







MEMORY MODEL



**PROCESS**  
**A**

► Memory is split into pages (each 4KiB on x86)

▶ The kernel *maps* its own memory into each process



▶ This “kernel” memory is only accessible by the kernel











Userspace







Kernel space









b) and c) are completely different scenarios

process could try to access the pages via a printer

► c) Kernel modules are marked "kernel only" but the



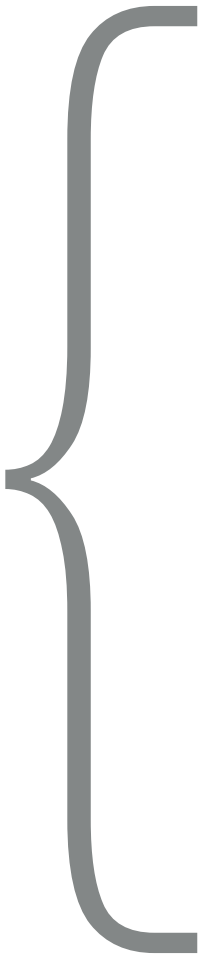
▶ **b)** Process **B** has no possibility to even *describe* the address













PROCES













OxOODOO.















