

QCVN Group Meeting

Haar measure & Unitary t-design

- Towards URB for qutrit

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① Unitary randomized benchmarking

- Generate RB sequences

- Measure survival probability & fit results

- Intuition behind URB

- Mathematical motivation

② Haar measure

- The motivation of measure theory

- The Haar measure

- Haar measure in $U(2)$

- Haar measure in $U(N)$

③ Unitary t-design

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- Motivation: spherical t-design

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Complete protocol of URB

A standard unitary randomized benchmarking consists of

- 1 Generate RB sequences

$$\mathcal{S}_{i_m} = \bigodot_{j=1}^{m+1} (\Lambda_{i_j, j} \odot C_{i_j}) \quad (1)$$

- 2 For each sequence, calculate survival probability

$$\text{Tr}[E_\psi \mathcal{S}_{i_m}(\rho_\psi)] \quad (2)$$

- 3 Average over random realizations to find the *average* fidelity

$$F_{\text{seq}}(m, \psi) = \text{Tr}(E_\psi \mathcal{S}_m(\rho_\psi)) \quad (3)$$

- 4 Fit the results to the model of exponential decay.

Generate RB sequences

Definition

A sequence of $m + 1$ quantum operations with the first m operations chosen *uniformly* at random from some group $\mathcal{G} \in U(d)$ and the final operation ($m + 1$) chosen so that the net sequence is the identity operation.

- We primarily focus on $C_3 \in U(3^n)$, because they can be realized efficiently on both quantum & classical hardware.
- For each length m , we choose K_m RB sequences. Each sequence contains m random element C_{i_j} sampled uniformly from C_3 .
- The $C_{i_{m+1}}$ element is defined as $(C_{i_1} \cdots C_{i_m})^{-1}$.

Measure survival probability

The survival probability is defined by

$$\text{Tr}[E_\psi \mathcal{S}_{i_m}(\rho_\psi)] \quad (4)$$

where ρ_ψ is the initial state (SPAM absorbed) and E_ψ is the POVM element (having off-diagonal non-zero entries). If noise-free,

$$\rho_\psi = E_\psi = |\psi\rangle\langle\psi|$$

In practice, survival probabilities are probabilities that the qutrit go back to the ground state $|0\rangle$. For example, in the noise-free situation,

$$\text{Tr}[E_\psi \mathcal{S}_{i_m}(\rho_\psi)] = p(0) \quad (5)$$

Measure survival probability

Theorem

The survival probability of the qutrit is the probability we obtain the initial state (or ground state if we prepare $|0\rangle$).

Proof. Suppose we prepare a statistical ensemble ρ , taking into account the SPAM error. Then the evolution of such system is governed by

$$\rho_\psi \rightarrow \sum_i p_i U |\psi_i\rangle \langle \psi_i| U^\dagger \quad (6)$$

The probability of getting result m is

$$p(m) = \sum_i p(m|i)p_i \quad (7)$$

$$= \sum_i \langle \psi_i | U^\dagger E_{|\psi\rangle} U | \psi_i \rangle p_i \quad (8)$$

$$= \text{Tr} \left[E_{|\psi\rangle} \sum_i p_i U | \psi_i \rangle \langle \psi_i | U^\dagger \right] \quad (9)$$

$$= \text{Tr} [E_{|\psi\rangle} \mathcal{S}_{i_m}(\rho_\psi)] \quad (10)$$

If we intended to prepare an initial ground state, then

$$\text{Tr} [E_{|0\rangle} \mathcal{S}_{i_m}(\rho_{|0\rangle})] = p(0) \quad \square$$

Average sequence fidelity and fit results

Average over K_m random realizations of the sequence to find the average sequence fidelity

$$F_{seq}(m, |\psi\rangle) = \text{Tr}[E_\psi S_{K_m}(\rho_\psi)] \quad (11)$$

$$= \text{Tr}\left[E_\psi \frac{1}{K_m} \sum_{i_m} S_{i_m}(\rho_\psi)\right] \quad (12)$$

Repeat for different values of m and fit the results for the averaged sequence fidelity to the model

$$F^{(0)}(m, |\psi\rangle) = A_0 \alpha^m + B_0 \quad (13)$$

where A_0 and B_0 absorb SPAM error. α is called the *Error per Clifford (EPC)*, relating to the average error-rate.

