

LINKED LISTS

→ Each node contains

- Data
- A pointer to next node in sequence

* * * Current != null vs Current.next != null * * *

→ $Current \neq null$: when you want to traverse the entire linked list, including the last node (start → end)

→ $Current.next \neq null$: when you need to stop right before the last node (when inserting at end or removing last node)

* * Remember * *

→ Update head

→ Don't access $Current.next$ when current is null

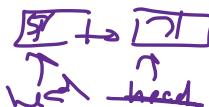
→ maintain references

→ advance current pointer

SLL

→ Operations on SLL

1) $addToFront() - O(1)$



• newNode's next to head

• $newNode \rightarrow head (head = newNode)$
~~newNode = head~~

2) $addToBack() - O(n)$

• Iterate until current's next is null NOT until current is null

• General Case:

1) $current \rightarrow head$

3) If next is null, you are sitting at the last node & all you need to do is set the next pointer to new node

• Edge Case: if head is null, point head to newNode

→ To add a node to empty list

1) create a new node containing data

2) set that newNode's next to current head

3) set head to newNode

→ To Iterate through SLL

1) create current node

2) $current = head$

3) $while (current \neq null) \{$

...
 $current = current.next;$

$\}$

4) To Remove from back



• $current = head$

• $while (current.next.next \neq null) \{ current = current.next; \}$
 $current.next = null;$

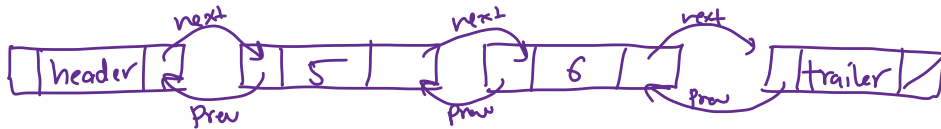
3) To Remove from Front



• Save data from head for returning

• $head = head.next$

Doubly LL



- Size: 0 \rightarrow both head & tail point to null
- Size: 1 \rightarrow both point to single node

\rightarrow To add to front

- 1) create newNode
- 2) newNode.next = head
- 3) head.prev = newNode
- 4) head = newNode

\rightarrow To add to back

- 1) newNode.prev = trailer.prev
- 2) trailer.prev.next = newNode
- 3) trailer.prev = newNode
- 4) newNode.next = trailer

\rightarrow To remove from back



it's okay if lastNode points to things, as long as nothing is pointing to it &

- 1) trailer.prev = trailer.prev.prev
- 2) trailer.prev.next = trailer

\rightarrow To remove from front



head = head.next

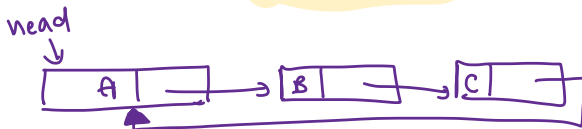
head.prev = null

Circular Linked Lists

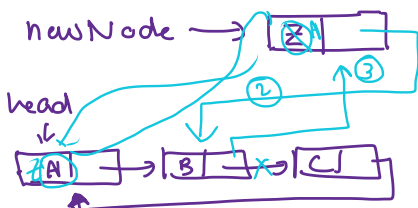
\rightarrow last node points back to head

\rightarrow can no longer use current == null to check if we've reached end of our list

\rightarrow must use current == head to terminate our loop



\rightarrow Adding to the front in O(1)



- 1) create a new empty node
- 2) newNode.next = head.next
- 3) head.next = newNode
- 4) newNode.data = head.data
- 5) head.data = newNode.data
- 6) head.data = newNode.data

Don't do:
create a new node, point to head & move head then reset last node to point to new head cause O(n) &

\rightarrow Adding to back in O(1)

1) all steps performed to add to front

2) head = head.next

Stacks

- push(e) - adds element e to top of stack
- pop() - removes + returns top element
- top() - returns not removes the top element of stack
→ peek()
- Allows elements stored in the stack to belong to any object type $\langle E \rangle$
- Array-based stack
 - all operations are $O(1)$
 - Stack $\langle \text{Integer} \rangle$ $S = \text{new ArrayStack} \langle \rangle();$

