[b]-1;}
77
    void update(int a,int b,int val){
78
79
      while(ltop[a]!=ltop[b]){
          if(deep[ltop[a]] < deep[ltop[b]])swap(a,b);</pre>
80
81
          tree.update(mp[ltop[a]],mp[a],val);
82
          a=fa[ltop[a]];
83
        if(deep[a]>deep[b])swap(a,b);
84
        tree.update(mp[a],mp[b],val);
85
86
87
    int query(int a,int b,int k){
88
      int sum=0;
89
      while(ltop[a]!=ltop[b]){
          if(deep[ltop[a]] < deep[ltop[b]]) swap(a,b);</pre>
90
91
          sum+=tree.query(mp[ltop[a]],mp[a],k);
92
          a=fa[ltop[a]];
93
94
        if(deep[a]>deep[b])swap(a,b);
95
        sum+=tree.query(mp[a],mp[b],k);
96
        return sum;
    }
97
```

98|}g;

2.15 树-动态点分治

```
1|//动态点分治(点分树)
  template < const int N > class Graph { public :
    vector<int>g[N+10];
4
    bool vis[N+10];
    int all,sz[N+10],root,maxt,father[N+10];
5
 6
    <u>int</u> _deep[N+10],_dfn[N+10],_cnt;
    \underline{int} son [N+10]
         lag[N+10][int(log(N+10)/log(2))+1];
    int logn[N+10];
10
    void add(int a,int b){
11
      g[a].emplace_back(b);g[b].emplace_back(a);
12
13
    int dfs_root(int now,int fa){
14
      int cnt=1;
15
      son[now]=0;
      for(auto to:g[now])if(to!=fa&&!vis[to]){
16
17
        int ch=dfs_root(to,now);
        son[now] = max(son[now],ch);
18
19
        cnt+=ch;
20
21
      son[now] = max(son[now],all-cnt);
|22|
      if(son[now] < son[root]) root=now;</pre>
23
      return sz[now]=cnt;
|24|
25
    void dfs_lca(int now,int fa){
26
      _dfn[now] =++_cnt;
       _lca[_cnt][0]=_deep[now]=_deep[fa]+1;
27
28
      for(auto to:g[now]) {
        if(to==fa) continue;
29
30
        dfs_lca(to,now);
31
        _lca[++_cnt][0]=_deep[now];
32
    }
33
34
    void cal_st(){
35
      \log n[0] = -1;
      rep(i,1,2e5+10) logn[i]=logn[i>>1]+1;
36
37
      rep(j,1,logn[_cnt])rep(i,1,_cnt-(1<<j)+1)
38
         _lca[i][j]=min(_lca[i][j-1],_lca[i+(1<<(j-1))][j-1]);
39
    int _dis(int a,int b){
40
      a=_dfn[a],b=_dfn[b];
41
42
      if(a>b) swap(a,b);
      int len=logn[b-a+1];
43
44
      <u>return</u> min(_lca[a][len],_lca[b-(1<<len)+1][len]);
45
46
    <u>int</u> dis(<u>int</u> a,<u>int</u> b){//两点距离
47
      int res=_deep[a]+_deep[b]-2*_dis(a,b);
48
      return res;
|49|
50
    void dfs_cal(int now,int fa){
51
      for(auto to:g[now]){
        if(to==fa||vis[to]) continue;
52
53
        dfs_cal(to,now);
54
55
56
    void dfs_dv(int now,int fa){
57
      father [now] = fa; vis [now] = 1;
58
      dfs_cal(now,0);
59
      for(auto to:g[now]){
60
        if(vis[to]) continue;
61
        all=sz[to],root=0;
62
        dfs_root(to,0);
63
        int tmp=root;
64
        dfs_dv(root,now);
65
66
67
    void init(int n){//初始化
68
       cnt=0;
69
      dfs_lca(1,0);
```

```
cal_st();
      son[0]=n;
71
72
      all=n;root=0;
73
      dfs_root(1,0);
74
      dfs_dv(root,0);
75
76
    void update(int pos){
77
      for(int now=pos;father[now];now=father[now]){
78
79
80
    }
81| };
82 Graph<1e6> tree;
```

2.16 最大流-ISAP 算法

```
template<typename T>class mxf{public:
     struct node{int to,next;T cap;}e[maxm<<1];</pre>
     int cur[maxn],head[maxn],dis[maxn],gap[maxn];
    int nume=1,s,t,tot;
5
    void init(int n){
6
      rep(i,0,n) head[i]=gap[i]=dis[i]=0;
8
9
    void add(int a,int b,T c){
10
      e[++nume]={b,head[a],c};head[a]=nume;
      e[++nume]={a,head[b],0};head[b]=nume;
11
12
    T dfs(int now,T flow=INF){
13
|14|
       if (now==t||!flow) return flow;
15
      T use=0,tmp;
16
      int d=dis[now]-1,to;
17
      for (int &i=cur[now];i;i=e[i].next) {
18
        <u>if</u>(dis[to=e[i].to]==d&&(tmp=e[i].cap)){
19
          e[i].cap-=(tmp=dfs(to,min(flow-use,tmp)));
20
          e[i^1].cap+=tmp;
21
          if((use+=tmp)==flow) return use;
22
23
24
      if (!--gap[dis[now]]) dis[s]=tot+1;
25
      ++gap[++dis[now]];
      cur[now] = head[now];
26
27
      return use;
28
29
    T getflow(<u>int</u> ss,<u>int</u> tt,<u>int</u> n,T ans=0){
30
      tot=n; s=ss; t=tt; gap[0]=tot;
31
      memcpy(cur,head,(tot+1)<<2);</pre>
32
      while(dis[s]<=tot) ans+=dfs(s);</pre>
33
      return ans;
34
35|};
36 mxf<int> net;
```

2.17 最大流-Dinic 算法

```
1 template < typename T > class mxf {public:
    struct node{int to,next;T cap;}e[maxm<<1];</pre>
    int cur[maxn],head[maxn],que[maxn],dis[maxn];
    int nume=1,s,t,tot,tp,ed;
    inline void adde(int a,int b,T c){e[++nume]={b,head[a],c};head[a]=nume;}
    inline void add(int a,int b,T c){adde(a,b,c);adde(b,a,0);}
    void init(int n=maxn-1){memset(head,0,(n+1)<<2);nume=1;}</pre>
    bool bfs(){
9
      rep(i,0,ed) dis[que[i]]=0;
10
      dis[t]=1,que[0]=t;
11
      tp=0,ed=1;
12
      cur[t]=head[t];
13
      int now, to;
|14|
      while(tp!=ed) for(int i=head[now=que[tp++]];i;i=e[i].next)
15
        if(dis[to=e[i].to]==0\&\&e[i^1].cap>0){
16
          cur[to] = head[to];
17
          dis[to]=dis[now]+1;
18
          if((que[ed++]=to)==s) return true;
19
```

```
<u>return</u> <u>false</u>;
21
    T dfs(<u>int</u> now,T flow=INF){
22
23
       if(now==t||flow==0) return flow;
24
       int to,d=dis[now]-1;
25
      T use=0,tmp;
      for(int &i=cur[now];i;i=e[i].next){
26
27
        if(dis[to=e[i].to]!=d||!(tmp=e[i].cap))continue;
28
        e[i].cap-=(tmp=dfs(to,min(tmp,flow-use)));
29
        e[i^1].cap+=tmp,use+=tmp;
30
         if(use==flow) return use;
31
32
       \underline{if}(use==0) dis[now]=-1;
33
      return use;
34
35
    T getflow(int ss,int tt,int n,T ans=0){
36
       s=ss,t=tt,tot=n;
37
       while(bfs())ans+=dfs(s);
38
       return ans;
39
40 };
41 mxf<<u>int</u>> net;
```

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2.18 网络流-上下界可行流-Dinic

```
11 //上下界可行流
  template<typename T>class mxf{public:
    struct node{int id,to,rev;T cap;}w[maxm<<1],e[maxm<<1];</pre>
     int cur[maxn],head[maxn],dis[maxn],que[maxn];
    T up[maxm<<1],in[maxn],sum,ans[maxm<<1];</pre>
5
 6
     int num[maxn], numv, nume, s, t, tot, tp, ed, fr[maxm<<1];</pre>
    bool bfs(){
      rep(i,0,ed) dis[que[i]]=0;
9
      dis[t]=1,que[0]=t;
10
      tp=0,ed=1;
11
       cur[t]=head[t];
12
       int now,to;
13
      while(tp!=ed) for(int i=head[now=que[tp++]];i<=num[now];++i)</pre>
14
        <u>if</u>(dis[to=e[i].to]==0&&e[e[i].rev].cap>0){
          cur[to]=head[to];
15
16
          dis[to]=dis[now]+1;
17
          if((que[ed++]=to)==s) return true;
18
19
      <u>return</u> <u>false</u>;
20
21
    T dfs(int now,T flow=0x3f3f3f3f){
22
       if(now==t||flow==0) return flow;
23
       int to,d=dis[now]-1;
24
      T use=0,tmp;
25
      for(int &i=cur[now];i<=num[now];++i){</pre>
26
        if(dis[to=e[i].to]!=d||!(tmp=e[i].cap))continue;
27
        e[i].cap-=(tmp=dfs(to,min(tmp,flow-use)));
28
        e[e[i].rev].cap+=tmp,use+=tmp;
29
        if(use==flow) return use;
30
31
      if(use==0) dis[now]=-1;
32
      return use;
33
34
    T getflow(<u>int</u> ss,<u>int</u> tt,<u>int</u> n,T ans=0){
35
       s=ss,t=tt,tot=n;
36
       while(bfs())ans+=dfs(s);
37
      return ans;
38
39
    void init(int n){
40
      rep(i,0,n) num[i]=in[i]=head[i]=dis[i]=0;
41
      nume=0;tot=n;sum=0;
42
43
    void add(int a,int b,T c,int id){
44
      w[++nume] = (node) \{id,b,0,c\}; ++num[a],fr[nume] = a;
45
      w[++nume] = (node) \{0,a,0,0\}; ++num[b], fr[nume] = b;
46
47
    void addbound(int a,int b,T c,T d,int id){
48
      add(a,b,d-c,id);
```

```
up[id]=d,in[b]+=c,in[a]-=c;
50
51
    bool fesbflow(int n){
52
      s=n+1, t=n+2; numv=n; tot=t;
53
      rep(i,1,numv){
54
         <u>if</u>(in[i]>0) add(s,i,in[i],0),sum+=in[i];
         if(in[i]<0) add(i,t,-in[i],0);</pre>
55
56
57
      head[1]=1;
58
      rep(i,2,tot) head[i]=head[i-1]+num[i-1];
59
      rep(i,1,tot-1) num[i]=head[i+1]-1;
60
      num[tot]=nume;
61
      rep(i,1,nume){
62
        e[head[fr[i]]+cur[fr[i]]++]=w[i];
63
        <u>if</u>(!(i%2)){
64
          e[head[fr[i]]+cur[fr[i]]-1].rev=head[w[i].to]+cur[w[i].to]-1;
          e[head[w[i].to]+cur[w[i].to]-1].rev=head[fr[i]]+cur[fr[i]]-1;
65
66
67
68
      T flow=getflow(s,t,t);
69
       if(flow<sum) return 0;</pre>
70
      rep(i,1,nume){
71
        node &x=e[i];
72
        \underline{if}(x.id) ans [x.id] = up[x.id] - x.cap;
73
74
      return 1;
    }
75
76 };
77 mxf<<u>int</u>> net;
```

2.19 网络流-上下界网络流-ISAP

```
template<typename T>class mxf{public:
     struct node{int to,rev;T cap;}w[maxm<<1],e[maxm<<1];</pre>
3
     int cur[maxn],head[maxn],dis[maxn],gap[maxn];
     int num[maxn],numv,nume,s,t,tot,last,fr[maxm<<1];</pre>
5
     T in[maxn],sum;
    T dfs(int now,T flow=INF){
 6
       if (now==t||!flow) return flow;
8
      T use=0,tmp;
9
      int d=dis[now]-1,to;
10
      for (int &i=cur[now];i<=num[now];++i) {</pre>
        <u>if</u>(dis[to=e[i].to]==d&&(tmp=e[i].cap)){
11
12
          e[i].cap-=(tmp=dfs(to,min(flow-use,tmp)));
          e[e[i].rev].cap+=tmp;
13
          <u>if</u>((use+=tmp)==flow) <u>return</u> use;
14
15
16
      }
17
      if (!--gap[dis[now]]) dis[s]=tot+1;
18
      ++gap[++dis[now]];
19
       cur[now] = head[now];
20
      return use;
21
22
    T getflow(<u>int</u> ss,<u>int</u> tt,<u>int</u> n,T ans=0){
23
      rep(i,0,n)dis[i]=gap[i]=0;
24
      tot=n; s=ss; t=tt; gap[0]=tot;
25
      memcpy(cur,head,(tot+1)<<2);</pre>
26
      while(dis[s]<=tot) ans+=dfs(s);</pre>
27
      return ans;
28
29
    void init(int n){
      rep(i,0,n) num[i]=in[i]=head[i]=dis[i]=0;
30
31
      nume=0;tot=n;sum=0;
32
33
    void add(int a,int b,T c){
      w[++nume] = (node) \{b,0,c\}; ++num[a], fr[nume] = a;
34
35
      w[++nume] = (node)\{a,0,0\}; ++num[b],fr[nume] = b;
36
37
    void addbound(int a,int b,T c,T d){
38
      add(a,b,d-c);
39
      in[b]+=c,in[a]-=c;
40
    void makeflow(int n){
41
```

```
s=n+1, t=n+2; numv=n; tot=t;
43
      rep(i,1,numv){
        <u>if</u>(in[i]>0) add(s,i,in[i]),sum+=in[i];
|44|
45
        if(in[i]<0) add(i,t,-in[i]);</pre>
46
47
      head[1]=1;
      rep(i,2,tot) head[i]=head[i-1]+num[i-1];
48
49
      rep(i,1,tot-1) num[i]=head[i+1]-1;
50
      num[tot] = nume;
51
      rep(i,1,nume){
52
        e[head[fr[i]]+cur[fr[i]]++]=w[i];
53
        <u>if</u>(!(i%2)){
54
          e[head[fr[i]]+cur[fr[i]]-1].rev=head[w[i].to]+cur[w[i].to]-1;
          e[head[w[i].to]+cur[w[i].to]-1].rev=head[fr[i]]+cur[fr[i]]-1;
55
56
57
      }
58
    T fesbflow(int n){
59
60
      makeflow(n);
61
      T flow=getflow(s,t,t);
62
      if(flow!=sum) return -1;
63
      return flow;
64
65
    T fesbflow(<u>int</u> ss, <u>int</u> tt, <u>int</u> n){
66
      add(tt,ss,INF);
67
      makeflow(n);
68
      rep(i,head[tt],num[tt])
69
        <u>if</u>(e[i].to==ss&&e[i].cap==INF) {
70
          last=i;
71
          break;
72
        }
73
      T flow=getflow(s,t,t);
74
      if(flow!=sum) return -1;
75
      return flow;
76
77
    T maxflow(int ss,int tt,int n){
78
       if(fesbflow(ss,tt,n)==-1) return -1;
79
      return getflow(ss,tt,n+2);
80
    T minflow(int ss,int tt,int n){
81
82
      if(fesbflow(ss,tt,n)==-1) return -1;
83
      node &x=e[last];
84
      T ans=INF-x.cap;
      x.cap=e[x.rev].cap=0;
85
86
       return ans-getflow(tt,ss,n+2);
87
88|};
89 mxf<int> net;
```

2.20 网络流-SW 全局最小割算法

```
11 //全局最小割
 2 #include <cstdio>
 3 <u>#include</u> <<u>iostream</u>>
 4 #include <cstring>
 5 #include <algorithm>
 6 #include <queue>
   #include <numeric>
   typedef long long LL;
   const int MAXV = 3010;
10 | \underline{\text{const}} | \underline{\text{int}} | \text{MAXE} = 100010 * 2;
11 const int INF = 0x3f3f3f3f;
12 int head [MAXV], val [MAXV], ecnt;
13 int to[MAXE], next[MAXE], weight[MAXE];
14 bool vis[MAXV];
15 int fa[MAXV], link[MAXV];
16 <u>int</u> n, m;
17 <u>void</u> init() {
     memset(head + 1, -1, \underline{\text{sizeof}}(\underline{\text{int}}) * n);
18
19
     memset(link + 1, -1, \underline{\text{sizeof}}(\underline{\text{int}}) * n);
20
     for (int i = 1; i <= n; ++i)</pre>
21
        fa[i] = i;
22
     ecnt = 0;
```

```
|23| }
24 void add_edge(int u, int v, int w) {
     to[ecnt] = v; weight[ecnt] = w; next[ecnt] = head[u]; head[u] = ecnt++;
25
     to[ecnt] = u; weight[ecnt] = w; next[ecnt] = head[v]; head[v] = ecnt++;
26
27 }
28 int findset(int u) {
    return u == fa[u] ? u : fa[u] = findset(fa[u]);
29
30|}
31 void merge(int u, int v) {
32
     int p = u;
33
     while (~link[p]) p = link[p];
34
     link[p] = v;
35
     fa[v] = u;
36|}
37 <u>int</u> MinimumCutPhase(<u>int</u> cnt, <u>int</u> &s, <u>int</u> &t) {
38
    memset(val + 1, 0, \underline{sizeof}(\underline{int}) * n);
39
    memset(vis + 1, 0, \underline{sizeof}(\underline{bool}) * n);
40
     std::priority_queue<std::pair<int, int>> que;
     t = 1;
41
     while (--cnt) {
42
43
       vis[s = t] = \underline{true};
       \underline{\text{for}} (\underline{\text{int}} u = s; ~u; u = link[u]) {
44
45
         for (int p = head[u]; ~p; p = next[p]) {
46
           int v = findset(to[p]);
47
           <u>if</u> (!vis[v])
48
             que.push(std::make_pair(val[v] += weight[p], v));
49
         }
       }
50
       t = 0;
51
52
       while (!t) {
53
         if (que.empty()) return 0;
54
         auto pa = que.top(); que.pop();
55
         if (val[pa.second] == pa.first) t = pa.second;
56
57
     return val[t];
58
59 }
60 int StoerWagner() {
61
     int res = INF;
     \underline{for} (\underline{int} i = n, s, t; i > 1; --i) {
|62|
       res = <u>std</u>::min(res, MinimumCutPhase(i, s, t));
63
       if (res == 0)
64
65
         break;
66
       merge(s, t);
67
     return res;
68
69|}
70 int main() {
     while (scanf("%d%d", &n, &m) != EOF) {
71
72
       init();
73
       \underline{\text{for}} \ (\underline{\text{int}} \ i = 0, u, v, w; i < m; ++i) {
         scanf("%d%d%d", &u, &v, &w);
74
75
         add_edge(u, v, w);
76
77
       printf("%d\n", StoerWagner());
78
79| \}
```

2.21 网络流-最大密度子图

```
1 const double eps=1e-8;
 2 template < typename T > class mxf {public:
    struct node{int to,next;T cap;}e[maxm<<1];</pre>
    int cur[maxn],head[maxn],que[maxn],dis[maxn],nume=1,s,t;
    inline void adde(int a,int b,T c){
      e[++nume]={b,head[a],c};head[a]=nume;
 6
    inline void add(int a, int b, T c) {adde(a,b,c); adde(b,a,0);}
9
    void init(int n=maxn-1){memset(head,0,(n+1)<<2);nume=1;}</pre>
10
    bool bfs(){
11
      memset(dis,-1,(t+1)<<2);
12
      dis[t]=0,que[0]=t;
13
      int tp=0,ed=1;
```

