### Loading...

### IDA Pro crash course

(don't tell pancake)

### Introduction to unpacking using r2

### Agenda

Introduction

Useful r2 commands for unpacking

Techniques and samples

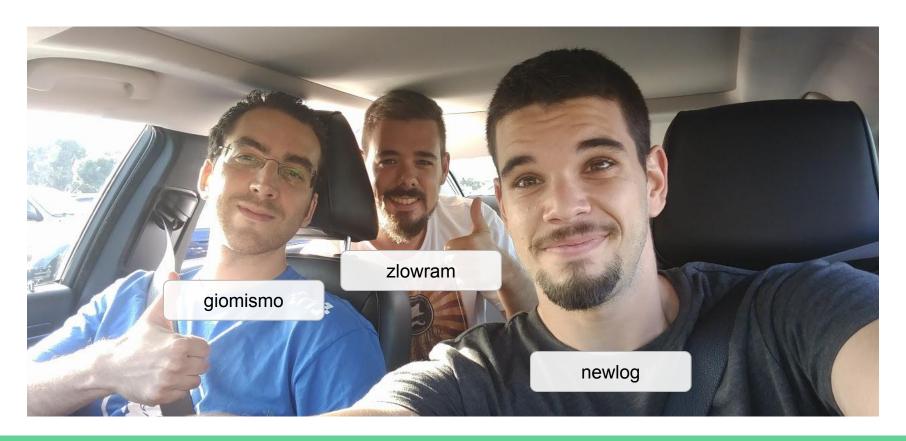
Dealing with the IAT

### Introduction

### Introduction

- About Us
- About You
- About Packers
- About PEs

### About Us

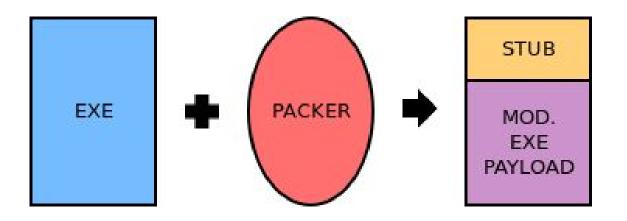


### **About You**

- 1. Who's dealt with malware unpacking?
- 2. Who's has unpacked malware with r2?
- 3. Who has successfully unpacked malware with r2?

