Knowledge, action, and the frame problem

Reasoning about Actions

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Outline

Situation calculus provides a framework for reasoning about actions.
This work presents an expansion to handle the *knowledge* possessed or acquired by the agent, and allow it to shape the agent's decisions.

- Knowledge is represented by one additional fluent
- Uniform axiomatization with the rest of sitcalc
- Ordinary actions and knowledge-producing ones are strictly separated
- Easy expansion of regression as defined in [Reiter2001]
- Desirable properties are direct consequences of the axiomatization (e.g. knowledge persistence / memory)

The K fluent

$$K(s', s)$$

Defines an accessibility relation between situations.

(Informal) definition

$K(s', s)$ is true if an agent in situation s could mistake the current situation for the other s', given its current knowledge.

Knowledge

Definition of knowledge

A fluent is known to be true (false) in a situation s if and only if it is true (false) in all situations accessible from s.

Shorthand notation: $\textbf{Knows}(\phi, s) \stackrel{\text{def}}{=} \forall s' \, \mathrm{K}(s', s) \rightarrow \phi(s')$

Knowledge-producing actions

Actions that have an effect on the agent's knowledge

SENSE actions

Learn the truth value of a formula. Example: check if a door is open or closed.

$$\textbf{Knows}(\mathrm{P}, \textsc{do}(\textsc{sense}_{\mathrm{P}}, s)) \vee \textbf{Knows}(\neg\mathrm{P}, \textsc{do}(\textsc{sense}_{\mathrm{P}}, s))$$

READ actions

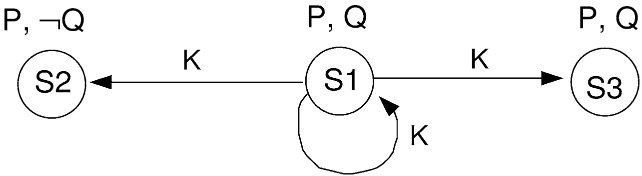
Learn the value of a term. Example: read a number on a sheet of paper.

$$\exists x\, \textbf{Knows}(\tau = x, \textsc{do}(\textsc{read}_{\tau}, s))$$

Knowledge effects

In order to complete the specification of the K fluent, we need to define its successor state axiom, determining how ordinary actions and knowledge-producing actions affect it.

Consider this case study with three accessible situations. The agent is in S1.



$$\textbf{Knows}(\mathrm{P}, S1) \wedge \neg\textbf{Knows}(\mathrm{Q}, S1)$$

Ordinary actions

Assume the agent performs a $\textsc{drop}$ action.

Informal idea

The agent cannot distinguish the current situation s from all the other s' accessible from it. Therefore, after executing the action, the agent may believe to be in any situation resulting from any s' after executing $\textsc{drop}$.

Axiomatization

$$\mathrm{K}(s'', \textsc{do}(\textsc{drop}, s)) \equiv \exists s' \, (\textsc{Poss}(\textsc{drop}, s') \land \mathrm{K}(s', s) \land s'' = \textsc{do}(\textsc{drop}, s'))$$

Ordinary actions

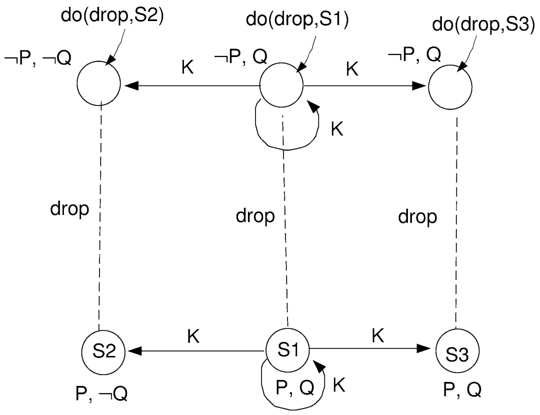
The **only knowledge gained is that the** DROP **action has been performed**, as well as anything that can be derived from the action effects.
For example, if DROP makes P false:

$$\mathrm{P}(\mathrm{DO}(a, s)) \equiv a \neq \mathrm{DROP} \wedge \mathrm{P}(s)$$

then

$$\textbf{Knows}(\neg \mathrm{P}, \mathrm{DO}(\mathrm{DROP}, S1))$$

but no extra knowledge is gained about Q.



