C++ Software Engineering

for engineers of other disciplines

Module 10 "C++ Parallelism" 3rd Lecture: Multithreading



CPU & Threads



- CPU sockets are divided into cores and cores are divided into threads.
- Numbers of CPUs is "socket * cores * threads".
- To get the best performance, it is ideal to launch a thread per CPU even better if threads are not siblings.

mrz@vu:~\$ cat /sys/devices/system/cpu/cpu0/topology/thread_siblings_list
0

 Threads which are in the same cores are called siblings and in most of the cases have inferior performance compared to threads on separate cores.

```
Architecture:
                   x86 64
CPU op-mode(s): 32-bit, 64-bit
Byte Order:
            Little Endian
CPU(s):
On-line CPU(s) list: 0-3
Thread(s) per core: 1
Core(s) per socket:
Socket(s):
NUMA node(s):
Vendor ID:
                    GenuineIntel
CPU family:
Model:
                    158
Model name:
                    Intel(R) Core(TM) i7-8850H CPU @ 2.60GHz
```

std::thread



- A class representing a single thread of execution.
- Begin execution upon construction.
- Threads can communicate via:
 - std::promise
 - std::condition variable
- Synchronizations on shared resources can happen using:
 - std::mutex
 - std::atomic
- Threads should either be joined or detached

Observers	
joinable	checks whether the thread is joinable, i.e. potentially running in parallel context (public member function)
get_id	returns the <i>id</i> of the thread (public member function)
native_handle	returns the underlying implementation-defined thread handle (public member function)
hardware_concurrency [static]	returns the number of concurrent threads supported by the implementation (public static member function)
Operations	
join	waits for a thread to finish its execution (public member function)
detach	permits the thread to execute independently from the thread handle (public member function)
swap	swaps two thread objects (public member function)

https://en.cppreference.com/w/cpp/thread/thread

Undetached threads which are not joined cause segfault!

```
void foo()
    // simulate expensive operation
    std::this thread::sleep for(std::chrono::seconds(1));
void bar()
    // simulate expensive operation
    std::this thread::sleep for(std::chrono::seconds(1));
int main()
    std::cout << "starting first helper...\n";</pre>
    std::thread helper1(foo);
    std::cout << "starting second helper...\n";</pre>
    std::thread helper2(bar);
    std::cout << "waiting for helpers to finish..." << std::endl;</pre>
    helper1.join();
    helper2.join();
    std::cout << "done!\n";</pre>
```

https://en.cppreference.com/w/cpp/thread/thread/join

```
void independentThread()
    std::cout << "Starting concurrent thread.\n";</pre>
    std::this thread::sleep for(std::chrono::seconds(2));
    std::cout << "Exiting concurrent thread.\n";</pre>
void threadCaller()
    std::cout << "Starting thread caller.\n";</pre>
    std::thread t(independentThread);
    t.detach();
    std::this thread::sleep for(std::chrono::seconds(1));
    std::cout << "Exiting thread caller.\n";</pre>
int main()
    threadCaller();
    std::this thread::sleep for(std::chrono::seconds(5));
```

https://en.cppreference.com/w/cpp/thread/thread/detach

Mutual Exclusion



- STL provides mutual exclusion:
 - lock in std::mutex
 - Different means for acquiring a lock:
 - std::lock gaurd
 - std::shared lock
 - std::unique lock
 - std::scoped_lock
 - Different strategies for locking.

Locking

lock	locks the mutex, blocks if the mutex is not available (public member function)
try_lock	tries to lock the mutex, returns if the mutex is not available (public member function)
unlock	unlocks the mutex (public member function)

https://en.cppreference.com/w/cpp/thread/mutex

Туре	Effect(s)
defer_lock_t	do not acquire ownership of the mutex
try_to_lock_t	try to acquire ownership of the mutex without blocking
adopt_lock_t	assume the calling thread already has ownership of the mutex

https://en.cppreference.com/w/cpp/thread/lock_tag

- Reentrant mutex to avoid dead lock by std::recursive_mutex
- std::lock avoid locks any given number of *lockables* using an algorithm which never deadlocks yet:

RAII wrapper for this function, and is generally preferred to a naked call to std::lock.