### **About Packers**

What a packer is



### **About Packers**

Packer transformations

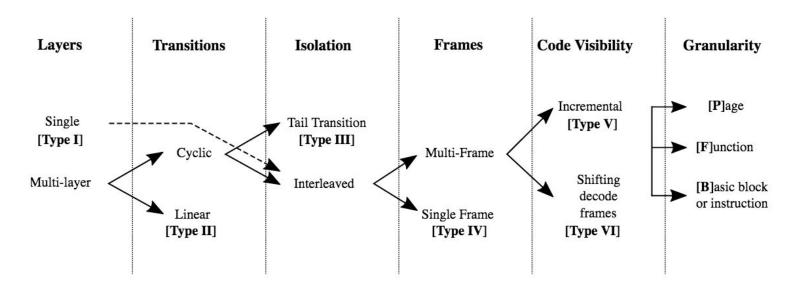
Compressors

Crypters

**Protectors** 

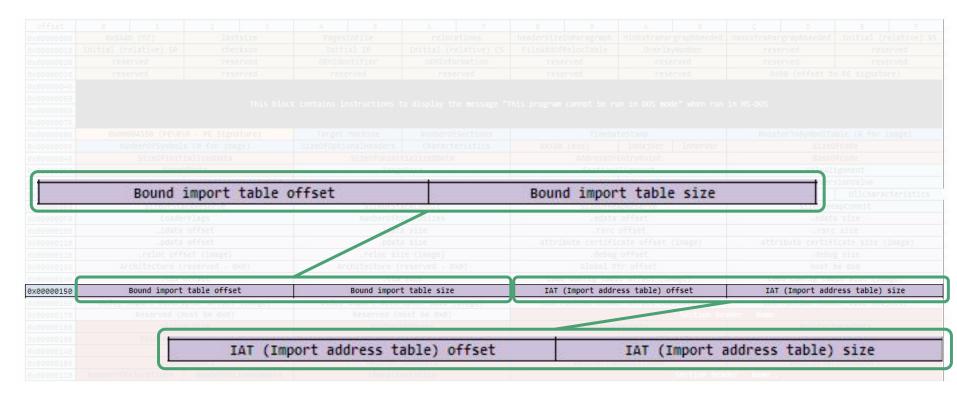
### **About Packers**

### Packer complexity types



img source: www.packerinspector.com/reference

### **About PEs**



img source: stackoverflow

### **About PEs**

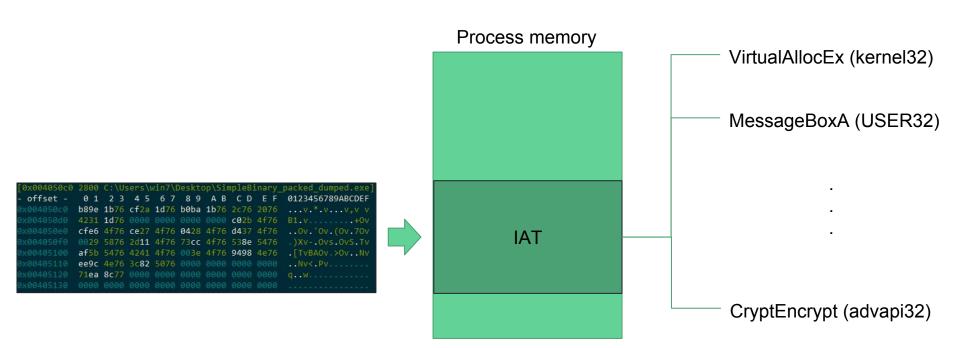
### **Import Table (IT)**

It is a structure defined within the PE header that defines which modules and functions will be imported, and thus used by the process.

### **Import Address Table (IAT)**

It is a table that the Windows Loader fills up with the virtual addresses where the API functions used by the PE reside into memory (they will not be mapped always to the same addresses). It can be found within the process mapped memory.

### **About PEs**



### **Follow Execution Flow**

r2 -d <file\_path>: opens file in r2 for debugging

dc: continues execution

ds: single step

dcr: continues until stack frame return

dr <register>: shows the value of specified register

### **Breakpoints**

**db** <address>: sets software breakpoint on specified address

**dbc** <address> <r2\_cmd>: assigns r2 command to be executed when bp on address is triggered

drx <number> <address> <len> <rwx>: sets hardware breakpoint on specified address range (by address and len) when accessed for read, write or execute

**db-** <address>: removes software breakpoint on specified address

drx- <number>: removes hardware breakpoint on specified address

### **Memory Maps**

**dm**: lists memory map of target process

dmi <address | libname>: lists loaded DLL symbols

### **Memory Dumping**

wtf <filename> <size> @<starting\_address>: saves memory region to file

dmd <filename> @<address>: dumps memory map at given address to a file

### **Misc**

/A <opcode>: searches for specified opcode on current section

**pf p4** @<address>: retrieves 4 bytes of content from pointer at given address

\*<address>: equivalent to pf p4 @ <address>

"e cmd.vprompt=px@esp": In visual mode, it shows the stack on top of the view

### **R2 Command Parsing**

"<text>:<row>[<column>]: Filters lines matching text, then selects the given row and finally returns the specified column

<cmd>j: Several r2 commands support this ending character to output data in json

# Techniques and samples

### **Useful Steps to Unpack Locky**

- Debug execution flow commands
- Hardware breakpoints
- Regular breakpoints
- List loaded DLL imports

### **Useful Commands to Unpack Locky**

- dcr
- drx 0 <address> 4 w
- db `<symbol\_from\_loaded\_DLL>`
- dmi kernel32~<symbol>:0[1]
- ds

### **VirtualAlloc**

### <u>Idea</u>

In order to write the PE into memory, Locky will reserve memory space. To do so, it uses the Windows API function **VirtualAlloc**.