Knowledge-producing actions

Consider an action $\textsc{sense}_{\mathrm{Q}}$ that provides information on whether Q is true or false. We define a **sensing result function** to represent the signal received by the agent in response:

Sensing result function

$$\textsc{sr}(\textsc{sense}_{\mathrm{Q}}, s) = r \equiv (r = \text{``}\textsc{yes}\text{''} \wedge \mathrm{Q}(s)) \vee (r = \text{``}\textsc{no}\text{''} \wedge \neg \mathrm{Q}(s))$$

Knowledge-producing actions

When the agent executes SENSE_Q, what are the accessible situations afterwards?

Informal definition

Initially, the agent cannot distinguish the current situation s from all the other s' accessible from it. Therefore, after executing SENSE_Q, the agent may believe to be in any situation that:

- results from any s' after executing the action,
- **AND** would yield the same sensing result as the one that has been observed.

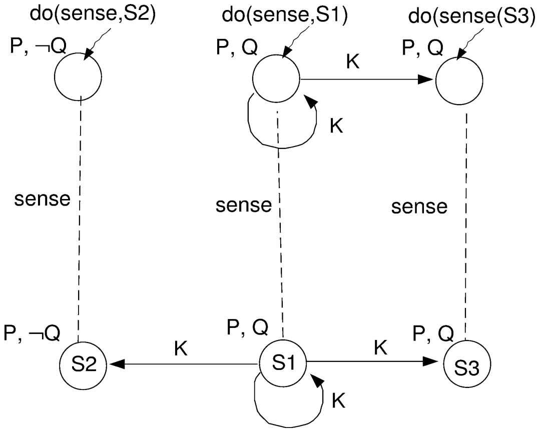
Axiomatization

$$\begin{aligned}\text{K}(s'', \text{DO}(\text{SENSE}_\text{Q}, s)) \equiv \exists s' \, (&\text{Poss}(\text{SENSE}_\text{Q}, s') \land \text{K}(s', s) \land \\ & s'' = \text{DO}(\text{SENSE}_\text{Q}, s') \land \text{SR}(\text{SENSE}_\text{Q}, s) = \text{SR}(\text{SENSE}_\text{Q}, s'))\end{aligned}$$

Knowledge-producing actions

After executing $\textsc{sense}_{\mathrm{Q}}$, **only situations with the same truth value for Q are accessible**.
Thus, in addition to knowing that $\textsc{sense}_{\mathrm{Q}}$ has been performed, the agent now knows the truth value of Q as well, by definition.

$$\textbf{Knows}(\mathrm{Q}, \textsc{do}(\textsc{sense}_{\mathrm{Q}}, S1))$$



Sensing results in general

The concept of sensing result extends to all types of action, allowing for a uniform axiomatization.

Ordinary actions

$$\mathrm{SR}(\mathrm{DROP}, s) = r \equiv r = \text{“OK”}$$

SENSE-type knowledge-producing actions

$$\mathrm{SR}(\mathrm{SENSE}_{\mathrm{Q}}, s) = r \equiv (r = \text{“YES”} \wedge \mathrm{Q}(s)) \vee (r = \text{“NO”} \wedge \neg \mathrm{Q}(s))$$

READ-type knowledge-producing actions

$$\mathrm{SR}(\mathrm{SENSE}_{\tau}, s) = r \equiv r = \tau(s)$$

The successor state axiom for K

Putting it all together, the definitive form is as follows:

Successor state axiom for the K fluent

$$\mathrm{K}(s'', \textsc{do}(a, s)) \equiv \exists s' \, (\textsc{Poss}(a, s') \wedge \mathrm{K}(s', s) \wedge \\ s'' = \textsc{do}(a, s') \wedge \textsc{sr}(a, s) = \textsc{sr}(a, s'))$$

What about...

