

Program Representation for General Intelligence

The Reduct Toolkit for Program Normalization

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1 Introduction

- Program Space
- What are Programs in Normal Forms?
- Effects of Reducing Programs in Normal Forms

2 The Reduct Toolkit

- Recall of Combo
- How it works
- Demo...
- What Remains to Be Implemented

Outline

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Introduction

- Program learning is a very useful skill
 - But program space is complex
 - Understanding and exploiting its properties
- ⇒ Reducing programs in normal form

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Program learning is **HARD!**

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Program space is:

- Vast \Rightarrow grows exponentially with program size



Program learning is **HARD!**

Program space is:

- Vast \Rightarrow **grows exponentially** with program size



- Chaotic \Rightarrow small syntactic variations can lead to **huge semantic variations**



Program learning is **HARD!**

Program space is:

- Vast \Rightarrow **grows exponentially** with program size



- Chaotic \Rightarrow small syntactic variations can lead to **huge semantic variations**



- Over-represented \Rightarrow **many programs with same semantics**

$$x \wedge y \equiv x \wedge y \wedge x \equiv \neg(\neg(x) \vee \neg(y)) \equiv \dots$$

less diversity packed into the program space

Over-representation

Different syntax, same semantics

x	y	$x \wedge y$
0	0	0
0	1	0
1	0	0
1	1	1

Over-representation

Different syntax, same semantics

x	y	$x \wedge y$	$\neg(\neg y \vee \neg x)$
0	0	0	0
0	1	0	0
1	0	0	0
1	1	1	1

Over-representation

Different syntax, same semantics

x	y	$x \wedge y$	$\neg(\neg y \vee \neg x)$	$x \wedge x \wedge y$
0	0	0	0	0
0	1	0	0	0
1	0	0	0	0
1	1	1	1	1

Over-representation

Different syntax, same semantics

x	y	$x \wedge y$	$\neg(\neg y \vee \neg x)$	$x \wedge x \wedge y$	$x \wedge y \wedge (\neg z \vee z)$
0	0	0	0	0	0
0	1	0	0	0	0
1	0	0	0	0	0
1	1	1	1	1	1

Over-representation

Different syntax, same semantics

x	y	$x \wedge y$	$\neg(\neg y \vee \neg x)$	$x \wedge x \wedge y$	$x \wedge y \wedge (\neg z \vee z)$...
0	0	0	0	0	0	...
0	1	0	0	0	0	...
1	0	0	0	0	0	...
1	1	1	1	1	1	...

Syntactic vs Semantic De-correlation

$$f_1 = x \wedge (y \vee z)$$

x	y	z	f_1
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	1

$$f_2 = \neg(x \wedge (y \vee z))$$

x	y	z	f_2
0	0	0	1
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	0

Syntactic vs Semantic De-correlation

$$f_1 = x \wedge (y \vee z)$$

x	y	z	f_1
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	1

$$f_2 = \neg(x \wedge (y \vee z))$$

x	y	z	f_2
0	0	0	1
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	0

- Syntactic distance between f_1 and f_2 :

Syntactic vs Semantic De-correlation

$$f_1 = x \wedge (y \vee z)$$

x	y	z	f_1
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	1

$$f_2 = \neg(x \wedge (y \vee z))$$

x	y	z	f_2
0	0	0	1
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	0

- Syntactic distance between f_1 and f_2 : 1

Syntactic vs Semantic De-correlation

$$f_1 = x \wedge (y \vee z)$$

x	y	z	f_1
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	1

$$f_2 = \neg(x \wedge (y \vee z))$$

x	y	z	f_2
0	0	0	1
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	0

- Syntactic distance between f_1 and f_2 : 1
- Semantic distance between f_1 and f_2 :

Syntactic vs Semantic De-correlation

$$f_1 = x \wedge (y \vee z)$$

x	y	z	f_1
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	1

$$f_2 = \neg(x \wedge (y \vee z))$$

x	y	z	f_2
0	0	0	1
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	0

- Syntactic distance between f_1 and f_2 : 1
- Semantic distance between f_1 and f_2 : 1

Syntactic vs Semantic De-correlation

$$f_1 = x \wedge (y \vee z)$$

x	y	z	f_1
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	1

$$f_2 = \neg(x \wedge (y \vee z))$$

x	y	z	f_2
0	0	0	1
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	0

- Syntactic distance between f_1 and f_2 : 1
- Semantic distance between f_1 and f_2 : 2