std::lock gaurd



- RAII-style mechanism for owning a mutex for duration of scoped block, to avoid hassle of lock and release.
- std::lock_gaurd is blocking and non-copyable.
- std::unique lock is a general mutex wrapper.

```
void transfer(Box &from, Box &to, int num)
{
    // don't actually take the locks yet
    std::unique_lock<std::mutex> lock1(from.m, std::defer_lock);
    std::unique_lock<std::mutex> lock2(to.m, std::defer_lock);

    // lock both unique_locks without deadlock
    std::lock(lock1, lock2);

    from.num_things -= num;
    to.num_things += num;

    // 'from.m' and 'to.m' mutexes unlocked in 'unique_lock' dtors
}
```

```
int g_i = 0;
std::mutex g_i_mutex; // protects g_i

void safe_increment()
{
    const std::lock_guard<std::mutex> lock(g_i_mutex);
    ++g_i;

    std::cout << std::this_thread::get_id() << ": " << g_i << '\n';

    // g_i_mutex is automatically released when lock
    // goes out of scope
}</pre>
```

```
Defined in header <mutex>
template < class Mutex >
class unique_lock;

Defined in header <shared_mutex>

template < class Mutex >
class lock_guard;

Defined in header <shared_mutex>
template < class Mutex >
class shared_lock;

Defined in header <mutex>
```

(since C++17)

template< class... MutexTypes >

class scoped lock;

std::condition variable

- Another mean of inter-thread communication.
- One thread needs to acquire a lock, change a value, and notify others.

```
std::mutex m;
std::condition variable cv;
std::string data;
bool ready = false;
bool processed = false;
int main()
    std::thread worker(worker thread);
    data = "Example data";
    // send data to the worker thread
        std::lock guard<std::mutex> lk(m);
        ready = true;
        std::cout << "main() signals data ready for processing\n";</pre>
    cv.notify one();
    // wait for the worker
        std::unique lock<std::mutex> lk(m);
        cv.wait(lk, []{return processed;});
    std::cout << "Back in main(), data = " << data << '\n';</pre>
    worker.join();
```

Notification notifie

notify_one	notifies one waiting thread (public member function)
notify_all	notifies all waiting threads (public member function)
Waiting	
wait	blocks the current thread until the condition variable is woken up (public member function)
wait_for	blocks the current thread until the condition variable is woken up or after the specified timeout duration (public member function)
wait_until	blocks the current thread until the condition variable is woken up or until specified time point has been reached (public member function)
Native handle	
native_handle	returns the native handle

```
void worker_thread()
{
    // Wait until main() sends data
    std::unique_lock<std::mutex> lk(m);
    cv.wait(lk, []{return ready;});

    // after the wait, we own the lock.
    std::cout << "Worker thread is processing data\n";
    data += " after processing";

    // Send data back to main()
    processed = true;
    std::cout << "Worker thread signals data processing completed\n";

    // Manual unlocking is done before notifying, to avoid waking up
    // the waiting thread only to block again (see notify_one for details)
    lk.unlock();
    cv.notify_one();
}</pre>
```

Output

main() signals data ready for processing
Worker thread is processing data
Worker thread signals data processing completed
Back in main(), data = Example data after processing

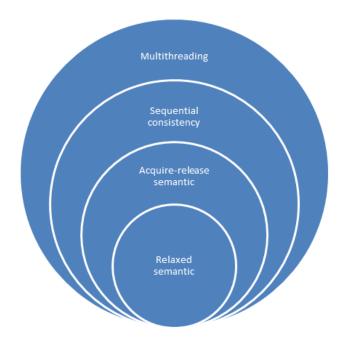
Memory Modeling Contracts

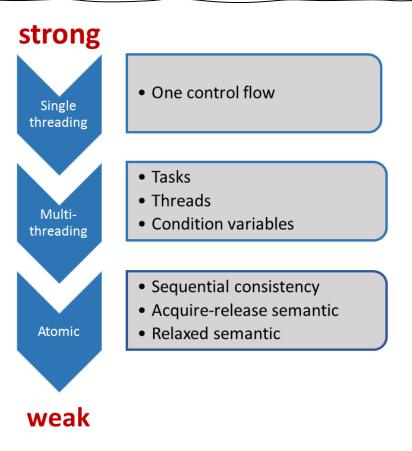


• The weaker the memory model, the higher optimization potential!

In computing, a memory model describes the interactions of threads through memory and their shared use of the data. A memory model allows a compiler to perform many important optimizations. https://en.wikipedia.org/wiki/Memory model (programming)

Expert levels





Atomic

• An area in memory where there is no race! If one thread writes to an atomic object while another thread reads

from it, the behavior is well-defined.

```
int cnt = 0;
auto f = [&]{cnt++;};
std::thread t1{f}, t2{f}, t3{f}; // undefined behavior

std::atomic<int> cnt{0};
auto f = [&]{cnt++;};
std::thread t1{f}, t2{f}, t3{f}; // OK
```

https://en.cppreference.com/w/cpp/language/memory model

 Memory order specifies how memory (including regular nonatomic accesses) are to be ordered around an atomic operation.

```
inline constexpr memory_order memory_order_relaxed = memory_order::relaxed;
inline constexpr memory_order memory_order_consume = memory_order::consume;
inline constexpr memory_order memory_order_acquire = memory_order::acquire;
inline constexpr memory_order memory_order_release = memory_order::release;
inline constexpr memory_order memory_order_acq_rel = memory_order::acq_rel;
inline constexpr memory_order memory_order_seq_cst = memory_order::seq_cst;
```

https://en.cppreference.com/w/cpp/atomic/memory_order



• User datatypes could be used as templated parameter for std::atomic but it is not guaranteed that the operations would be lock-free.The only guaranteed lock-free data type is std::atomic_flag.

