Advanced ML course

Lecture 5: Word representations and Recurrent Neural Networks

Neychev Radoslav

ML Instructor (MIPT, HSE, Harbour.Space, BigData Team)

Research Scientist, MIPT

30.10.2019, Moscow, Russia

Outline

- 1. Classic word representations
- 2. Word embeddings
- 3. RNN intuitions
- 4. LSTM
- 5. Relations to CNN
- 6. Names generation from scratch
- 7. Q & A

Optional: Small attention outro

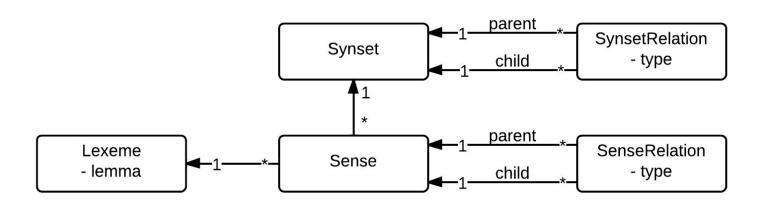
Different layers

Layers

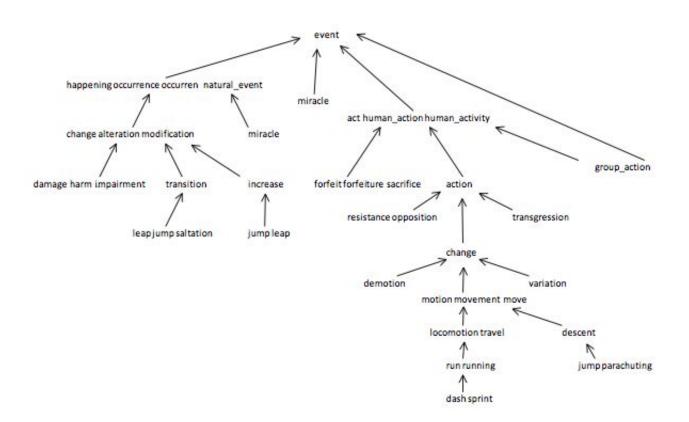
- a. Dense layer (done)
- b. Convolutional layer (next lecture)
- c. Pooling layer (next lecture)
- d. Dropout layer (done)
- e. Batchnorm layer (batch normalization) (done)
- f. Embeddings (aka word2vec, GloVe) (today)
- g. Recurrent layers (today)

How to represent text in a computer?

Use a taxonomy like WordNet that has hypernyms (is-a) relationships and synonym sets



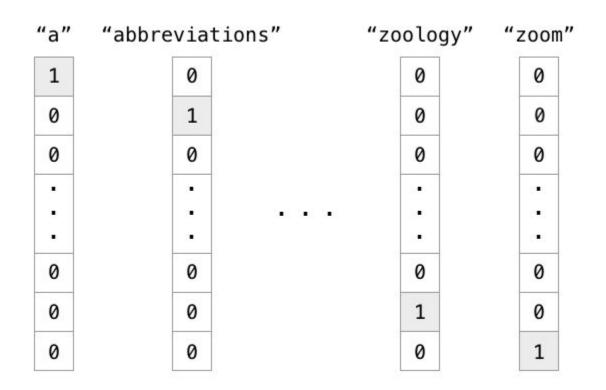
How to represent text in a computer: WordNet



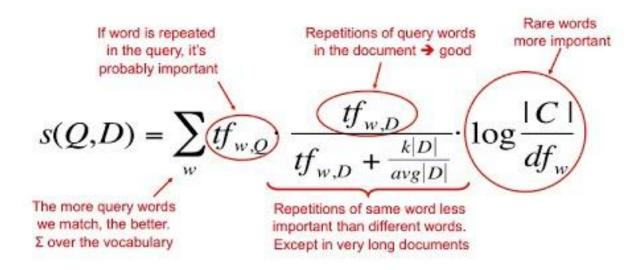
Discrete representations: problems

- Missing new words
- Subjective
- Requires human labor to create and adapt
- Hard to compute accurate word similarity

Discrete representations: one-hot encoding



Making it better: TF-iDF



TF - term frequency

iDF - Inversed Document Frequency

TF-iDF: an example

$$ext{tf("this",}\,d_1)=rac{1}{5}=0.2$$
 $ext{tf("this",}\,d_2)=rac{1}{7}pprox 0.14$ $ext{idf("this",}\,D)=\log\Bigl(rac{2}{2}\Bigr)=0$

$\operatorname{tfidf}("this",d_1,D) = 0.2 imes 0 = 0$
$ ext{tfidf}(" ext{this}",d_2,D)=0.14 imes0=0$