Syntactic vs Semantic De-correlation

$$f_1 = x \wedge (y \vee z)$$

x	y	z	f_1
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	1

$$f_2 = \neg(x \wedge (y \vee z))$$

x	y	z	f_2
0	0	0	1
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	0

- Syntactic distance between f_1 and f_2 : 1
- Semantic distance between f_1 and f_2 : 3

Syntactic vs Semantic De-correlation

$$f_1 = x \wedge (y \vee z)$$

x	y	z	f_1
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	1

$$f_2 = \neg(x \wedge (y \vee z))$$

x	y	z	f_2
0	0	0	1
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	0

- Syntactic distance between f_1 and f_2 : 1
- Semantic distance between f_1 and f_2 : 4

Syntactic vs Semantic De-correlation

$$f_1 = x \wedge (y \vee z)$$

x	y	z	f_1
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	1

$$f_2 = \neg(x \wedge (y \vee z))$$

x	y	z	f_2
0	0	0	1
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	0
1	0	1	0
1	1	0	0
1	1	1	0

- Syntactic distance between f_1 and f_2 : 1
- Semantic distance between f_1 and f_2 : 5

Syntactic vs Semantic De-correlation

$$f_1 = x \wedge (y \vee z)$$

x	y	z	f_1
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	1

$$f_2 = \neg(x \wedge (y \vee z))$$

x	y	z	f_2
0	0	0	1
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	0

- Syntactic distance between f_1 and f_2 : 1
- Semantic distance between f_1 and f_2 : 6

Syntactic vs Semantic De-correlation

$$f_1 = x \wedge (y \vee z)$$

x	y	z	f_1
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	1

$$f_2 = \neg(x \wedge (y \vee z))$$

x	y	z	f_2
0	0	0	1
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	0

- Syntactic distance between f_1 and f_2 : 1
- Semantic distance between f_1 and f_2 : 7

Syntactic vs Semantic De-correlation

$$f_1 = x \wedge (y \vee z)$$

x	y	z	f_1
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	1

$$f_2 = \neg(x \wedge (y \vee z))$$

x	y	z	f_2
0	0	0	1
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	0

- Syntactic distance between f_1 and f_2 : 1
- Semantic distance between f_1 and f_2 : 8

Syntactic vs Semantic De-correlation

$$g_1 = \text{false}$$

x	y	z	g_1
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	0
1	1	0	0
1	1	1	0

$$g_2 = x \wedge y \wedge z$$

x	y	z	g_2
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	0
1	1	0	0
1	1	1	1

Syntactic vs Semantic De-correlation

$$g_1 = \text{false}$$

x	y	z	g_1
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	0
1	1	0	0
1	1	1	0

$$g_2 = x \wedge y \wedge z$$

x	y	z	g_2
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	0
1	1	0	0
1	1	1	1

- Syntactic distance:

Syntactic vs Semantic De-correlation

$$g_1 = \text{false}$$

x	y	z	g_1
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	0
1	1	0	0
1	1	1	0

$$g_2 = x \wedge y \wedge z$$

x	y	z	g_2
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	0
1	1	0	0
1	1	1	1

- Syntactic distance: 1

Syntactic vs Semantic De-correlation

$$g_1 = \text{false}$$

x	y	z	g_1
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	0
1	1	0	0
1	1	1	0

$$g_2 = x \wedge y \wedge z$$

x	y	z	g_2
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	0
1	1	0	0
1	1	1	1

- Syntactic distance: 2

Syntactic vs Semantic De-correlation

$$g_1 = \text{false}$$

x	y	z	g_1
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	0
1	1	0	0
1	1	1	0

$$g_2 = x \wedge y \wedge z$$

x	y	z	g_2
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	0
1	1	0	0
1	1	1	1

- Syntactic distance: 3

Syntactic vs Semantic De-correlation

$$g_1 = \text{false}$$

x	y	z	g_1
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	0
1	1	0	0
1	1	1	0

$$g_2 = x \wedge y \wedge z$$

x	y	z	g_2
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	0
1	1	0	0
1	1	1	1

- Syntactic distance: 4

Syntactic vs Semantic De-correlation

$$g_1 = \text{false}$$

x	y	z	g_1
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	0
1	1	0	0
1	1	1	0

$$g_2 = x \wedge y \wedge z$$

x	y	z	g_2
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	0
1	1	0	0
1	1	1	1

- Syntactic distance: 5

Syntactic vs Semantic De-correlation

$$g_1 = \text{false}$$

x	y	z	g_1
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	0
1	1	0	0
1	1	1	0

$$g_2 = x \wedge y \wedge z$$

x	y	z	g_2
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	0
1	1	0	0
1	1	1	1

