

# Queues

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→ first in, first out

→ enqueue (insert an element at end of queue)

→ dequeue (removes & returns the element at front of queue)

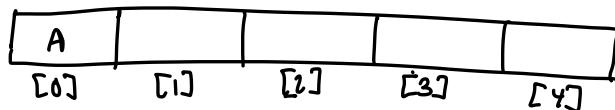
Public interface Queue<E> {

```
int size();
boolean isEmpty();
E first();
void enqueue(E e);
E dequeue();
```

## → Array Linear Queues

→ front of queue is fixed at first slot in the array whereas the rear of the queue moves down with each Enqueue. Now we dequeue the front element in the queue.

(a) queue.enqueue('A')

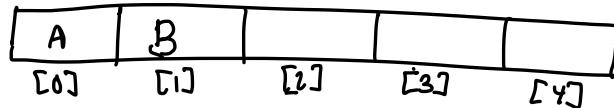


front = 0  
rear = 0

Enqueue: O(1)

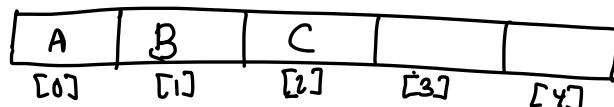
Dequeue: O(1)

(b) queue.enqueue('B')



front = 0  
rear = 1

(c) queue.enqueue('C')



front = 0  
rear = 2

\* In dequeue, instead of moving elements up to beginning of array, it merely just increments the index of front indicator to next slot

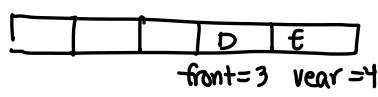
(d) queue.enqueue(item)



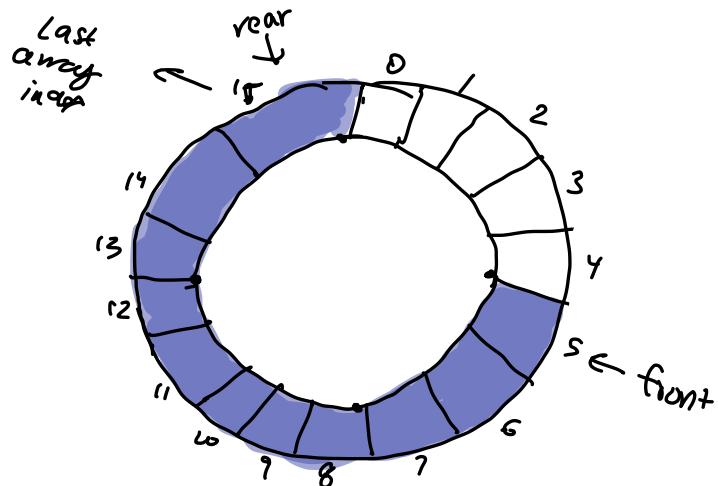
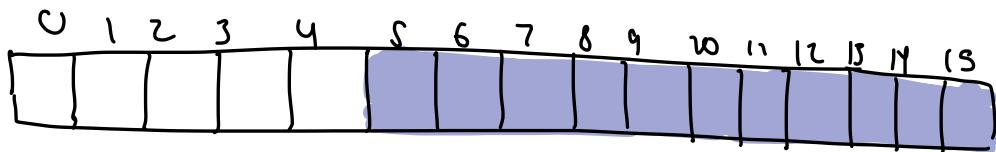
front = 1  
rear = 2