Results

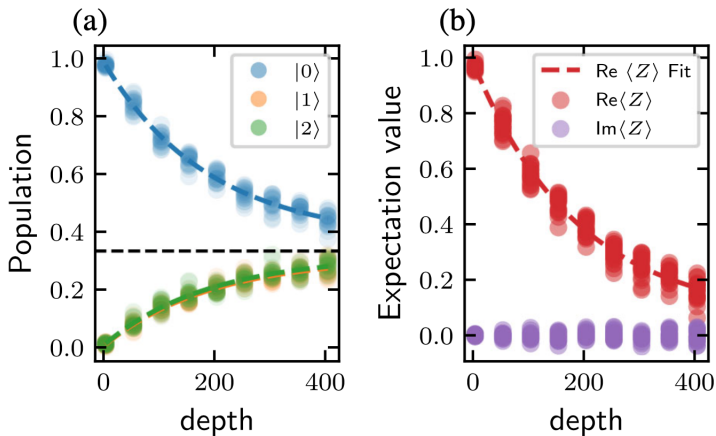


Figure: Qutrit randomized benchmarking [Morvan *et al.* Phys. Rev. Lett.126.210504 (2021)].

Intuition behind URB

- By sampling uniformly noisy gates, we create a *depolarizing channel*, characterized by the probability α of the qutrit not being turned into a maximally mixed state $I/3$.
- After a sequence of m gates, where the error per gate rate is α , then the resulting density matrix is

$$\rho_f^m = \alpha^m \rho_i + (1 - \alpha^m)/3 \cdot I \quad (14)$$

- Suppose we start with $|0\rangle$, and the entire sequence is equivalent to the identity operator. The probability of successfully measuring $|0\rangle$ is

$$p(0) = \begin{cases} \alpha^m [\rho_i]_{00} + 1 - \alpha/3 \\ \alpha^m + 1 - \alpha^m/3 = 2/3\alpha^m + 1/3 \end{cases} \quad (15)$$

The mathematical theory behind

- The exponential decay model is not a result of repeating gate in a sequential manner, but *uniformly randomized gates from the Clifford group*. These Clifford gates are noisy, thus $C_{ij} \odot \Lambda$, where Λ constitutes a depolarizing channel.
- Formally speaking, taking an average over a finite group C_3 of a quantum channel Λ constitutes a twirl.
- Twirling over $U(3)$ yields exactly the same result as the Clifford group, because the Clifford group $C_3 \in U(3)$ is a *unitary 2-design* of the unitary group.
- Buzzing words: *uniformly-randomized, unitary t -design, twirling over a finite group*.

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The motivation of measure theory

The theory of random matrices involves extensively the generalised concept of *measure*.

- In $\mathbb{R}^1, \mathbb{R}^2, \mathbb{R}^3$, think about length, area, and volume.
- What about higher dimensions?
- What about other spaces?
- What about vector spaces, e.g $\mathcal{M}_{3 \times 3}$.

Definition (Loosely speaking)

Measure μ tells us how mathematical objects (points, matrices, functions, etc.) are distributed in a mathematical set or space—some place is immensely dense, some is not.

Example 1: Triple integral of volume V

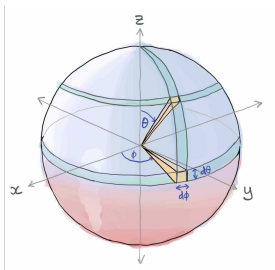


Figure: The measure μ takes into account that these infinitesimal volume is not uniformly distributed over \mathbb{R}^3 .

The volume of the sphere is *properly measured* when including the measure factor μ

$$V = \int_0^r \int_0^{2\pi} \int_0^\pi \mu d\rho d\phi d\theta, \quad (16)$$

where $\mu = \rho^2 \sin \theta$ weights portions of the sphere differently depending on where they are in the space.

The Haar measure

- Similar to points in spherical coordinates uniquely defined by (ρ, ϕ, θ) , unitary matrices are uniquely defined by three parameters, e.g (θ, ϕ, λ) for every element in $U(2)$.
- For every dimension N , the unitary matrices of size $N \times N$ constitute the unitary group $U(N)$. Operations on $U(N)$ requires proper measure μ , or the *Haar measure*.
- For an N -dimensional system, the Haar measure tells us how to weight the elements of $U(N)$. As an example, suppose f is a function acts on $V \in U(N)$, the integral over the group is

$$\int_{V \in U(N)} f(V) d\mu_N(V) \quad (17)$$

where the analytical expression of μ_N is desired.

Haar measure in $U(2)$

Problem (Sampling)

Given the structure of the group $U(2)$, sample elements of the unitary group $U(2)$ in a properly uniform manner.