Document 1

Term	Term Count
this	1
is	1
a	2
sample	1

Document 2

Term	Term Count
this	1
is	1
another	2
example	3

Word 'this' is not very informative

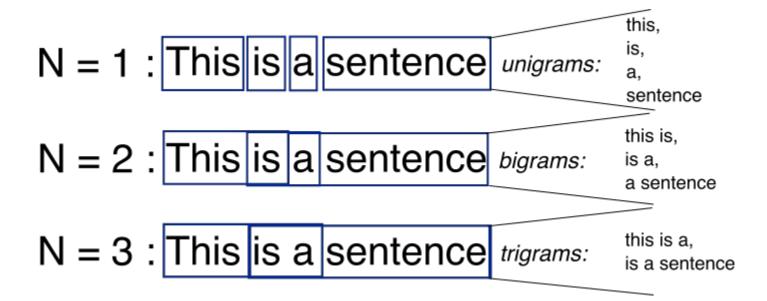
Words cooccurrences

One of the most successful ideas of statistical NLP:

"You shall know a word by the company it keeps"

(J. R. Firth 1957: 11)

Words cooccurrences: n-gramms



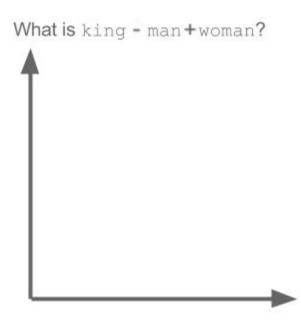
N-gramm approach problems

- Increase in size with vocabulary
- Very high dimensional: require a lot of storage
- Subsequent classification models have sparsity issues

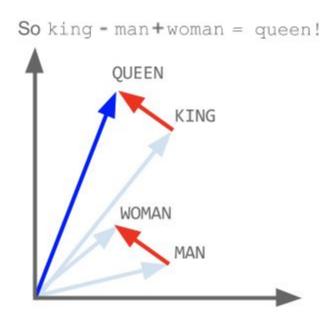


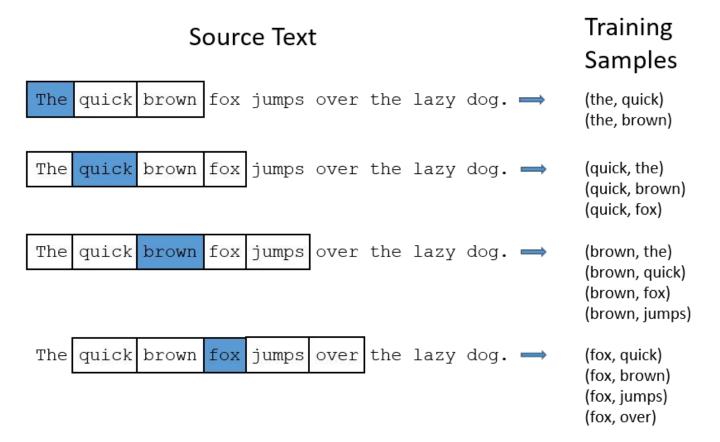
Models are less robust

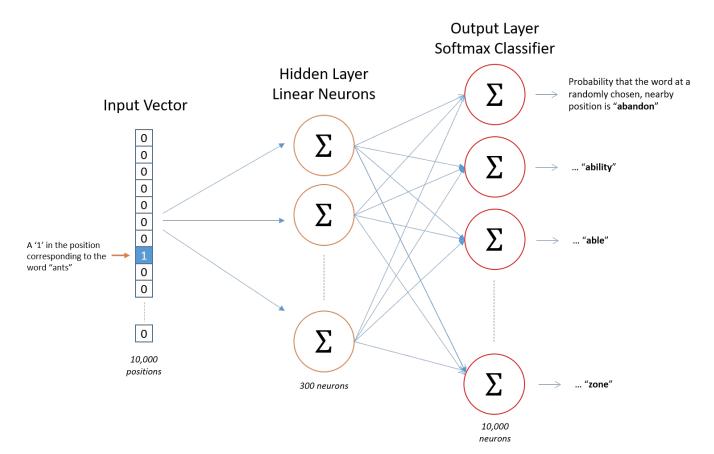
Embeddings: intuition

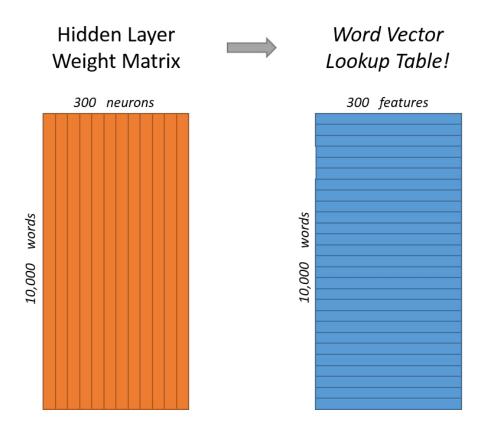


Embeddings: intuition









- Word vectors with 300 components
- Vocabulary of 10,000 words.
- Weight matrix with 300 x 10,000 = 3 million weights each!

Training is too long and computationally expensive

How to fix this?

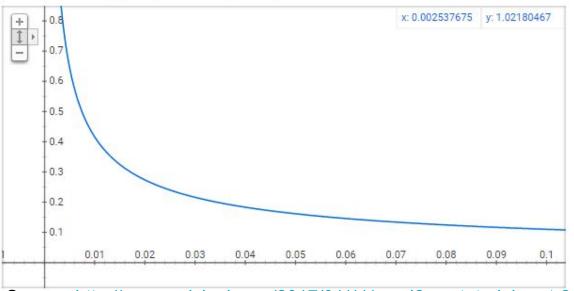
Basic approaches:

- 1. Treating common word pairs or phrases as single "words" in their model.
- 2. Subsampling frequent words to decrease the number of training examples.
- 3. Modifying the optimization objective with a technique they called "Negative Sampling", which causes each training sample to update only a small percentage of the model's weights.