→ Queue is full when if  $\text{rear} == \text{Max-Items} - 1$

But when removing elements, we have empty slots at the low array index



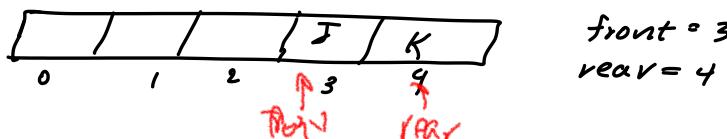
## → Circular Array Queue



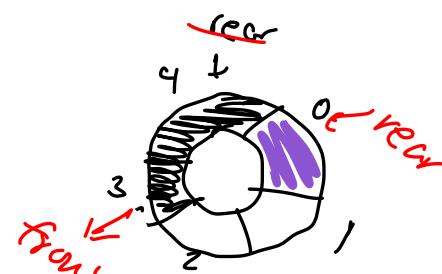
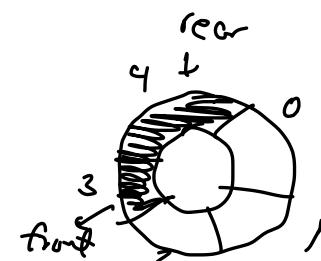
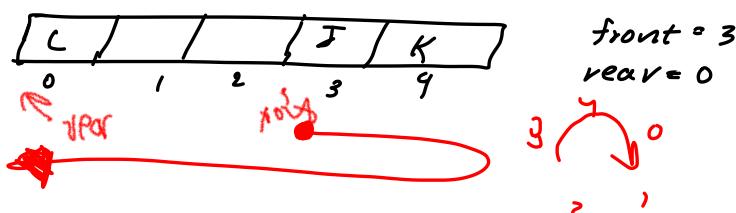
→ When rear reaches the last array index, move rear to under 0.

$$\text{rear} = \underline{\underline{(\text{rear} + 1) \% \text{ maxQue}}};$$

(a) queue.enqueue('L')

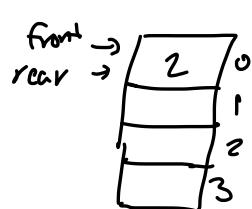


(b) can wrap the queue around to top of array

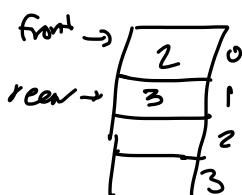


## → Example

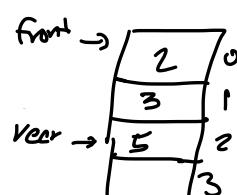
• q.enq(2)



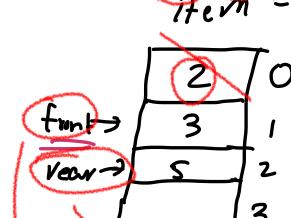
• q.enq(3)



• q.enq(5)



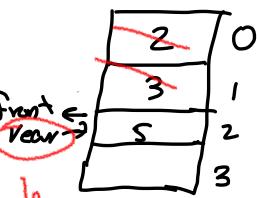
• q.deq(item)



→ front moves to next spot & 2 is deleted

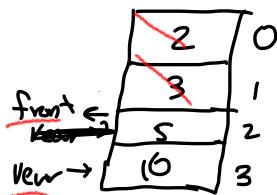
deq move

•  $q.\text{deq}(\text{item})$   
 $\text{item} = 3$

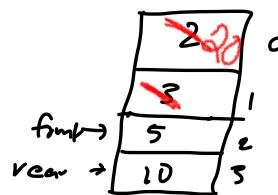


↓  
 doesn't move

•  $q.\text{enq}(10)$



•  $q.\text{enq}(20)$  ??

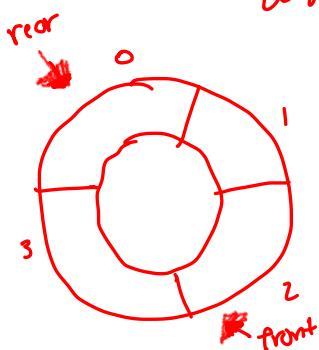
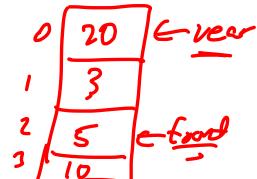


let queue elements wrap around  
 if ( $\text{rear} == \text{maxQue} - 1$ )  
 $\text{rear} = 0;$   
 else  
 $\text{rear} = \text{rear} - 1;$

or

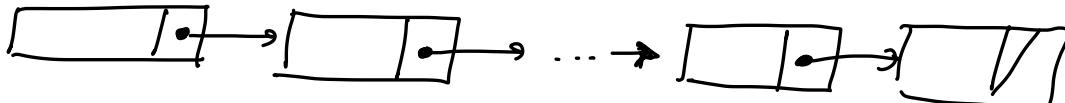
$\text{rear} = (\text{rear} + 1) \% \text{maxQue};$

wrap arrd:



