Static Analysis of Bottom Tier Sensor Networks

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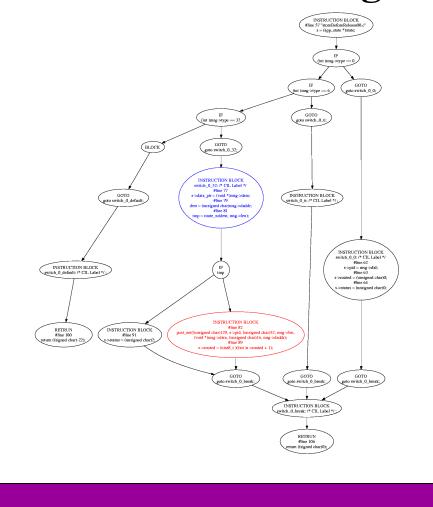
Motivation: Bringing Accessible Static Analysis To Sensor Network Development

Static Analysis is Cool

- Has the potential to find bugs before they arise
- Streamlines the development process
- Lowers learning curve for new developers

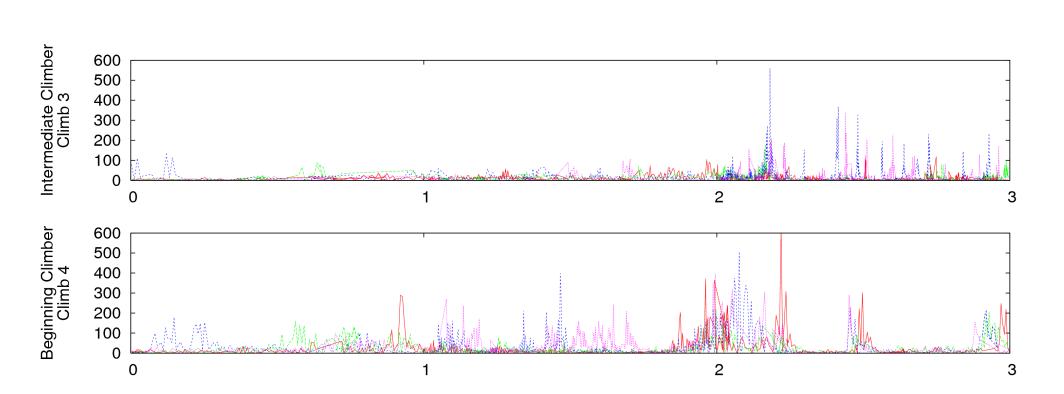
Bottom Tier Sensing Devices are Cool

- Developer can look at the entire system
- All-in-one hardware is great for hobbyist and students
- Show you amazing things about the world we live in

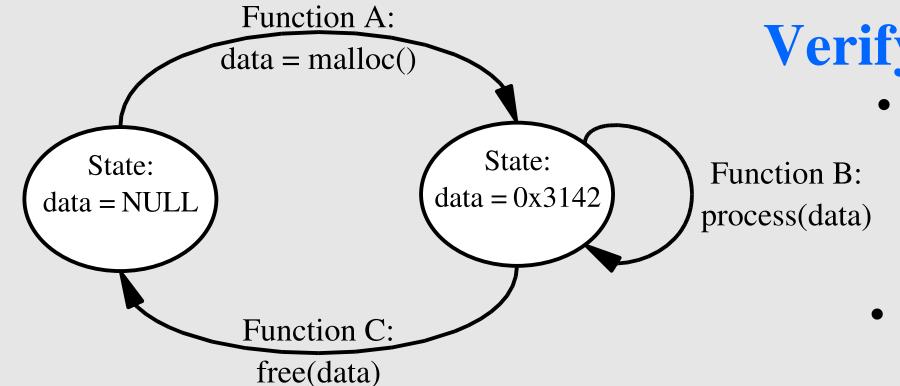


Combination of These?

- Difficulties developing sensor network applications
 - Limited insight into a running system
 - Common failure mode is simply "not working"
 - Failures are not well contained
- Desire gaining insight into bottom tier sensing devices by applying static analysis



Problem: Mapping Traditional Analysis Onto Very Reactive Systems

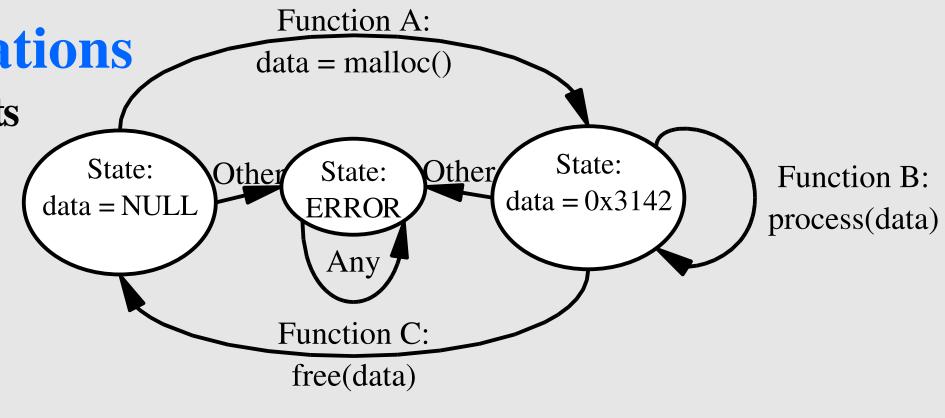


Verifying Sensor Network Applications

Event handlers react to external events

Sensor detects an event of significance

- Message is received over the radio
- Timer injects event into the system
- Naïve model assumes any event can occur at any time