```
while(tp!=ed){
        int now=que[tp++]; if(tp==maxn) tp=0;
15
16
        for(int i=head[now];i;i=e[i].next){
17
          int to=e[i].to;
          <u>if</u>(dis[to]==-1&&e[i^1].cap>0){
18
19
            dis[to]=dis[now]+1;
20
            if(to==s) return true;
21
            que[ed++]=to;
            \underline{if}(ed==maxn) ed=0;
23
        }
24
      }
25
26
      <u>return</u> <u>false</u>;
27
28
    T dfs(<u>int</u> now,T flow=1e9){
      if(now==t||flow==0) return flow;
29
30
      T use=0;
31
      for(int &i=head[now];i&&use!=flow;i=e[i].next){
32
        int to=e[i].to;
33
        if(dis[to]+1!=dis[now])continue;
34
        T tmp=dfs(to,min(e[i].cap,flow-use));
        e[i].cap-=tmp,e[i^1].cap+=tmp,use+=tmp;
35
36
      if(use==0) dis[now]=-1;
37
38
      return use;
39
    T getflow(int ss,int tt){
40
41
      s=ss,t=tt;T ans=0;
      memcpy(cur,head,(t+1)<<2);
42
43
      while(bfs()){
44
        ans+=dfs(s);
        memcpy(head, cur, (t+1) << 2);
45
46
47
      return ans;
48
49 };
50 mxf<double> net;
  const int maxn2=500;
51
52 int mt[maxn2] [maxn2];
53 double d[maxn2];
54 int val[maxn2], tag[maxn2];
55 void init(int n){rep(i,1,n)rep(j,i+1,n) mt[i][j]=mt[j][i]=0;}
56 void adde(int a, int b, int v) {mt[a][b]=mt[b][a]=v;}
  <u>const</u> <u>double</u> all=400*2200;//点权和 + 边权和 *2
57
58 bool check(double mid){
59
    int s=n+1,t=n+2;
60
    net.init(n+3);
61
    double f=0;
    rep(i,1,n)
62
63
      d[i]=0.0;
64
      rep(j,1,n){
65
        if(i==j||!mt[i][j])continue;
66
        d[i]+=mt[i][j];
  //如果公式计算出来, 边权跟 mid 有关, 就要加上相应的 mid
67
68
        net.add(i,j,mt[i][j]);
69
70
71
    rep(i,1,n){
72
      if(tag[i]){
73
        f+=all+2*mid-d[i];
74
        net.add(s,i,INF);
75
      }else {
76
        net.add(s,i,all);
77
        net.add(i,t,all+2*mid-d[i]);
   //有点权的话,这个 2*mid 还要再乘那个点权
78
79
      }
80
81
    double x=net.getflow(s,t);
82
    double ans=(all*n-f-x)*0.5;
83
    return ans>eps;
84|}
85 <u>int</u> main(){IO;
```

```
cin>>casn;
87
      cout << fixed << setprecision(10);</pre>
88
     rep(kase,1,casn){
89
       \underline{cin} >> n;
90
       init(n);
91
       rep(i,1,n)<u>cin</u>>>val[i];
92
       rep(i,1,n)rep(j,i+1,n)
93
         if(val[j]<val[i]) adde(i,j,1);</pre>
94 //
         rep(i,1,n) tag[i]=0;//是否必须用
95
       double l=0,r=n,mid;
96
       while(r-l>=eps){
         mid=(1+r)*0.5;
97
98
          if(check(mid)) l=mid;
99
         else r=mid;
100
       cout << "Case #"<<kase<<": "<<(1+r)*0.5<<'\n';</pre>
101
102
103
     return 0;
104|}
```

2.22 费用流-Dijkstra

```
1|//原始对偶算法,dijkstra 寻找增广路,单路增广
 21 // 无优化空间
  template<typename T1, typename T2, const int N>class mcf{public:
  #define pdi pair<T2,int>
    priority_queue<pdi, vector<pdi>, greater<pdi>>que;
5
    struct node{int to;T1 cap;T2 cost;int rev;};
6
    int prev[N+10],pree[N+10],numv;
8
    T2 dis[N+10],h[N+10];
9
    vector<node> g[N+10];
10
    void init(int n){
11
      numv=n;
|12|
      rep(i,0,n) g[i].clear();
13
|14|
    inline void add(int from,int to,T1 cap,T2 cost){
15
      g[from].push_back({to,cap,cost,(int)g[to].size()});
16
      g[to].push_back({from,0,-cost,(<u>int</u>)g[from].size()-1});
17
18
    pair<T1,T2>getcost(<u>int</u> s,<u>int</u> t,<u>int</u> n){
19
      numv=n:
20
      T1 flow=0; T2 cost=0;
21
      fill_n(h,numv+1,0);
22
      \underline{\text{while}}(1)
23
        fill_n(dis,numv+1,INF);
24
        dis[s]=0;que.push(make_pair(0,s));
25
        while(!que.empty()){
          auto now=que.top();que.pop();
26
          if(dis[now.second] < now.first) continue;</pre>
27
28
          int x=now.second;
29
          int cnt=0;
30
          for(auto &i:g[x])
            if(i.cap>0&&dis[i.to]>dis[x]+h[x]-h[i.to]+i.cost){
31
32
              dis[i.to]=dis[x]+i.cost+h[x]-h[i.to];
              prev[i.to]=x;
33
34
              pree[i.to]=cnt++;
              que.push(make_pair(dis[i.to],i.to));
35
36
            }else cnt++;
37
38
        if(dis[t] == INF) break;
39
        rep(i,0,numv) h[i]+=dis[i];
40
        T1 d=INF;
41
        for(int now=t;now!=s;now=prev[now])
42
          d=min(d,g[prev[now]][pree[now]].cap);
43
        if(d==INF)break;
44
        flow+=d; cost+=d*h[t];
45
        for(int now=t;now!=s;now=prev[now]){
          node &e=g[prev[now]][pree[now]];
46
47
          e.cap-=d,g[now][e.rev].cap+=d;
48
49
50
      return make_pair(flow,cost);
51
```

```
52 };
53 mcf<<u>int</u>,<u>int</u>,(<u>int</u>)1e4>net;
```

```
费用流-SPFA
 1 //zkw 费用流, 单路增广
21 //可将单路增广改为多路增广 + 当前弧优化,但提升不大
3 <u>template</u><<u>typename</u> T1, <u>typename</u> T2><u>class</u> mcf{<u>public</u>:
     int nume=1,s,t,numv,head[maxn],pre[maxn];
    bool vis[maxn];
    queue<<u>int</u>>q;
    T1 flow[maxn], mflow;
8
    T2 dis[maxn], mcost;
9
    struct node{int to,next;T1 cap;T2 cost;}e[maxm<<1];</pre>
10
    void init(int n=maxn-10){
11
      numv=n;
12
      fill(head,head+n+2,0);nume=1,mflow=mcost=0;
13
|14|
    <u>inline void</u> add(<u>int</u> from, <u>int</u> to, <u>int</u> cap, T2 cost){
15
      e[++nume] = {to,head[from],cap,cost};head[from] = nume;
      e[++nume]={from,head[to],0,-cost};head[to]=nume;
16
17
    bool spfa(){
18
19
      fill(dis,dis+2+numv,INF);
20
      fill(vis, vis+2+numv, false);
21
      dis[s]=0;flow[s]=INF;q.push(s);
22
      while (!q.empty()){
23
        int now=q.front();q.pop();
24
        vis[now] = false;
25
        for (int i=head[now];i;i=e[i].next){
26
           int to=e[i].to;
27
          T2 cost=e[i].cost;
          if (e[i].cap&&dis[now]+cost<dis[to]){</pre>
28
29
            dis[to]=dis[now]+cost;
30
            flow[to] = min(flow[now], e[i].cap);
31
            pre[to]=i;
32
            <u>if</u> (!vis[to]){
33
              vis[to]=true;
              q.push(to);
34
35
          }
36
37
        }
38
39
      return dis[t] < INF;</pre>
40
41
    void dfs(){
42
      int x=t;
      while (x!=s){
43
        int i=pre[x];
44
        e[i].cap-=flow[t];
45
46
        e[i^1].cap+=flow[t];
47
        x=e[i^1].to;
      }
48
49
      mflow+=flow[t];
50
      mcost+=(T2)flow[t]*dis[t];
51
52
    pair<T1,T2> getcost(<u>int</u> ss,<u>int</u> tt){
53
      s=ss,t=tt;
54
      while (spfa())dfs();
55
       return make_pair(mflow,mcost);
56
57 };
58 mcf<int,int>net;
```

2.24 杂项-欧拉路径

```
int vis[maxn];
int cnt;
struct node {
   int to,flag,id,next;
}e[maxm];
int head[maxn],nume,deg[maxn];
inline void _add(int a,int b,int c){
   e[++nume]=(node){b,1,c,head[a]};
```

```
head[a]=nume;
10 }
  inline void add(int a,int b,int c){
11
12
    _add(a,b,c);_add(b,a,-c);
13 }
14 vector < int > ans [maxn];
15 void dfs(int now){
16
    vis[now]=1;
17
    for(int i=head[now];i;i=e[i].next){
18
      if(!e[i].flag) continue;
      e[i].flag=e[i^1].flag=0;
19
20
      dfs(e[i].to);
21
      ans[cnt].push_backe(-e[i].id);
22
23 }
24 void solve(){
25
    rep(i,1,n){
26
       <u>if</u>(!vis[i]&&deg[i]&1) {
27
        cnt++;
28
        dfs(i);
29
30
31
    rep(i,1,n){
32
      <u>if</u>(!vis[i]&&deg[i]){
33
        cnt++:
34
        dfs(i);
35
36
37 }
```

2.25 杂项-三元环计数

```
1 vector<pii>g[maxn];
  int deg[maxn],a[maxn],b[maxn],cnt[maxn],pos[maxn],v[maxn];
   int main() {
     \underline{\text{while}}(\underline{\text{cin}}>>n>>m)
5
      rep(i,1,n){
6
         g[i].clear();
         v[i]=deg[i]=pos[i]=0;
8
9
      rep(i,1,m){
10
         <u>cin</u>>>a[i]>>b[i];
11
         deg[a[i]]++,deg[b[i]]++;
12
13
      rep(i,1,m){
         cnt[i]=0;
14
15
         if(deg[a[i]] < deg[b[i]])g[a[i]].emplace_back(b[i],i);</pre>
16
         else if(deg[a[i]]>deg[b[i]])g[b[i]].emplace_back(a[i],i);
17
18
           <u>if</u>(a[i] < b[i]) g[a[i]].emplace_back(b[i],i);
19
           else g[b[i]].emplace_back(a[i],i);
20
      }
21
22
      rep(i,1,m){
23
         <u>int</u> u=a[i],to=b[i];
24
         for(auto j:g[u]) pos[j.fi]=j.se,v[j.fi]=i+1;
25
         for(auto j:g[to]){
           int t=j.fi;
26
27
           if(v[t]==i+1){
28
             cnt[i]++
29
             cnt[pos[t]]++;
30
             cnt[j.se]++;
31
         }
32
33
34
      ll ans=0;
35
      rep(i,1,m) ans+=111*cnt[i]*(cnt[i]-1)/2;
36
       cout << ans << endl;
37
38|}
```

3 Geometry

```
1| //点类与基础
2 #define db double
3 #define pb push_back
 4 | \underline{\text{const}} | \text{db eps} = 1\text{e-7};
 5 \mid \underline{\text{const}} \mid \text{db pi} = \text{acos}(-1);
 6 const db inf = 1e9;
 7 int sign(db x) { if (fabs(x) < eps) return 0; return x > 0 ? 1 : -1; }
8 int cmp(db k1, db k2) { return sign(k1-k2); }
9|//k1 在 k2、k3 之间:
10 bool inmid(db k1, db k2, db k3) { return sign(k2-k1)*sign(k3-k1) \leq 0; }
11 //区间相交判定,区间 1 在区间 2 前:
12 bool intersect(db 11,db r1,db 12,db r2){
    <u>if</u>(11>r1) swap(11,r1); <u>if</u>(12>r2) swap(12,r2);
13
    <u>return</u> cmp(r1,12)!=-1 && cmp(r2,11)!=-1;
14
15|}
16 struct point{
    db x, y;
point(){}
17
18
19
    point(db k1, db k2){ x = k1, y = k2; }
    //向量加法、点 + 向量 = 点:*/
21
    point operator + (const point &k1) const { return point(x+k1.x, y+k1.y); }
22
    //向量减法、点-点 = 向量:*/
    point operator - (const point &k1) const { return point(x-k1.x, y-k1.y); }
23
    //向量数乘:*/
24
25
    point operator * (db k1) const { return (point){x*k1, y*k1}; }
26
    //向量数除:*/
27
    point operator / (db k1) const { return (point) {x/k1, y/k1}; }
28
    //比较两个点(向量)是否相同:*/
29
    bool operator == (const point &k1) const {
      \underline{\text{return}} \text{ cmp}(x,k1.x) == 0 \&\& \text{ cmp}(y,k1.y) == 0;
30
31
    //逆时针旋转:*/
32
33
    point turn(db k1){
34
      \underline{\text{return}} \text{ (point)} \{x*\cos(k1)-y*\sin(k1), x*\sin(k1)+y*\cos(k1)\};
35
36
    //逆时针旋转 90 度:*/
37
    point turn90(){return (point){-y, x};}
38
    //比较两个点(向量)的大小:
    //x 越小则点越小,若 x 相等,则 y 越小点越小. 可以实现按点的坐标排序*/
39
    bool operator < (const point k1) const{</pre>
40
41
      \underline{int} a = cmp(x, k1.x);
      \underline{if}(a == -1) \underline{return} 1;
42
43
      else if(a == 1) return 0;
      else return cmp(y,k1.y)==-1;
44
45
46
    //向量模长:
47
    db len(){ return sqrt(x*x+y*y); }
    //向量模长的平方:
48
    db len2(){ <u>return</u> x*x+y*y; }
49
50
    //单位向量:
51
    point unit(){ return (*this)/(*this).len(); }
    //向量的极角:
52
53
    db angle() { return atan2(y, x); }
54
    //将点放入第一象限:
55
    //当横坐标为负时,或横坐标为 0 纵坐标为负时,将点按原点做对称角度是 |-/2,/2|
56
    point getdel(){
      if (sign(x) = -1 | (sign(x) = -0 \& sign(y) = -1)) return (*this)*(-1);
57
58
      else return (*this);
59
    //判断点是否在 1 2 象限,或者在 x 的负半轴上角度是(0, ]
60
    bool getp() const {return sign(y)==1 || (sign(y)==0&&sign(x)==-1); }
61
    void scan(){cin>>x>>y;}
    void print(){cout<<x<' '<<y<<'\n'; }</pre>
63
64 };
65 // 判断 k1 在 [k2,k3] 内:
66 bool inmid(point k1, point k2, point k3){
    <u>return</u> inmid(k1.x,k2.x,k3.x) && inmid(k1.y,k2.y,k3.y);
68|}
69 //得到两点中点:
70 point midpo(point k1, point k2){ return (k1+k2)/2; }
```

```
71 //两点距离的平方
72 db dis2(point k1, point k2){
    \underline{\text{return}} \ (k1.x-k2.x)*(k1.x-k2.x) + (k1.y-k2.y)*(k1.y-k2.y);
73
74 }
75 db dis(point k1, point k2){ return sqrt(dis2(k1, k2)); }
76 // 义乘:
77| db cross(point k1, point k2) { return k1.x*k2.y - k1.y*k2.x; }
78 //点乘:
79 db dot(point k1, point k2) { return k1.x*k2.x + k1.y*k2.y; }
80 //向量夹角:
81 db rad(point k1, point k2){
    return acos(dot(k1,k2)/k1.len()/k2.len());
82
83
  //return atan2(cross(k1,k2), dot(k1,k2));
84|}
85| //极角排序,[-,]:
86 bool compareangle (point k1, point k2){
    return k1.getp()<k2.getp() ||</pre>
87
88
          (k1.getp()==k2.getp() && sign(cross(k1,k2))>0);
89|}
90| //k1 k2 k3 逆时针 1 顺时针-1 否则 0:
91 int clockwise(point k1,point k2,point k3){return sign(cross(k2-k1,k3-k1));}
```

儿何类-直线类与线段类

```
1 //直线与线段
 2| //直线类
3 struct line{
 4 //方向为 p[0]->p[1]
    point p[2];
    line()\{\}
    line(db x1,db y1,db x2,db y2){p[0]=point(x1,y1),p[1]=point(x2,y2);}
    line(point k1, point k2) {p[0]=k\bar{1}; p[\bar{1}]=k2;}
    point& operator [] (int k){return p[k];}
10| //点在直线左侧的判定:
  //沿着 p0->p1 的左侧为 1, 右侧为 0
12
    bool include(point k){
      return sign(cross(p[0]-k,p[1]-k))>0;
13
14
15 //方向向量:
16
    point dir(){return p[1]-p[0];}
17 //向外 (左) 平移 eps
18
    line push(){
19
      point delta=(p[1]-p[0]).turn90().unit()*eps;
20
      return {p[0]-delta, p[1]-delta};
21
22|};
23 //线段类:
24 struct segment{
25
    point p[2];
26
    segment(){}
27
    segment(db x1,db y1,db x2,db y2) \{p[0]=point(x1,y1),p[1]=point(x2,y2);\}
    segment(point a, point b) { p[0] = a, p[1] = b; }
28
    point dir(){return p[1]-p[0];}
point& operator [] (int k){ return p[k]; }
29
30
31|};
32 //q 到直线 k1,k2 的投影:
33 point proj(point q, point k1, point k2){
34
    point k=k2-k1;
35
    \underline{\text{return}} \ k1+k*(\text{dot}(q-k1,k)/k.len2());
36|}
37| //q 关于直线 k1, k2 的对称点:
38 point reflect(point q, point k1, point k2){
39
    return proj(q,k1,k2)*2-q;
40|}
41  // 点在线投上的判定:
42 bool checkons(point q,point k1,point k2){
    \underline{\text{return}} \text{ inmid}(q,k1,k2) \&\& \text{ sign}(\text{cross}(k1-q, k2-k1)) == 0;
|43|
44|}
45| //点在直线上的判定:
46 bool checkonl(point q,point k1,point k2){
   return sign(cross(k1-q, k2-k1))==0;
```