- **...mixing ordinary and knowledge effects?**
  We assume that ordinary and knowledge actions are disjoint: each action is going to be axiomatized as either affecting *only* the K fluent or as not affecting it at all. This does not cause loss of generality.
- **...knowledge of arbitrary formulae?**
  They already work within this system.
  Example: **Knows**$(\forall x(\textsc{man}(x) \rightarrow \textsc{mortal}(x)) \wedge \textsc{man}(Socrates))$

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Theorems

Theorem 1 (Knowledge-producing effects). *For all situations* s*, all fluents* P *(other than* K*) and knowledge-producing action terms* α*, if* $\mathrm{P}(s)$ *then* $\mathrm{P}(\textsc{do}(\alpha, s))$.

Theorem 2 (Default persistence of ignorance). *For an action* α *and a situation* s*, if* $\neg\mathbf{Knows}(\mathrm{P}, s)$ *holds and the axiomatization entails*

$$\forall s\ \mathrm{P}(s) \equiv \mathrm{P}\big(\textsc{do}(\alpha, s)\big)$$

and

$$\forall y\ \neg\mathbf{Knows}\big(\big(\textsc{Poss}(\alpha) \wedge \textsc{sr}(\alpha) = y\big) \rightarrow \mathrm{P}, s\big)$$

then

$$\neg\mathbf{Knows}\big(\mathrm{P}, \textsc{do}(\alpha, s)\big)$$

holds as well.

Theorems

Theorem 3 (Knowledge incorporation). *For a knowledge-producing action α, a fluent or the negation of a fluent* F*, a fluent or the negation of a fluent* P*, and a situation s, if the axiomatization entails*

$$\exists y\ \mathbf{Knows}\big(\mathrm{F} \equiv \textsc{sr}(\alpha) = y, s\big)$$

and also

$$\mathrm{F}(s), \quad \textsc{Poss}(\alpha, s),$$

and

$$\mathbf{Knows}(\mathrm{F} \rightarrow \mathrm{P}, s)$$

hold, then

$$\mathbf{Knows}\big(\mathrm{P}, \textsc{do}(\alpha, s)\big)$$

holds as well.

Theorems

Theorem 4 (Memory). *For all fluents* P *and situations* s*, if* **Knows**(P, s) *holds then* **Knows**$(\mathrm{P}, \textsc{do}(\alpha, s))$ *holds as long as the axiomatization entails*

$$\forall s\ \mathrm{P}(s) \equiv \mathrm{P}\big(\textsc{do}(\alpha, s)\big)$$

Theorem 5 (Knowledge of effects of actions). *If* α *is an ordinary (not a knowledge-producing) action, and if the axiomatization entails*

$$\forall s\ \phi[s] \rightarrow \mathrm{P}\big(\textsc{do}(\alpha, s)\big)$$

where ϕ *is an arbitrary formula with situation terms suppressed*[8] *and* P *is a fluent or its negation, then the following is also entailed:*

$$\mathbf{Knows}\big((\textsc{Poss}(\alpha) \wedge \phi), s\big) \rightarrow \mathbf{Knows}\big(\mathrm{P}, \textsc{do}(\alpha, s)\big)$$

Reasoning Tasks

Projection Task

Determing if a sentence G is true in the situation resulting from the execution of an action sequence is represented by the query:

$$F \models G(\mathrm{DO}([\alpha_1, \ldots, \alpha_n], S_0))$$

Reasoning Tasks

Projection Task

Determing if a sentence G is true in the situation resulting from the execution of an action sequence is represented by the query:

$$F \models G(\mathrm{DO}([\alpha_1, \ldots, \alpha_n], S_0))$$

Legality Task

The situation term

$$\mathrm{DO}(\alpha_m, \mathrm{DO}(\alpha_{m-1} \ldots \mathrm{DO}(\alpha_1, S_0) \ldots))$$

is a legal situation iff

$$[\alpha_1, \ldots, \alpha_m]$$

is a legal action sequence

Regression

- Reduce Reasoning about future situations to reasoning about the initial situation;
- Sound and Complete Reasoner.

