# Spelling Correction: Edit Distance

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I am writing this email on behaf of ... The user typed 'behaf'.

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- behave
- ....

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#### Isolated word error correction

Pick the one that is closest to 'behaf'

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- Pick the one that is closest to 'behaf'
- How to define 'closest'?

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- Pick the one that is closest to 'behaf'
- How to define 'closest'?
- Need a distance metric

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#### Isolated word error correction

- Pick the one that is closest to 'behaf'
- How to define 'closest'?
- Need a distance metric
- The simplest metric: edit distance

## Edit Distance

The minimum edit distance between two strings

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- Is the minimum number of editing operations

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- The minimum edit distance between two strings
- Is the minimum number of editing operations
  - Insertion
  - Deletion
  - Substitution

## Example

Edit distance from 'intention' to 'execution'

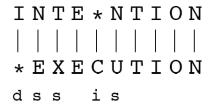
What is the minimum number of edit operations needed to go from 'intention' to 'execution'?

### Example

Edit distance from 'intention' to 'execution'

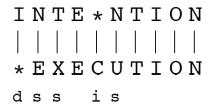


- 1. Delete I
- 2. Substitute N with E
- 3. Substitute T with X
- 4. Insert C
- 5. Substitute N with U

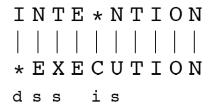


One delete (d), one insert (i), 3 substitution (s) operations.

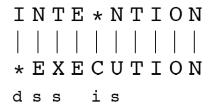
What is the cost of each operation?



- If each operation has a cost of 1 (Levenshtein)
  - Distance between these is 5



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Weighted versions of edit distance also exist, where different operations, especially different substitutions, are weighted differently — look up details

Searching for a path (sequence of edits) from the *start string* to the *final string*:

Initial state: the word we are transforming

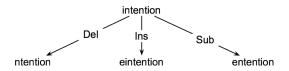
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- Operators: insert, delete, substitute
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A naive method — try all possible edit operations starting with the first character of initial state

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- Lot of distinct paths end up at the same state
- Don't have to keep track of all of them Ideally, we want to keep track of only those paths that take us 'towards' the goal state.

# How to navigate?

- The space of all edit sequences is huge
- Lot of distinct paths end up at the same state
- Don't have to keep track of all of them
- Keep track of the shortest path to each state

# Defining Minimum Edit Distance Matrix

## For two strings

- X of length n
- Y of length m

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## For two strings

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- Y of length m

# We define D(i,j)

- the edit distance between X[1..i] and Y[1..j]
- i.e., the first i characters of X and the first j characters of Y

Thus, the edit distance between X and Y is D(n,m)

## Dynamic Programming

• A tabular computation of D(n,m)

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# Computing Minimum Edit Distance

#### Dynamic Programming

- A tabular computation of D(n,m)
- Solving problems by combining solutions to subproblems
- Bottom-up
  - Compute D(i,j) for small i,j
  - Compute larger D(i,j) based on previously computed smaller values
  - Compute D(i,j) for all i and j till you get to D(n,m)

## Dynamic Programming Algorithm

#### Initialization

$$D(i,0) = i$$
  
 $D(0,j) = j$ 

#### Recurrence Relation:

For each 
$$i = 1...M$$
  
For each  $j = 1...N$   

$$D(i,j) = \min \begin{cases} D(i-1,j) + 1 & \text{The } X[i] \text{ character is deleted} \\ D(i,j-1) + 1 & \text{The } Y[j] \text{ character is inserted} \\ D(i-1,j-1) + 2; & \text{if } X(i) \neq Y(j) \\ \text{substitute } X[i] \text{ by } Y[j] \\ \text{if they are not same} \end{cases}$$
Fermination:

Termination:

D(N,M) is distance

N	9									
0	8									
I	7									
Т	6									
N	5									
Е	4									
Т	3									
N	2									
I	1									
#	0	1	2	3	4	5	6	7	8	9
	#	Е	Χ	Е	С	U	Т	I	0	N

N	9									
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	#	Е	Χ	Е	С	U	Т	I	0	N

$$\begin{split} D(\textit{i,j}) = \min \quad \begin{cases} D(\textit{i-1,j}) + 1 \\ D(\textit{i,j-1}) + 1 \\ D(\textit{i-1,j-1}) + \\ 0; \ \ \textit{if} \ S_1(\textit{i}) \neq S_2(\textit{j}) \end{cases} \end{split}$$

N	9	8	9	10	11	12	11	10	9	8
0	8	7	8	9	10	11	10	9	8	9
I	7	6	7	8	9	10	9	8	9	10
Т	6	5	6	7	8	9	8	9	10	11
N	5	4	5	6	7	8	9	10	11	10
Е	4	3	4	5	6	7	8	9	10	9
Т	3	4	5	6	7	8	7	8	9	8
N	2	3	4	5	6	7	8	7	8	7
I	1	2	3	4	5	6	7	6	7	8
#	0	1	2	3	4	5	6	7	8	9
	#	Е	Χ	Е	С	U	Т	I	0	N

Edit distance: 8

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• Computing edit distance may not be sufficient for some applications

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  - We often need to align characters of the two strings to each other

Which parts of the two strings are aligned?