## → Linked List Based Implementation

$q.\text{Front}$



$q.\text{Rear}$



•  $q.\text{enqueue}(\text{Item} \& \text{newItem})$

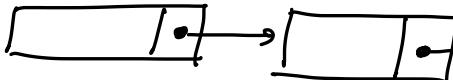
• function: adds newItem to rear of queue

• Precond: queue has been initialized & is not full

→  $\text{rear}.\text{next} = \text{newNode}$

⇒  $\text{rear} = \text{newNode}$

$q.\text{Front}$



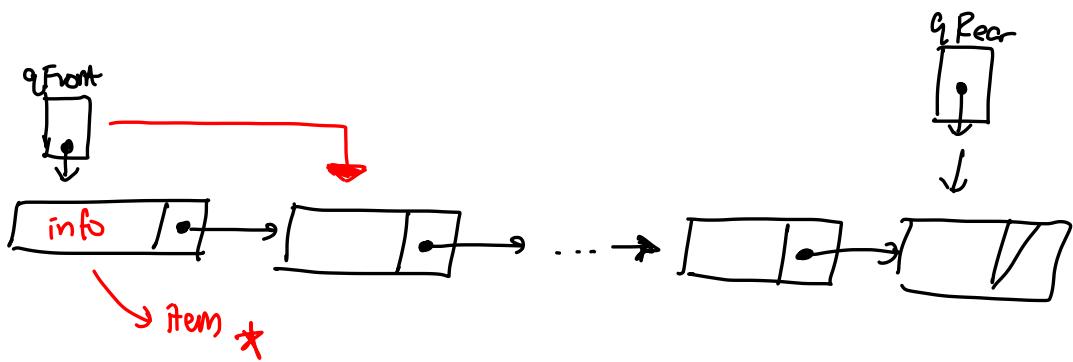
$q.\text{Rear}$



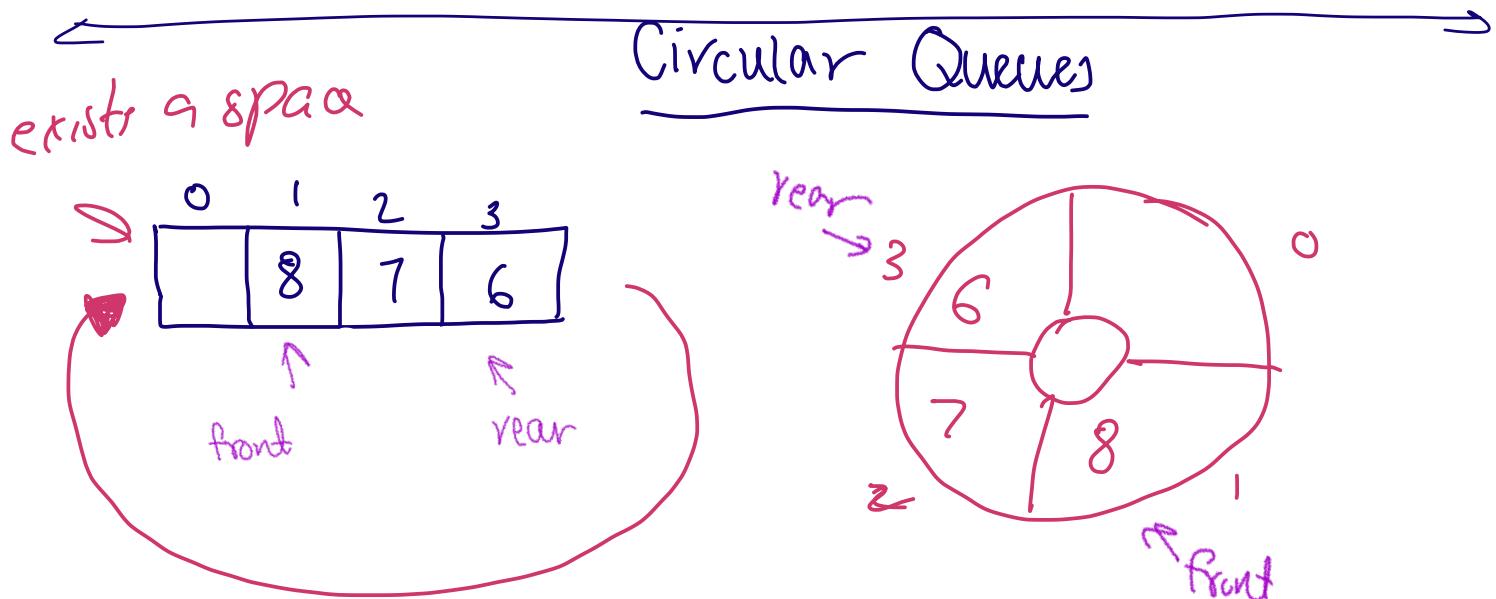
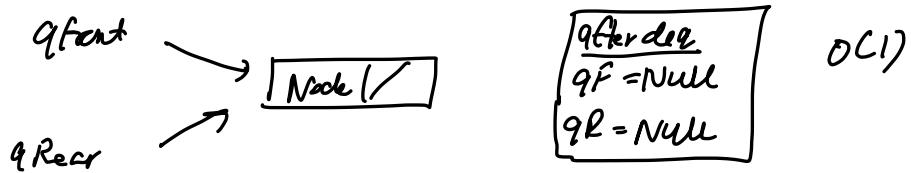
$\text{newNode}$