- The mathematical space we're dealing with is $U(2)$.
- The operation is sampling.
- An alternative angle of looking: sampling quantum *states* uniformly at random.

$$\rho_{|0\rangle} \rightarrow U_{\mu_N} |\psi(\theta, \varphi)\rangle \langle \psi(\theta, \varphi)| U_{\mu_N}^\dagger \quad (18)$$

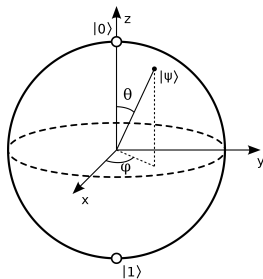


Figure:
 $SU(2) \subseteq U(2) \cong SO(3)$.

Visualization of state sampling

- Note that the density matrix of a single-qubit state can be expanded in the form

$$\rho = \frac{1}{2}(\hat{I} + \mathbf{r} \cdot \hat{\sigma}), \quad (19)$$

where $\mathbf{r}(r_x, r_y, r_z)$ is the Bloch vector, uniquely defines a mixed state. The Bloch vector component is $r_\alpha = \text{Tr}(\sigma \cdot \rho) \in \mathbb{R}^3$.

- The explicit matrix representation of any arbitrary single-qubit operator $U(2)$ is

$$U(\theta, \phi, \lambda) = \begin{pmatrix} \cos(\theta/2) & -e^{i\lambda} \sin(\theta/2) \\ e^{i\phi} \sin(\theta/2) & e^{i(\lambda+\phi)} \cos(\theta/2) \end{pmatrix}, \quad (20)$$

hence each element of $U(2)$ is effectively defined by θ, ϕ, λ .

Visualization of state sampling

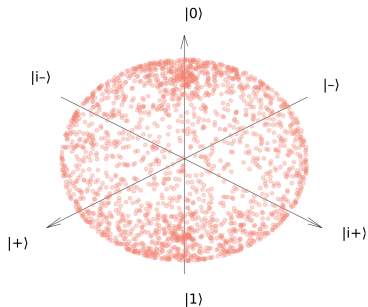


Figure: $d\mu_2 = d\theta d\phi d\lambda$. Not Haar-random.

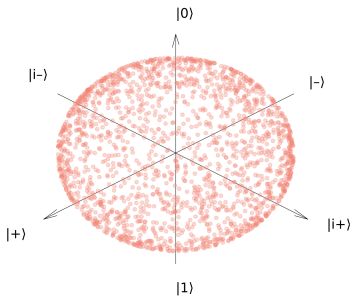
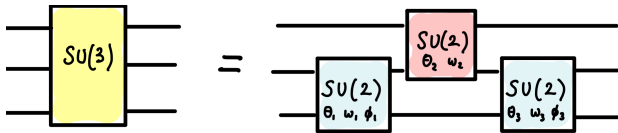


Figure: $d\mu_2 = \sin \theta d\theta d\phi d\lambda$. Haar-random.

Haar measure in $U(N)$

- It's hard to visualize when $N = d^n$. We thus remain focus on 3^1 .
- From the study of photonics, we know that we can decompose any $SU(N)$ operation recursively by sandwiching an $SU(2)$ between two $SU(N-1)$ [H. de Guise *et al.*, Phys. Rev. A 97 022328 (2018)].
- For $N = 3^1$, the exact Haar measure $d\mu_3$ follows accordingly



$$d\mu_3 = \sin \theta_1 \sin \theta_2 \sin^2(\theta_2/2) \sin \theta_3 \prod_{i=1}^3 d\theta_i d\phi_i d\lambda_i \quad (21)$$

Final note on Haar measure

- Unlike $SU(2)$, sampling uniformly over $SU(3)$ requires us to sample θ_i from *different distribution*.
- This suggests the higher dimension of the unitary group, the more parameters we need to care about. Generally speaking an N -dimensional unitary requires at least $N^2 - 1$ parameters.
- Other methods of sampling uniformly are thus desired, e.g. *Haar-random matrices from QR decomposition* [F. Mezzaddri, arXiv:math-ph/0609050v2 (2007)].
- Haar measure is invariant under unitary transformation, i.e.

$$\begin{aligned}\int_{V \in U(N)} f(WV) d\mu_N(V) &= \int_{V \in U(N)} f(VW) d\mu_N(V) \\ &= \int_{V \in U(N)} f(V) d\mu_N(V)\end{aligned}$$

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The problem of efficient randomization

Sampling from N -dimensional unitary group requires us to keep track of $N^2 - 1$ parameters, or $d^{2n} - 1$.

	1	2	3	...	n
Qubit ($d = 2$)	2	15	63	...	$2^{2n} - 1$
Qutrit ($d = 3$)	8	80	728	...	$3^{2n} - 1$

Table: The number of parameters we need to specify a $SU(d^n)$ operator. The exponential scaling hinders us from scalability.

Problem (Efficient sampling)

How to scalably and efficiently sample random unitaries in a uniform manner?