We will put a breakpoint on this function and check if its parameters are appropriate to fit lengthy data (size).

For each VirtualAlloc "suspicious" call, we will set a (write) hardware breakpoint on its return memory address. We will check if the PE header (starting with "MZ") is written starting at that address.

```
dc
db `dmi kernel32~VirtualAlloc:0[1]`
dc _______ 10 times
[0x76202fb6]> pf p4 @ esp + 8^[1]
0x0001ae00 — VirtualAlloc Size
[0x76202fb6] > *esp + 16
0x0000004 — VirtualAlloc Permissions
dcr
S
drx 0 `dr eax` 4 w _____ eax contains VirtualAlloc returned memory address
dc
```

```
orne cates to labor to wheel and
on at hereits we want to be
```

### **Useful Steps to Unpack Dridex Dropper**

- Debug execution flow commands
- Regular breakpoints
- List loaded DLL imports

### **Useful Commands to Unpack Dridex Dropper**

- dc
- db `<symbol\_from\_loaded\_DLL>`
- dmi ntdll~<symbol>:0[1]
- \*<register+offset>

### RtIDecompressBuffer

### <u>Idea</u>

As with Locky, the Dridex Dropper uses VirtualAlloc in order to reserve memory space where the code will write the PE.

However the Windows API function **RtIDecompressBuffer** is used to fill the buffer that will contain the PE, so there is no need to search on all memory regions reserved by VirtualAlloc as shown with Locky.

This procedure is valid for unpacking Locky too.

```
dc
db `dmi ntdll ~RtlDecompressBuffer:0[1]`
dc
dcr
*esp+8
*esp+12
wtf rtl_dump.exe `*esp+12`@`*esp+8`
```

```
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```

#### **Useful Steps to Unpack UPX #1**

- Context switching instructions
- Hardware breakpoints on read

#### **Useful Commands to Unpack UPX #1**

- drx 0 <address> 4 r
- db <address>
- drx-
- dc

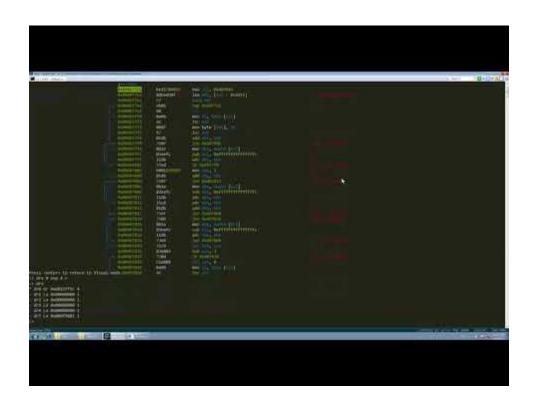
# pushal

#### <u>Idea</u>

Having a context switching instruction (i.e. **pushal**) at the beginning of the stub code is suspicious.

The **pushal** instruction stores all the registers into the stack so probably, at some point near the original code, it will restore those values back into the registers.

dc db entry0 dc ds drx 0 esp 4 r dc db 0x00407953 drxdc ds



#### **Useful Hints to Unpack UPX #2**

- Long jumps (not always though)
- Jumps with nonsense code afterwards

#### **Useful Commands to Unpack UPX #2**

- /x <bytes>
- e search.in = io.section

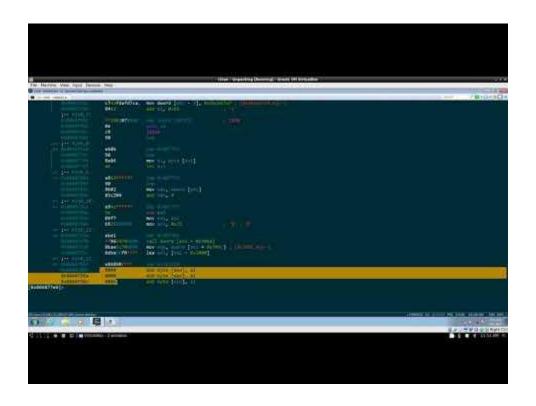
# Spotting the tail jump

#### <u>Idea</u>

The tail jump is the last instruction (a jump) executed within the stub before executing the original decompressed code.

Usually, the jump to the original code will be a long jump, and probably it will have a bunch of nonsense code afterwards, since it will never be executed.