```
49| //点在射线 k1->k2 上的判定:
50 bool checkonr(point q, point k1, point k2){
    \underline{\text{return}} \ \text{sign}(\text{cross}(q-k1, k2-k1)) == 0 \&\& \ \text{sign}(\text{dot}(q-k1, k2-k1)) >= 0;
51
52|}
53| //直线平行判定, 可以重合:
54 bool parallel(line k1, line k2) { return sign(cross(k1.dir(), k2.dir()))==0; }
55| //直线同向判定:
56 bool samedir(line k1, line k2){
    return parallel(k1,k2)&&sign(dot(k1.dir(),k2.dir()))==1;
57
58| }
59|//直线的比较,极角排序,范围是 [- ;
60 bool operator < (line k1, line k2){
    if (samedir(k1,k2)) return k2.include(k1[0]);
62
    return compareangle(k1.dir(),k2.dir());
|63| }
64 //直线相交判定:
65| //叉积计算面积,两直线不平行必相交 (除去重合的情况),平行时,三角形面积相等:
66 bool checkll(point k1, point k2, point k3, point k4) {
67
     \underline{\text{return}} \text{ cmp}(\text{cross}(k3-k1,k4-k1),\text{cross}(k3-k2,k4-k2))!=0;
|68|
69| //直线相交判定:
70 bool checkll(line k1, line k2){
    <u>return</u> checkll(k1[0],k1[1],k2[0],k2[1]);
72| \}
73| //直线交点:
74 point getll(point k1, point k2, point k3, point k4) {
    db w1=cross(k1-k3,k4-k3), w2=cross(k4-k3,k2-k3);
76
     <u>return</u> (k1*w2+k2*w1)/(w1+w2);
77|}
78| //直线交点:
79 point getll(line k1, line k2){
     <u>return</u> getll(k1[0],k1[1],k2[0],k2[1]);
80
81|}
82| //直线与线段相交判定:
83| //线段的两端点在直线的两侧
84 bool checkls(point k1, point k2, point k3, point k4){
85
     \underline{\text{return}} \ \text{sign}(\text{cross}(k1-k3, k2-k3)) * \underline{\text{sign}}(\text{cross}(k1-k4, k2-k4)) <= 0;
86|}
87 // 线段相交判定:
88 bool checkss(point k1, point k2, point k3, point k4) {
     return intersect(k1.x,k2.x,k3.x,k4.x)&&intersect(k1.y,k2.y,k3.y,k4.y) &&
89
90
           sign(cross(k3-k1,k4-k1))*sign(cross(k3-k2,k4-k2)) <= 0 &&
91
           sign(cross(k1-k3,k2-k3))*sign(cross(k1-k4,k2-k4))<=0;
92|}
93| // 线段相交判定:
94 bool checkss(segment k1, segment k2){
95
     return checkss(k1[0], k1[1], k2[0], k2[1]);
96|}
97 //线段规范相交判定:
98| //端点相交不算
99 bool strictcheckss(point k1, point k2, point k3, point k4){
100
    <u>return</u> sign(cross(k3-k1,k4-k1))*sign(cross(k3-k2,k4-k2))<0 &&
101|
           sign(cross(k1-k3,k2-k3))*sign(cross(k1-k4,k2-k4))<0;
102| \}
103| // 线段规范相交判定:
104 bool strictcheckss(segment k1, segment k2){
     <u>return</u> strictcheckss(k1[0], k1[1], k2[0], k2[1]);
106|}
107 //点到直线的距离:
108 db displ(point q, point k1, point k2){
    if(k1 == k2) return dis(q, k1);
     return fabs(cross(k2-k1, q-k1)) / (k2-k1).len();
110
111 }
112| //点到直线的距离:
|113| db displ(point q, line l){
     <u>return</u> displ(q, 1[0], 1[1]);
115|}
117 db disps(point q,point k1,point k2){
```

```
118
     point k3 = proj(q,k1,k2);
     if (inmid(k3,k1,k2)) return dis(q, k3);
119
120|
     else return min(dis(q, k1),dis(q, k2));
121|}
122| //点到线段的距离:
123 db disps(point q, segment k1){
124
    <u>return</u> disps(q, k1[0], k1[0]);
125|}
126| //线段到线段间的距离:
127 db disss(point k1, point k2, point k3, point k4) {
128
     if (checkss(k1,k2,k3,k4)) return 0;
129
     else return min(min(disps(k3,k1,k2),disps(k4,k1,k2)),
130
                   min(disps(k1,k3,k4),disps(k2,k3,k4)));
131| }
132 //线段到线段间的距离:
133 db disss(segment k1, segment k2){
134
    <u>return</u> disss(k1[0], k1[1], k2[0], k2[1]);
135|}
```

3.3 几何类-圆类

```
1 // 圆类
2 struct circle{
3
    point o; db r;
    circle(){}
    circle(point _o, db _r)\{ o = _o, r = _r; \}
5
6 //点在圆内判定:
7
    bool include(point k){ return cmp(dis(o, k), r) <= 0; }</pre>
8|};
91//求直线与圆的交点沿着 k2->k3 方向给出,相切给出两个:
10 vector point petcl (circle k1, point k2, point k3) {
    point k=proj(k1.o,k2,k3);
11
12
    db d=k1.r*k1.r-(k-k1.o).len2();
13
    if (sign(d)==-1) return {};
    point del=(k3-k2).unit()*sqrt(max((db)0.0, d));
14
15
    return {k-del, k+del};
16|}
17 // 返回两个圆的公切线数量:
18
  int checkposcc(circle k1,circle k2){
19
    <u>if</u> (cmp(k1.r,k2.r)==-1) swap(k1,k2);
    db d=dis(k1.o,k2.o); <u>int</u> w1=cmp(d,k1.r+k2.r),w2=cmp(d,k1.r-k2.r);
|20|
    if (w1>0) return 4; //相离:
21
22
    <u>else if</u> (w1==0) <u>return</u> 3; //相切:
    <u>else</u> <u>if</u> (w2>0) <u>return</u> 2; //相交:
23
24
    <u>else if</u> (w2==0) <u>return</u> 1; //内切:
25
    else return 0; //内含:
|26| }
27| //求两圆交点沿圆 k1 逆时针给出, 相切给出两个:
28 <u>vector</u><point> getcc(circle k1,circle k2){
29
    int pd=checkposcc(k1,k2)
30
    <u>if</u>(pd==0||pd==4) <u>return</u> {};
    db a=(k2.o-k1.o).len2();
31
32
    db cosA=(k1.r*k1.r+a-k2.r*k2.r)/(2*k1.r*sqrt(max(a,(db)0.0)));
33
    db b=k1.r*cosA:
34
    db c=sqrt(max((db)0.0,k1.r*k1.r-b*b));
351
    point k=(k2.o-k1.o).unit(), m=k1.o+k*b, del=k.turn90()*c;
    return {m-del, m+del};
36
37|}
38| //过圆外一点作圆的切线的切点:
39| //沿圆 k1 逆时针给出
40 vector <point > tangentcp(circle k1, point k2){
    db = (k2-k1.0).len(), b=k1.r*k1.r/a, c=sqrt(max((db)0.0, k1.r*k1.r-b*b));
41
42
    point k=(k2-k1.0).unit(),m=k1.0+k*b,del=k.turn90()*c;
    return {m-del, m+del};
43
44|}
45| //求两圆的外切线:
46 vector < line > tangentoutcc (circle k1, circle k2) {
    int pd=checkposcc(k1,k2); if (pd==0) return {};
|47|
48 //内含,返回一条切线
    <u>if</u> (pd==1){
49
      point p1=getcc(k1,k2)[0]; point p2=p1+((p1-k1.o).turn90()/(p1-k1.o).len());
50
```

```
<u>return</u> {(line){p1,p2}};
52
53
     if (cmp(k1.r,k2.r)==0){
      point del=(k2.o-k1.o).unit().turn90().getdel();
|54|
55
      return {(line){k1.o-del*k1.r,k2.o-del*k2.r},(line){k1.o+del*k1.r,k2.o+del*k2.r}};
    } else {
56
      point p=(k2.0*k1.r-k1.0*k2.r)/(k1.r-k2.r);
57
58
      vector<point>A=tangentcp(k1,p),B=tangentcp(k2,p);
59
      vector<line>ans; for (int i=0;i<A.size();i++) ans.push_back((line){A[i],B[i]});</pre>
60
       return ans;
     }
61
62| \}
63| //求两圆的内切线:
64 vector < line > tangentincc (circle k1, circle k2) {
     int pd=checkposcc(k1,k2); if (pd<=2) return {};</pre>
65
66
     <u>if</u> (pd==3){
67
      point p1=getcc(k1,k2)[0]; point p2=p1+((p1-k1.o).turn90()/(p1-k1.o).len());
      <u>return</u> {(line){p1,p2}};
68
69
70
    point p=(k2.0*k1.r+k1.0*k2.r)/(k1.r+k2.r);
71
     vector<point>A=tangentcp(k1,p),B=tangentcp(k2,p);
72
     vector<line>ans; for (int i=0;i<A.size();i++) ans.push_back((line){A[i],B[i]});</pre>
73
     return ans;
74 }
75 //求两圆所有切线:
76 vector<line> tangentcc(circle k1,circle k2){
77
     int flag=0; if (k1.r<k2.r) swap(k1,k2),flag=1;</pre>
78
     vector<line>A=tangentoutcc(k1,k2),B=tangentincc(k1,k2);
79
     for (line k:B) A.push_back(k);
     if (flag) for (line &k:A) swap(k[0],k[1]);
80
81
     return A;
82|}
83| // 圆 k1 与三角形 k2 k3 k1.o 的有向面积交:
84 db circleinsarea(circle k1, point k2, point k3){
    point k=k1.o; k1.o=k1.o-k; k2=k2-k; k3=k3-k;
     int pd1=k1.include(k2),pd2=k1.include(k3);
87
     vector<point>A=getcl(k1,k2,k3);
88 //有一个点在圆内或圆上:
    \underline{if} (pd1>=0){
89
90| //三角形整个落在圆内, 返回三角形的有向面积
91
      <u>if</u> (pd2>=0) <u>return</u> cross(k2,k3)/2;
92| //三角形的一个点落在圆内, 一个点落在圆外
      else return k1.r*k1.r*rad(A[1],k3)/2+cross(k2,A[1])/2;
93
94| //三角形的一个点落在圆内, 一个点落在圆外
    } else if (pd2>=0){
96
       return k1.r*k1.r*rad(k2,A[0])/2+cross(A[0],k3)/2;
97 //否则, 三角形的两个点都落在圆外:
    }else {
98
99
       <u>int</u> pd=cmp(k1.r,disps(k1.o,k2,k3));
100| //返回一个扇形面积:
       <u>if</u> (pd<=0) <u>return</u> k1.r*k1.r*rad(k2,k3)/2;
101|
102 //返回两个扇形加一个三角形:
      else return cross(A[0],A[1])/2+k1.r*k1.r*(rad(k2,A[0])+rad(A[1],k3))/2;
103
104
105|}
106| //以 k1k2 为直径的圆:
107 circle getexcir2(point k1, point k2){
108
    point c = midpo(k1, k2);
109
     db r = dis(c, k1);
110
    return circle(c, r);
111|}
112| //三角形的外接圆:
113 circle getexcir3(point k1, point k2, point k3){
114
    point c = getll( midpo(k1, k2),
115|
                    (k1-k2).turn90()+midpo(k1, k2),
116
                    midpo(k1, k3)
                    (k1-k3).turn90()+midpo(k1, k3));
117
    db r = dis(c, k1);
118
119
     return circle(c, r);
120 }
121| //最小圆覆盖:
```

```
122 circle mincircover(<u>vector</u><point>A){
       random_shuffle(A.begin(), A.end()); int n = A.size();
123
124
       circle now = circle(A[0], 0);
       for(int i = 0; i < n; i++)if(!now.include(A[i])){
125
126
         now = circle(A[i], 0);
127
         \underline{\text{for}}(\underline{\text{int}} \ j = 0; \ j < i; \ j++)\underline{\text{if}}(!\text{now.include}(A[j]))\{
128
           now = getexcir2(A[i], A[j]);
129
           \underline{\text{for}}(\underline{\text{int}} \ k = 0; \ k < j; \ k++) \ \underline{\text{if}}(!\text{now.include}(A[k]))
130
              now = getexcir3(A[i], A[j], A[k]);
131|
      }
132
133
       return now;
134|}
```

3.4 几何函数-多边形与半平面

```
1 //多边形函数
 2| //三角形面积:
3 db tarea(point a, point b, point c){
     \underline{\text{return}} \text{ fabs}((b.x-a.x)*(c.y-a.y)-(b.y-a.y)*(c.x-a.x))/2;
5|}
6| //多边形面积:
71//多边形用 vector<point> 表示, 逆时针
8 db polyarea(<u>vector</u><point>A){
    db ans = 0;
9
     sort(all(A),compareangle);
10
     \underline{\text{for}}(\underline{\text{int}} \ \text{i=0}; \text{i<A.size}(); \text{i++}) \text{ ans } += \text{cross}(A[i], A[(i+1)\%A.\text{size}()]);
11
     return fabs(ans/2);
|12|
13|}
14| //多边形周长:
15 db polyperimeter(<u>vector</u><point>&A){
16
    db ans = 0;
     \underline{\text{for}}(\underline{\text{int}} \ i = 0; \ i < A.size(); \ i++) \ \text{ans} += dis(A[i], \ A[(i+1)\%A.size()]);
17
18
     return ans;
19|}
20 //多边形重心:
21 point polyfocus(vector<point>&A){
     int n = A.size();
23
     db sumx= 0, sumy = 0, sumarea = 0, area;
     for(int i = 1; i+1 < n; i++){
25
       area = cross(A[i]-A[0], A[i+1]-A[0])/2.0;
26
       sumarea += area;
       sumx += (A[0].x+A[i].x+A[i+1].x)*area;
27
       sumy += (A[0].y+A[i].y+A[i+1].y)*area;
28
29
     return point(sumx/sumarea/3.0, sumy/sumarea/3.0);
30
31|}
32| //点与多边形的位置关系:
33 // 2 内部 1 边界 0 外部
34 <u>int</u> contain(<u>vector</u><point>&A, point q){
     int pd=0; A.push_back(A[0]);
36
     for (int i=1;i<A.size();i++){</pre>
37
       point u=A[i-1], v=A[i];
38
       if (checkons(q,u,v)) return 1;
39
       if (cmp(u.y,v.y)>0) swap(u,v);
40
       \underline{if} (cmp(u.y,q.y)>=0||cmp(v.y,q.y)<0) \underline{continue};
41
       <u>if</u> (sign(cross(u-v,q-v))<0) pd^=1;
42
43
     return pd<<1;</pre>
44
     \underline{int} wn = 0;
     \underline{int} n = A.size();
45
     for(int i = 0; i < n; i++){
46
       if(checkons(q, A[i], A[(i+1)%n])); return -1;//onside
47
       \underline{int} k = sign(cross(A[(i+1)%n]-A[i], q-A[i]));
48
49
       \underline{int} d1 = sign(A[i].y-q.y);
50
       \underline{int} d2 = sign(A[(i+1)\%n].y-q.y);
51
       if(k > 0 \&\& d1 \le 0 \&\& d2 > 0) wn++;
       if (k < 0 && d2 <= 0 && d1 > 0) wn--;
52
53
54
     if(wn != 0) return 1;//inside
55
     return 0;//outside
56|}
```

```
57 //逆时针凸包判定:
58 int checkconvex(vector<point>&A){
     <u>int</u> n=A.size(); A.pb(A[0]); A.pb(A[1]);
59
     for (int i=0;i<n;i++) if(sign(cross(A[i+1]-A[i],A[i+2]-A[i]))==-1) return 0;</pre>
60
61
62 }
63| //求凸包:
64 //flag=0 不严格 flag=1 严格
65 vector <point > convexhull(vector <point > A, int flag=1){
     int n=A.size(); vector<point>ans(n*2);
sort(A.begin(), A.end());
66
67
     \underline{int} now=-1;
68
   //下凸壳
69
70
     for(int i=0;i<n;i++){</pre>
       while(now>0 && sign(cross(ans[now]-ans[now-1], A[i]-ans[now-1])) < flag) now--;</pre>
71
72
       ans[++now]=A[i];
73
74
     int pre=now;
   //上凸壳
75
     for(int i=n-2;i>=0;i--){
76
77
        while(now>pre && sign(cross(ans[now]-ans[now-1], A[i]-ans[now-1])) < flag) now--;</pre>
       ans [++now] = A[i];
781
79
   //因为 A[O] 会被算两次,所以舍弃最后一次的 A[O]
80
81
     ans.resize(now);
82
     return ans;
83|}
84| //切割凸包:
85 //保留直线左边的所有点
86 vector<point> convexcut(vector<point>A,point k1,point k2){
87
     int n=A.size(); A.push_back(A[0]); vector<point>ans;
88
     for(int i=0;i<n;i++){</pre>
89
       int w1=clockwise(k1,k2,A[i]), w2=clockwise(k1,k2,A[i+1]);
90
           (w1>=0) ans.push_back(A[i]);
91
       if (w1*w2<0) ans.push_back(getll(k1,k2,A[i],A[i+1]));
92
93
     return ans;
94 }
95| //凸包最近点对:
96| //先要按照 x 坐标排序
97 bool _cmp(point k1, point k2) {return k1.y<k2.y;}
98 db closestpoint(vector<point>&A, int 1, int r){
99
     if (r-1<=5){
100| //当点数小于等于 5 时,暴力计算:
101
       db ans=1e20;
       \underline{\text{for}} \ (\underline{\text{int}} \ \text{i=l;i<=r;i++}) \ \underline{\text{for}} \ (\underline{\text{int}} \ \text{j=i+1;j<=r;j++}) \ \text{ans=min(ans,dis(A[i],A[j]));}
102
103
       return ans;
104
105|
     int mid=l+r>>1; db ans=min(closestpoint(A,1,mid),closestpoint(A,mid+1,r));
106
     \underline{\text{vector}} < \text{point>B}; \underline{\text{for}} (\underline{\text{int}} \ i=1; i\leq r; i++) \underline{\text{if}} (abs(A[i].x-A[mid].x) < = ans) B.push\_back(A[i]);
     sort(B.begin(),B.end(),_cmp);
107
     for (int i=0;i<B.size();i++) for (int j=i+1;j<B.size()&&B[j].y-B[i].y<ans;j++)</pre>
108
109
       ans=min(ans,dis(B[i],B[j]));
     return ans;
110
111|}
112| //凸包的直径 (最远点对):
113| //旋转卡壳,得到的答案为最远距离的平方
114 db convexdiameter(vector < point > &A) {
115
     \underline{int} n = A.size();
116
      \underline{int} now = 1;
     db res = 0;
117
     for(int i = 0; i < n; i++){
118
119
       while(1){
         db x=cross(A[i]-A[(i+1)%n],A[i]-A[(now+1)%n]);
120|
121
         db y=cross(A[i]-A[(i+1)%n],A[i]-A[now]);
122
         if(x<y) break;</pre>
123
         now=(now+1)%n;
124
125
       res = max(res, dis2(A[now], A[i]));
126
127
     return res;
```

```
128| \}
129| //点集中的最大三角形:
130 db maxtriangle(<u>vector</u><point>&A){
131
     int m = A.size();
     int a = 1, b = 2;
132
133
     db res = 0;
     for(int i = 0; i < m; i++){</pre>
134
       \frac{\text{while}(\text{cross}(A[a]-A[i], A[(b+1)\%m]-A[i])}{\text{cross}(A[a]-A[i], A[b]-A[i])}
135
136
         b = (b + 1) \% m;
137
       res = \max(\text{res}, \text{cross}(A[a]-A[i], A[b]-A[i]) / 2.0);
       \frac{\text{while}(\text{cross}(A[(a+1)\%m]-A[i], A[b]-A[i])}{\text{cross}(A[a]-A[i], A[b]-A[i])}
138
139
         a = (a + 1) \% m;
140
       res = \max(\text{res}, \text{cross}(A[a]-A[i], A[b]-A[i]) / 2.0);
141
142
     return res;
143 }
144| //凸包间的最小距离:
145 db mindisbetconvex(vector<point>&A, vector<point>&B){
146
     int = A.size(), m = B.size();
|147|
     if(n < 3 \&\& m < 3){
148
       if(n == 1)
149
         <u>if</u>(m == 1) <u>return</u> dis(A[0], B[0]);
150
         else return disps(A[0], B[0], B[1]);
       }
151
152
       else{
153
         <u>if</u>(m == 1) <u>return</u> disps(B[0], A[0], A[1]);
154
         else return disss(A[0], A[1], B[0], B[1]);
155
156
     int ai = 0, bi = 0;
157
     for(int i = 0; i < n; i++) if(A[i].y < A[ai].y){ ai = i;</pre>
158
     for(int i = 0; i < m; i++) if(B[i].y > A[bi].y){ bi = i; }
159
     db ans = 1e18;
160
161
     for(int i = 0; i < n; i++){
162
       db ck;
       while (ck = sign(cross(B[(bi+1)\%m]-B[bi], A[(ai+1)\%n]-A[ai])) < 0) bi = (bi+1)\%m;
163
164
       if(ck == 0) ans = min(ans, disss(A[(ai+1)%n], A[ai], B[(bi+1)%m], B[bi]));
165
       else ans = min(ans, disps(B[bi], A[(ai+1)\%n], A[ai]));
       ai = (ai+1)%n;
166
167
168
     return ans;
169 }
170| //最小正方形覆盖:
171 db minsquarecover(<u>vector</u><point>&A, db rad){
     db minx = inf, maxx = -inf, miny = inf, maxy = -inf;
172
     for(int i = 0; i < A.size(); i++){</pre>
173
174
       point p = A[i].turn(rad);
       minx = min(minx, p.x);
175
       miny = min(miny, p.y);
176
177
       maxx = max(maxx, p.x);
178
       maxy = max(maxy, p.y);
179
     return max(maxx-minx, maxy-miny);
180
181|}
182| //三分--最小正方形覆盖:
183 db t_divide(<u>vector</u><point>&A, db 1, db r){
184
     db m,rm,eps=1e-8;
185
     \underline{\text{while}}(r-1) = \exp s
       m=1+(r-1)/3;
186
187
       rm=r-(r-1)/3;
       if(minsquarecover(A,m)>minsquarecover(A,rm)) l=m;
188
189
       else r=rm;
190
     return minsquarecover(A, (m+rm)/2);
191
192|}
193| //求半平面交:
194| //半平面是逆时针方向, 输出按照逆时针
195 <u>vector</u><point> gethalf(<u>vector</u><line> L){
     int n = L.size();
196
197
     sort(L.begin(), L.end());
     int first = 0, last = 0;
198
199 //双端队列指针
```

30

31

c2.r = 0.5*((1/(d01-r1))-(1/(d01+r1)))*r0*r0;

db d02 = r0*r0/(d01+r1)+c2.r;

