

5. Arrays

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<https://docs.oracle.com/javase/7/docs/api/java/util/Arrays.html>

<https://docs.oracle.com/javase/7/docs/api/java/lang/String.html>

- Organizes items sequentially in memory
index $0 \rightarrow n$

- Good for iterating if you know the indices
- * optimal for access, indexing (access element by index number), $O(1)$
 - insert/delete/search is $O(n)$

// ArrayList push is $O(1)$

- Treat String as array of characters

2 Types

1) Static - fixed size

2) Dynamic - reserved size for additional elements and can grow in size

String reversal practice

Input = ["H", "E", "L", "O", "O"]

Output = ["O", "O", "L", "E", "H"]

Approach:

2 pointer approach

needs: temp value for storage,
divide loop length by 2



```
class Solution {  
    public void reverseString(char[] s) {  
        for (int i = 0; i < s.length / 2; i++) {  
            char temp = 'a'; // Saves placeholder value for swapping  
            temp = s[i];
```

```
s[i] = s[s.length - (1 + i)];  
s[s.length - (1 + i)] = temp;
```

Time: $O(n/2) \rightarrow O(n)$, because we have 1 for loop that is dependent on size of input

Space: $O(1)$, because we create a constant space char

Merge Sorted Arrays Practice

Input:

- 2 integer arrays, `nums1` and `nums2`, sorted in non-decreasing order
- 2 integers, `m` and `n`, representing number of elements in `nums1` and `nums2` respectively

Output: Merge `nums1` and `nums2` into a single array sorted in non-decreasing order

Assume: final sorted array should be stored inside `nums1`. `nums1` has a length of `m + n`, where `m` are the elements that should be merged, and the last `n` elements are arbitrary and set to 0. `nums2` is length of `n`

Example 1:

Input: `nums1 = [1, 2, 3, 0, 0, 0]` `m = 3`

`nums2 = [2, 5, 6]`, `n = 3`

output: `nums1 = [1, 2, 2, 3, 5, 6]`

Approach 1:

create temp array with size M , do sort & insert
of $nums1[i]$ and $nums2[i]$ into temp array then copy into $nums1$
for loop

Time : $O(n)$, b/c 1 for loop

Time: $O(n)$, b/c 1 for loop
 size: $O(n)$, b/c we're creating an array dependent on input size

Approach 2: fill values starting the back of nums1 to avoid shifting elements

3-pointer approach, each pointer at the end of both arrays,
compare values at pointers and insert the larger value,
decrement the pointer that was inserted

Edge case: pointers < 0 , to avoid out of bounds error if any of the arrays are empty or when pointer goes to end of array

e.g. input: $\text{nums}_1 = [1]$, $m=1$ | $\text{nums}_2 = []$, $n=0$

Ans 4: 0

output: 0
 input: nums1 = [0], m = 0 | nums2 = [1], n = 1

output A: [1]

Input:

$\text{nums1} = [1, 2, 3, 0, 0, 0]$, $m = 3$
 $\text{nums2} = [2, 5, 6]$, $n = 3$

nums 2 = [2, 5, 6]

, n = 3

— [1, 2, 8, 0, 0, 0] | [1, 2, 3, ³0, ⁵0, ⁶0]
[2, 8, 6]

```
class Solution {
public void merge(int[] nums1, int m, int[] nums2, int n) {
    // 3 Pointer Approach
    int pointer1 = m-1;
    int pointer2 = n-1;
    int index = m + n -1;

    while (index >= 0) {
        // Edge case check if one of the arrays are empty or when
        // pointer goes to the end of an array
        if (pointer1 < 0) {
            nums1[index] = nums2[pointer2--];
        }
        else if (pointer2 < 0){
            nums1[index] = nums1[pointer1--];
        }

        // Comparison because we are not the end of either arrays
        else{
            if (nums1[pointer1] > nums2[pointer2]){
                nums1[index] = nums1[pointer1--];
            }
            else{
                nums1[index] = nums2[pointer2--];
            }
        }
        // Decrement after checks
        index--;
    }
}

// O(n+m) time, because it is dependent if n>m or n<m, O(max(n,m))

// O(1) space, because space created is constant
```