Mockingbird Semantics

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1 MbFilter

1.1 Grammar

```
\langle breakMethodG \rangle = 1G \mid 2\_GP
 \langle \mathit{breakMethodS} \rangle = \mathtt{1S} \mid \mathtt{16^2\_SP}
\langle \mathit{breakMethodH} \rangle = \mathtt{1H}
                   \langle expr \rangle = \langle expr \rangle > \langle expr \rangle
                                  | fix x.\langle expr\rangle
                                  | x
                                  | test\{\langle expr \rangle\}
                                  hit
                                  Shade
                                  | if |g| > n then \langle expr \rangle else \langle expr \rangle
                                  | BreakMapReduceG(\langle breakMethodG \rangle){\langle expr \rangle}
                                  | BreakMapReduceS(\langle breakMethodS \rangle){\langle expr \rangle}
                                     BreakMapReduceH(\langle breakMethodH \rangle) \{\langle expr \rangle\}
         \langle sizeRange \rangle = [a, b] \mid [a+]
       \langle entity\text{-}type \rangle = \texttt{GeoSamp} \langle sizeRange \rangle \langle sizeRange \rangle
                                  | Hit \(\sizeRange\)
                                  \mid Frag \langle sizeRange \rangle
\langle transition\text{-}type \rangle = \langle entity\text{-}type \rangle \rightarrow \langle entity\text{-}type \rangle
     \langle breakG\text{-}type \rangle = (\langle sizeRange \rangle \rightarrow \langle sizeRange \rangle) \langle sizeRange \rangle
      \langle breakS\text{-}type \rangle = \langle sizeRange \rangle \ (\langle sizeRange \rangle \rightarrow \langle sizeRange \rangle)
     \langle breakH\text{-}type \rangle = \langle sizeRange \rangle \rightarrow \langle sizeRange \rangle
```

1.2 Typing Rules

$$\frac{\Gamma \vdash e_1 : \tau \to \tau' \quad \Gamma \vdash e_2 : \tau' \to \tau''}{\Gamma \vdash e_1 >> e_2 : \tau \to \tau''}$$
(1)

$$\frac{\Gamma, x : \phi \vdash e : \phi}{\Gamma \vdash \text{fix } x.e : \phi} \tag{2}$$

$$\frac{\Gamma(x) = \phi}{\Gamma \vdash \mathsf{call}\ x : \phi} \tag{3}$$

$$\frac{\Gamma \vdash e : \texttt{GeoSamp} \ \sigma_1 \ \sigma_2 \to \tau}{\Gamma \vdash \texttt{test}\{e\} : \texttt{GeoSamp} \ \sigma_1 \ \sigma_2 \to \tau} \tag{4}$$

$$\frac{\cdot}{\Gamma \vdash \mathtt{hit} : \mathtt{GeoSamp} \; [1,1] \; [1,1] \to \mathtt{Hit}} \tag{5}$$

$$\frac{\cdot}{\Gamma \vdash \mathtt{Shade} : \mathtt{Hit} \ [1,1] \to \mathtt{Frag}} \tag{6}$$

$$\frac{\Gamma \vdash e_1 : \texttt{GeoSamp} \ [n+1,b] \ \sigma \to \tau \quad \Gamma \vdash e_2 : \texttt{GeoSamp} \ [a,n] \ \sigma \to \tau}{\Gamma \vdash \texttt{if} \ |g| > n \ \texttt{then} \ e_1 \ \texttt{else} \ e_2 : \texttt{GeoSamp} \ [a,b] \ \sigma \to \tau} \tag{7}$$

$$\frac{\Gamma \vdash e : \texttt{GeoSamp} \to \tau}{\Gamma \vdash \texttt{BreakMapReduceG}(d)\{e\} : \texttt{GeoSamp} \to \tau} \tag{8}$$

$$\frac{\Gamma \vdash e : \texttt{GeoSamp} \to \tau}{\Gamma \vdash \texttt{BreakMapReduceS}(d)\{e\} : \texttt{GeoSamp} \to \tau} \tag{9}$$

$$\frac{\Gamma \vdash e : \mathtt{Hit} \to \tau}{\Gamma \vdash \mathtt{BreakMapReduceH}(d)\{e\} : \mathtt{Hit} \to \tau} \tag{10}$$