Proposed Solution: Drive Static Analysis with Ordering Specifications

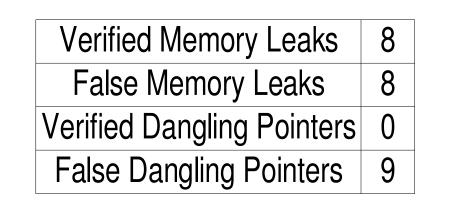
Lighthouse

• Simple exclusive resource ownership model

- All heap memory is owned by exactly one module
- Owner is responsible for releasing the memory or transferring it to another module
- Modules may not free or transfer ownership of heap data that they do not own

• Uses standard dataflow techniques to verify each function

- May / must pointer analysis used to resolve aliases
- Function summaries used to handle function calls



Warnings generated from the analysis of 213 unique versions of SOS modules totalling 28042 SLOC

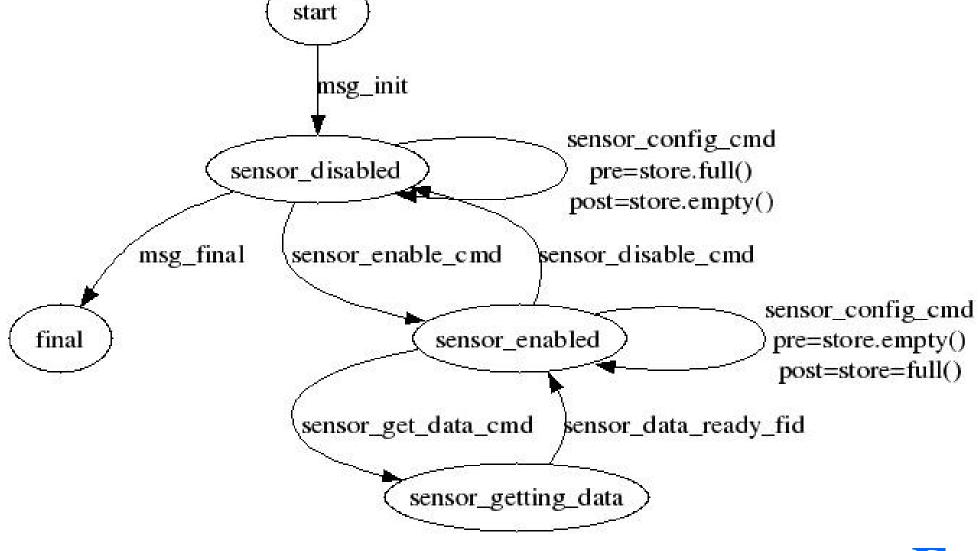
Verified Memory Leaks 2
False Memory Leaks 22
Verified Dangling Pointers 0
False Dangling Pointers 11

Warnings generated from the analysis of 40 source files making up the SOS kernel totalling 9223 SLOC

mod_op = (sos_module_op_t*) ker_msg_take_data(msg); if(mod_op == NULL) return -ENOMEM; if (mod_op->op == MODULE_OP_INSMOD) { existing_module = ker_get_module(mod_op->mod_id); if(existing_module != NULL) { uint8_t ver = sos_read_header_byte(...); if (ver < mod_op->version) { ker_unload_module(...); else { return SOS_OK; ret = fetcher_request(KER_DFT_LOADER_PID, mod_op->mod_id, mod_op->version, entohs(mod_op->size), msg->saddr); $s->pend = mod_op;$ ker_led(LED_RED_TOGGLE); return SOS_OK; return SOS_OK;

Error: Expression mod_op is not stored after
 instruction #line 125
mod_op = (sos_module_op_t*)
 ker_msg_take_data((unsigned char)18, msg);

Checkpoint



- Static analysis driven by finite state machine specification
- Specification provided by the developer describes correct execution of a program
 - Set of arbitrary states associated with specific source code
 - Set of events that cause transitions between states
 - Persistent store pre- and post- conditions for stores when a specific event causes a transition between two states
- Critical for the reactive sensing environments
- Checkpoint used to refine the Lighthouse analysis

Future Work

- Specification inference
 - Infer memory manipulation within individual functions
 - Infer interfunction per- / post- state dependencies

Finite state machine specification

- What is the best way to describe the FSM?
- Transitioning from functional specification to slice-oriented specification