Subsampling frequent words.

 w_i is the word, $z(w_i)$ is the fraction of this word in the whole text,

Graph for (sqrt(x/0.001)+1)*0.001/x



 $P(w_i)$ is the probability of *keeping* the word:

$$P(w_i) = (\sqrt{\frac{z(w_i)}{0.001}} + 1) \cdot \frac{0.001}{z(w_i)}$$

Source: http://mccormickml.com/2017/01/11/word2vec-tutorial-part-2-negative-sampling/

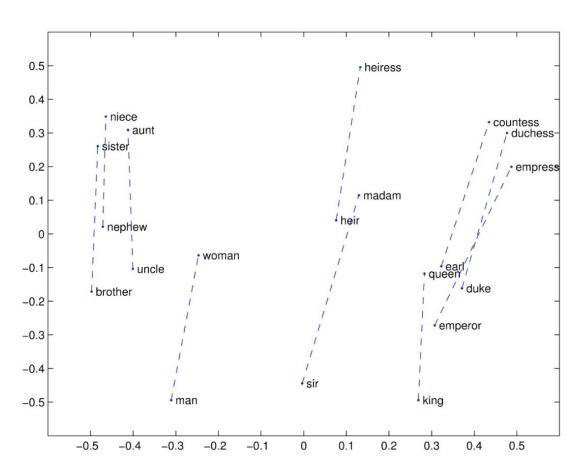
Embeddings: negative sampling

Negative Sampling idea: only few words error is computed. All other words has zero error, so no updates by the backprop mechanism.

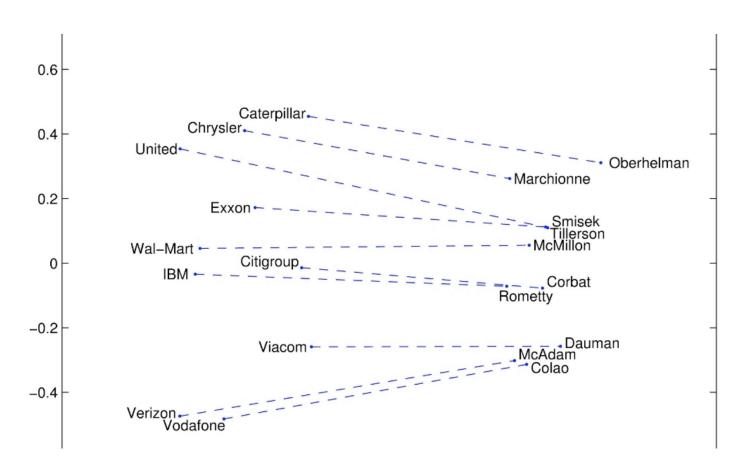
More frequent words are selected to be negative samples more often. The probability for a selecting a word is just it's weight divided by the sum of weights for all words.

$$P(w_i) = \frac{f(w_i)^{3/4}}{\sum_{j=0}^{n} (f(w_j)^{3/4})}$$

GloVe Visualizations



GloVe Visualizations: Company - CEO



Small outro about word representations

Word vectors are simply vectors of numbers that represent the meaning of a word

Approaches:

- One-hot encoding
- Bag-of-words models
- Counts of word / context co-occurrences
- TF-IDF
- Predictions of context given word (skip-gram neural network models, e.g. word2vec)

RNNs generating...

Shakespeare

PANDARUS:

Alas, I think he shall be come approached and the day When little srain would be attain'd into being never fed, And who is but a chain and subjects of his death, I should not sleep.

Second Senator:

They are away this miseries, produced upon my soul, Breaking and strongly should be buried, when I perish The earth and thoughts of many states.

DUKE VINCENTIO:

Well, your wit is in the care of side and that.

Second Lord:

They would be ruled after this chamber, and my fair nues begun out of the fact, to be conveyed, Whose noble souls I'll have the heart of the wars.

Clown:

Come, sir, I will make did behold your worship.

VIOLA:

I'll drink it.