Regression

Definition of Regressor Operator through:

- Ordinary Actions
- Knowledge-Producing Actions

Regression Theorem

Theorem

For any ground situation term

$$F \models G(s_{gr}) \quad iff \quad F - F_{SS} \models R^*_\theta[G(s_{gr})]$$

Proof.

It suffices to show the preservation of logical equivalence

$$F \models G(s_{gr}) \equiv \mathrm{R}^*_\theta[G(s_{gr})]$$

$\square$

Regression Theorem

Theorem (Regression-Ordinary Actions)

$$R^*_\theta[\textbf{\textit{Knows}}(W, DO(a, s))] \equiv$$

$$\textbf{\textit{Knows}}(POSS(a) \rightarrow R^*_\theta[W[DO(a, s')]], s)$$

Regression-Ordinary Actions

Proof.

1. **Knows**(W, DO(a, s))
2. by the definition of **Knows**:

$$\forall s'' \ \mathsf{K}(s'', \mathsf{DO}(a, s)) \ \equiv \rightarrow \mathsf{W}[s'']$$

$\square$

Regression-Ordinary Actions

Proof.

3. by Successor State Axiom:

$$\mathsf{K}(s'', \mathsf{DO}(a, s)) \equiv \exists s'(\mathsf{K}(s', s) \land \mathsf{POSS}(a, s) \land s'' = \mathsf{DO}(a, s'))$$

$$\forall s''(\exists s'(\mathsf{K}(s', s) \land \mathsf{POSS}(a, s) \land s'' = \mathsf{DO}(a, s')) \rightarrow \mathsf{W}[s''])$$

4. the axiomatization entails:

$$\forall s, s' \mathsf{SR}(a, s) = \mathsf{SR}(a, s')$$

$$\forall s'(\mathsf{K}(s', s) \land \mathsf{POSS}(a, s') \rightarrow W[\mathsf{DO}(a, s')])$$

$\square$

Regression-Ordinary Actions

Proof.

5. Inductive Hypothesis:

$$\forall s'(\mathsf{K}(s', s) \wedge (\mathsf{POSS}(a, s'))) \rightarrow \mathsf{R}^*_\theta[W[\mathsf{DO}(a, s')]]$$

$$\forall s'(\mathsf{K}(s', s) \wedge (\mathsf{POSS}(a, s'))) \rightarrow \mathsf{R}^*_\theta[W[\mathsf{DO}(a, s')]]^{-1}$$

6. again by definition of **Knows**:

$$\textbf{Knows}(\mathsf{POSS}(a) \rightarrow \mathsf{R}^*_\theta[W[\mathsf{DO}(a, s')]]^{-1}, s)$$

$\square$

Regression Theorem

Theorem (Regression- Knowledge Producing Actions)

$$\mathcal{R}^*_\theta[\textbf{\textit{Knows}}(W, DO(SENSE_i, s))] \equiv$$

$$\exists y \; SR(SENSE_i, s) = y \wedge$$

$$\textbf{\textit{Knows}}((POSS(a) \wedge SR(SENSE_i) = y) \rightarrow \mathcal{R}^*_\theta[W[DO(SENSE_i)]]^{-1}, s)$$

Regression- Knowledge Producing Actions

(vi) $\quad \mathbf{Knows}\big(W, \textsc{do}(\textsc{sense}_i, s)\big)$

by the definition of **Knows**

$$\forall s''\ \mathrm{K}\big(s'', \textsc{do}(\textsc{sense}_i, s)\big) \rightarrow W[s'']$$

by the successor state axiom for K (sentence (18)), and also and the inductive hypothesis

$$\begin{aligned}&\forall s'\ \big(\mathrm{K}(s', s) \wedge \\ &\qquad \textsc{Poss}(\textsc{sense}_i, s') \wedge \textsc{sr}(\textsc{sense}_i, s) = \textsc{sr}(\textsc{sense}_i, s')\big) \rightarrow \\ &\qquad\quad \mathcal{R}_\Theta\big[W[\textsc{do}(\textsc{sense}_i, s')]\big]\end{aligned}$$