```
dc
db entry0
dc
/A jmp
pd 4 @ hit*
<look for a jump with nonsense opcodes after>
```



#### **Useful Hints to Unpack UPX #3**

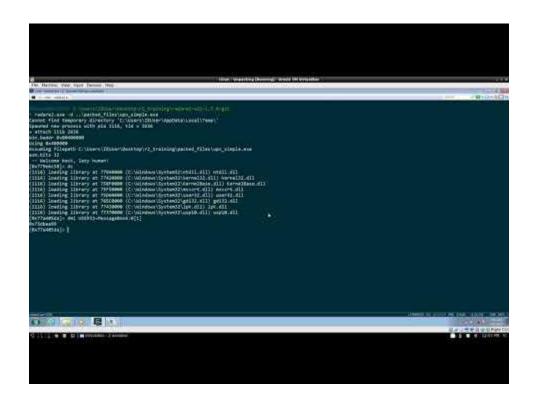
- API functions called by the original malware
- Breakpoints on API functions of loaded DLLs
- Function prologue

#### **Useful Commands to Unpack UPX #3**

- db <address>
- dmi <dll\_name>~<symbol>:0[1]
- dcr

# Breaking on common imports

```
dc
db `dmi USER32~MessageBox:0[1]`
dc
dcr
ds
<spot the function prologue>
```



Dealing with the IAT

# Dealing with the IAT

#### **Problem**

UPX and other packers stubs do not decompress the whole PE into memory. Instead they just decompressed all the sections of the PE, and they resolve all the API calls the original code used (i.e. they load the necessary modules, fix the references. etc).

To obtain a valid PE with the "original" code we need to:

- Find the IAT
- Dump the process
- Fix the dump (Fix IAT and Rebuild PE)

# Finding the IAT

Once the OEP has been spotted, the IAT will be somewhere in the memory.

Since the IAT contains references to API functions, the following steps could be followed:

- Get all the modules that the process has loaded (and the addresses of the functions they expose)
- Find an API pointer within the code (either call or jmp instructions)
- We now have a pointer to an API reference within the IAT
- Scan backwards and forward until no more API references are found, and we've found the start and end of the IAT (and also the size).

## Dumping the process

Once the IAT was found, we need to dump the process memory to disk.

The dump has to be a PE (not necessarily valid yet), so we can use the original PE structure, and dump the memory into it.

Alternatively, we could create a new PE if we have that capability.

# Fixing the dump

Once we have a PE, we need to make it valid:

- Update all references on the code to API functions to point to the new IAT (Fix IAT)
- Realign (Rebuild PE)
- Rebase (Rebuild PE)
- Update header checksum (Rebuild PE)

All this process is tedious and would be repetitive, so tools like Scylla now come handy

# Scylla

It is well known import reconstruction tool for x86/x64 (https://github.com/NtQuery/Scylla)

It automatizes all the process already described, and it's GPLv3 licensed!

However, it's full GUI based and does not integrate very well with the radare2 workflow.

# Scylla

# **R2SCYLLA**

#### R2SCYLLA

#### Scylla plugin for radare2

It allows full use of Scylla without the need of the Scylla GUI, integrating with the radare2 workflow.

Written in Python, uses the DLL compiled version of Scylla via CTypes.

#### R2SCYLLA

#### <u>Install</u>

git clone <a href="https://github.com/zlowram/r2scylla">https://github.com/zlowram/r2scylla</a>

#### <u>Usage</u>

#!pipe python r2scylla.py fulldump -h

usage: r2scylla.py fulldump [-h] pid oep base\_address binary\_path

#### R2SCYLLA

