

12. Sorting

Wednesday, June 1, 2022 5:30 PM

<https://www.toptal.com/developers/sorting-algorithms>

- Matters for big data sets
- Focusing on Bubble, Selection, Insertion, Merge, Quick, Heap, Radix, Counting
badgoodnon-compare
- Arrays.sort uses TimSort for object arrays (stable)
QuickSort for primitive arrays
Insertion for small arrays
Merge for mostly sorted arrays
Dual pivot QuickSort for everything else

<http://hg.openjdk.java.net/jdk8/jdk8/jdk/file/46c727d6ecc2/src/share/classes/java/util/DualPivotQuicksort.java>

Tradeoffs (when to use which)

- Bubble Sort
 - Never use, $O(n^2)$ average
 - double for loop, comparing $O(i)$ and $O(i+1)$ n^2 times

```
void bubbleSort(int arr[])  
{  
    int n = arr.length;  
    for (int i = 0; i < n - 1; i++)  
        for (int j = 0; j < n - i - 1; j++)  
            if (arr[j] > arr[j + 1]) {  
                // swap arr[j+1] and arr[j]  
                int temp = arr[j];  
                arr[j] = arr[j + 1];  
                arr[j + 1] = temp;  
            }  
}
```

<https://www.geeksforgeeks.org/bubble-sort/>

- Selection Sort
 - slightly better than Bubble
 - $O(n^2)$ average, but selection has less swaps
 - Scans for smallest item, then swap the index

```
void sort(int arr[])  
{  
    int n = arr.length;  
  
    // One by one move boundary of unsorted subarray
```

<https://www.geeksforgeeks.org/selection-sort/>

```

for (int i = 0; i < n-1; i++)
{
    // Find the minimum element in unsorted array
    int min_idx = i;
    for (int j = i+1; j < n; j++)
        if (arr[j] < arr[min_idx])
            min_idx = j;

    // Swap the found minimum element with the first
    // element
    int temp = arr[min_idx];
    arr[min_idx] = arr[i];
    arr[i] = temp;
}
}

```

Good sorts

Tradeoffs (when to use which)

<https://www.geeksforgeeks.org/insertion-sort/>

- Insertion sort

- Good for small data sets and partially sorted data

- $O(n)$ best case, $O(n^2)$ avg/worse

- Split array into sorted and unsorted part, values from unsorted are picked and place in correct order in sorted part

GFG code

first pass:
 index 0 is unsorted section
 1 is sorted section

```

void sort(int arr[])
{
    int n = arr.length;
    for (int i = 1; i < n; ++i) {
        int key = arr[i]; // first unsorted element
        int j = i - 1;

```

$O(n^2)$ { for () {
 while () {
 }
 }
 }

repeat
 1) select first unsorted element
 2) swap other element to right
 to create correct pos &
 shift unsorted element

```

/* Move elements of arr[0..i-1], that are
greater than key, to one position ahead
of their current position */
while (j >= 0 && arr[j] > key) {
    arr[j + 1] = arr[j]; // swap
    j = j - 1; // go back index
}
arr[j + 1] = key;
}

```

- Shifts all element to the right
 to create pos for unsorted element
 - inserts unsorted element at correct pos

3) advance pointer to right one.

- Divide and conquer (Merge sort and Quick sort)
 - Merge: $O(n \log n)$ time, $O(n)$ space, stable sorting
 - Quick: $O(n \log n)$ time | $O(n^2)$ worst time, $O(\log n)$ space, not stable

<https://www.geeksforgeeks.org/merge-sort/>

Merge sort

- Divides array into 2 halves, calls itself for two halves (until it is 1 item), then merge sorted halves
- Stable, if there are equivalent elements, the original order is preserved
- compares local list to each other instead of every element
- $O(n)$ space b/c we need to create divided arrays

```
// Merges two subarrays of arr[].
// First subarray is arr[l..m]
// Second subarray is arr[m+1..r]
void merge(int arr[], int l, int m, int r)
{
    // Find sizes of two subarrays to be merged
    int n1 = m - l + 1;
    int n2 = r - m;

    /* Create temp arrays */
    int L[] = new int[n1];
    int R[] = new int[n2];

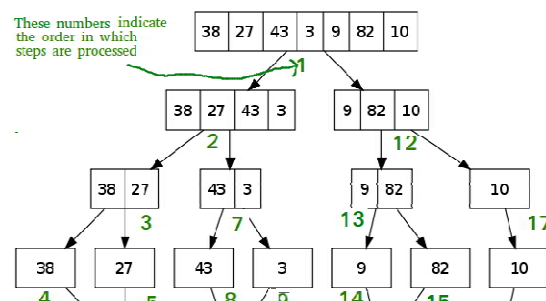
    /* Copy data to temp arrays */
    for (int i = 0; i < n1; ++i)
        L[i] = arr[l + i];
    for (int j = 0; j < n2; ++j)
        R[j] = arr[m + 1 + j];

    /* Merge the temp arrays */

    // Initial indexes of first and second subarrays
    int i = 0, j = 0;

    // Initial index of merged subarray array
    int k = l;
    while (i < n1 && j < n2) {
        if (L[i] <= R[j]) {
            arr[k] = L[i];
            i++;
        }
        else {
            arr[k] = R[j];
            j++;
        }
        k++;
    }
}
```