by the definition of equality and the existential quantifier

$$\begin{aligned}&\forall s'\ \big(\mathrm{K}(s', s) \wedge \\ &\quad \textsc{Poss}(\textsc{sense}_i, s') \wedge \exists y\ \textsc{sr}(\textsc{sense}_i, s) = y \\ &\qquad \wedge\ \textsc{sr}(\textsc{sense}_i, s') = y\big) \rightarrow\ \mathcal{R}_\Theta\big[W[\textsc{do}(\textsc{sense}_i, s')]\big]\end{aligned}$$

Regression- Knowledge Producing Actions

by the definition of the connectives and quantifiers

$$\begin{aligned}&\forall y\ \mathrm{SR}(\mathrm{SENSE}_i, s) = y \ \rightarrow \\ &\quad \forall s' \left(\mathrm{K}(s', s) \wedge \mathrm{Poss}(\mathrm{SENSE}_i, s') \wedge \mathrm{SR}(\mathrm{SENSE}_i, s') = y\right) \rightarrow \\ &\qquad \mathcal{R}_\Theta\left[W[\mathrm{DO}(\mathrm{SENSE}_i, s')]\right]\end{aligned}$$

by the definition of the connectives, quantifiers, and the fact that there can only be one denotation of $\mathrm{SR}(\mathrm{SENSE}_i, s)$

$$\begin{aligned}&\exists y\ \mathrm{SR}(\mathrm{SENSE}_i, s) = y \ \wedge\ \forall s' \left(\mathrm{K}(s', s) \rightarrow\right. \\ &\quad \left.\mathrm{Poss}(\mathrm{SENSE}_i, s') \ \wedge\ \mathrm{SR}(\mathrm{SENSE}_i, s') = y\right) \rightarrow \\ &\qquad \mathcal{R}_\Theta\left[W[\mathrm{DO}(\mathrm{SENSE}_i, s')]\right]\end{aligned}$$

by the definitions of **Knows**

$$\begin{aligned}&\exists y\ \mathrm{SR}(\mathrm{SENSE}_i, s) = y \wedge \\ &\quad \mathbf{Knows}\left(\left(\mathrm{Poss}(\mathrm{SENSE}_i) \wedge \mathrm{SR}(\mathrm{SENSE}_i) = y\right)\right. \\ &\qquad \left.\rightarrow \mathcal{R}_\Theta\left[W[\mathrm{DO}(\mathrm{SENSE}_i, s')]\right]^{-1}, s\right) \qquad \square\end{aligned}$$

Legality Testing

Method

Given n actions, the action precondition axioms are:

$$\forall x_n \ \mathsf{POSS}(A_n(x_n), s) \equiv \Pi A_n(x_n, s)$$

A sequence is a legal action sequence iff:

$$F \models \Pi(\alpha_1)[s \rightarrow S_0] \wedge \Pi(\alpha_2)[s \rightarrow \mathsf{DO}(\alpha_1, S_0)] \wedge \ldots$$

$$\wedge \Pi(\alpha_m)[s \rightarrow \mathsf{DO}([\alpha_1, \ldots, \alpha_{m-1}], S_0)]$$

So iff, by Regression:

$$F - F_{SS} \models \Pi(\alpha_1)[s \rightarrow S_0] \wedge \mathsf{R}^*_\theta[\Pi(\alpha_2)[s \rightarrow \mathsf{DO}(\alpha_1, S_0)]] \wedge \ldots$$

$$\wedge \mathsf{R}^*_\theta[\Pi(\alpha_m)[s \rightarrow \mathsf{DO}([\alpha_1, \ldots, \alpha_{m-1}], S_0)]]$$

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The Problem

- A robot has to manage **n plants** in a garden.
- The robot performs an action on one plant at a time.
- **A plant can be watered only if it is dry and the temperature is known.**
- The robot has access to a watering can with unlimited capacity.
- The robot performs action on an object only if it is near it.
- The robot can hold only one object at a time.