- Syntactic distance: 6

Syntactic vs Semantic De-correlation

$$g_1 = \text{false}$$

x	y	z	g_1
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	0
1	1	0	0
1	1	1	0

$$g_2 = x \wedge y \wedge z$$

x	y	z	g_2
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	0
1	1	0	0
1	1	1	1

- Syntactic distance: 6
- Semantic distance:

Syntactic vs Semantic De-correlation

$$g_1 = \text{false}$$

x	y	z	g_1
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	0
1	1	0	0
1	1	1	0

$$g_2 = x \wedge y \wedge z$$

x	y	z	g_2
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	0
1	1	0	0
1	1	1	1

- Syntactic distance: 6
- Semantic distance: 1

Syntactic vs Semantic De-correlation

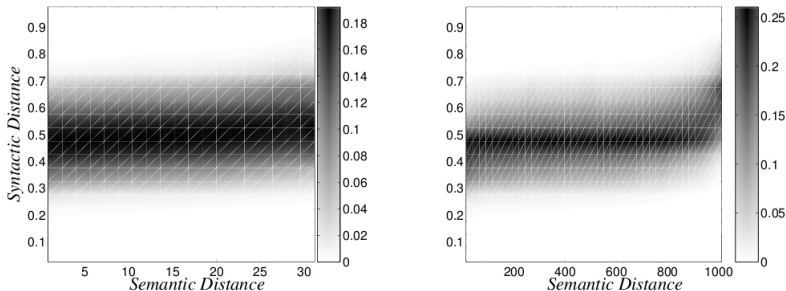


Figure: Syntactic vs semantic distance for random formulae with arities five (left) and ten (right). *Extracted from Moshe Looks PhD thesis.*

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Programs in Normal Forms

- 1 unique representation (possibly the shortest)

$$x \wedge y \equiv x \wedge y \wedge x \equiv \neg(\neg(x) \vee \neg(y)) \equiv \dots$$

- 2 interesting properties \Rightarrow program space more regular



Outline

1 Introduction

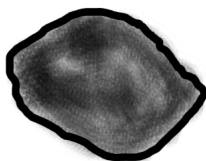
- Program Space
- What are Programs in Normal Forms?
- **Effects of Reducing Programs in Normal Forms**

2 The Reduct Toolkit

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Avoiding over-representation

Before reduction:



Program space

$$x \wedge y \equiv x \wedge y \wedge x \equiv \neg(\neg(x) \vee \neg(y)) \equiv \dots$$

$$x \vee y \equiv x \vee y \vee x \equiv \neg(\neg(x) \wedge \neg(y)) \equiv \dots$$

$$x \wedge y \wedge z \equiv x \wedge y \wedge z \wedge x \wedge y \wedge z \equiv \dots$$

$$x \wedge y \wedge \neg(z) \equiv x \wedge y \wedge \neg(z) \wedge x \wedge y \wedge \neg(z) \equiv \dots$$

...

Avoiding over-representation

After reduction:



Program space

$$\cancel{x \wedge y} \equiv \cancel{x \wedge y \wedge x} \equiv \cancel{\neg(\neg(x) \vee \neg(y))} \equiv \dots$$

$$\cancel{x \vee y} \equiv \cancel{x \vee y \vee x} \equiv \cancel{\neg(\neg(x) \wedge \neg(y))} \equiv \dots$$

$$\cancel{x \wedge y \wedge z} \equiv \cancel{x \wedge y \wedge z \wedge x \wedge y \wedge z} \equiv \dots$$

$$\cancel{x \wedge y \wedge \neg(z)} \equiv \cancel{x \wedge y \wedge \neg(z) \wedge x \wedge y \wedge \neg(z)} \equiv \dots$$

...

Increase Syntactic vs Semantic Correlation

Before reduction

$$f_1 = x \wedge (y \vee z)$$

x	y	z	f_1
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	1

$$f_2 = \neg(x \wedge (y \vee z))$$

x	y	z	f_2
0	0	0	1
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	0

- Syntactic distance between f_1 and f_2 : 1
- Semantic distance between f_1 and f_2 : 8

Increase Syntactic vs Semantic Correlation

After reduction

$$f_1 = x \wedge (y \vee z)$$

x	y	z	f_1
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	1

$$f_2 = \neg x \vee (\neg y \wedge \neg z)$$

x	y	z	f_2
0	0	0	1
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	0

- Syntactic distance between f_1 and f_2 : ?
- Semantic distance between f_1 and f_2 : 8