1.3 Denotational Semantics

$$|e_1 >> e_2| = |e_2| \circ |e_1| \tag{11}$$

$$|\operatorname{fix} x.e| = |[(\operatorname{fix} x.e)/(\operatorname{call} x)]e| \tag{12}$$

$$|\mathtt{test}\{e\}| = |e| \circ id_{\breve{0}} \cup (\lambda(G, S).\{(\bot, k) : (s, k) \in S\}) \circ id_{\breve{0}} \tag{13}$$

$$|\mathtt{hit}| = \{((\{g\}, \{(s,k)\}), \{(h,k)\}) : h = isect(g,s)\}$$
 (14)

$$|Shade| = \{(\{(s,k)\}, \{(f,k)\}) : f = shade(h)\}$$
(15)

$$|\text{if } |g| > n \text{ then } e_1 \text{ else } e_2| = |e_1| \circ id_{|g| > n} \cup |e_2| \circ id_{|g| \le n}$$
 (16)

$$|\mathtt{BreakMapReduceG}(d)\{e\}| = \begin{cases} \{((G,S), \{F_1 \oplus \cdots \oplus F_n\}) \\ : ((G_i,S),f_i) \in |e|, (G,\{G_1,...,G_n\}) \in |d| \} \end{cases} \tag{17}$$

$$|\texttt{BreakMapReduceS}(d)\{e\}| = \begin{cases} \{((G,S), \{F_1 \cup \dots \cup F_n\}) \\ : ((G,S_i), f_i) \in |e|, (S, \{S_1, ..., S_n\}) \in |d| \} \end{cases} \tag{18}$$

$$|\mathtt{BreakMapReduceH}(d)\{e\}| = \begin{cases} (H, \{F_1 \cup \dots \cup F_n\}) \\ : (H_i, F_i) \in |e|, (H, \{H_1, \dots, H_n\}) \in |d| \} \end{cases} \tag{19}$$

1.4 Theorems

- 1. For any expression e, if $((G, S), F) \in |e|$, then keys(S) = keys(F).
- 2. For any expression $e: \texttt{GeoSamp} \to \texttt{Frag}, \text{ if } ((G,S),F) \in |e|, \text{ then for all } (s,k) \in S, (\bigoplus_{g \in G} shade(isect(g,s)), k) \in F.$

2 MbOrder

2.1 Grammar

```
e = e >> e
      | fix x. e
      \mid x
      hit
      \mid test\{e\}
      \mid filt\{e\}
      \mid ifsizeG(>n) e else e
      \mid buildG\{d_G\}
      \mid buildS\{d_S\}
      \mid \ \mathtt{unboundG}\{e\}
       \mid unboundS\{e\}
      | \ \mathtt{mmrG}\{e\}
       | \mathtt{mmrS}\{e\}
d_{lpha}=\mathtt{flat}
     |p_{\alpha}>=>d_{\alpha}
      | fix x. d_{\alpha}
      \mid x
      | bound(d_{\alpha})
      | ifsize\alpha(>n) d_{\alpha} else d_{\alpha}
p_G = 1 \mathrm{G} \mid 2 \mathrm{\_GP}
p_S = 1S | 16^2_SP
  \tau = \sigma \to \sigma
  \sigma = {\tt GeoSamp} \ \delta \ \delta
      \mid Hit \delta
  \delta = *
      \mid [\delta]
      \mid \#\delta
      \delta + \delta
      | fix x. \delta
      \mid x
```

2.2 Typing Rules

$$\frac{\Gamma \vdash e_1 : \sigma \to \sigma' \quad \Gamma \vdash e_2 : \sigma' \to \sigma''}{\Gamma \vdash e_1 >> e_2 : \sigma \to \sigma''}$$
(20)

$$\frac{\Gamma, x: \tau \vdash e: \tau}{\Gamma \vdash \text{fix } x.e: \tau} \tag{21}$$

$$\frac{\Gamma(x) = \tau}{\Gamma \vdash x : \tau} \tag{22}$$

$$\frac{\Gamma \vdash e : \texttt{GeoSamp} \ \# \delta_1 \ \# \delta_2 \to \sigma}{\Gamma \vdash \texttt{test}\{e\} : \texttt{GeoSamp} \ \# \delta_1 \ \# \delta_2 \to \sigma} \tag{23}$$

$$\frac{\cdot}{\Gamma \vdash \mathtt{hit} : \mathtt{GeoSamp} \ * \ * \to \mathtt{Hit} *} \tag{24}$$

$$\frac{\Gamma \vdash e : \texttt{GeoSamp} \ \delta_1 \ \delta_2 \to \texttt{Hit*}}{\Gamma \vdash \texttt{mmrG}(d)\{e\} : \texttt{GeoSamp} \ [\delta_1] \ \delta_2 \to \texttt{Hit*}} \tag{25}$$