Motivation: spherical t -design

Problem (Averaging P over a sphere)

What is the average A of a d variables polynomial over the surface of a real d -dimensional unit sphere $S(R^d)$?

$$A = \int_{S(R^d)} P_t(u) d\mu(u) \quad (22)$$

- This integration—in principle—can be numerically calculated, provided we have the proper measure $d\mu(u)$.
- What if one could alternatively approximate A by uniformly sampling a sufficiently large number of points on the sphere?
- What if one could alternatively **calculate exactly** A by uniformly sampling a sufficiently large number of points on the sphere?

Motivation: spherical t -design

Definition (Spherical t -design)

Let $P_t : \mathcal{S}(R^d) \rightarrow \mathbb{R}$ be a polynomial in d variables with all terms homogeneous in degree at most t . A set $X = \{x | x \in \mathcal{S}(R^d)\}$ is a spherical t -design if

$$\frac{1}{|X|} \sum_{x \in X} P_t(x) = \int_{\mathcal{S}(R^d)} P_t(u) d\mu(u) \quad (23)$$

holds for all possible P_t , where $d\mu$ is the uniform, normalized spherical measure. A spherical t -design is also a k -design for all $k < t$.

- ① One natural question is how do we find X , provided f ?
- ② The other question concerns with the structure of the set (group) X/\mathcal{G} itself.

Unitary t -design

Instead of averaging polynomials P_t over spheres, we consider polynomials that are functions of the entries of $U \in U(N)$.

Definition (Unitary t -design)

Let $P_{t,t}(U)$ be a polynomial with homogeneous degree at most t in d variables in the entries of a unitary matrix U , and degree t in the complex conjugates of those entries. A unitary t -design is a set of K unitaries $\{U_k\}$ such that

$$\frac{1}{K} \sum_{k=1}^K P_{t,t}(U_k) = \int_{U(d)} P_{t,t}(U) d\mu(U) \quad (24)$$

holds for all possible $P_{t,t}$ and where $d\mu$ is the uniform Haar measure.

Unitary t -design in action: Haar measure

Let us revisit our efficient randomization problem. The ultimate goal of randomization is to calculate the fidelity. More precisely, the *average fidelity* of a quantum channel can be calculated by twirling over the Haar measure of the respective Lie group.

Definition (Twirling \mathcal{E} over a Haar-measure)

Suppose a quantum channel \mathcal{E} . The average fidelity with respect to the full set of Haar-random states $U_{\mu N} \rho U_{\mu N}^\dagger$ is *twirling channel* \mathcal{E} ,

$$\bar{F}_{\mathcal{E}} = \int_{U(d)} d\mu(U) U^\dagger \mathcal{E}(U \rho U^\dagger) U \quad (25)$$

Unitary t -design in action: Fidelity

Using the result of [J. Emerson *et al.*, arXiv:quant-ph/0606161v2 (2012)], we conclude that twirling any quantum channel over a unitary t -design is equivalent to twirling over a Haar-measure unitary group.

$$\frac{1}{K} \sum_{j=1}^K U_j^\dagger \mathcal{E}(U_j \rho) U_j = \int_{U(d)} d\mu(U) U^\dagger \mathcal{E}(U \rho U^\dagger) U \quad (26)$$

- Note that the inner product involves two U and two U^\dagger , thus implies this unitary t -design is a *unitary 2-design*.
- What is the unitary 2-design then?

Twirling over Clifford group

Theorem (\mathcal{C}_3^n)

The multi-qutrit Clifford group \mathcal{C}_3^n forms a unitary 2-design.

A result from [M. Nielsen, arXiv:quant-ph/0205035v2 (2002)] prove that twirling a quantum channel \mathcal{E} over a Haar measure yields a depolarizing channel.

Corollary

Twirling over the multi-qutrit Clifford group yields a depolarizing channel.

From this, we informally confirm the aforementioned results of URB,

$$\mathcal{E}_{dep}(\rho_i) = p\rho_i + (1-p)\frac{I}{3^n} \quad (27)$$

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- 1 The metrics we concern is the average fidelity of a quantum channel \mathcal{E} . This average fidelity can be calculated by twirling that channel over Haar-random unitary group [J. Emerson *et al.*, arXiv.quant-ph/0503243].
- 2 This is better than QPT, but still not scalable and efficient. We can however get around by using the unitary t -design.
- 3 Twirling over a unitary t -design is equivalent to twirling over Haar-random unitary group, and yields a depolarizing channel.
- 4 The depolarizing channel has an exponential decay with respect to gate length m .

Next week: The Clifford group \mathcal{C}_3 on single-qutrit.