Increase Syntactic vs Semantic Correlation

After reduction

$$f_1 = x \wedge (y \vee z)$$

x	y	z	f ₁
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	1

$$f_2 = \neg x \vee (\neg y \wedge \neg z)$$

x	y	z	f ₂
0	0	0	1
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	0

- Syntactic distance between f_1 and f_2 : 1
- Semantic distance between f_1 and f_2 : 8

Increase Syntactic vs Semantic Correlation

After reduction

$$f_1 = x \wedge (y \vee z)$$

x	y	z	f ₁
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	1

$$f_2 = \neg x \vee (\neg y \wedge \neg z)$$

x	y	z	f ₂
0	0	0	1
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	0

- Syntactic distance between f_1 and f_2 : 2
- Semantic distance between f_1 and f_2 : 8

Increase Syntactic vs Semantic Correlation

After reduction

$$f_1 = x \wedge (y \vee z)$$

x	y	z	f ₁
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	1

$$f_2 = \neg x \vee (\neg y \wedge \neg z)$$

x	y	z	f ₂
0	0	0	1
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	0

- Syntactic distance between f_1 and f_2 : 3
- Semantic distance between f_1 and f_2 : 8

Increase Syntactic vs Semantic Correlation

After reduction

$$f_1 = x \wedge (y \vee z)$$

x	y	z	f ₁
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	1

$$f_2 = \neg x \vee (\neg y \wedge \neg z)$$

x	y	z	f ₂
0	0	0	1
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	0

- Syntactic distance between f_1 and f_2 : 4
- Semantic distance between f_1 and f_2 : 8

Increase Syntactic vs Semantic Correlation

After reduction

$$f_1 = x \wedge (y \vee z)$$

x	y	z	f ₁
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	1

$$f_2 = \neg x \vee (\neg y \wedge \neg z)$$

x	y	z	f ₂
0	0	0	1
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	0

- Syntactic distance between f_1 and f_2 : 5
- Semantic distance between f_1 and f_2 : 8

Increase Syntactic vs Semantic Correlation

Before and after reduction

$$g_1 = \text{false}$$

x	y	z	g_1
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	0
1	1	0	0
1	1	1	0

$$g_2 = x \wedge y \wedge z$$

x	y	z	g_2
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	0
1	1	0	0
1	1	1	1

- Syntactic distance: 6
- Semantic distance: 1

Increase Syntactic vs Semantic Correlation

Before reduction

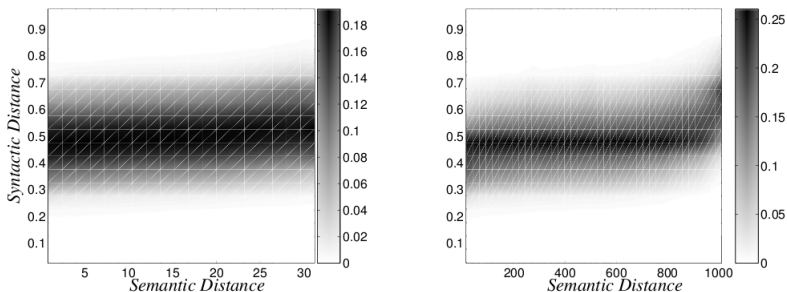


Figure: Syntactic vs semantic distance for random formulae with arities five (left) and ten (right). *Extracted from Moshe Looks PhD thesis.*

Increase Syntactic vs Semantic Correlation

After reduction

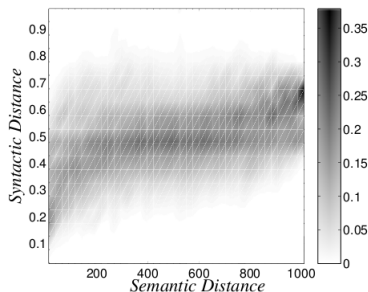
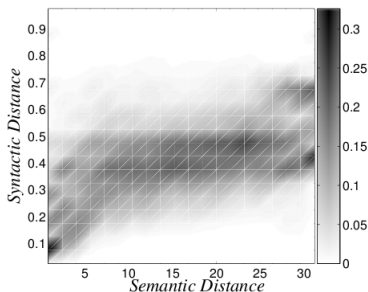
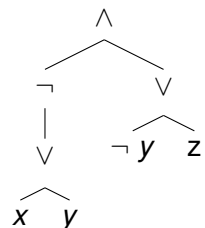



Figure: Syntactic vs semantic distance for random formulae in normal form with arities five (left) and ten (right). *Extracted from Moshe Looks PhD thesis.*