• dequeue:

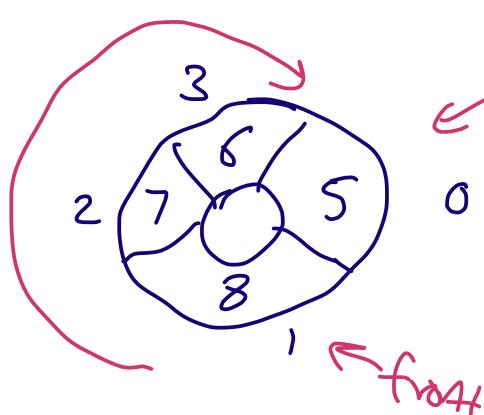
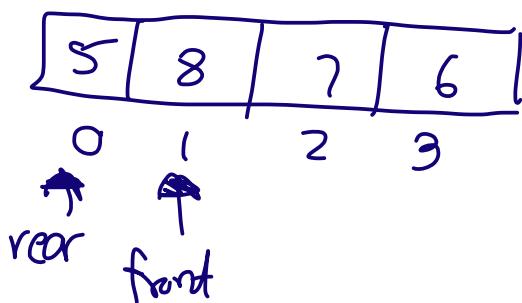
frontpointer = front . next



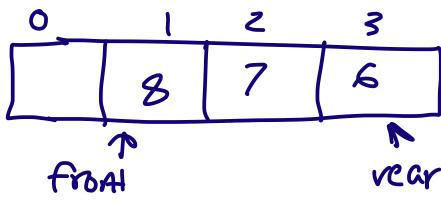
\* Special Case: Queue contains only one element



enqueue(5);



rear (still ahead of front)  
value of rear is < value of front, but position is head

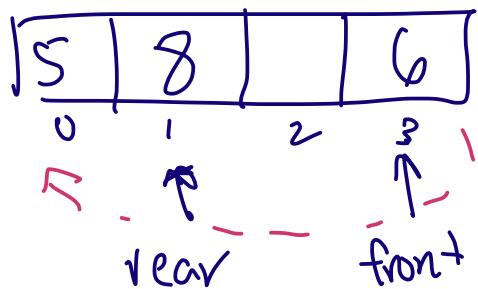


Add(s)

$$\text{rear} = 4 \% 4 \\ \text{rear} = 0$$

$$\text{rear}++ = 4$$

$$\leftarrow \text{rear} = \text{rear} + 1 \% \text{length}$$



Delete from front()

$$\text{front}++ = 4$$

$$\text{front} = \text{front} + 1 \% \text{length}$$

Code:

void addrear (int element) {

$$\text{if}(\text{rear} - \text{front} \geq \text{arr.length} - 1)$$

$$\text{if}(\text{front} - \text{rear} = 1) \{$$

// Queue is full

return;