```
Page 43
200|
    line *q = \underline{new} line[n];
201| //双端队列
202
    point *p = new point[n];
203| //p[i]  为  1[i]  和  1[i+1]  的交点
204 q[last] = L[0];
205| //初始化为一个半平面
206
    for(int i = 0; i < n; i++){
      while(first < last && !L[i].include(p[last-1])) last--;</pre>
207
208
      while(first < last && !L[i].include(p[first])) first++;</pre>
      q[++last] = L[i];
209|
210
      if(samedir(q[last], q[last-1])) last--;
211
      if(first < last) p[last-1] = getll(q[last], q[last-1]);</pre>
212
213
    while(first < last && !q[first].include(p[last-1])) last--;</pre>
214
    vector<point>ans;
    if(last - first <= 1) return ans;</pre>
215
216
    p[last] = getll(q[last], q[first]);
    for(int i = first; i <= last; i++) ans.pb(p[i]);</pre>
217
218
     <u>return</u> ans;
219 }
220 int checkpos(line k1,line k2,line k3){return k3.include(getll(k1,k2));}
221| //求半平面交:
222 //半平面是逆时针方向,输出按照逆时针
223 <u>vector</u><line> gethl(<u>vector</u><line> L){
224
    sort(L.begin(),L.end()); deque<line> q;
225
    <u>for</u> (<u>int</u> i=0;i<(<u>int</u>)L.size();i++){
      226
      while (q.size()>1&&!checkpos(q[q.size()-2],q[q.size()-1],L[i])) q.pop_back();
227
      while (q.size()>1&&!checkpos(q[1],q[0],L[i])) q.pop_front();
228
      q.push_back(L[i]);
229
230
231
    while (q.size()>2&&!checkpos(q[q.size()-2],q[q.size()-1],q[0])) q.pop_back();
232
    while (q.size()>2&&!checkpos(q[1],q[0],q[q.size()-1])) q.pop_front();
    vector<line>ans; for (int i=0;i<q.size();i++) ans.push_back(q[i]);</pre>
233
234
    return ans;
235|}
       几何函数-圆的反演
  3.5
 1 /*
 2|一、反演的概念
 3|设在平面内给定一点 O 和常数 k(k 不等于零),对于平面内任意一点 A,
 4] 确定 A, 使 A 为直线 OA 上一点, 并且有向线段 OA 与 OA 满足 OA \cdot OA = k, 我们称这种变换是以 O 为反演中心,
 5|以 k 为反演幂的反演变换,简称反演. 称 A 为 A 关于 O(r) 的互为反演点.
 6 二、作已知点的反演点的方法
```

```
71 给出反演极 O 和反演幂 k>O, 作点 A 的反演点 A.
8| 令 k=r^2, 作出反演基圆 O(r),
9 1) 若点 A 在 O(r) 外,则过点 A 作圆的切线 (两条),两个切点相连与 OA 连线交点就是点 A.
10|2) 若点 A 在 O(r) 内,则把上述过程逆过来:
11 连结 OA, 过点 A 作直线垂直于 OA, 直线与 O(r) 的交点处的切线的交点就是点 A.
12|3) 若点 A 在 O(r) 上, 反演点 A 就是点 A 自身.
13|4)0 没有反演点
14| 三、圆的反演变换
15| 圆在不同情形下的反演成像:
    当圆不经过反演中心,它的反演图形仍旧是个不过反演中心的圆,并且反演中心为这两个互为反形的圆的位似中心:
16 1.
    当圆与反演圆相交,交点是保持不变的;
17|2.
    当圆在反演圆的外面的时候, 反演成像位于圆的内部; 反之, 当圆位于反演圆的内部, 反演成像位于圆的外部.
18 3.
    当圆经过反演中心,它的反演图形是一条直线.
19|4.
    反之,任意一条不过反演中心的直线,其反演成像是一个经过反演中心的圆.
20
21 5. 相切两圆反向任相切,且切点不变,若切点是反演中心,
    则其反象是两条平行直线;两圆相切,若反演中心在某圆上,则为反形为相切的直线与圆;
22
23 */
24 //c1 关于 c0 的反演圆:
25 circle getinvertcir(circle c1, circle c0){
26
   circle c2;
  db x0 = c0.o.x, y0 = c0.o.y, r0 = c0.r,
27
28
  db x1 = c1.o.x, y1 = c1.o.y, r1 = c1.r;
29
  db \ d01 = dis(c0.o, c1.o);
```

```
//db _d02 = r0*r0/(d01-r1)-c2.r;
    c2.o.x = x0 + d02/d01*(x1-x0);
33
    c2.o.y = y0 + d02/d01*(y1-y0);
34
35
    return c2;
36 }
37| //直线 k1 关于 c0 的反演圆:
38 circle getinvertcir(line k1, circle c0){
|39|
    point a = proj(c0.o, k1[0], k1[1]);
    db oa = dis(c0.o, a);
40
41
    db ob = c0.r*c0.r/oa;
    point v = a-c0.0; v = v/v.len();
42
    point b = c0.o+v*ob;
43
44
    circle res;
45
    res.o = midpo(c0.o, b);
46
    res.r = ob/2;
47
    return res;
48 }
49| //c1 关于 c0 的反演直线:
50 line getinvertline(circle c1, circle c0){
511
    point v = c1.0 - c0.0;
    v = v / v.len();
52
   db d = c0.r*c0.r / (2*c1.r);
53
    point k1 = v * d + c0.o;
54
55
    v = v.turn90();
    point k2 = k1 + v;
56
    return line(k1, k2);
57
58|}
59|//p 关于 c 的反演点:
60 point getinvertpoint(point p, circle c){
61
    point v = (p-c.o).unit();
62|
    db len = c.r*c.r/dis(c.o, p);
63
    return c.o+v*len;
64 }
```

几何函数-圆上整点

```
1 struct point{
     11 x,y;//两圆上整数点对
     point(11 _x=0,11 _y=0){x=_x,y=_y;}
     void print2(){printf("%lld %lld\n",x,y);}
     void print1(){printf("%lld %lld ",x,y);}
5
6
    bool operator == (const point&other)const{
7
       return x==other.x&&y==other.y;
8
9
    bool operator<(const point&other)const{</pre>
10
       if(x==other.x) return y<other.y;</pre>
11
       return x<other.x;</pre>
12
13|};
14 ll dis(point a, point b){
15
     \underline{\text{return}} (a.x-b.x)*(a.x-b.x)+(a.y-b.y)*(a.y-b.y);
16
17 bool check(ll n){ll x=sqrt(n); return x*x==n;}
18 <u>void</u> solve(11 rr,11 r,11 d,<u>vector</u><point>&A){
19
     for (ll i=1;i*i<=rr;++i){</pre>
20
       ll t=rr-i*i;
21
       <u>if</u>(!check(t))<u>continue</u>;
22
       11 j=sqrt(t);
23
       if(i>=j)break;
24
       \underline{if}(\underline{gcd(i*i,t)}==1){
25
         11 x=i*j*d;
26
         11 y=sqrt(r*r-x*x);
         A.pb(point(x,y));
27
28
         A.pb(point(-x,y));
29
         A.pb(point(x,-y));
30
         A.pb(point(-x,-y));
31
32
33|}
34 <u>void</u> deal(<u>vector</u><point>&A,ll r){
35
    A.pb(point(0,r));
36
    A.pb(point(0,-r));
    A.pb(point(-r,0));
```

```
A.pb(point(r,0));
39
    r<<=1
    for (ll d=1;d*d<=r;d++){</pre>
40
      if(r%d!=0) continue;
41
42
      solve(r/d,r/2,d,A);
43
       if(d*d==r) break;
44
      solve(d,r/2,r/d,A);
45
46|}
47
  vector<point>veca, vecb;
48
  struct node{
49
    point A,B;
    node(point _A,point_B){A=_A;B=_B;}
50
    bool operator<(const node&other)const{</pre>
51
52
       if(A==other.A)return B<other.B;</pre>
53
       else return A<other.A;</pre>
54
55|};
56
  vector<node>ans;
57
  int main(){
    int T;ll a,b,c;
58
    scanf("%d",&T);
59
    while (T--){
60
61
      veca.clear();vecb.clear();ans.clear();
      scanf("%11d%11d%11d",&a,&b,&c);
62
63
      deal(veca,a);deal(vecb,b);
64
      for(int i=0;i<veca.size();i++)</pre>
65
        for(int j=0;j<vecb.size();j++)</pre>
66
          if(dis(veca[i],vecb[j])==c*c)
67
            ans.pb(node(veca[i],vecb[j]));
68
      sort(ans.begin(),ans.end());
69
      printf("%d\n",ans.size());
70
      for(int i=0;i<ans.size();++i){</pre>
        ans[i].A.print1();ans[i].B.print2();
71
72
73
74 }
```

3.7 几何函数-多边形面积并

```
多边形的交, 多边形的边一定是要按逆时针方向给出
   * 还要判断是凸包还是凹包,调用相应的函数
3
   * 面积并,只要和面积减去交即可
   */
5
6 #include <bits/stdc++.h>
7
  using namespace std;
  const int maxn = 300;
  const double eps = 1e-8;
10 int dcmp(double x){
11
       <u>if</u>(x > eps) <u>return</u> 1;
12
      <u>return</u> x < -eps ? -1 : 0;
13 }
  struct Point{
|14|
      double x, y;
15
16|};
  <u>double</u> cross(Point a, Point b, Point c) ///叉积
17
  {
18
19
      \underline{\text{return}} \quad (a.x-c.x)*(b.y-c.y)-(b.x-c.x)*(a.y-c.y);
20| \}
21 Point intersection(Point a, Point b, Point c, Point d) {
22
      Point p = a;
      \frac{\text{double}}{\text{double}} \ t = ((a.x-c.x)*(c.y-d.y)-(a.y-c.y)*(c.x-d.x))/((a.x-b.x)*(c.y-d.y)-(a.y-b.y)*(c.x-d.x));
23
24
      p.x += (b.x-a.x)*t;
25
      p.y +=(b.y-a.y)*t;
26
      return p;
27 }
28| //计算多边形面积
29|
  double PolygonArea(Point p[], int n){
30
      if(n < 3) return 0.0;
      <u>double</u> s = p[0].y * (p[n - 1].x - p[1].x);
31
      p[n] = p[0];
32
33
      for(int i = 1; i < n; ++ i)
```

```
s += p[i].y * (p[i - 1].x - p[i + 1].x);
35
       return fabs(s * 0.5);
36|}
37
   double CPIA(Point a[], Point b[], int na, int nb)//ConvexPolygonIntersectArea
38 {
39
       Point p[20], tmp[20];
40
       int tn, sflag, eflag;
41
       a[na] = a[0], b[nb] = b[0];
42
       memcpy(p,b,sizeof(Point)*(nb + 1));
43
       \underline{for}(\underline{int} \ i = 0; \ i < na \&\& \ nb > 2; \ i++){}
           sflag = dcmp(cross(a[i + 1], p[0],a[i]));
44
45
           \underline{\text{for}}(\underline{\text{int}} \ j = \text{tn} = 0; \ j < \text{nb}; \ j++, \ \text{sflag} = \text{eflag})\{
                if(sflag>=0) tmp[tn++] = p[j];
|46|
                eflag = dcmp(cross(a[i + 1], p[j + 1], a[i]));
47
                if((sflag - eflag) == -2)
48
                    tmp[tn++] = intersection(a[i], a[i + 1], p[j], p[j + 1]); //求交点
49
50
51
           memcpy(p, tmp, sizeof(Point) * tn);
           nb = tn, p[nb] = p[0];
52
53
54
       if(nb < 3) return 0.0;
55
       return PolygonArea(p, nb);
56|}
57| <u>double SPIA(Point a[], Point b[], <u>int</u> na, <u>int</u> nb)//SimplePolygonIntersectArea 调用此函数</u>
58| {
59
       <u>int</u> i, j
       Point t1[4], t2[4];
60
       double res = 0, num1, num2;
61
       a[na] = t1[0] = a[0], b[nb] = t2[0] = b[0];
62
       for(i = 2; i < na; i++){</pre>
63
           t1[1] = a[i-1], t1[2] = a[i];
64
           num1 = dcmp(cross(t1[1], t1[2],t1[0]));
65
            if (num1 < 0) swap(t1[1], t1[2]);
66
67
           for(j = 2; j < n\bar{b}; j++)
68
69
                t2[1] = b[j - 1], t2[2] = b[j];
                num2 = dcmp(cross(t2[1], t2[2],t2[0]));
70
71
                if(num2 < 0) swap(t2[1], t2[2]);
72
                res += CPIA(t1, t2, 3, 3) * num1 * num2;
73
74
75
       return res;
76 }
77 Point p1[maxn], p2[maxn];
78 <u>int</u> n1, n2;
79
   int main(){
       freopen("in.txt", "r", stdin);
80
81
       \underline{\text{while}}(\underline{\text{cin}} >> \text{n1} >> \text{n2}) \{
           for(int i = 0; i < n1; i++) scanf("%lf%lf", &p1[i].x, &p1[i].y);
for(int i = 0; i < n2; i++) scanf("%lf%lf", &p2[i].x, &p2[i].y);</pre>
82
83
            double Area = SPIA(p1, p2, n1, n2);
84
85
           cout << Area << endl;</pre>
86
       return 0;
87
88|}
```

3.8 几何函数-辛普森积分

```
11/*自适应辛普森积分*/
2 db f(db x){
      /*积分表达式,或积分微元*/
3
4 }
5 db simpson(db 1,db r){return (r-1)*(f(1)+f(r)+4*f((1+r)/2))/6;}/*辛普森积分公式*/
6 db asr(db 1,db r,db eps,db s){
      double mid=(1+r)/2, ls=simpson(1,mid),rs=simpson(mid,r);
8
      if(fabs(ls+rs-s) \le eps*15){
9
         return ls+rs+(ls+rs-s)/15.0;
10
      return asr(1, mid, eps/2, ls) + ars(mid, r, eps/2, rs);
11
12 }
13 <u>int</u> main(){
      db ans = asr();
14
|15| }
```

几何函数-最小圆覆盖

```
1 #include <algorithm>
  #include <iostream>
3 #include <cstring>
  #include <cstdio>
5 #include <cmath>
7 using namespace std;
9 struct vec
10 {
|11|
      double x, y;
12
      vec (const double& x0 = 0, const double& y0 = 0) : x(x0), y(y0) {}
13
      vec operator + (const vec& t) const {return vec(x+t.x, y+t.y);}
      vec operator - (const vec& t) const {return vec(x-t.x, y-t.y);}
14
15
      vec operator * (const double& t) const {return vec(x*t, y*t);}
16
      vec operator / (const double& t) const {return vec(x/t, y/t);}
17
      const double len2 () const {return x*x + y*y;}
      const double len () const {return sqrt(len2());}
18
19
       vec norm() const {return *this/len();}
20
       vec rotate_90_c () \{return \ vec(y, -x);\}
21 };
23 double dot(const vec& a, const vec& b) {return a.x*b.x + a.y*b.y;}
24 double crs(const vec& a, const vec& b) {return a.x*b.y - a.y*b.x;}
26 vec lin_lin_int(const vec& p0, const vec& v0, const vec& p1, const vec& v1)
27 {
28
      \underline{\text{double}} \ \text{t = crs(p1-p0, v1) / crs(v0, v1)};
29
      return p0 + v0 * t;
30| \}
32 vec circle(const vec& a, const vec& b, const vec& c)
33| {
      return lin_lin_int((a+b)/2, (b-a).rotate_90_c(), (a+c)/2, (c-a).rotate_90_c());
34
35|}
37 int n;
38 vec pot[100005];
40 <u>int</u> main()
41| {
42
      scanf("%d", &n);
43
      for(int i=1; i<=n; i++) scanf("%lf%lf", &pot[i].x, &pot[i].y);</pre>
44
      random_shuffle(pot+1, pot+n+1);
45
      vec o;
46
      double r2 = 0;
47
      for(int i=1; i<=n; i++)</pre>
48
          <u>if</u>((pot[i]-o).len2() > r2)
49
50
51
              o = pot[i], r2 = 0;
52
              <u>for</u>(<u>int</u> j=1; j<i; j++)
53
54
                  \underline{if}((pot[j]-o).len2() > r2)
55
56
                      o = (pot[i]+pot[j])/2, r2 = (pot[j]-o).len2();
57
                      <u>for</u>(<u>int</u> k=1; k<j; k++)
58
59
                          \underline{if}((pot[k]-o).len2() > r2)
60
                             o = circle(pot[i], pot[j], pot[k]), r2 = (pot[k]-o).len2();
61
62
                     }
63
                  }
64
              }
65
66
67
      printf("%.101f\n%.101f \%.101f\n", sqrt(r2), o.x, o.y);
68
69
      return 0;
70| \}
```

3.10 几何函数-最小球覆盖

```
1 \mid \underline{\text{double}} \quad \text{cx=0,cy=0,cz=0};
  double x[maxn],y[maxn],z[maxn];
3 double dis(int now){
     return sqrt((x[now]-cx)*(x[now]-cx)+
5
           (y[now]-cy)*(y[now]-cy)+
6
           (z[now]-cz)*(z[now]-cz));
7
  }
8
   int main() {
9
     cin>>n;
10
     rep(i,1,n) <u>cin</u>>>x[i]>>y[i]>>z[i];
11
     int pos=1;
     double dmax=1e4,ans=1e18;
12
     while(dmax>1e-7){
13
14
       rep(i,1,n) <u>if</u>(dis(i)>dis(pos)) pos=i;
15
       double d=dis(pos);
16
       ans=min(ans,d);
       cx+=(x[pos]-cx)/d*dmax;
17
18
       cy+=(y[pos]-cy)/d*dmax;
19
       cz+=(z[pos]-cz)/d*dmax;
20
       dmax*=0.98;
21
    printf("%.6f \n",ans);
22
23
     return 0;
|24| }
```

4 Game

4.1 SG 函数