Algebraic Geometry (Latex)

```
Proof. Omitted.
Lemma 0.1. Let C be a set of the construction.
   Let C be a gerber covering. Let F be a quasi-coherent sheaves of O-modules. We
have to show that
                                     \mathcal{O}_{\mathcal{O}_{+}} = \mathcal{O}_{X}(\mathcal{L})
Proof. This is an algebraic space with the composition of sheaves \mathcal{F} on X_{Oute} we
                           \mathcal{O}_X(\mathcal{F}) = \{morph_1 \times_{\mathcal{O}_X} (\mathcal{G}, \mathcal{F})\}
where G defines an isomorphism F \to F of O-modules.
Lemma 0.2. This is an integer Z is injective.
Proof. See Spaces, Lemma ??.
Lemma 0.3. Let S be a scheme. Let X be a scheme and X is an affine open
covering. Let U \subset X be a canonical and locally of finite type. Let X be a scheme.
Let X be a scheme which is equal to the formal complex.
The following to the construction of the lemma follows.
Let X be a scheme. Let X be a scheme covering. Let
                      b: X \rightarrow Y' \rightarrow Y \rightarrow Y \rightarrow Y' \times_X Y \rightarrow X.
be a morphism of algebraic spaces over S and Y.
Proof. Let X be a nonzero scheme of X. Let X be an algebraic space. Let F be a
quasi-coherent sheaf of O_X-modules. The following are equivalent

 F is an algebraic space over S.

    (2) If X is an affine open covering.
Consider a common structure on X and X the functor O_X(U) which is locally of
finite type.
```

Linux kernel (source code)

```
* If this error is set, we will need anything right after that BSD.
static void action new function(struct s stat info *wb)
 unsigned long flags;
 int lel idx bit = e->edd, *sys & -((unsigned long) *FIRST COMPAT);
 buf[0] = 0xFFFFFFFF & (bit << 4);
 min(inc, slist->bytes);
 printk(KERN WARNING "Memory allocated %02x/%02x, "
   "original MLL instead\n"),
   min(min(multi run - s->len, max) * num data in),
   frame pos, sz + first seg);
 div u64 w(val, inb p);
 spin unlock(&disk->queue lock);
 mutex unlock(&s->sock->mutex);
 mutex unlock(&func->mutex);
 return disassemble(info->pending bh);
```

Proof. Omitted.

Lemma 0.1. Let C be a set of the construction.

Let C be a gerber covering. Let F be a quasi-coherent sheaves of O-modules. We have to show that

$$\mathcal{O}_{\mathcal{O}_X} = \mathcal{O}_X(\mathcal{L})$$

.

Proof. This is an algebraic space with the composition of sheaves F on $X_{\acute{e}tale}$ we have

$$\mathcal{O}_X(\mathcal{F}) = \{morph_1 \times_{\mathcal{O}_X} (\mathcal{G}, \mathcal{F})\}$$

where G defines an isomorphism $F \to F$ of O-modules.

Lemma 0.2. This is an integer Z is injective.

Proof. See Spaces, Lemma ??.

Lemma 0.3. Let S be a scheme. Let X be a scheme and X is an affine open covering. Let $U \subset X$ be a canonical and locally of finite type. Let X be a scheme. Let X be a scheme which is equal to the formal complex.

The following to the construction of the lemma follows.

Let X be a scheme. Let X be a scheme covering. Let

$$b: X \to Y' \to Y \to Y \to Y' \times_X Y \to X.$$

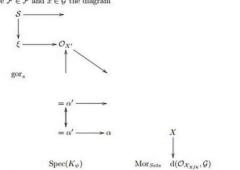
be a morphism of algebraic spaces over S and Y.

Proof. Let X be a nonzero scheme of X. Let X be an algebraic space. Let \mathcal{F} be a quasi-coherent sheaf of \mathcal{O}_X -modules. The following are equivalent

- F is an algebraic space over S.
- (2) If X is an affine open covering.

Consider a common structure on X and X the functor $\mathcal{O}_X(U)$ which is locally of finite type.

This since $\mathcal{F} \in \mathcal{F}$ and $x \in \mathcal{G}$ the diagram



is a limit. Then G is a finite type and assume S is a flat and F and G is a finite type f_* . This is of finite type diagrams, and

- the composition of G is a regular sequence,
- O_{X'} is a sheaf of rings.

Proof. We have see that $X = \operatorname{Spec}(R)$ and $\mathcal F$ is a finite type representable by algebraic space. The property $\mathcal F$ is a finite morphism of algebraic stacks. Then the cohomology of X is an open neighbourhood of U.

Proof. This is clear that G is a finite presentation, see Lemmas ??. A reduced above we conclude that U is an open covering of C. The functor F is a "field

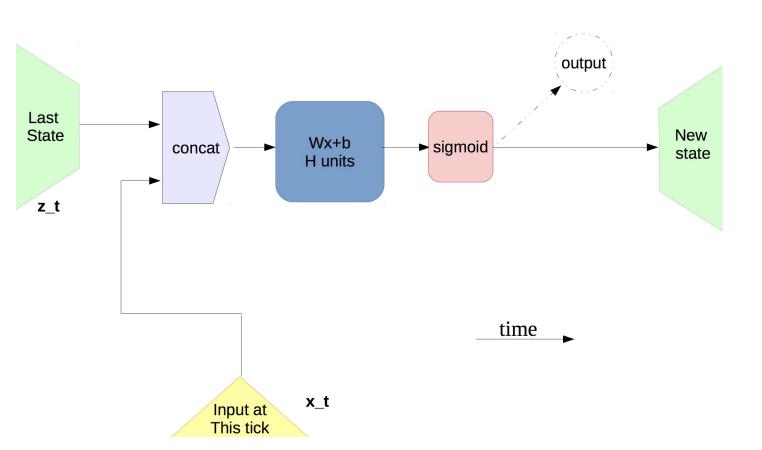
$$\mathcal{O}_{X,x} \longrightarrow \mathcal{F}_{\overline{x}} -1(\mathcal{O}_{X_{\ell tale}}) \longrightarrow \mathcal{O}_{X_{\ell}}^{-1}\mathcal{O}_{X_{\lambda}}(\mathcal{O}_{X_{\eta}}^{\overline{v}})$$

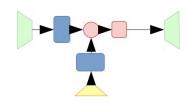
is an isomorphism of covering of \mathcal{O}_{X_i} . If \mathcal{F} is the unique element of \mathcal{F} such that X is an isomorphism.