What edit operations were required to go from one string to the other?

- Computing edit distance may not be sufficient for some applications
  - We often need to align characters of the two strings to each other
- We do this by keeping a "backtrace"
- Every time we enter a cell, remember where we came from
- When we reach the end,
  - Trace back the path from the upper right corner to read off the alignment

N	9									
0	8									
т	7									
1	/									
Т	6									
N	5									
Е	4									
Т	3									
N	2									
I	1									
#	0	1	2	3	4	5	6	7	8	9
	#	Е	Χ	Е	C	U	Т	I	0	N

$$D(i,j) = \min \left\{ \begin{cases} D(i-1,j) + 1 \\ D(i,j-1) + 1 \\ D(i-1,j-1) + \end{cases} \left[ \begin{array}{c} 2; & \text{if } S_1(i) \neq S_2(j) \\ 0; & \text{if } S_1(i) = S_2(j) \end{array} \right. \right.$$

							_			
N	9									
0	8									
Ι	7									
Т	6									
N	5									
Е	4	3	4							
Т	3	4	5							
N	2	3	4							
I	1	2	3							
#	0	1	2	3	4	5	6	7	8	9
	#	Е	Χ	Е	С	U	Т	I	0	N

$$D(i,j) = \min \begin{bmatrix} D(i-1,j) + 1 \\ D(i,j-1) + 1 \\ D(i-1,j-1) + \\ 0; & \text{if } S_1(i) \neq S_2(j) \\ 0; & \text{if } S_1(i) = S_2(j) \end{bmatrix}$$

As we are filling up each cell, remember which operation was used at each step (the lowest-cost operation)

### Minimum Edit with Backtrace

n	9	↓ 8	<u></u> ∠←↓9	<b>∠</b> ←↓ 10	∠←↓ 11	∠←↓ 12	↓ 11	↓ 10	↓9	<b>/8</b>	
0	8	↓ 7	<b>∠</b> ←↓8	∠←↓ 9	<b>∠</b> ←↓ 10	∠←↓ 11	↓ 10	↓ 9	∠ 8	← 9	
i	7	↓ 6	∠←↓ 7	∠-↓8	∠←↓ 9	∠←↓ 10	↓9	∠ 8	← 9	← 10	
t	6	↓ 5	∠←↓ 6	∠-↓7	∠<∪↓ 8	∠←↓ 9	∠ 8	← 9	← 10	<b>←</b> ↓ 11	
n	5	↓ 4	<b>∠</b> ←↓ 5	∠←↓ 6	∠←↓ <b>7</b>	∠←↓ <b>8</b>	<b>∠</b> ←↓9	<b>∠</b> ←↓ 10	∠←↓ 11	<b>∠</b> ↓ 10	
e	4	∠3	← 4	<b>∠</b> ← 5	← 6	← 7	<b>←</b> ↓ 8	<b>∠</b> ←↓9	∠←↓ 10	↓9	
t	3	∠-↓4	∠← <b>↓</b> 5	∠-↓6	∠-↓7	∠←↓ 8	∠ 7	←↓ 8	∠←↓ 9	↓ 8	
n	2	∠ <b></b>	∠←↓ 4	∠←J 5	∠←↓ 6	∠←↓ 7	∠←↓ 8	↓ 7	∠←↓ 8	Z 7	
i	1	∠←↓ 2	∠←↓3	∠←↓ 4	∠<↓ 5	∠<↓ 6	∠←↓ 7	∠ 6	← 7	← 8	
#	0	1	2	3	4	5	6	7	8	9	
	#	e	X	e	c	u	t	i	0	n	

Left arrow implies insertion Down arrow implies deletion Diagonal arrow implies substitution

## Adding Backtrace to Minimum Edit

Base conditions:

$$D(i,0) = i$$
  $D(0,j) = j$ 

Termination:

D(0,j) = j D(N,M) is distance

Recurrence Relation:

Time

Time

O(nm)

Space

Time		
O(nm)		
Space		
Space O(nm)		
Backtrace		

Time

O(nm)

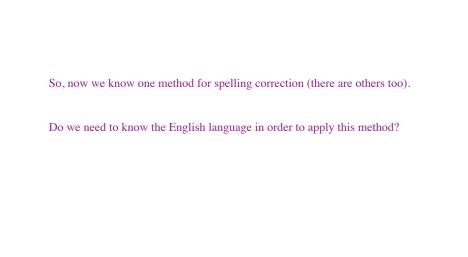
Space

O(nm)

**Backtrace** 

O(n+m)

Worst case complexity (all left arrows, followed by all down arrows)



So, now we know one method for spelling correction (there are others too).

Do we need to know the English language in order to apply this method?

Are there any situations where this method will fail, but a knowledge of the English language could have helped to identify spelling mistakes?