\}

$$\text{if}(\text{front} == -1) \{$$

front = 0;

rear = 0;

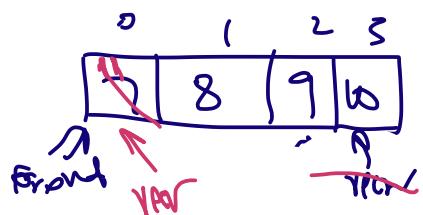
arr[0] = element;

return;

\} (rear + 1)

$$\text{rear} = (\text{rear} + 1) \% \text{arr.length};$$

arr[rear] = element;



add(11)

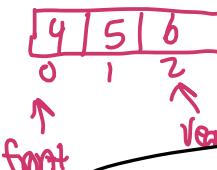
so

$$\text{rear} = 4 \% 4 = 0$$

Note

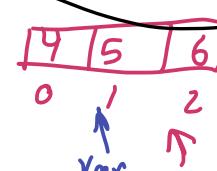
→ 2 instances when queue full:

①



$$\bullet \text{rear} - \text{front} = \text{len} - 1$$

②



$$\bullet \text{front} - \text{rear} = 1$$



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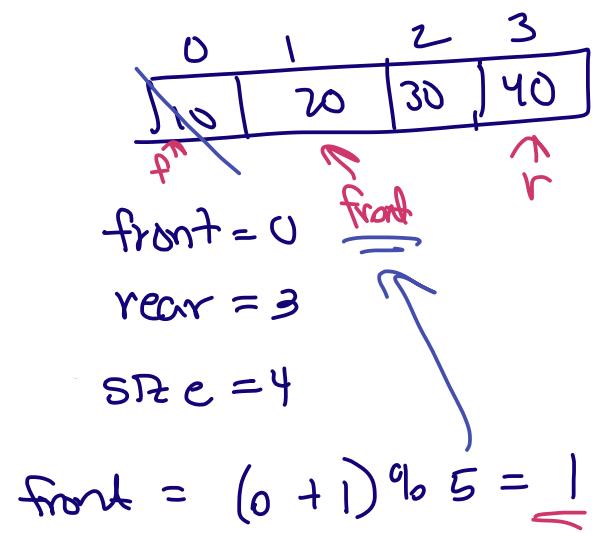
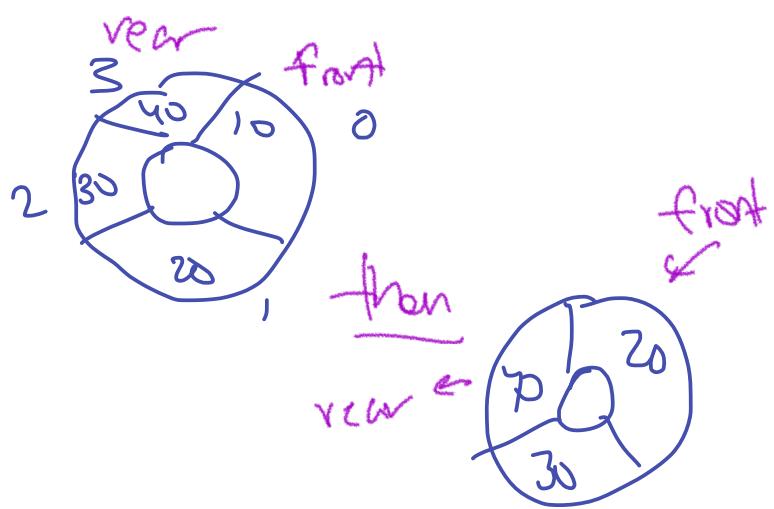
int deletefront() {
    if (front == -1)
        return -1;
    }

    if (rear == front) {
        // last element
        int temp1 = arr[rear];
        front = -1;
        rear = -1;
        return temp1;
    }

    int temp = arr[front];
    front = (front + 1) % arr.length();
    return temp;
}

```

↑  
front rear



temp = 5  
 front = -1  
 rear = -1

Conventions to indicate queue is empty