The property \mathcal{F} is a disjoint union of Proposition ?? and we can filtered set of presentations of a scheme \mathcal{O}_{X} -algebra with \mathcal{F} are opens of finite type over S. If \mathcal{F} is a scheme theoretic image points.

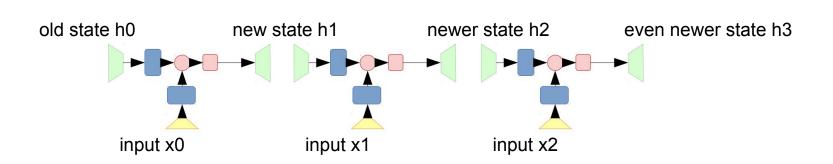
If \mathcal{F} is a finite direct sum $\mathcal{O}_{X_{\lambda}}$ is a closed immersion, see Lemma ??. This is a sequence of \mathcal{F} is a similar morphism.

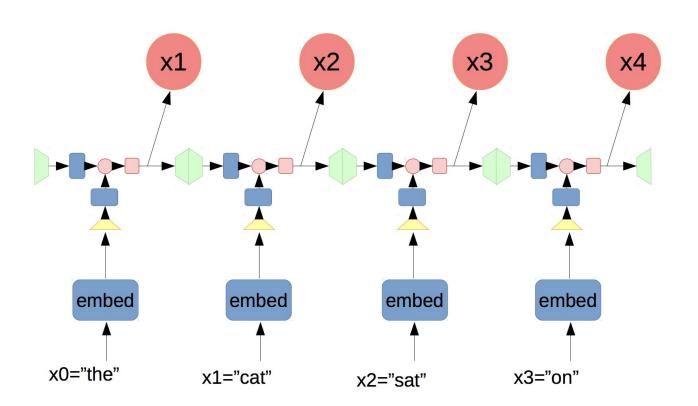
```
#include <asm/io.h>
#include <asm/prom.h>
#include <asm/e820.h>
#include <asm/system info.h>
#include <asm/setew.h>
#include <asm/pgproto.h>
#define REG PG vesa slot addr pack
#define PFM NOCOMP AFSR(0, load)
#define STACK DDR(type) (func)
#define SWAP ALLOCATE(nr)
                            (e)
#define emulate sigs() arch get unaligned child()
#define access rw(TST) asm volatile("movd %%esp, %0, %3" :: "r" (0)); \
 if ( type & DO READ)
static void stat PC SEC read mostly offsetof(struct seg argsqueue, \
         pC>[1]);
static void
os prefix(unsigned long sys)
#ifdef CONFIG PREEMPT
 PUT PARAM RAID(2, sel) = get state state();
  set_pid_sum((unsigned long)state, current_state_str(),
          (unsigned long)-1->lr full; low;
```





We use same weight matrices for all steps





Now with formulas

$$h_{0} = \bar{0}$$

$$h_{1} = \sigma(\langle W_{\text{hid}}[h_{0}, x_{0}] \rangle + b)$$

$$h_{2} = \sigma(\langle W_{\text{hid}}[h_{1}, x_{1}] \rangle + b) = \sigma(\langle W_{\text{hid}}[\sigma(\langle W_{\text{hid}}[h_{0}, x_{0}] \rangle + b, x_{1}] \rangle + b)$$

$$h_{i+1} = \sigma(\langle W_{\text{hid}}[h_{i}, x_{i}] \rangle + b)$$

$$P(x_{i+1}) = \operatorname{softmax}(\langle W_{\text{out}}, h_{i} \rangle + b_{\text{out}})$$

RNNs generating...

Shakespeare

PANDARUS:

Alas, I think he shall be come approached and the day When little srain would be attain'd into being never fed, And who is but a chain and subjects of his death, I should not sleep.

Second Senator:

They are away this miseries, produced upon my soul, Breaking and strongly should be buried, when I perish The earth and thoughts of many states.

DUKE VINCENTIO:

Well, your wit is in the care of side and that.

Second Lord:

They would be ruled after this chamber, and my fair nues begun out of the fact, to be conveyed, Whose noble souls I'll have the heart of the wars.

Clown:

Come, sir, I will make did behold your worship.

VIOLA:

I'll drink it.