```
1 int sg[maxn] [maxn];
   int getmex(bool *vis,int n){
     int t=0;
3
|4|
    while(vis[t]) ++t;
    return t;
6 }
   int getsg(int n,int m){
    if(~sg[n][m]) return sg[n][m];
9
     if(n<=m) return sg[n][m]=1;</pre>
10
    \underline{bool} \ vis[n+1] = \{0\};
|11|
    rep(i,1,m)vis[getsg(n-i,m)]=1;
12
     return sg[n][m]=getmex(vis,n);
13|}
14 <u>int</u> main() {
15
    memset(sg,-1,sizeof sg);
16
     cin>>casn;
17
    while(casn--){
18
       \underline{cin} >> n >> m;
       if(getsg(n,m)) cout<<"first\n";</pre>
19
20
       else cout<<"second\n";</pre>
21
22
    return 0;
```

4.2 威佐夫游戏

```
int main(){
    double k=(1+sqrt(5.0))/2;
    while(scanf("%d%d",&n,&m)!=EOF) {
        if (n>m) swap(n,m);
        int t=m-n;
        if (n==(int)((double)t*k)) printf("0\n");
        else printf("1\n");
}
```

4.3 K 倍取石子博弈

```
#include <iostream>
#include<cstdio>
#include<cstring>
#include<algorithm>
using namespace std;
const int maxn=2000000;
int a[maxn],b[maxn];
```

```
9 <u>int</u> main()
10 {
11
        int t,n,k;
12
        cin>>t;
13
        for(int cas=1;cas<=t;cas++)</pre>
|14|
            <u>cin</u>>>n>>k; //n 个石子每次拿前一次的最多 k 倍
15
            printf("Case %d: ",cas);
16
17
            \underline{if} (n<=k+1)
18
                 printf("lose\n");
19
20
                 continue;
21
22
            a[0]=b[0]=1;
23
            <u>int</u> i=0, j=0;
24
            \underline{\text{while}}(a[i] < n)
25
            {
26
                 i++;
27
                 a[i]=b[i-1]+1;
28
                 while(a[j+1]*k<a[i]) j++;</pre>
29
                 <u>if</u>(a[j]*k<a[i]) b[i]=b[j]+a[i];
30
                 else b[i]=a[i];
            }
31
            \underline{if}(a[i]==n)printf("lose \n");
33
34
            <u>else</u>
35
36
                 int ans=0;
37
                 while(n)
                 {
38
39
                      \underline{if}(n)=a[i]
40
                          n-=a[i];
41
                          ans=a[i];
42
43
44
                      i-
45
46
                 cout << ans << endl;</pre>
47
            }
        }
48
49 }
```

Math

数论-基础函数

```
1 ll p=1e9+7;
2 11 gcd(11 a,11 b) { return b?gcd(b,a%b):a; }
3 11 lcm(11 a,11 b) {\underline{\text{return}} a*gcd(a,b)/b;}
4 | 11 exgcd(11 a,11 b,11 &x,11 &y) {
    \underline{if}(b==0) \underline{return} (x=1,y=0,a);
     <u>if</u>(a==0) <u>return</u> (x=0,y=1,b);
    11 r=exgcd(b,a\%b,y,x);
8
     y=(a/b)*x;
9
     return r;
10|}
11 | 11 | 1cm_mod(11 a,11 b,11 c=p) {
12
    return (a/gcd(a,b)*b)%c;
13
14 | 11 pow_mod(ll a, ll b, ll c=p, ll ans=1) {
    while(b) {
15
16
       if(b&1) ans=(a*ans)%c;
       a=(a*a)%c,b>>=1;
17
18
19
    return ans;
20| \}
21| //long double 1e9 以下表现良好, 不会出现误差
22 11 mul_mod_2(11 a,11 b,11 m){
23 ll c=\bar{a}*b-\bar{(11)}((\underline{long double})a*b/m+0.5)*m;
   return c<0?c+m:c;</pre>
24
25|
26| //不丢失精度的快速乘
27 | 11 mul_mod(11 a,11 b,11 c){return (__int128)a*b%c;}
28 11 pow_mul_mod(l1 a, l1 b, l1 c=p, l1 ans=1){
```

```
while(b){
30
       if(b&1)res=mul_mod(res,a,c);
31
       a=mul_mod(a,a,c),b>>=1;
32
33
     return res;
34|}
35 | 11 inv_gcd(11 a,11 c=p){
36
    a%=c;
|37|
     if(a<0)a+=c;
38
     11 b=c,u=0,v=1;
39
     \underline{\text{while}}(a) {
40
       11 t=b/a;b=t*a;
41
       swap(a,b);
42
       u-=t*v;
43
       swap(u,v);
44
45
     if(u<0)u+=c;
46
     return u;
47 }
```

数论-欧拉函数线性筛

```
1 / / \text{ #x that } x \le n \&\& \gcd(x,n) == 1
 2 int euler_phi(int n) {
     \underline{int} m = (\underline{int}) \operatorname{sqrt}(n+0.5);
     int ans = n;
     for (int i = 2; i <= m; ++ i) if (n % i == 0) {
        ans = ans / i * (i-1);
        while (n\%i == 0) n /= i;
 8
 Q.
     <u>if</u> (n > 1) ans = ans / n * (n-1);
10
     return ans;
11|}
12 int phi[maxn];
13 void phi_table(int n) {
|14|
     for (int i = 2; i <= n; ++ i) phi[i] = 0;</pre>
15
     phi[1] = 1;
     <u>for</u> (<u>int</u> i = 2; i <= n; ++ i) {
16
17
        <u>if</u> (!phi[i]) {
18
          <u>for</u> (<u>int</u> j = i; j <= n; j += i) {
            <u>if</u> (!phi[j]) phi[j] = j;
phi[j] = phi[j] / i * (i-1);
19
20
21
22
23
        phi[i] += phi[i-1];
24
25| \}
```

数论-莫比乌斯函数线性筛

```
1 int mu[maxn],prime[maxn],sum[maxn],nump;
2 bool isp[maxn];
3 void getmu(){
    mu[1]=1,nump=0;
5
    int n=maxn-10;
6
    rep(i,2,n){
       <u>if</u>(!isp[i])    prime[++nump]=i,mu[i]=-1;
      for(int j=1;j<=nump&&prime[j]*i<=n;j++){</pre>
8
9
        isp[i*prime[j]]=1;
10
        <u>if</u>(i%prime[j]==0) mu[i*prime[j]]=0,j=nump+10;
11
        else mu[i*prime[j]]=-mu[i];
|12|
13
|14|
    rep(i,1,n){
15
      sum[i]=sum[i-1]+mu[i];
16
17| \}
```

数论-Miller Rabin 素数判定

```
1| //两种米勒罗宾素数筛实现
21//第一种速度稍快,但实现麻烦,使用前需要先初始化
3|//单次复杂度约为 1-maxpe 中的素数个数 *logn, maxpe 不要小于 50
4 const int maxpe=100;
5 auto randint=bind(uniform_int_distribution<int>(1,1e9),mt19937(rand()));
```

```
6 | 11 mul_mod(11 a,11 b,11 c){return (__int128)a*(__int128)b%(__int128)c;}
  int mul_mod(int a,int b,int c){return (11)a*(11)b%c;}
  template<typename T> int pow_mod(int a,int b,int c){
     int res=1;
10
     while(b){
11
       if(b&1)res=mul_mod(res,a,c);
|12|
      a=mul_mod(a,a,c);
13
    }
14
15
     return res;
16|}
17 <u>template<typename</u> T> <u>class</u> miller_rabin{<u>public</u>:
    T prime[maxpe],cntp;
18
19
     void init(int n=maxpe-5){
20
       cntp=0;
21
      bool vis[n+1];
22
      rep(i,2,n)vis[i]=1;
      rep(j,2,n)<u>if</u>(vis[j]==1)
23
        for(int m=2;j*m<=n;++m)vis[j*m]=0;</pre>
|24|
25
      rep(i,2,n)<u>if</u>(vis[i]==1) prime[cntp++]=i;
26
27
     bool _test(T n,T a,T d) {
28
       if(n==2||n==a) return true;
       if((n&1)==0) return false;
29
      while(!(d&1))d>>=1;
30
31
      T t=pow_mod(a,d,n);
32
       while (d!=n-1&&t!=1&&t!=n-1) {
33
        t=mul_mod(t,t,n);
34
         d <<=1;
35
36
      \underline{\text{return}} (t==n-1||(d&1)==1);
37
38
    bool test(T n) {
       if(n<2||n%2==0) return false;</pre>
39
40
      rep(i,0,cntp-1) <u>if(!_test(n,prime[i],n-1)) return false;</u>
41
       <u>return true;</u>
42
43|};
44 miller_rabin<int> miller;
45| //第二种实现
46| //速度稍慢, 复杂度约为 logn * time, time 不要低于 10
47
  const int test_time=20;
48
  template<typename T> bool miller_rabin(T n) {
     if(n<3)return n==2;
49
50
     T = n-1, b=0;
     <u>while</u> (a\%2==0) a/=2,++b;
51
     for(int i=1,j; i<=test_time;++i){</pre>
52
       T x=randint()\%(n-2)+2, v=pow_mod(x,a,n);
53
54
       \underline{if}(v==1||v==n-1) \underline{continue};
      for(j=0;j<b;++j){</pre>
55
56
        v=mul_mod(v,v,n);
57
        if(v==n-1) break;
58
59
      if(j>=b)return 0;
60
61
     return 1;
62| \}
```

5.5 数论-Pollard Rho 因子分解

```
1|//得到 n 的一个随机因子,包括自身和 1
2|//复杂度 O(n^{1/4})
3|//要保证 randint 的随机范围大于等于测试数字
5 int mul_mod(int a, int b, int c){return (11)a*(11)b%c;}
  auto randint=bind(uniform_int_distribution<11>(1e9,1e18),mt19937(rand()));
6
  template<typename T> T pollard_rho(T n,T c=randint()) {
7
   T i=1,k=2,x=randint()\%(n-1)+1,y=x,d;
   \underline{\text{while}}(1)
10
     i++;
11
     x=(mul_mod(x,x,n)+c)%n;
     d=_gcd(n,y-x);
12
13
     if(d>1&&d<n)return d;</pre>
```

```
if(y==x)return n;
      if(i==k){
15
        k<<=1;
16
17
        y=x;
18
19
    }
20 }
21 //分解因子,map 中即为素因子从小到大,first 为因子,second 为次幂
22 //分解 1000 个 1e18 的数字约为 600ms,10000 个 1e9 的数字约为 200ms
23 map<11, int > factor;
24 template typename T>void get_factor(T n,T c=randint()) {
25
    if(n==1) return;
    if(miller_rabin(n)) {
|26|
27
      factor[n]++;
28
      return;
29
    T p=n;
30
31
    while(p>=n) p=pollard_rho(p,c--);
32
    get_factor(p,c);
33
    get_factor(n/p,c);
34|}
36 //新版实现常数降低
  template<typename T>class Pollard_rho{public:
37
38
    ull s1,s2,s3,s4;
39
    Pollard_rho(){
40
      srand(time(0)+rand());
41
      s1=(ull)rand()*rand()*rand();
      s2=(ull)rand()*rand()*rand();
42
43
      s3=(ull)rand()*rand()*rand();
44
      s4=(ull)rand()*rand()*rand();
45
46
    ull getinteger() {
      ull t=s1^(s1<<11);
47
48
      s1=s2;s2=s3;s3=s4;
49
      <u>return</u> s4=s4^(s4>>19)^t^(t>>8);
50
    T randinteger(T 1,T r){return T(getinteger()%((ull)r-l+1)+1);}
51
52
    T get_factor(T n,T c) {
53
    T i=1,k=2,x=randinteger(1,1e18)\%(n-1)+1,y=x,d;
54
    \underline{\text{while}}(1)
55
      i++;
56
      x=(mul_mod(x,x,n)+c)%n;
57
      d=_gcd(n,y-x);
58
      if(d>1&&d<n)return d;</pre>
59
      if(y==x)return n;
      \underline{if}(i==k){
60
61
        k <<=1;
62
        y=x;
63
64
65|}
66
67 Pollard_rho<ll> rho;
       多项式-拉格朗日插值
```

```
1 class polysum {public:
     ll a[maxn],f[maxn],g[maxn],p[maxn],p1[maxn],p2[maxn],b[maxn],h[maxn][2],C[maxn];
     11 \operatorname{calcn}(\underline{\operatorname{int}} d, 11 *a, 11 n) {//len=d get(an)}
       if (n<=d) return a[n];</pre>
       p1[0]=p2[0]=1;
6
       rep(i,0,d) {
7
         11 t=(n-i+mod) mod;
8
         p1[i+1]=p1[i]*t%mod;
9
10
       rep(i,0,d) {
11
         11 t=(n-d+i+mod)\%mod;
         p2[i+1]=p2[i]*t%mod;
12
13
|14|
       ll ans=0;
15
       rep(i,0,d) {
16
         11 t=g[i]*g[d-i]%mod*p1[i]%mod*p2[d-i]%mod*a[i]%mod;
```

24

37

41

47

```
if ((d-i)&1) ans=(ans-t+mod)%mod;
        else ans=(ans+t)%mod;
18
19
20
      return ans;
    }
21
22
    void init(int maxm) {//init
      f[0]=f[1]=g[0]=g[1]=1;
23
      rep(i,2,maxm+4) f[i]=f[i-1]*i%mod;
25
      g[maxm+4] = pow_mod(f[maxm+4], mod-2);
26
      per(i,1,maxm+3) g[i]=g[i+1]*(i+1)%mod;
27
    ll polysum(ll n,ll *a,ll m){//a[i]会被修改
28
      // 初始化预处理阶乘和逆元 (取模乘法)a[0].. a[m] \sum_{i=0}^{n-1} a[i]
29
      // len=m,psum_n
30
31
      a[m+1] = calcn(m,a,m+1);
      rep(i,1,m+1) a[i]=(a[i-1]+a[i])%mod;
32
33
      return calcn(m+1,a,n-1);
34
35
    ll qpolysum(ll R,ll n,ll *a,ll m) { // a[0].. a[m] \sum_{i=0}^{n-1} a[i] *R^i
36
      if (R==1) return polysum(n,a,m);
      a[m+1]=calcn(m,a,m+1);
38
      ll r=pow_mod(R,mod-2),p3=0,p4=0,c,ans;
39
      h[0][0]=0;
40
      h[0][1]=1;
      rep(i,1,m+1) {
42
        h[i][0]=(h[i-1][0]+a[i-1])*r\mod;
43
        h[i][1]=h[i-1][1]*r\mod;
44
      rep(i,0,m+1) {
45
        ll t=g[i]*g[m+1-i]%mod;
46
        if (i&1) p3=((p3-h[i][0]*t)\(mod+mod)\(mod,p4=((p4-h[i][1]*t)\)\(mod+mod)\(mod;\)
48
        <u>else</u> p3=(p3+h[i][0]*t)\mod, p4=(p4+h[i][1]*t)\mod;
49
      c=pow_mod(p4,mod-2)*(mod-p3)%mod;
50
      rep(i,0,m+1) h[i][0]=(h[i][0]+h[i][1]*c)\mbox{mod};
51
      rep(i,0,m+1) C[i]=h[i][0]
52
      ans=(calcn(m,C,n)*pow_mod(R,n)-c)\mod;
53
      if (ans<0) ans+=mod;</pre>
54
55
      return ans;
56
57|}
```

多项式-快速傅立叶变换

```
1| const double pi=acos(-1.0);
  struct cp{double x,y;};
3 cp operator*(cp a,cp b){return {a.x*b.x-a.y*b.y,a.x*b.y+a.y*b.x};}
4 cp operator + (cp a, cp b) {return {a.x+b.x,a.y+b.y};}
5 cp operator (cp a,cp b) {return {a.x-b.x,a.y-b.y};}
6 #define carr vector <cp>
  <u>const</u> <u>int</u> maxl=6e4+10;//卷积单个数组的最大长度
7
8
  class fourier{public:
    int rev[max1<<2],len,pw;</pre>
10
    vector<cp> wt;
    void init_0(int ml=maxl){
11
12
      for(int mid=1;mid<2*ml;mid<<=1){</pre>
13
        wt.resize(mid*2+1);
14
        cp wn={cos(pi/mid),sin(pi/mid)};
        wt[mid] = (cp) \{1,0\};
15
16
        rep(j,1,mid-1) wt[mid+j]=wt[mid+j-1]*wn;
17
    }
18
19
    void init(int n){
20
      len=1,pw=0;
21
      while(len<=n) len<<=1,++pw;--pw;
22
      rep(i,0,len-1) rev[i]=rev[i>>1]>>1|(i&1)<<pw;
23
24
    void transform(carr &a, int flag){
25
      if(a.size()!=len) a.resize(len,(cp){0,0});
26
      rep(i,0,len-1) <u>if</u>(i<rev[i]) swap(a[i],a[rev[i]]);
27
      for(int mid=1;mid<len;mid<<=1){</pre>
28
        for(int r=mid<<1,j=0;j<len;j+=r){</pre>
29
          for(int k=0;k<mid;++k){</pre>
```

```
cp wn=wt[mid+k];
31
            if(flag==-1) wn.y=-wn.y;
            cp y=wn*a[mid+j+k];
32
33
            a[j+k+mid]=a[j+k]-y,a[j+k]=a[j+k]+y;
34
        }
35
36
37
      <u>if</u>(flag==-1) rep(i,0,len-1) a[i].x/=len;
38
    }//会破坏掉 a,b 数组, 视情况可以去掉引用
39
    carr mul(carr &a,carr &b){
40
      int la=a.size(),lb=b.size();
41
      <u>if</u>((11)1a*1b<=1000){
42
        carr c(la+lb,(cp)\{0,0\});
43
        rep(i,0,la-1) rep(j,0,lb-1) c[i+j]=c[i+j]+a[i]*b[j];
44
        return c;
45
46
      init(la+lb);
      carr c(len,(cp){0,0});
47
48
      transform(a,1); transform(b,1);
49
      rep(i,0,len-1)c[i]=a[i]*b[i];
50
      transform(c,-1);
51
      return c;
52
53 }fft;
```

5.8 多项式-快速傅立叶变换-数组实现

```
1 const double pi=acos(-1.0);
  struct cp{double x,y;};
3 cp operator*(cp a,cp b){return {a.x*b.x-a.y*b.y,a.x*b.y+a.y*b.x};}
4 cp operator + (cp a, cp b) {return {a.x+b.x,a.y+b.y};}
5 cp operator-(cp a,cp b){return {a.x-b.x,a.y-b.y};}
 6 const int max1=2e6+10;
  class fourier{public:
    int rev[max1<<2],len,pw;</pre>
9
    void init(int n){
10
      len=1,pw=0;
       \underline{\text{while}}(\text{len} \leq n) \text{ len} \leq 1, ++pw; --pw;
11
12
      rep(i,0,len-1) rev[i]=rev[i>>1]>>1|(i&1)<<pw;
13
14
    cp c1[max1<<2],c2[max1<<2];</pre>
15
    vector<cp> wt;
16
    void init_0(){
17
      for(int mid=1;mid<2*maxl;mid<<=1){</pre>
18
        wt.resize(mid*2+1);
19
        cp wn{cos(pi/mid),sin(pi/mid)};
20
        wt[mid] = (cp) \{1,0\}
        rep(j,1,mid-1) wt[mid+j]=wt[mid+j-1]*wn;
21
22
    }
23
24
    void transform(cp*a, int flag){
      rep(i,0,len-1) if(i<rev[i]) swap(a[i],a[rev[i]]);
25
26
      for(int mid=1;mid<len;mid<<=1){</pre>
27
        for(int r=mid<<1,j=0;j<len;j+=r){</pre>
28
          for(int k=0;k<mid;++k){</pre>
29
            cp wn=wt[mid+k];
30
            if(flag==-1) wn.y=-wn.y;
31
            cp y=wn*a[mid+j+k];
32
            a[j+k+mid]=a[j+k]-y,a[j+k]=a[j+k]+y;
33
34
        }
35
36
      <u>if</u>(flag==-1) rep(i,0,len-1) a[i].x/=len;
37
38
    void fix(vector<int>&a){
39
      while(!a.empty()&&!a.back())a.pop_back();
40
     }//传入整数,进行卷积
41
    void mul(vector<int> &a, vector<int> &b){
42
       int la=a.size(),lb=b.size();
43
      init(la+lb);
      rep(i,0,la-1) c1[i]=(cp){a[i],0};
44
45
      rep(i,la,len-1) c1[i]=(cp){0,0};
46
      rep(i,0,1b-1) c2[i]=(cp){b[i],0};
```