(Non)Fluents

Fluents

- $\textsc{Near}(x, s) \rightarrow$ Robot is near object x in situation s
- $\textsc{Holding}(x, s) \rightarrow$ Robot is holding object x in situation s
- $\textsc{Moist}(p, s) \rightarrow$ Plant p is moist in situation s
- $\textsc{Temperature}(p, s) \rightarrow$ Temperature value of the spot near plant p
- $\textsc{MoistPlants}(s) \rightarrow$ Number of moist plants

Non-Fluents

- $\textsc{WateringCan}(x) \rightarrow$ Object x is a watering can
- $\textsc{Thermometer}(x) \rightarrow$ Object x is a thermometer
- $\textsc{Moisturemeter}(x) \rightarrow$ Object x is a moisturemeter

Actions

All actions are to be axiomatized as affecting only either the K fluent or other fluents.

Normal

- $\textsc{goto}(x) \rightarrow$ Go to object x
- $\textsc{water}(p) \rightarrow$ Water plant p
- $\textsc{pickup}(x) \rightarrow$ Pick up object x
- $\textsc{putdown}(x) \rightarrow$ Put down object x

Knowledge

- $\textsc{checkmoisture}(p) \rightarrow$ Check moisture of plant p
- $\textsc{checktemperature}(p) \rightarrow$ Check the temperature of the spot near plant p

Effects

Sensing Result Axioms

- $\text{SR}(\text{GOTO}(x), s) = r \equiv r = \text{“OK”}$
- $\text{SR}(\text{CHECKMOISTURE}(p), s) = r \equiv (r = \text{“YES”} \land \text{Moist}(p, s)) \lor (r = \text{“NO”} \land \neg\text{Moist}(p, s))$
- $\text{SR}(\text{CHECKTEMPERATURE}(p), s) = r \equiv r = \text{Temperature}(p, s)$

Knowledge Action Effects

- **Kwhether**$(\text{Moist}(p), \text{DO}(\text{CHECKMOISTURE}(p), s))$
- **Kref**$(\text{Temperature}(p), \text{DO}(\text{CHECKTEMPERATURE}(p), s))$

Preconditions

- $\textsc{Poss}\,(\textsc{water}(p), s) \equiv \textsc{Near}(p, s) \land \textsc{Holding}(x, s) \land \textsc{WateringCan}(x) \land$ $\textbf{Knows}(\neg\textsc{Moist}(p), s) \land$ $\textbf{Kref}(\textsc{Temperature}(p), s)$
- $\textsc{Poss}\,(\textsc{pickup}(x), s) \equiv \textsc{Near}(x, s) \land \neg\exists y.\textsc{Holding}(y, s)$
- $\textsc{Poss}\,(\textsc{putdown}(x), s) \equiv \textsc{Holding}(x, s)$
- $\textsc{Poss}\,(\textsc{checkmoisture}(p), s) \equiv$ $\textsc{Near}(p, s) \land \textsc{Holding}(x, s) \land \textsc{MoistureMeter}(x)$
- $\textsc{Poss}\,(\textsc{checktemperature}(p), s) \equiv$ $\textsc{Near}(p, s) \land \textsc{Holding}(x, s) \land \textsc{Thermometer}(x)$

Successor State Axioms

In general $F(x, \textsc{do}(\alpha, s)) \equiv \Phi_F^+(x, a, s) \lor (F(x, s) \land \neg\Phi_F^-(x, a, s))$