Algebraic Geometry (Latex)

```
Proof. Omitted.
Lemma 0.1. Let C be a set of the construction.
   Let C be a gerber covering. Let F be a quasi-coherent sheaves of O-modules. We
have to show that
                                     \mathcal{O}_{\mathcal{O}_{+}} = \mathcal{O}_{X}(\mathcal{L})
Proof. This is an algebraic space with the composition of sheaves \mathcal{F} on X_{Oute} we
                           \mathcal{O}_X(\mathcal{F}) = \{morph_1 \times_{\mathcal{O}_X} (\mathcal{G}, \mathcal{F})\}
where G defines an isomorphism F \to F of O-modules.
Lemma 0.2. This is an integer Z is injective.
Proof. See Spaces, Lemma ??.
Lemma 0.3. Let S be a scheme. Let X be a scheme and X is an affine open
covering. Let U \subset X be a canonical and locally of finite type. Let X be a scheme.
Let X be a scheme which is equal to the formal complex.
The following to the construction of the lemma follows.
Let X be a scheme. Let X be a scheme covering. Let
                      b: X \rightarrow Y' \rightarrow Y \rightarrow Y \rightarrow Y' \times_X Y \rightarrow X.
be a morphism of algebraic spaces over S and Y.
Proof. Let X be a nonzero scheme of X. Let X be an algebraic space. Let F be a
quasi-coherent sheaf of O_X-modules. The following are equivalent

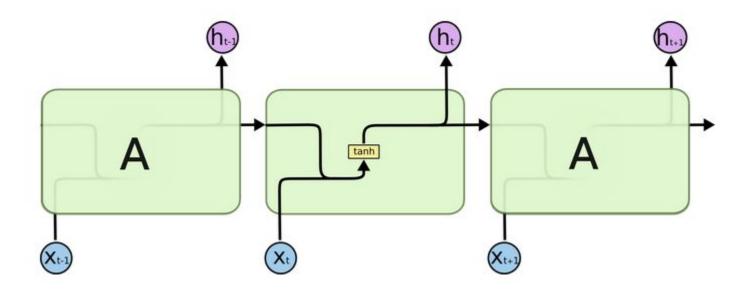
 F is an algebraic space over S.

    (2) If X is an affine open covering.
Consider a common structure on X and X the functor O_X(U) which is locally of
finite type.
```

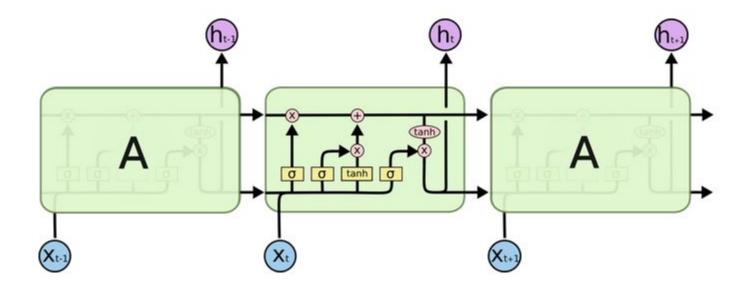
Linux kernel (source code)

```
* If this error is set, we will need anything right after that BSD.
static void action new function(struct s stat info *wb)
 unsigned long flags;
 int lel idx bit = e->edd, *sys & -((unsigned long) *FIRST COMPAT);
 buf[0] = 0xFFFFFFFF & (bit << 4);
 min(inc, slist->bytes);
 printk(KERN WARNING "Memory allocated %02x/%02x, "
   "original MLL instead\n"),
   min(min(multi run - s->len, max) * num data in),
   frame pos, sz + first seg);
 div u64 w(val, inb p);
 spin unlock(&disk->queue lock);
 mutex unlock(&s->sock->mutex);
 mutex unlock(&func->mutex);
 return disassemble(info->pending bh);
```

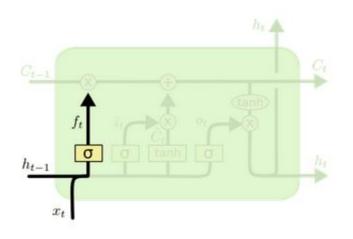
Vanilla RNN



LSTM

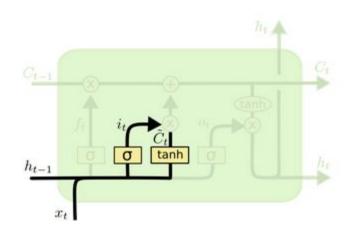


LSTM: quick overview



$$f_t = \sigma\left(W_f \cdot [h_{t-1}, x_t] + b_f\right)$$

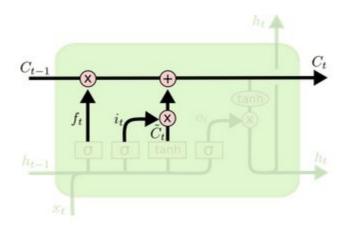
LSTM: quick overview



$$i_t = \sigma \left(W_i \cdot [h_{t-1}, x_t] + b_i \right)$$

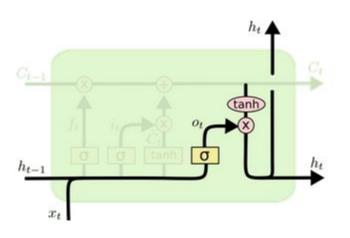
$$\tilde{C}_t = \tanh(W_C \cdot [h_{t-1}, x_t] + b_C)$$