```
rep(i,lb,len-1) c2[i]=(cp){0,0};
      transform(c1,1); transform(c2,1);
48
49
      rep(i,0,len-1) c1[i]=c1[i]*c2[i];
50
      transform(c1,-1);
      a.resize(len);
51
52
      rep(i,0,len-1) a[i] = \underline{int}(c1[i].x+0.5)?1:0;
53
      fix(a);
54
55
    void sqr(vector<int> &a){
56
      int la=a.size(); init(2*la);
57
      rep(i,0,la-1)c1[i]=(cp){a[i],0};
58
      rep(i,la,len-1)c1[i]=(cp)\{0,0\};
59
      transform(c1,1);
      rep(i,0,len-1) c1[i]=c1[i]*c1[i];
60
61
      transform(c1,-1);
62
      a.resize(len);
63
      rep(i,0,len-1) a[i]=<u>int</u>(c1[i].x+0.5)?1:0;
64
      fix(a);
65
66
    vector<int> pow(vector<int>a,int k){
67
      fix(a);
68
      vector<int>ret;
69
      while(k){
70
        <u>if</u>(k&1){
71
          if(ret.empty()) ret=a;
72
          else mul(ret,a);
73
74
        sqr(a);
75
        k/=2;
76
77
      return ret;
78
79 }
```

多项式-杜教 BM

```
namespace bm{
     const int maxl=1e4+10;
     11 res[maxl],base[maxl],_c[maxl],_md[maxl];
     vector<11> md;
 5
     ll inv(ll a,ll c=mod) {
 6
       a\%=c; \underline{if}(a<0)a+=c;
 7
       ll b=c,u=0,v=1;
 8
       while(a) {
 9
         11 t=b/a;b=t*a;
10
         swap(a,b);u-=t*v;
11
         swap(u,v);
12
13
       if(u<0)u+=c;
|14|
       return u;
15
16
     void mul(ll *a,ll *b,int k) {
       for(int i=0;i<k+k;i++) _c[i]=0;
for(int i=0;i<k;i++) if (a[i])</pre>
17
18
19
         for(int j=0;j<k;j++) _c[i+j]=(_c[i+j]+a[i]*b[j])%mod;</pre>
       for (ll i=k+k-1;i>=k;i--) if (_c[i])
20
21
         for(int j=0;j<md.size();j++)</pre>
22
             c[i-k+md[j]]=(_c[i-k+md[j]]-_c[i]*_md[md[j]])%mod;
23
       for(int i=0;i<k;i++) a[i]=_c[i];</pre>
24
25
     int solve(ll n, vector<11> a, vector<11> b) {
26
     //a 系数 b 初值 b[n+1]=a[0]*b[n]+...
     //求出的是第 n+1 项
27
28
       11 ans=0,pnt=0;
29
       ll k=a.size();
       for(int i=0;i<k;i++) _md[k-1-i]=-a[i];_md[k]=1;</pre>
30
31
       md.clear();
32
       \underline{\text{for}}(\underline{\text{int}} \ i=0; i< k; i++) \ \underline{\text{if}} \ (\underline{\text{md}}[i]!=0) \ \text{md.push\_back}(i);
33
       for(int i=0;i<k;i++) res[i]=base[i]=0;</pre>
34
       res[0]=1;
35
       <u>while</u> ((111<<pnt)<=n) pnt++;
       for (11 p=pnt;p>=0;p--) {
36
         mul(res,res,k);
37
```

```
38
        if ((n>>p)&1) {
39
          for (ll i=k-1;i>=0;i--) res[i+1]=res[i];res[0]=0;
          for(int j=0;j<md.size();j++)</pre>
40
            res[md[j]]=(res[md[j]]-res[k]*_md[md[j]])%mod;
41
42
43
      for(int i=0;i<k;i++) ans=(ans+res[i]*b[i])%mod;</pre>
44
45
      if (ans<0) ans+=mod;</pre>
46
      return ans;
47
    vector<ll> init(vector<ll> s) {
48
49
      <u>vector</u><ll> coe(1,1),base(1,1);
50
       <u>int</u> len=0,m=1,b=1;
51
      for(int n=0;n<s.size();n++) {</pre>
52
        11 d=0;
53
        for(int i=0;i<len+1;i++) d=(d+(l1)coe[i]*s[n-i])%mod;</pre>
54
        if (d==0) ++m;
55
        else if (2*len<=n) {</pre>
          vector<11> tmp=coe;
56
57
          11 c=mod-d*inv(b)%mod;
          while (coe.size() < base.size() + m) coe.push_back(0);</pre>
58
59
          for(int i=0;i<base.size();i++) coe[i+m]=(coe[i+m]+c*base[i])%mod;</pre>
60
          len=n+1-len; base=tmp; b=d; m=1;
61
        } else {
          11 c=mod-d*inv(b)%mod;
62
63
          while (coe.size() < base.size() + m) coe.push_back(0);</pre>
64
          for(int i=0;i<base.size();i++) coe[i+m]=(coe[i+m]+c*base[i])%mod;</pre>
65
66
67
68
      return coe;
69
70
    vector<ll> c,a;
71
    void inita(vector<11> _a){
72
73
       c=init(a);c.erase(c.begin());
74
      for(auto &i:c) i=(mod-i)%mod;
75
76
    int get(ll n) {
      return solve(n,c,vector<11>(a.begin(),a.begin()+c.size()));
77
78
79
    int get(vector<11> a,11 n) {
80
      vector<ll> c=init(a);
81
      c.erase(c.begin());
82
      for(int i=0;i<c.size();i++) c[i]=(mod-c[i])%mod;</pre>
83
       return solve(n,c,vector<11>(a.begin(),a.begin()+c.size()));
84
85|};
```

5.10 多项式-快速数论变换

```
1 const 11 mod=998244353, modg=3, modi=332748118;
   int a[maxn],b[maxn];
 3 <u>int</u> pow_mod(<u>int</u> a, <u>int</u> b){
     int ans=1;
 5
     while(b){
 6
        if(b&1) ans=(ll)ans*a%mod;
 7
       a=(11)a*a\mod,b>>=1;
 8
     return ans;
10|}
11 int add(int a,int b){
12
     a+=b; if (a>mod) return a-mod;
     return a;
13
14|}
15 \mid \underline{int} \quad sub(\underline{int} \quad a, \underline{int} \quad b) 
16
     a-=b; if(a<0) return a+mod;
17
     return a;
18| }
19 #define arr vector<int>
20 <u>const</u> <u>int</u> maxl=2e4+10;//卷积单个数组的最大长度
21 arr operator *(arr&,arr&b){
22 <u>int</u> len=a.size();
```

```
arr c(len));
24
    rep(i,0,len-1) c[i]=a[i]*b[i];
25
    return c;
26 }
27
  class nubmer{public:
28
     int rev[max1<<2],len,pw;</pre>
29
    <u>int</u> wt[2][max1<<2];
30
    void init_0(){
31
       <u>int</u> len=1;//最开始的初始化,整个程序一次就够
      while(len<=maxl) len<<=1;</pre>
32
33
      for(int mid=1;mid<len;mid<<=1){</pre>
34
        11 wn1=pow_mod(modg,(mod-1)/(mid<<1));</pre>
35
        11 wn2=pow_mod(modi,(mod-1)/(mid<<1));</pre>
36
        wt[0][mid]=wt[1][mid]=1;
37
        ll wt1=wn1,wt2=wn2;
38
        rep(j,1,mid){
39
          wt[0][mid+j]=wt1;wt1=wt1*wn1%mod;
40
          wt[1][mid+j]=wt2;wt2=wt2*wn2%mod;
41
42
43
    }
44
    void init(int n){
45
      len=1,pw=0;
46
       \underline{\text{while}}(\text{len} \leq n) \text{ len} \leq 1, ++pw; --pw;
47
      rep(i,0,len-1) rev[i]=rev[i>>1]>>1|(i&1)<<pw;
48
49
    void transform(arr &a, int flag){
50
      11* f=flag==1?wt[0]:wt[1];
51
       if(a.size()!=len) a.resize(len);
      rep(i,0,len-1) <u>if</u>(i<rev[i]) swap(a[i],a[rev[i]]);
52
53
       for(int mid=1;mid<len;mid<<=1){</pre>
54
        for(int r=mid<<1,j=0;j<len;j+=r){</pre>
55
          11 *p=f+mid;
           for(int k=0;k<mid;++k,++p){</pre>
56
57
             <u>int</u> x=a[j+k],y=(*p)*a[j+k+mid]%mod;
            a[j+k+mid]=sub(a[j+k],y),a[j+k]+=y;
58
59
             \underline{if}(a[j+k]>mod) a[j+k]-=mod;
60
61
        }
62
63
       <u>if</u>(flag==-1) {
64
        11 inv=pow_mod(len,mod-2);
65
        rep(i,0,len-1){
66
          a[i]=a[i]*inv%mod;
67
          if(a[i]<0)a[i]+=mod;</pre>
68
      }
69
70
    }
    void fix(arr &a){//分治 ntt 优化
71
72
      while(!a.empty()&&!a.back()) a.pop_back();
73
     }//会破坏掉 a,b 数组, 视情况可以去掉引用
74
     arr mul(arr &a,arr &b){
75
       int la=a.size(),lb=b.size();
76
      <u>if</u>(la*lb<=1000){
77
        arr c(la+lb,0);
78
        rep(i,0,la-1) rep(j,0,lb-1)
79
           c[i+j]=(c[i+j]+(l1)a[i]*b[j])%mod;
80
        return c;
81
82
      int n=la+lb;
83
      init(n);
      transform(a,1);transform(b,1);
84
85
       arr c=a*b;
86
      rep(i,0,len-1)c[i]=(ll)a[i]*b[i]%mod;
87
      transform(c,-1);
88
       return c;
89
91| }ntt;
```

5.11 多项式-快速沃尔什变换

1 #define add(a,b) ((a+=b)>=mod?a-=mod:a)

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```
2 class walsh{public:
3
     void transform_or(ll *a, int len, int flag){
       for(int i=1;i<len;i<<=1)</pre>
|4|
5
         for(int p=i<<1,j=0;j<len;j+=p)</pre>
           for(int k=0;k<i;++k)</pre>
6
7
             add(a[i+j+k],flag==1?a[j+k]:mod-a[j+k]);
8
9
     void transform_and(ll *a,int len,int flag){
       for(int i=1;i<len;i<<=1)</pre>
10
11
         for(int p=i<<1,j=0;j<len;j+=p)</pre>
12
           for(int k=0;k<i;++k)</pre>
             add(a[j+k],flag==1?a[i+j+k]:mod-a[i+j+k]);
13
14
15
     void transform_xor(ll *a, int len, int flag){
       for(int i=1;i<len;i<<=1)</pre>
16
17
         for(int p=i<<1,j=0;j<len;j+=p)</pre>
18
           for(int k=0;k<i;++k){</pre>
19
             <u>int</u> x=a[j+k],y=a[i+j+k];
20
             a[j+k]=(x+y) \text{mod}, a[i+j+k]=(x+mod-y) \text{mod};
21
             if(flag==-1)
22
               a[j+k]=a[j+k]*inv2\mod,a[i+j+k]=111*a[i+j+k]*inv2\mod;
23
     }
24
25|}fwt;
26
   int main(){
     \underline{\operatorname{cin}} >> n; n = (1 << n);
27
    rep(i,0,n-1) \underline{cin} >> a[i];
28
    rep(i,0,n-1) <u>cin</u>>>b[i];
29
30
    fwt.transform_or(a,n,1),fwt.transform_or(b,n,1);
    rep(i,0,n-1) ans[i]=a[i]*b[i]%mod;
31
32
    fwt.transform_or(a,n,-1),fwt.transform_or(b,n,-1);
33
    fwt.transform_or(ans,n,-1);
34
    rep(i,0,n-1) <u>cout</u><<ans[i]<<' ';<u>cout</u><<<u>endl</u>;
    fwt.transform_and(a,n,1),fwt.transform_and(b,n,1);
|35|
    rep(i,0,n-1) ans[i]=a[i]*b[i]%mod;
36
37
    fwt.transform_and(a,n,-1),fwt.transform_and(b,n,-1);
38
    fwt.transform_and(ans,n,-1);
    rep(i,0,n-1) cout << ans[i] << ' '; cout << endl;</pre>
39
40
     fwt.transform_xor(a,n,1),fwt.transform_xor(b,n,1);
41
     rep(i,0,n-1) ans[i]=a[i]*b[i]%mod;
42
     fwt.transform_xor(a,n,-1),fwt.transform_xor(b,n,-1);
43
     fwt.transform_xor(ans,n,-1);
     rep(i,0,n-1) cout << ans[i] << ' '; cout << endl;
44
45|}
```

5.12 线性代数-异或线性基

```
<u>template</u><<u>typename</u> T, <u>const</u> <u>int</u> len> <u>class</u> lbass{<u>public</u>:
     T d[len+1]; int cnt; bool flag;
     void init(){flag=0;memset(d,0,sizeof d);}
     bool insert(T x){
5
       for(int i=len;x&&i>=0;--i)
 6
         \underline{if}((T)1 << i\&x){
           if(!d[i]) {d[i]=x;return true;}
8
           else x^=d[i];
9
         }
10
       flag=1;
11
       <u>return</u> <u>false</u>;
12
     //线性基和 x 异或的最值
13
     T querymax(T x=0){
14
15
       per(i,0,len)x=max(x,x^d[i]);
16
       <u>return</u> x;
17
18
     T querymin(T x=0){
19
       per(i,0,len)x=min(x,x^d[i]);
20
       return x;
21
22
     //求所有异或值去重后的第 k 大,需要预处理
23
     T p[len+1];
24
     void makekth(){
25
       cnt=0;
       memset(p,0,sizeof p);
26
```

```
27
      T ans=0;
28
      per(i,0,len)per(j,0,i-1)
        if((T)1<<j&d[i]) d[i]^=d[j];</pre>
29
30
      rep(i,0,len) <u>if</u>(d[i]) p[cnt++]=d[i];
31
32
    T querykth(T k){
      <u>if</u>(flag)--k;//包含零
33
       if(!k) return 0;
34
35
      T res=0;
36
      if(k>=(T)1<<cnt) return -1;
37
      per(i,0,len) <u>if</u>((T)1<<i&k)res^=p[i];
38
      return res;
    }
39
40|};
41 template typename T, const int len lbass T, len merge (const lbass T, len &a, const lbass T, len &b) {
42
    lbass<T,len>res=a;
43
    rep(i,0,len) res.insert(b.d[i]);
44
    return res;
45| \}
```

5.13 线性代数-线段树维护区间线性基

```
1 class segtree{public:
   #define nd node[now]
   #define ndl node[now<<1]</pre>
   #define ndr node[now<<1|1]</pre>
5
     struct segnode{
       int l,r,flag,val;
6
7
       int d[32];
8
       inline void init(){val=flag=0;memset(d,0,sizeof d);}
       inline void insert(ll x){
10
         <u>for</u>(<u>register</u> <u>int</u> i=30;x&&i>=0;--i)
11
             if(x&(111<<i)){
12
               <u>if</u>(!d[i]) {d[i]=x;<u>return</u>;}
13
               else x^=d[i];
|14|
15
16
       \underline{int} count(){\underline{int} ans=0;per(i,0,30) \underline{if}(d[i])ans++; \underline{return} ans;}
       void update(int x){val^=x;flag^=x;}
17
18
     node[maxn << 2|3];
19
     <u>inline</u> segnode marge(segnode &a, segnode b) const {
20
       segnode ans;ans.init();
21
       per(i,0,30) ans.insert(a.d[i]),ans.insert(b.d[i]);
22
       ans.insert(a.val^b.val);
23
       ans.val=a.val;
24
       ans.l=a.l,ans.r=b.r;
25
       return ans;
26
27
     <u>inline</u> <u>void</u> down(<u>int</u> now){
28
       if(nd.flag){
29
         ndl.update(nd.flag);ndr.update(nd.flag);
30
         nd.flag=0;
31
32
33
     void maketree(int s,int t,int now=1){
34
       nd.l=s,nd.r=t;nd.init();
35
       <u>if</u>(s==t) {<u>cin</u>>>nd.val;<u>return</u> ;}
36
       maketree(s,(s+t)/2,now << 1);
37
       maketree((s+t)/2+1,t,now << 1|1);
38
       nd=marge(ndl,ndr);
39
40
     void update(int s,int t,int x,int now=1){
41
       if(s<=nd.l&&t>=nd.r) {nd.update(x);return;}
42
       down(now);
43
       if(s<=ndl.r) update(s,t,x,now<<1);</pre>
       if(t>ndl.r) update(s,t,x,now<<1|1);</pre>
44
45
       nd=marge(ndl,ndr);
46
47
     segnode query(<u>int</u> s,<u>int</u> t,<u>int</u> now=1){
       if(s<=nd.l&&t>=nd.r) {
48
49
         if(s==nd.1) {
50
           segnode x;x.init();
51
           return marge(x,nd);
```

```
}else return nd;
53
54
       down(now);
55
       segnode ans;ans.init();
56
       if(s<=ndl.r) ans=marge(ans,query(s,t,now<<1));</pre>
       if(t>ndl.r) ans=marge(ans,query(s,t,now<<1|1));</pre>
57
58
       nd=marge(ndl,ndr);
59
       return ans;
     }
601
61| }tree;
62
  int main() {
63
       <u>cin</u>>>n>>m;
64
       register int a,b,c,d;
65
       tree.maketree(1,n);
66
       \underline{\text{while}}(m--) {
67
           cin>>a>>b>>c;
68
           if(a==1)cin>>d;tree.update(b,c,d);
69
           else cout<<(1<<tree.query(b,c).count())<<endl;</pre>
70
71 }
```