is_lock_free	checks if the atomic object is lock-free (public member function)
store	atomically replaces the value of the atomic object with a non-atomic argument (public member function)
load	atomically obtains the value of the atomic object (public member function)
operator T	loads a value from an atomic object (public member function)
exchange	atomically replaces the value of the atomic object and obtains the value held previously (public member function)
compare_exchange_weak compare_exchange_strong	atomically compares the value of the atomic object with non-atomic argument and performs atomic exchange if equal or atomic load if not (public member function)
wait (C++20)	blocks the thread until notified and the atomic value changes (public member function)
notify_one(C++20)	notifies at least one thread waiting on the atomic object (public member function)
notify_all(C++20)	notifies all threads blocked waiting on the atomic object (public member function)
	https://en.cppreference.com/w/cpp/atomic/atomic

https://en.cppreference.com/w/cpp/atomic/atomic flag test and set





CPU affinity



Where to have mutex?



```
//std::lock_guard<std::mutex> guard(g_pages_mutex);
size_t i = std::rand()%5;
std::cout << "Sleeping for " << i << " second in thread "
std::this_thread::sleep_for(std::chrono::seconds(i));
std::string result = "fake content";

std::lock_guard<std::mutex> guard(g_pages_mutex);
g_pages[url] = result;
```

Reenterant Mutex



```
class X {
    std::recursive mutex m;
    std::string shared;
  public:
    void fun1() {
      std::lock guard<std::recursive mutex> lk(m);
      shared = "fun1";
      std::cout << "in fun1, shared variable is now " << shared << '\n';</pre>
    void fun2() {
      std::lock guard<std::recursive mutex> lk(m);
      shared = "fun2";
      std::cout << "in fun2, shared variable is now " << shared << '\n';</pre>
      fun1(); // recursive lock becomes useful here
      std::cout << "back in fun2, shared variable is " << shared << '\n';</pre>
```

Scoped Lock



```
use std::scoped lock to acquire two locks without worrying about
// other calls to assign lunch partner deadlocking us
  and it also provides a convenient RAII-style mechanism
std::scoped lock lock(e1.m, e2.m);
  Equivalent code 1 (using std::lock and std::lock guard)
// std::lock(e1.m, e2.m);
// std::lock guard<std::mutex> lk1(e1.m, std::adopt lock);
// std::lock guard<std::mutex> lk2(e2.m, std::adopt lock);
  Equivalent code 2 (if unique locks are needed, e.g. for condition variables)
// std::unique lock<std::mutex> lk1(e1.m, std::defer lock);
// std::unique lock<std::mutex> lk2(e2.m, std::defer lock);
 ' std::lock(lk1, lk2);
```

Threads are not FAIR!



alice and bob are waiting for locks
christina and alice are waiting for locks
dave and bob are waiting for locks
christina and bob are waiting for locks
alice and bob got locks
dave and bob got locks
christina and bob got locks
christina and alice got locks
Employee alice has lunch partners: bob christina
Employee bob has lunch partners: alice dave christina
Employee christina has lunch partners: bob alice
Employee dave has lunch partners: bob

alice and bob are waiting for locks
dave and bob are waiting for locks
christina and bob are waiting for locks
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dave and bob got locks
christina and bob got locks
christina and alice are waiting for locks
christina and alice got locks
Employee alice has lunch partners: bob christina
Employee bob has lunch partners: alice dave christina
Employee christina has lunch partners: bob alice
Employee dave has lunch partners: bob

```
std::unique_lock<std::mutex> lk(cv_m);
std::cerr << "Waiting... \n";
cv.wait(lk, []{return i == 1;});
std::cerr << "...finished waiting. i == 1\n";</pre>
```

```
std::this_thread::sleep_for(std::chrono::seconds(1));
    std::lock guard<std::mutex> lk(cv m);
    std::cerr << "Notifying...\n";</pre>
cv.notify all();
std::this thread::sleep for(std::chrono::seconds(1));
    std::lock guard<std::mutex> lk(cv m);
    i = 1:
    std::cerr << "Notifying again...\n";</pre>
//cv.notify all();
```