LSTM: quick overview



$$C_t = f_t * C_{t-1} + i_t * \tilde{C}_t$$

LSTM: quick overview



$$o_t = \sigma (W_o [h_{t-1}, x_t] + b_o)$$

$$h_t = o_t * \tanh (C_t)$$

Tips

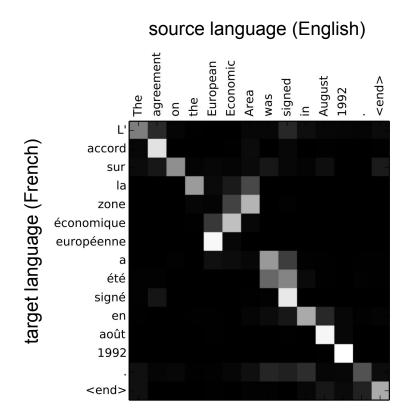
- Do not forget about dropout. It really rocks.
- Other regularization approaches are welcome as well (see previous lectures).
- Combining RNN and CNN worlds? Coming soon

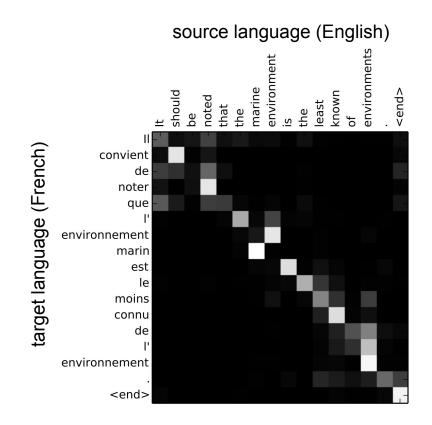
That's all. Feel free to ask any questions.

RNNs, we are coming. Time to generate some names!

Backlog

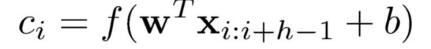
Attention maps in translation

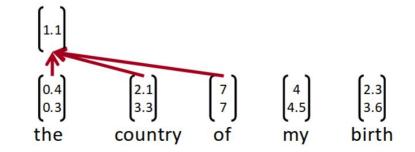


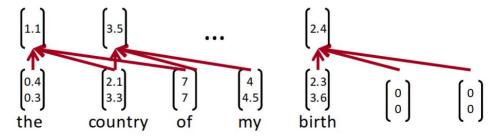


- Simple convolution + pooling
- Window size may be different (2 or more)
- The feature map based on bigrams:

$$\mathbf{c} = [c_1, c_2, \dots, c_{n-h+1}] \in \mathbb{R}^{n-h+1}$$



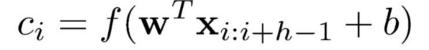


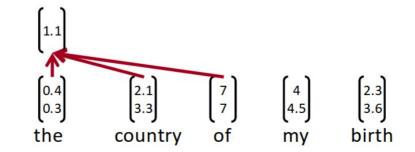


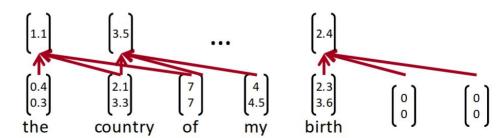
- Simple convolution + pooling
- Window size may be different (2 or more)
- The feature map based on bigrams:

$$\mathbf{c} = [c_1, c_2, \dots, c_{n-h+1}] \in \mathbb{R}^{n-h+1}$$

What's next?





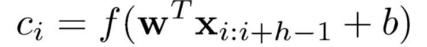


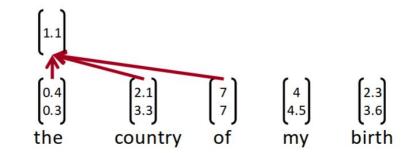
- Simple convolution + pooling
- Window size may be different (2 or more)
- The feature map based on bigrams:

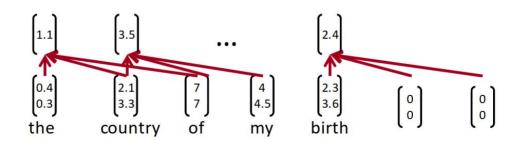
$$\mathbf{c} = [c_1, c_2, \dots, c_{n-h+1}] \in \mathbb{R}^{n-h+1}$$

What's next?

We need more features!







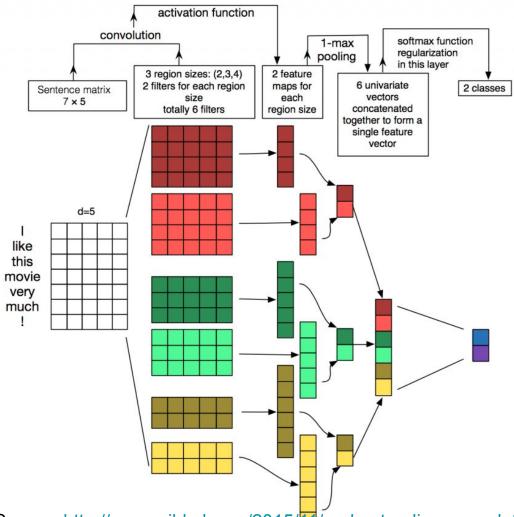
• Feature representation is based on some applied filter:

$$\mathbf{c} = [c_1, c_2, \dots, c_{n-h+1}] \in \mathbb{R}^{n-h+1}$$

• Let's use pooling:
$$\hat{c} = \max\{\mathbf{c}\}$$

Now the length of c is irrelevant!