5.14 线性代数-矩阵运算

```
class matrix{public://mod
                     int a,b;
                    vector<vector<11> > x;
                    matrix(int a=1,int b=1){
                             \underline{\text{this}}->a=a,\underline{\text{this}}->b=b;x.resize(a);
   6
                             for(auto &i:x){i.resize(b);std::fill(all(i),0);}
   7
   8
                    \underline{\text{void}} \ \ e(\underline{\text{int}} \ n) \{a=b=n; x=\text{matrix}(n,n).x; \underline{\text{for}}(\underline{\text{int}} \ i=0; i < n; i++)x[i][i]=1; \}
   9
                    \underline{\text{void}} \text{ fill(1l } xx=0) \{ \underline{\text{for}} (\underline{\text{int}} \text{ i=0; i<a; i++)} \underline{\text{for}} (\underline{\text{int}} \text{ j=0; j<b; j++)} x[i][j] = xx; \}
10
                    void fill(vector<vector<1l>> &y){x=y;a=y.size();b=y[0].size();}
11
                    matrix operator *(matrix &m){
12
                            matrix ans(a,m.b);
13
                            \frac{\mathbf{for}(\mathbf{int} \ \mathbf{i=0}; \mathbf{i<a}; \mathbf{i++})\mathbf{for}(\mathbf{int} \ \mathbf{j=0}; \mathbf{j<m.b}; \mathbf{j++})\mathbf{for}(\mathbf{int} \ \mathbf{k=0}; \mathbf{k<b}; \mathbf{k++})\mathbf{if}(\mathbf{x[i][k]\&\&m.x[k][j]})}{\mathbf{for}(\mathbf{int} \ \mathbf{i=0}; \mathbf{i<a}; \mathbf{i=0}; \mathbf{i=0
                                     ans.x[i][j] = (mod+ans.x[i][j] + (x[i][k]*m.x[k][j]+mod)%mod)%mod;
14
15
                             return ans;
16
17
                    matrix operator +(matrix &m){
18
                            matrix ans(a,m.b);
                            for(int i=0;i<a;i++)for(int j=0;j<b;j++)ans.x[i][j]=(x[i][j]+m.x[i][j]+mod)%mod;</pre>
19
20
                            return ans;
21
22
                    matrix operator -(matrix &m){
23
                            matrix ans(a,m.b);
                             for(int i=0;i<a;i++)for(int j=0;j<b;j++)ans.x[i][j]=(x[i][j]-m.x[i][j]+mod)%mod;</pre>
24
25
                             return ans;
26
27
                   matrix pow(ll p){
28
                            matrix ans;ans.e(a);matrix t;t.fill(x);
29
                             while(p){if(p&1) ans=t*ans;t=t*t;p>>=1;}return ans;
30
31|};
```

5.15 杂项-高精度整数类

```
1 namespace bignumbers{
  const int maxd=9999,dlen=4;
  class bignumber {public:
    int len,a[maxn];
    bignumber(){len=1;memset(a,0,sizeof(a));}
5
    bignumber(const 11);
 6
    bignumber(const char*);
    bignumber(const bignumber &);
    bignumber & operator = (const bignumber &);
10
    friend istream& operator>>(istream&, bignumber&);
|11|
    bignumber operator + (const bignumber &) const;
12
    bignumber operator - (const bignumber &) const;
    bignumber operator*(const bignumber &)const;
13
    bignumber operator/(const int &)const;
14
15
    bignumber operator (const int &) const;
16
    11 operator%(const 11 &)const;
17
    bool operator>(const bignumber &T)const;
```

```
bool operator > (const int &t)const;
19
    void print();
20 };
21 bignumber::bignumber(const 11 b) {
22
    int c,d=b;len=0;
23
    memset(a,0,sizeof(a));
24
    while(d>maxd){
      c=d-(d/(maxd+1))*(maxd+1);
26
      d=d/(maxd+1);
27
      a[len++]=c;
28
29
    a[len++]=d;
30|}
31 bignumber::bignumber(<u>const</u> <u>char</u> *s) {
32
    int t,k,index,L;
33
    memset(a,0,sizeof(a));
34
    L=strlen(s);
35
    len=L/dlen;
36
    if(L%dlen)len++;
37
    index=0;
    for(int i=L-1;i>=0;i-=dlen) {
38
39
      t=0,k=i-dlen+1;
40
      if(k<0)k=0;
      rep(j,k,i)t=t*10+s[j]-'0';
41
42
      a[index++]=t;
43
44|}
45| bignumber::bignumber(const bignumber &T):len(T.len) {
|46|
    memset(a,0,<u>sizeof</u>(a));
47
    rep(i,0,len-1)a[i]=T.a[i];
48|}
49 bignumber &bignumber::operator=(const bignumber &n) {
50
    memset(a,0,sizeof(a));
51
    rep(i,0,n.len-1)a[i]=n.a[i];
52
    len=n.len;
53
    return *this;
54|}
55
  char ch[maxn*dlen];
56
  <u>istream</u>& <u>operator</u>>>(<u>istream</u> &in, bignumber &b) {
57
    in>>ch:
    int L=strlen(ch);
58
59
    int count=0,sum=0;
60
    for(int i=L-1;i>=0;) {
       int t=1;sum=0;
61
62
      for(int j=0;j<4&&i>=0;j++,i--,t*=10)
        sum+=(ch[i] - '0')*t;
63
64
      b.a[count++]=sum;
65
66
    b.len=count++;
    return in;
67
68
69 bignumber bignumber::operator+(const bignumber &T)const {
70
    bignumber t(*this);
     int big=T.len>len?T.len:len;
71
    rep(i,0,big-1){
72
      t.a[i]+=T.a[i]
73
74
      <u>if</u>(t.a[i]>maxd)t.a[i+1]++,t.a[i]-=maxd+1;
75
76
    <u>if</u>(t.a[big]!=0) t.len=big+1;
77
    else t.len=big;
78
    return t;
79
80 bignumber bignumber::operator-(const bignumber &T)const {
81
    bool flag;
82
    bignumber t1,t2;
83
    \underline{if}(*\underline{this}>T) {
      t1=*<u>this</u>;t2=T;flag=0;
84
85
     } else {
86
       t1=T;t2=*this;flag=1;
87
88
    int j,big=t1.len;
89
    rep(i,0,big-1){
      <u>if</u>(t1.a[i]<t2.a[i]) {
90
```

```
j=i+1;
         while(t1.a[j]==0) j++;
92
         t1.a[j--]--;
93
         while (j>i) t1.a[j--]+=maxd;
94
         t1.a[i]+=maxd+1-t2.a[i];
95
96
       } else t1.a[i]-=t2.a[i];
97
98
     t1.len=big;
99
     while (t1.a[t1.len-1] == 0 \& \& t1.len > 1)
100
       t1.len--,big--;
101
     <u>if</u>(flag) t1.a[big-1]=0-t1.a[big-1];
102
     return t1;
103 }
104 bignumber bignumber::operator*(const bignumber &T)const {
105
     bignumber ret;
106
     int up,temp,temp1;
107
     rep(i,0,len-1){
       up=0;
108
109|
       rep(j,0,T.len-1){
         temp=a[i]*T.a[j]+ret.a[i+j]+up;
110
         if(temp>maxd) {
111
           temp1=temp-temp/(maxd+1)*(maxd+1);
112
113
           up=temp/(maxd+1),ret.a[i+j]=temp1;
114
         } else up=0,ret.a[i+j]=temp;
115
       if(up!=0)ret.a[i+T.len]=up;
116
117
     ret.len=len+T.len;
118
119
     while(ret.a[ret.len-1] == 0&&ret.len>1)ret.len--;
120
     return ret;
121|}
122 bignumber bignumber::operator/(const int &b)const {
123
     bignumber ret;
124
     int down=0;
125
     per(i,0,len-1) {
       ret.a[i]=(a[i]+down*(maxd+1))/b;
126
127
       down=a[i]+down*(maxd+1) - ret.a[i]*b;
128
129
     ret.len=len:
     while(ret.a[ret.len-1] == 0&&ret.len>1) ret.len--;
130|
131
     return ret;
132|}
133 ll bignumber::<u>operator</u>%(<u>const</u> ll &b)<u>const</u> {
134
     int d=0;
135
     per(i,0,len-1)
       d=((d*(maxd+1))\%b+a[i])\%b;
136
137
     return d;
138|}
139 bignumber bignumber::operator (const int &n)const {
     if(n==0)return 1;
140
141
     if(n==1)return *this;
142
     bignumber t, ret(1);
143|
     int m=n,i;
     \underline{\text{while}}(m>1) {
144
       t=*this;
|145|
|146|
       for(i=1;(i<<1)<=m;i<<=1)t=t*t;</pre>
147
       m-=i,ret=ret*t;
148
       if(m==1)ret=ret*(*this);
149
150
     return ret;
151|}
152 bool bignumber::operator>(const bignumber &T)const {
153
     if(len>T.len)return true;
     else if(len==T.len) {
154
       int ln=len-1;
155
       while (a[ln] == T.a[ln] \&\& ln >= 0) ln --;
156
157
       158
       else return false;
159
     }else return false;
160|}
161 bool bignumber::operator>(const int &t)const {
162
     bignumber b(t);
     return *this>b;
163
```

```
164|}
165 void bignumber::print() {
     printf("%d",a[len-1]);
166|
     per(i,0,len-2)printf("%04d",a[i]);
167
     printf("\n");
168
169 }
170 \\ \lambda \text{using bignumbers::bignumber;}
```

5.16 杂项-分数类

```
1|//num 是分子,den 是分母,分母始终保持为正
  template<typename T>class farction{public:
    T num, den;
    farction(T num=0,T den=1) {
      if (den<0) {
6
        num=-num;
7
        den=-den;
8
      T g=_gcd(abs(num),den);
9
      this->num=num/g;
10
11
      this->den=den/g;
12
13
    farction operator +(const farction &o) const {
14
      return farction(num*o.den+den*o.num,den*o.den);
15
16
    farction operator -(const farction &o) const {
17
      return farction(num*o.den-den*o.num,den*o.den);
18
19
    farction operator *(const farction &o) const {
20
      return farction(num*o.num,den*o.den);
21
|22|
    farction operator /(const farction &o) const {
23
      return farction(num*o.den,den*o.num);
24
25
    bool operator <(const farction &o) const {</pre>
26
      return num*o.den<den*o.num;</pre>
27
28
    bool operator >(const farction &o) const {
29
      return num*o.den*o.num;
30
31
    bool operator ==(const farction &o) const {
32
      return num*o.den==den*o.num;
33
34|};
```

杂项-N 进制快速幂优化

```
1|//底数不变, 求值次数 1e5 以上的情况, 可以加速 10-100 倍
2| //1e9: maxp 1e3,maxv 1e9,100 倍
3 //1e18: maxp 32000, maxv 1e18,14 倍左右
4 const 11 maxp=32000, maxv=1e18;
  const ll maxw=log(maxv)/log(maxp)+1;
6 class basepow{public:
7
    11 pw[maxw][maxp];
8
    void init(ll base){
9
      base%=mod;
10
      rep(i,0,maxw-1)pw[i][0]=1;
11
      rep(i,1,maxp-1) pw[0][i]=(pw[0][i-1]*base)%mod;
12
      rep(i,1,maxw-1){
       pw[i][1]=pw[i-1][maxp-1]*pw[i-1][1]%mod;
13
       rep(j,2,maxp-1) pw[i][j]=pw[i][j-1]*pw[i][1]%mod;
14
15
16
17
    inline 11 getpow(11 b,11 res=1,int cnt=0){
      while(b){
18
19
       res=res*pw[cnt++][b%maxp]%mod;
20
       b/=maxp;
21
      return res;
22
23
    }
24 }p;
```

6 String

6.1 KMP 算法

```
1| //输入的数组从 O 开始, next 函数为 p, 从 1 开始
 2 template<typename T>class prefix{public:
    int p[maxn],lens;
    T *s;
5
    void init(T *_s,int _lens){
6
      s=_s,lens=_lens;
      rep(i,0,lens-1) p[i]=0;
      p[0] = -1;
9
      int now=0,pos=-1;
10
      while (now<lens)</pre>
11
        \underline{if}(pos==-1||s[now]==s[pos]) p[++now]=++pos;
12
        else pos=p[pos];
13
14
    vector<int> find(T *t,int lent){
15
      int now,pos=0;
16
      vector<int> ans;
17
      while(now<lent)</pre>
        <u>if</u>(pos==-1||t[now]==s[pos]) pos++,now++;
18
19
        else pos=p[pos];
20
        if(pos==lens) pos=p[pos],ans.push_back(now-lens);
21
22
      return ans;
23
24|};
25 prefix<char> kmp;
```

6.2 manacher 算法

```
1 class manacher
     char ma[maxn<<1]; int lenp[maxn<<1];</pre>
     void getp(char *s,int len){
       \underline{int} p=0;
       ma[p++]='$',ma[p++]='#';
forn(i,len) ma[p++]=s[i],ma[p++]='#';
 5
 6
 7
       <u>int</u> r=0,mid=0;
       forn(i,(len+1)<<1){
 8
 9
         lenp[i]=r>i?min(lenp[(mid<<1)-i],r-i):1;</pre>
         while(ma[i+lenp[i]]==ma[i-lenp[i]]) lenp[i]++;
10
11
         if(i+lenp[i]>r)r=i+lenp[i],mid=i;
12
13
14 }
```

6.3 后缀数组-倍增

```
1 namespace suffix{
     int tr[maxn],rank[maxn],sa[maxn],h[maxn],has[maxn],n;
     int cmp(int x,int y,int k){
       \underline{if}(x+k>n||y+k>n)\underline{return} 0;
5
       return rank[x] == rank[y] && rank[x+k] == rank[y+k];
6
     void getsa(char *s,int _n,int m=233){
8
       int i,cnt;n=_n;
9
       for(i=1;i<=n;i++)has[s[i]]=0;</pre>
       for(i=0;i<=m;i++)has[i]=0;</pre>
10
       for(i=1;i<=n;i++)has[s[i]]++;</pre>
11
12
       <u>for</u>(i=1,cnt=0;i<=m;i++)<u>if</u>(has[i])tr[i]=++cnt;
       for(i=1;i<=m;i++)has[i]+=has[i-1];</pre>
13
       for(i=1;i<=n;i++)rank[i]=tr[s[i]],sa[has[s[i]]--]=i;</pre>
14
15
       for(int k=1;cnt!=n;k<<=1){</pre>
16
         for(i=1;i<=n;i++)has[i]=0;
17
         for(i=1;i<=n;i++)has[rank[i]]++;</pre>
18
         for(i=1;i<=n;i++)has[i]+=has[i-1];
         \underline{for}(i=n;i>=1;i--)\underline{if}(sa[i]>k)tr[sa[i]-k]=has[rank[sa[i]-k]]--;
19
20
         for(i=1;i<=k;i++)tr[n-i+1]=has[rank[n-i+1]]--;</pre>
21
         for(i=1;i<=n;i++)sa[tr[i]]=i;</pre>
22
         for(i=1,cnt=0;i<=n;i++)tr[sa[i]]=cmp(sa[i],sa[i-1],k) ? cnt:++cnt;</pre>
23
         for(i=1;i<=n;i++)rank[i]=tr[i];</pre>
24
25
       fill_n(h,n+2,0);
       for(int i=1;i<=n;i++){</pre>
```

6.4 最小表示算法

```
1 int minrep(int *s, int n){
      <u>int</u> i=0,j=1,k=0,t;
      \underline{\text{while}}(i < n \& \& j < n \& \& k < n) \{
 4
        t=s[i+k>=n?i+k-n:i+k]-s[j+k>=n?j+k-n:j+k];
 5
        \underline{if}(!t) ++k;
 6
        else {
 7
           if(t>0) i+=k+1;
 8
           <u>else</u> j+=k+1;
 9
           <u>if</u>(i==j) ++j;
10
           k=0;
11
12
      }
13
      return min(i,j);
14| }
```

6.5 字符串哈希

```
1 const 11 decm=13331, modh=(11)1e17+13;
2 inline 11 mulh(11 a,11 b){
    ll tmp=(a*b-(ll)((long double)a/modh*b+1.0e-8)*modh);
4
    return tmp<0?tmp+modh:tmp;</pre>
5
6 inline 11 addh(11 a,11 b){11 x=a+b; if(x<0)return a+b+modh; else if(x>=modh)x-modh; else return x;}
  11 pw[maxn];
8 void init(int n=maxn-5){pw[0]=1;rep(i,1,n) pw[i]=mulh(pw[i-1],decm);
9 class strhash{public:
10
    11 hs[maxn],len;
    void calhs(string &s,int n){len=n;rep(i,0,len-1)hs[i+1]=addh(mulh(hs[i],decm),s[i]);}
11
12
    inline 11 geths(int a,int b){return addh(hs[b],-mulh(hs[a-1],pw[b-a+1]));}
13 }table;
```

6.6 自动机-tire 树

```
1 const int csize=26,minc='a';
 2 class trie {public:
3 #define nd node[now]
    struct tnode{
      int cnt,son[csize];
 6
      tnode(){
7
        cnt=0;
8
        memset(son,0,sizeof son);
9
10
    }node[maxn];
11
    int sz=0;
12
    void insert(char *s,int len){
13
      int now=0;
14
      rep(i,0,len-1){
15
        int ch=s[i]-minc;
        if(!nd.son[ch]) nd.son[ch]=++sz;
16
17
        now=nd.son[ch];
18
19
      node[now].cnt++;
20
21
    int find(char *s,int len){
22
      int ch,now=0;
23
      rep(i,0,len-1){
24
        ch=s[i]-minc;
25
        if(!nd.son[ch]) return -1;
26
        now=nd.son[ch];
27
28
      if(!nd.cnt) return -1;
29
      return 1;
30
31 }tree;
```

6.7 自动机-tire 图

```
1 #include < iostream >
2 #include <cstdio>
3 #include < cstring >
4 #include <vector>
5 #include <queue>
6 using namespace std;
8 const int N = 1000000 + 10, INF = 0x3f3f3f3f;
10 struct node {
    node *next[26];
11
    node *suff; //指向后缀节点
    bool flag;
13
14 } trie[N], *root;
15 <u>int</u> tot;
16 char ori[N];
17 node* node_init() {
18
    trie[tot].flag = false;
    trie[tot].suff = NULL;
19
   memset(trie[tot].next, 0, sizeof trie[tot].next);
20
21
    return trie + tot++;
22 }
23 void trie_insert(char *s) {
24
    node *p = root;
    for(int i = 0; s[i]; i++) {
  int j = s[i] - 'a';
25
26
      if(p->next[j] == NULL)
27
        p->next[j] = node_init();
28
29
      p = p-next[j];
30
31
    p->flag = <u>true</u>;
32|}
33 void trie_graph() {
    //trie[0] 为虚拟节点,root 为 trie[1],trie[0] 的所有边均指向 root,方便以后操作
34
35
    for(int i = 0; i < 26; i++)
      trie[0].next[i] = root;
36
37
    root->suff = trie + 0;
    trie[0].suff = NULL;
38
39
    queue<node*> que;
40
    que.push(root);
41
    while(! que.empty()) {
42
      node *p = que.front();
43
      que.pop();
      for(int i = 0; i < 26; i++)
44
45
        if(p-\text{next}[i]) {
          //查看父亲节点的后缀节点是否存在编号为 i 的边,若没有,就一直往上找,直到根节点
46
47
          node *ptr = p->suff;
          while(ptr && ! ptr->next[i])
48
            ptr = ptr->suff;
49
          //要么找到了,要么循环到了虚拟节点后,后缀节点被置为 root
50
51
          p->next[i]->suff = ptr->next[i];
          if(ptr->next[i]->flag)
52
53
            p->next[i]->flag = true;
54
          que.push(p->next[i]);
55
56
    }
57|}
58 bool trie_query(char *s) {
59
    node *p = root;
60
    \underline{for}(\underline{int} \ i = 0; \ s[i]; \ i++) \ \{
      <u>int</u> j = s[i] - 'a';
61
      while(true) {
62
        if(p->next[j] != NULL) {
63
          p = p-\sum_{j=1}^{n} (j);
64
65
          if(p->flag == true)
            return true;
66
67
          break;
68
        } else
          p = p \rightarrow suff;
69
70
    }
71
```

```
72
    <u>return</u> <u>false</u>;
73 }
74 int main() {
75
    int n;
    scanf("%d", &n);
76
    tot = 2;//tot+0 是虚拟节点, tot+1 是根节点, 故从 2 开始
77
78
    root = trie + 1;
    memset(trie[root].next, 0, sizeof trie[root].next);
79
80
    for(int i = 1; i <= n; i++) {</pre>
      scanf("%s", ori);
81
82
      trie_insert(ori);
83
    trie_graph();//建 trie 图
84
    scanf("%s", ori);
851
    puts(trie_query(ori) ? "YES" : "NO");
86
87
    return 0;
88|}
```