compare elements of
2 subarrays and merge



```

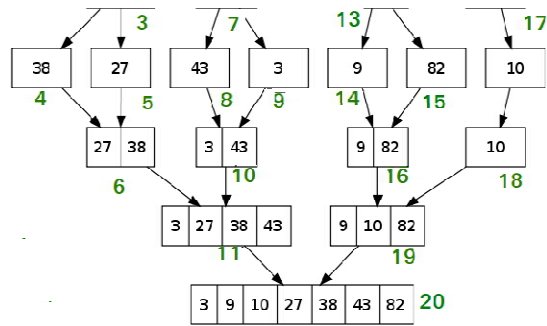
arr[k] = L[i];
i++;
}
else {
arr[k] = R[j];
j++;
}
k++;
}

/* Copy remaining elements of L[] if any */
while (i < n1) {
arr[k] = L[i];
i++;
k++;
}

/* Copy remaining elements of R[] if any */
while (j < n2) {
arr[k] = R[j];
j++;
k++;
}
}

```

compare elements of
2 subarrays and merge them



sort() recursively halves array
merge() combines array

```

// Main function that sorts arr[l..r] using
// merge()
void sort(int arr[], int l, int r)
{
if (l < r) {
// Find the middle point
int m = l + (r-l)/2;

// Sort first and second halves
sort(arr, l, m);
sort(arr, m + 1, r);

// Merge the sorted halves
merge(arr, l, m, r);
}
}

```

Main

int arr[] = {2, ..., 3}

← Merge sort of = new MergeSort();
ob.sort(arr, 0, arr.length - 1);

recursively splits until left of arr = right of arr
(only 1 element)
arr.length == 1

Quick sort

<https://www.geeksforgeeks.org/quick-sort/>

- picks an element as pivot (Many ways) and partitions the array around the pivot
- target of partition: given an element x as a pivot, place x in its correct position by putting all smaller elements before x and all larger elements after x

- $O(\log(n))$ space

- $O(n \log n)$ time average, $O(n^2)$ worse if bad pivot

// A utility function to swap two elements

static void swap(int[] arr, int i, int j)

```
{
    int temp = arr[i];
    arr[i] = arr[j];
    arr[j] = temp;
}
```

/* This function takes last element as pivot, places the pivot element at its correct position in sorted array, and places all smaller (smaller than pivot) to left of pivot and all greater elements to right of pivot */

static int partition(int[] arr, int low, int high)

```
{
    // pivot
    int pivot = arr[high];
```

```
    // Index of smaller element and
    // indicates the right position
    // of pivot found so far
    int i = (low - 1);
```

```
    for(int j = low; j <= high - 1; j++)
    {
```

```
        // If current element is smaller
        // than the pivot
        if (arr[j] < pivot)
        {
```

```
            // Increment index of
            // smaller element
            i++;
            swap(arr, i, j);
        }
    }
```

```
    swap(arr, i + 1, high);
    return (i + 1);
}
```

/* The main function that implements QuickSort

```
    arr[] --> Array to be sorted,
    low --> Starting index,
    high --> Ending index
```

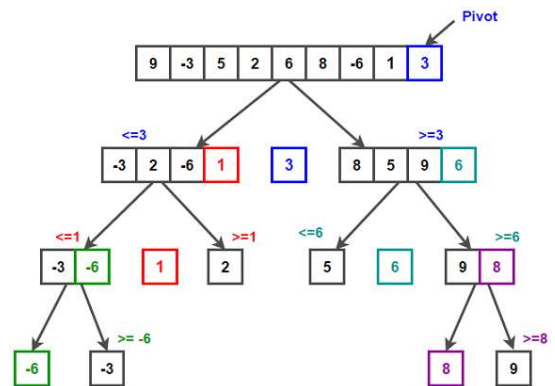
*/

static void quickSort(int[] arr, int low, int high)

```
{
    if (low < high)
    {
```

```
        // pi is partitioning index, arr[p]
        // is now at right place
        int pi = partition(arr, low, high);
```

```
        // Separately sort elements before
        // partition and after partition
        quickSort(arr, low, pi - 1);
        quickSort(arr, pi + 1, high);
    }
}
```



}
}

Merge vs Quick

- if worried about worst case : use merge
- if worried about memory : use quick
 - unless the sorting is external, then use merge
- merge sort is stable

Radix sort and counting sort (non-comparison sort)

- Only works with integers in a restricted range
- uses assumptions about how numbers are stored in memory

<https://opendatastructures.org/newhtml/ods/latex/sorting.html#tex2htm-121>