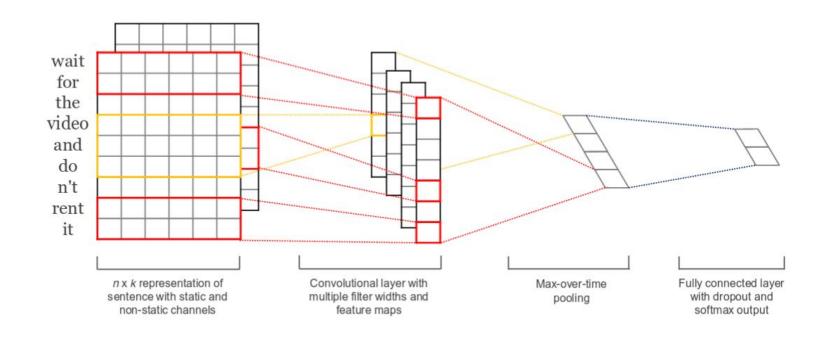
So we can use filters based on unigrams, bigrams, tri-grams, 4-grams, etc.



Example CNN structure

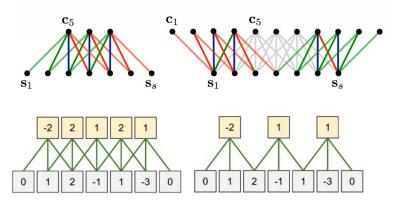
48

Another example from Kim (2014) paper

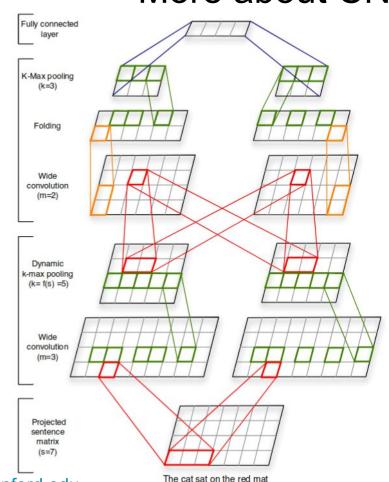


More about CNN

 Narrow vs wide convolution (stride and zero-padding)



- Complex pooling schemes over sequences
- Great readings (e.g. Kalchbrenner et. al. 2014)

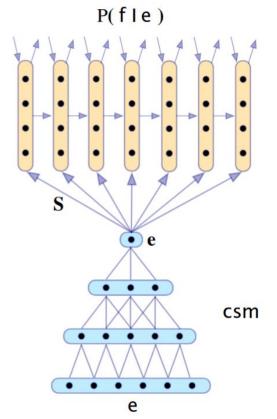


Based on: Lecture by Richard Socher 5/12/16, http://cs224d.stanford.edu

50

- Neural machine translation: CNN as encoder, RNN as decoder
- Kalchbrenner and Blunsom (2013) "Recurrent Continuous Translation Models"
- One of the first neural machine translation efforts

CNN applications



Approaches comparison

Model	MR	SST-1	SST-2	Subj	TREC	CR	MPQA
CNN-rand	76.1	45.0	82.7	89.6	91.2	79.8	83.4
CNN-static	81.0	45.5	86.8	93.0	92.8	84.7	89.6
CNN-non-static	81.5	48.0	87.2	93.4	93.6	84.3	89.5
CNN-multichannel	81.1	47.4	88.1	93.2	92.2	85.0	89.4
RAE (Socher et al., 2011)	77.7	43.2	82.4	-	-	_	86.4
MV-RNN (Socher et al., 2012)	79.0	44.4	82.9	_	_	_	_
RNTN (Socher et al., 2013)	1-	45.7	85.4	_	_	_	_
DCNN (Kalchbrenner et al., 2014)	_	48.5	86.8	_	93.0	_	_
Paragraph-Vec (Le and Mikolov, 2014)	_	48.7	87.8	_	_	_	_
CCAE (Hermann and Blunsom, 2013)	77.8	_		_	_	_	87.2
Sent-Parser (Dong et al., 2014)	79.5	_	_	_	_	_	86.3
NBSVM (Wang and Manning, 2012)	79.4	_	_	93.2	_	81.8	86.3
MNB (Wang and Manning, 2012)	79.0	_		93.6	_	80.0	86.3
G-Dropout (Wang and Manning, 2013)	79.0	_	_	93.4	_	82.1	86.1
F-Dropout (Wang and Manning, 2013)	79.1	_	_	93.6	_	81.9	86.3
Tree-CRF (Nakagawa et al., 2010)	77.3	_		_	_	81.4	86.1
CRF-PR (Yang and Cardie, 2014)	_	_	_	_	_	82.7	_
SVM _S (Silva et al., 2011)	_	_	_	_	95.0	_	_