- $\textsc{Near}(x, \textsc{do}(\alpha, s)) \equiv \alpha = \textsc{goto}(x) \lor (\textsc{Near}(x, s) \land \neg\exists y.\alpha = \textsc{goto}(y))$
- $\textsc{Holding}(x, \textsc{do}(\alpha, s)) \equiv \alpha =$
  $\textsc{pickup}(x) \lor (Holding(x, s) \land \neg\exists r.\alpha = \textsc{putdown}(x))$
- $\textsc{Moist}(p, \textsc{do}(\alpha, s)) \equiv \alpha = \textsc{water}(p) \lor (\textsc{Moist}(p, s) \land \neg\exists r.\alpha = \textsc{water}(p))$
- $\textsc{Temperature}(p, \textsc{do}(\alpha, s)) \equiv \textsc{Temperature}(p, s)$
- $\textsc{MoistPlants}(p, \textsc{do}(\alpha, s)) = n \equiv$
  $(\textsc{MoistPlants}(p, s) = n - 1 \land \alpha = \textsc{water}(p)) \lor (\textsc{MoistPlants}(p, s) = n)$

Theorems

1. Knowledge-producing actions do not change the state of the world. If $\mathrm{P}(s)$ then $\mathrm{P}(\textsc{do}(\alpha, s))$, given that α is knowledge-producing and P is not K
2. Nothing is learned about P by doing action α, as long as α does not affect P.
3. Agents know the consequences of knowledge acquired through knowledge-producing actions.
4. If the agent knows P at s, then P is also known at $\textsc{do}(\alpha, s)$ as long as the effect of α is not to make P false.
5. Agents know the effects of (ordinary) actions.

Initial Situation

- $\textsc{WateringCan}(c)$
- $\textsc{Thermometer}(t)$
- $\textsc{MoistureMeter}(m)$
- $\textsc{Plant}(p_1)$

  $\vdots$
- $\textsc{Plant}(p_n)$
- $\textsc{Near}(c)$

Goal

Have n moist plants, namely $\textsc{MoistPlants}(p, s) = n$

Solution

For every plant p, with the following situations

$$\mathrm{S}_1 =$$
$$\textsc{do}(\textsc{putDown}(m), \textsc{do}(\textsc{checkMoisture}(p),$$
$$\textsc{do}(\textsc{goTo}(p), \textsc{do}(\textsc{pickUp}(m), \textsc{do}(\textsc{goTo}(m), \mathrm{S}_0)))))$$
$$\mathrm{S}_2 =$$
$$\textsc{do}(\textsc{putDown}(t), \textsc{do}(\textsc{checkTemperature}(p),$$
$$\textsc{do}(\textsc{pickUp}(t), \textsc{do}(\textsc{goTo}(t), \mathrm{S}_1))))$$
$$\mathrm{S}_3 =$$
$$\textsc{do}(\textsc{putDown}(c), \textsc{do}(\textsc{water}(p),$$
$$\textsc{do}(\textsc{pickUp}(c), \textsc{do}(\textsc{goTo}(c), \mathrm{S}_2))))$$

In the end, we need to entail that $\textsc{MoistPlants}(p, s) = n$

We can use the successor state axioms for that

Legality testing (with a different S_0)

$$\textsc{do}(\textsc{water}(p), \textsc{do}(\textsc{pickup}(c), \textsc{do}(\textsc{goto}(c), \textsc{do}(\textsc{putdown}(t), \mathrm{S}_0))))$$

$$\Downarrow$$

We need to entail

$$\textsc{Near}(p, \mathrm{S}_0) \wedge \textsc{Holding}(t, \mathrm{S}_0) \wedge \textsc{Thermometer}(t, \mathrm{S}_0) \wedge \textbf{Knows}(\neg\textsc{Moist}(p, \mathrm{S}_0)) \wedge$$
$$\textbf{Kref}(\textsc{Temperature}(p), \textsc{do}(\textsc{checktemperature}(p), \mathrm{S}_0))$$

$$\Downarrow$$

$$\exists y. y = \textsc{Temperature}(p, s)$$

Golog

```
while ¬Knows(HealthyPlants = n) do
    (Πp) goto(p);
        checkMoisture(p);
        if Knows(¬Moist(p))
            goto(t);
                pickup(t);
                goto(p);
                checkTemperature(p);
                putdown(t);
            goto(c);
                pickup(c);
                goto(p);
                water(p);
                putdown(c);
end while
```

Knowledge, action, and the frame problem

Thank you for listening!

Any questions?