6.8 自动机-AC 自动机

```
1| const int csize=128,minc=0;
  class autom{public:
3 #define nd node[now]
    struct acnode{int id,son[csize],fail;}node[maxn];
    int sz=0;
    queue<<u>int</u>> que;
    void clear(int n=maxn-5){
8
      memset(node,0,sizeof(acnode)*n);
9
      sz=0;
    }
10
11
    void insert(int id, char *s, int len=0){
12
      if(!len)len=strlen(s);
13
      int now=0;
|14|
      rep(i,0,len-1){
15
        int ch=s[i]-minc;
16
        if(!nd.son[ch]) nd.son[ch]=++sz;
17
        now=nd.son[ch];
18
19
      nd.id=id;
20
21
    void init(){
22
      int now=0;
23
      rep(i,0,csize-1) <u>if</u>(nd.son[i])
24
        que.push(nd.son[i]);
25
      while(!que.empty()){
26
        now=que.front();que.pop();
27
        rep(i,0,csize-1){
28
          if(nd.son[i]) {
29
            node[nd.son[i]].fail=node[nd.fail].son[i];
30
            que.push(nd.son[i]);
31
          }else nd.son[i]=node[nd.fail].son[i];
32
        }
33
34
35
    vector<int> query(char *t,int len=0){
36
      if(!len)len=strlen(t);
37
      int now=0;
38
      vector<int> ans;
39
      rep(i,0,len-1) {
40
        now=nd.son[t[i]-minc];
41
        for(int j=now;j;j=node[j].fail)
42
          if(node[j].id)ans.push_back(node[j].id);
43
44
      return ans;
    }
45
46 }acam;
```

6.9 自动机-后缀自动机

```
1 const int maxn=1e5+10,maxsz=26+1,ch0='a';
2 class saffixam {public:
3  int fa[maxn<<1],son[maxn<<1][maxsz];
4  int last,cnt,len[maxn<<1];
5  void insert(int ch) {</pre>
```

```
int pre=last,now=++cnt;
|7|
      last=now,len[now]=len[pre]+1;
8
      for(; pre&&!son[pre][ch]; pre=fa[pre])son[pre][ch]=now;
9
      if(!pre)fa[now]=1;
10
      <u>else</u> {
11
        int q=son[pre][ch];
12
        if(len[pre]+1==len[q])fa[now]=q;
13
        else {
14
          int nq=++cnt;
15
          memcpy(son[nq],son[q],sizeof(son[0]));
16
          fa[nq]=fa[q];
17
          len[nq]=len[pre]+1;
18
          fa[q]=fa[now]=nq;
19
          for(; son[pre][ch]==q; pre=fa[pre])son[pre][ch]=nq;
20
21
      }
22
    }
23
    void init() {
24
      rep(i,0,cnt)
25
        memset(son[i],0,sizeof son[0]);
26
        fa[i]=len[i]=0;
27
28
      last=cnt=1;
29
    }
30
    void insert(string &s ) {for(auto i:s)insert(i-ch0);}
31
    bool find(string &s) {
32
      int now=1;
33
      for(auto i:s) if(!(now=son[now][i-ch0]))return 0;
34
35
36| } sam;
```

6.10 自动机-回文自动机

```
struct PAT {
     struct node {
       int len,num,fail,son[26];
     } t[maxn];
     int last,n,tot,s[maxn];
6
     void init() {
       memset(t,0,<u>sizeof</u>(t));
8
       tot=last=1;
9
       n=0;
10
       t[0].len=0;
11
       t[1].len=-1;
       t[0].fail=t[1].fail=1;
12
13
       s[0]=-1;
|14|
15
     int add(int c) {
16
       int p=last;
17
       s[++n]=c;
18
       \underline{\text{while}}(s[n]!=s[n-1-t[p].len])
19
         p=t[p].fail;
20
       <u>if</u>(!t[p].son[c]) {
21
         int v=++tot,k=t[p].fail;
22
         \underline{\text{while}}(s[n] != s[n-t[k].len-1])
           k=t[k].fail;
23
24
         t[v].fail=t[k].son[c];
25
         t[v].len=t[p].len+2;
26
         t[v].num=t[t[v].fail].num+1;
27
         t[p].son[c]=v;
28
29
       last=t[p].son[c];
30
       return t[last].num;
31
32| } T;
```

7 Others

7.1 常用头文件

```
#define forn(i,x,g,e) for(int i=g[x];i;i=e[i].next)
      #define IO ios::sync_with_stdio(false);cin.tie(0);cout.tie(0)
      #define ull unsigned long long
  8 #define fi first
  9 #define se second
10 #define mp make_pair
11 #define pii pair<ll,ll>
12 #define all(x) x.begin(),x.end()
13 #define show(x) cout<<#x<<"="<<x<<endl
14 <u>#define</u> showa(a,b) <u>cout</u><<#a<<'['<<b<<"]="<<a[b]<<endl
15 <u>#define</u> show2(x,y) <u>cout</u><<#x<<"="<<x<" "<<#y<<"="\leftarrow \cdot 
16 <u>#define</u> show3(x,y,z) <u>cout</u><<#x<<"="<<x<" "<<#y<<"="<<y<" "<<#z<<"="<<z<<endl
17 #define show4(w,x,y,z) cout<<#w<<"="<<w<" "<<#x<<"="<<x<" "<<#y<<"="<<y<" "<<#z<<"="<<z<<endl
18 #define show5(v,w,x,y,z) cout<<#v<<" "<<#w<<" "<<#x<<" "<<#x<<" "<<#y<<" "<<#z<<" "<<#z<<" "
               ="<<z<<end]
19| #define showa2(x,a,b) cout<<#x<<": ";rep(i,a,b) cout<<x[i]<<' ';cout<<endl
20 using namespace std;//head
                    <u>int</u> maxn=1e6+10,maxm=2e6+10;
22 const 11 INF=0x3f3f3f3f, mod=1e9+7;
23 int casn,n,m,k;
24 int main() {
28 }
29 //struct node{
30 // int x,y;
31 // friend bool operator<(const node&a,const node&b){
32 //
                   return a.x==b.x?a.y<b.y:a.x<b.x;
33 // }
34 //}a [maxn];
35 // auto _start=chrono::high_resolution_clock::now();
36 // auto
                            _end=chrono::high_resolution_clock::now();
37 // cerr<="elapsed time: "<<chrono::duration<double,milli>(_end-_start).count()<<" ms\n";
38 //int size=(64)<<20; //64MB
39 // __asm__("movq %0, %%rsp\n"::"r"((char*)malloc(size)+size));
                 快速读写
```

```
1 namespace io{
     const int L=(1<<21)+1;
     char ibuf [L],*iS,*iT,obuf [L],*oS=obuf,*oT=obuf+L-1,c,st[55];int f,tp;
     #define gc() (iS==iT?(iT=(iS=ibuf)+fread(ibuf,1,L,stdin),(iS==iT?EOF:*iS++)):*iS++)
     inline void flush(){fwrite(obuf,1,oS-obuf,stdout);oS=obuf;}
6
     <u>inline</u> <u>void</u> putc(<u>char</u> x){*oS++=x;<u>if</u>(oS==oT)flush();}
     template<class I>
8
     <u>inline</u> <u>void</u> gi(I&x){
9
       for(f=1,c=gc();c<'0'||c>'9';c=gc())if(c=='-')f=-1;
10
       for(x=0; c \le '9' \&\&c \ge '0'; c=gc())x=x*10+(c\&15); x*=f;
11
12
     template<class I>
13
     inline void print(I x){
14
       <u>if</u>(!x)putc('0');<u>if</u>(x<0)putc('-'),x=-x;
       while(\bar{x})st[++tp]=x%10+'\bar{0}',x/=10;
15
       while(tp)putc(st[tp--]);
16
17
     inline void gs(char*s,int&l){
18
       for(c=gc();c!='_'&&(c<'a'||c>'z');c=gc());
19
       for(1=0;c=='_'||c<='z'&&c>='a';c=gc())s[1++]=c;
20
21
22|};
```

7.3 快速读写加强版

```
namespacefastio{//支持读取整数,字符串,输出整数boolisdigit(char c){return c>=48&&c<=57;}</td>const int maxsz=1e7;class fast_iostream{public:char ch=get_char();bool endf=1,flag;char get_char(){static char buffer[maxsz],*a=buffer,*b=buffer;static char buffer[maxsz],*a=buffer,1,maxsz, stdin),b==a)?EOF:*a++;preturn b==a&&(b=(a=buffer)+fread(buffer,1,maxsz, stdin),b==a)?EOF:*a++;
```

```
template<typename type>bool get_int(type& tmp){
12
      flag=tmp=0;
13
      while(!isdigit(ch)&&ch!=EOF){flag=ch=='-';ch=get_char();};
14
       <u>if</u>(ch==EOF)<u>return</u> endf=0;
15
      do{tmp=ch-48+tmp*10;}while(isdigit(ch=get_char()));
16
       if(flag)tmp=-tmp;
17
      return 1;
18
    int get_str(char* str){
19
20
      char* tmp=str;
21
      while (ch== ' r' | ch== ' n' | ch== ' ') ch=get_char();
22
       if(ch==E0F)return(endf=0),*tmp=0;
       \underline{do}\{*(tmp++)=ch; ch=get\_char();\}\underline{while}(ch!='\r'\&\&ch!='\n'\&\&ch!=' '\&\&ch!=EOF);
23
24
       *(tmp++)=0;
25
      return(int)(tmp-str-1);
26
27
    fast_iostream& operator>>(char* tmp){get_str(tmp); return *this;}
28
    template<typename type>fast_iostream& operator>>(type& tmp){get_int(tmp);return *this;}
29
     operator bool() const {return endf;}
30|};
31 template < typename type > void put(type tmp) {
    if (tmp==0) {putchar(48); return;}
32
33
    static int top,stk[21];
34
    if (tmp<0){tmp=-tmp;putchar('-');}</pre>
|35|
    while(tmp)stk[++top]=tmp%10,tmp/=10;
36
    while(top)putchar(stk[top--]+48);
37
38 } fastio::fast_iostream io;
```

7.4 LIS

```
1 int num[maxn],dp[maxn],len;
   int lwb(int now){
 3
     int l=0,r=len,ans=0;
     while(l<r){</pre>
 5
       <u>int</u> mid=(l+r)>>1;
 6
       if(dp[mid]>=now) l=mid+1;
 7
       else r=mid;
 8
     return 1;
10|}
11 int lwb2(int now){
12
     int l=0,r=len;
13
     while(l<r){</pre>
       <u>int</u> mid=(l+r)>>1;
14
15
       if(dp[mid]>=now) r=mid;
       else l=mid+1;
16
17
     return 1;
18
19 }
20 <u>int</u> main(){
21
     cin>>n;
|22|
     <u>for</u>(<u>int</u> i=1;i<=n;i++) <u>cin</u>>>num[i];
23
     for(int i=1;i<=n;i++){</pre>
       if(dp[len]>=num[i]) dp[++len]=num[i];
24
25
       else dp[lwb(num[i])]=num[i];
26
27
     cout<<len<<endl;</pre>
28
|29|
     memset(dp,0,sizeof dp);
30
     for(int i=1;i<=n;i++){</pre>
31
       <u>if</u>(dp[len]<num[i]) dp[++len]=num[i];
32
       else dp[lwb2(num[i])]=num[i];
33
34
     cout << len << endl;</pre>
35|}
```

7.5 莫队算法

```
1 ll a[maxn],ans[maxn],sz,sum;
2 class block{public:
3     ll cnt[maxn];
4     struct node{ll l,r,id;
5     bool operator <(const node &b)const {
     if(1/sz!=b.1/sz) return l<b.1;
}</pre>
```

```
if((1/sz)&1) return r<b.r;</pre>
8
         return r>b.r;
9
10
     };
     void update(int pos,ll flag=1){
11
12
       sum+=flag*a[pos]*(cnt[a[pos]]*211+flag);
       cnt[a[pos]]+=flag;
13
|14|
15
  }ask;
16 <u>vector</u> <block::node > tab;
17
   int main() {
     \underline{cin} >> n >> m; sz = sqrt(n);
18
19
     rep(i,1,n) <u>cin</u>>>a[i];
20
     rep(i,1,m) {
21
       <u>int</u> a,b;<u>cin</u>>>a>>b;
22
       tab.push_back({a,b,i});
23
24
     sort(all(tab));
25
     int l=tab[0].1,r=tab[0].1-1;
26
     for(auto now:tab){
27
       while(1>now.1) ask.update(--1,1);
28
       while(l<now.l) ask.update(l++,-1);</pre>
29
       while(r<now.r) ask.update(++r,1);</pre>
       while(r>now.r) ask.update(r--,-1);
30
31
       ans[now.id] = sum;
32
33|}
```

7.6 大随机数

```
1 template < typename T > class random_xor {public:
    T s1,s2,s3,s4;
3 //
       random_xor(){
4 //
        srand(time(0)+rand());
5 //
        s1=(T)rand()*rand()*rand();
6 //
        s2=(T)rand()*rand()*rand();
7 //
8 //
      T getint(){
9 //
        T s3=s1,s4=s2;
10 //
        s1=s4;s3^=s3<<23;
        s2=s3^s4^(s3>>17)^(s4>>26);
11 //
12 //
        return s2+s4;
13 // }
14
    random_xor(){
      srand(time(0)+rand());
151
16
      s1=(T)rand()*rand()*rand();
      s2=(T)rand()*rand()*rand();
17
18
      s3=(T)rand()*rand()*rand();
19
      s4=(T)rand()*rand()*rand();
20
21
    T getint() {
|22|
      T t=s1^(s1<<11);
23
      s1=s2;s2=s3;s3=s4;
      return s4=s4^(s4>>19)^t^(t>>8);
24
25
26
    T get(T 1,T r){\underline{\text{return}} getint()%(r-1+1)+1;}
27
    T 	ext{ operator} ()(T 1,T r){return getint()%(r-1+1)+1;}
28
      operator ()(){return getint();}
29 };
30 random_xor<<u>unsigned</u> __int128> a;
31 srand(time(0));
32 //mt19937::result_type seed=time(0);
33 auto randint=bind(uniform int distribution<int>(1,r),mt19937(time(0)));
34 auto randfloat=bind(uniform_real_distribution<double>(1,r),mt19937(time(0)));
```

7.7 大质数

```
1 1e9+7,1e9+9,1e9+21,1e9+33,

2 1e9+87,1e9+93,1e9+97,1e9+103,

3 1e9+123,1e9+181,1e9+207,1e9+223,

4 1e9+241,1e9+271,1e9+289,1e9+297,

5 1e9+321,1e9+349,1e9+363,1e9+403,
```

7.8 编译器位操作

```
1 #pragma GCC target ("popcnt")
2| __builtin_parity(n);//1 的个数的奇偶性
3 __builtin_popcount(n);//1 个数
4 __builtin_ctz(n);//末尾 0 个数,n!=0
5 __builtin_clz(n); //前导 0 的个数, n!=0
6|__builtin_ffs(n);//最后 1 的位置,从 1 开始
7|__builtin_parity(n)//x 中 1 的奇偶性
8| uint32_t __builtin_bswap32(uint32_t x)//按字节翻转
```

最大子矩阵算法

```
1 /*
2 把二维看成一维
3| 先枚举行的起点和终点
 4| 再把起点行和终点行间每一列的数值压缩到每一个点上
5| 然后求一个最长连续子段和
6 复杂度 O(n^3)
7
  */
8
  #include <string.h>
  #include <stdio.h>
10
  #include <iostream>
11 #include <algorithm>
12 using namespace std;
13 const int maxn=1e3+10;
14 const int maxm=1e6+10;
15 const int INF=0x3f3f3f3f;
16 #define 11 long long
17 | \underline{int} | casn,n,m,k;
18 <u>int</u> smax(<u>int</u> a[], <u>int</u> len){
19
    int mx=0, sub=0;
20
    for(int i=1;i<=len;i++){</pre>
21
      sub=max(a[i],sub+a[i]);
22
      mx=max(sub,mx);
23
24
    return mx;
25|}
  int arr[maxn];
26
27 int dp[maxn] [maxn];
28 int main(){
29 #define test
30 #ifdef test
    freopen("in.txt", "r", stdin); freopen("out.txt", "w", stdout);
31
32 #endif
    while (\sim scanf("%d", \&n))
34
35
      for(int i=1;i<=n;i++){</pre>
        for(int j=1;j<=n;j++){
  scanf("%d",&dp[i][j]);</pre>
36
37
38
39
40
      int ans=-INF;
      for(int i=1;i<=n;i++){</pre>
41
42
        memset(arr,0,sizeof arr);
43
        <u>for</u>(<u>int</u> j=i;j<=n;j++){
          for(int k=1;k<=n;k++){</pre>
44
            arr[k]+=dp[j][k];
45
46
47
          ans=max(ans,smax(arr,n));
48
      }
49|
      printf("%d\n",ans);
50
51
52 #ifdef test
    fclose(stdin);fclose(stdout);system("out.txt");
53
54 #endif
55
    return 0;
56|}
```

7.10约瑟夫环算法

```
1 #include <bits/stdc++.h>
2 using namespace std;
3 typedef long long 11;
```

```
6|11 ysf(11 n,11 k,11 num) //n 为人数 k 为每次点名的数第 num 个出列的人的序号 0- n-1
7
8
      if(k==1)
9
          return num-1;
10
      ll ans=0;
11
      if(num<k)</pre>
12
13
          ans=(k-1)%(n-num+1);
14
          11 p=n-num+1;
15
          for(11 i=2; i<=num; i++)</pre>
16
17
              ans=(ans+k)%(++p);
          }
18
19
      }
                                          //num 很大时用乘法加速
      else
21
22
          ans=-1;
23
          for(ll i=n-num+1; i<=n; i++)</pre>
24
25
              ll j=min(n,(i-1)+((i-1)-ans+(k-2))/(k-1));
26
              ans=(ans+k*(j-i+1))%j,i=j;
27
          }
28
30
      return ans;
31 }
32
  int main()
  {
33
34
      11 T,T2;
      cin>>T;
35
36
      T2=T;
      while (T--)
37
38
39
          ll n, k, num;
40
          <u>cin</u>>>n>>num>>k;
          cout<<<"Case #"<<T2-T<<": "<<ysf(n,k,num)+1<<endl;</pre>
41
42
45 }
        java 高精度
  7.11
```

```
package acm;
   import java.math.BigInteger;
 3 import java.util.Scanner;
 4 public class Main{
     public static final int maxn=1000,maxm=200000;
     public static BigInteger gcd(BigInteger a, BigInteger b) {
       <u>if</u>(b.compareTo(BigInteger.ZERO)==0)<u>return</u> a;
 8
       else return gcd(b, a.remainder(b));
 9
10
     public static void main(String[] argc){
11
       int[]a=new int[maxn];
       for(int i=2;i<maxn;i++) a[i]=i;</pre>
12
13
       \underline{\text{for}}(\underline{\text{int}} \ i=2; i < \max; i++) \ \underline{\text{if}}(a[i]!=0)
14
         \underline{\text{for}}(\underline{\text{int}} \ j=i*2; j<\text{maxn}; j=j+i) \ a[j]=0;
15
       Scanner cin=new Scanner(System.in);
16
       int casn=cin.nextInt();
17
       int[] prime=new int[10000]
18
       for(int ii=1;ii<=casn;++ii) {</pre>
19
         BigInteger x=cin.nextBigInteger();
         BigInteger res=BigInteger.ONE;
20
21
         BigInteger now=BigInteger.ONE;
22
         <u>int</u> i=1;
23
         while(true) {
24
           i++;
25
           if(a[i]==0) continue;
           \underline{\textbf{if}}(res.multiply(BigInteger.valueOf(a[i])).compareTo(x)<=0) {
26
27
             res=res.multiply(BigInteger.valueOf(a[i]));
28
             now=now.multiply(BigInteger.valueOf(a[i]+1));
```