## **HOMEWORK 2**

The problems below are either restatements or slight adaptations of problems in Sipser.

- 1 Draw NFAs with the specified number of states recognizing the following languages over the binary alphabet:
  - **a.** the language  $\{0\}$ , with two states;
  - **b.** the language of strings that end in 00, with three states.
- 2 Draw NFAs for the following languages:
  - **a.** binary strings that begin with a 1 and end with a 0, or contain at least three 1s;
  - **b.** binary strings that contain the substring 1010 or do not contain the substring 110.
- Let  $L = \{w : w \text{ contains an even number of 0s and an odd number of 1s and does not contain the substring 01}. Draw a DFA with 5 states that recognizes <math>L$ .
- 4 Let L be the language of all strings over  $\{0, 1\}$  that do not contain a pair of 1s that are separated by an odd number of symbols. Draw a DFA with 5 states that recognizes L.
- Let  $M = (Q, \Sigma, \delta, q_0, F)$  be an NFA that recognizes a language L. Does the NFA  $(Q, \Sigma, \delta, q_0, Q \setminus F)$ , which is result of swapping the accept and reject states in M, necessarily recognize the complement of L? Prove or give a counterexample.
- 6 Let  $L_n$  be the language of all binary strings of the form 111...1 for some k that is a multiple of n. For each  $n \ge 1$ , construct a DFA or NFA that recognizes  $L_n$ .
- Let D be the language of binary strings that contain an equal number of occurrences of the substrings 01 and 10. Thus  $101 \in D$  because 101 contains a single 01 and a single 10, but  $1010 \notin D$  because 1010 contains two 10s and one 01. Construct a DFA or NFA for D.