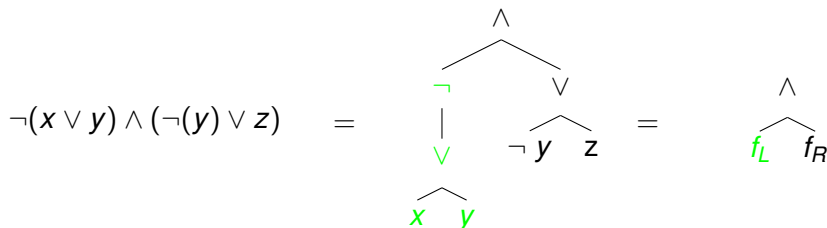
Improve Structure, interpretable, hierarchy

Before reduction:

$$\neg(x \vee y) \wedge (\neg(y) \vee z) =$$

$$=$$


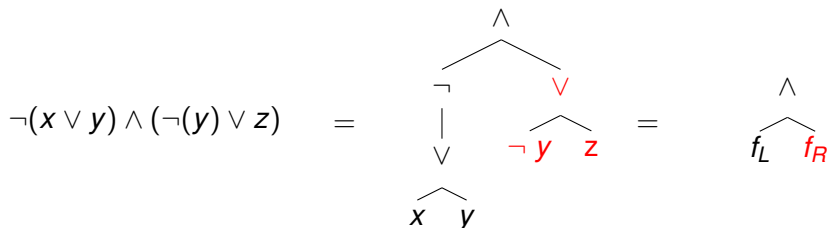
Improve Structure, interpretable, hierarchy

Before reduction:



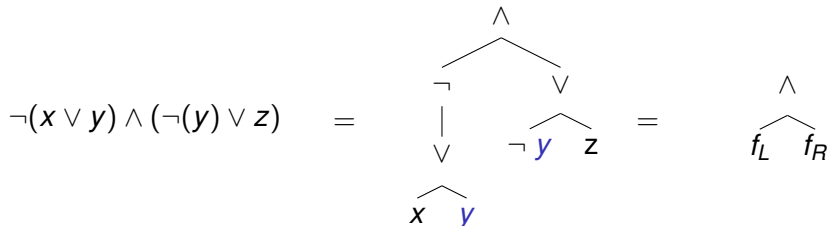
Improve Structure, interpretable, hierarchy

Before reduction:



Improve Structure, interpretable, hierarchy

Before reduction: f_L and f_R are not independent.



Improve Structure, interpretable, hierarchy

After reduction: g_L and g_R are **independent**.

$$\neg x \wedge \neg y \quad = \quad \begin{array}{c} \wedge \\ \swarrow \quad \searrow \\ \neg x \quad \neg y \end{array} \quad = \quad \begin{array}{c} \wedge \\ \swarrow \quad \searrow \\ g_L \quad g_R \end{array}$$

Improve Structure, interpretable, hierarchy

After reduction: g_L and g_R are **independent**.

$$\neg x \wedge \neg y \quad = \quad \begin{array}{c} \wedge \\ \swarrow \quad \searrow \\ \neg x \quad \neg y \end{array} \quad = \quad \begin{array}{c} \wedge \\ \swarrow \quad \searrow \\ g_L \quad g_R \end{array}$$

Benefice

- Easier to understand
- Better hierarchical structure

To sum up

Reduction allows us to:

- 1 Avoid **over-representation**
- 2 Increase **syntactic vs semantic correlation**
- 3 **Improve structure**
- 4 Smaller programs take **less memory** and usually run faster

Conclusion

Reduction is good!

But beware!

Reduction to normal form is:

- undecidable for recursive functions



But beware!

Reduction to normal form is:

- undecidable for recursive functions



- undecidable for primitive recursive functions



But beware!

Reduction to normal form is:

- undecidable for recursive functions



- undecidable for primitive recursive functions



- NP-hard for Boolean functions



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In practice

- **incomplete** reduction methods

But beware!

Reduction to normal form is:

- undecidable for recursive functions



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- NP-hard for Boolean functions



In practice

- **incomplete** reduction methods
- rely on **heuristics**

But beware!

Reduction to normal form is:

- undecidable for recursive functions



- undecidable for primitive recursive functions



- NP-hard for Boolean functions



In practice

- **incomplete** reduction methods
- rely on **heuristics**
- trade-off between **efficiency** and completeness

Outline

1 Introduction

- Program Space
- What are Programs in Normal Forms?
- Effects of Reducing Programs in Normal Forms

2 The Reduct Toolkit

- Recall of Combo
- How it works
- Demo...
- What Remains to Be Implemented

The Reduct Toolkit

- Part of the OpenCog framework

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- Coded in C++

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- Handles the following domains
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 - actions perceptions to control agent in virtual world

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Reduct Operates on Combo

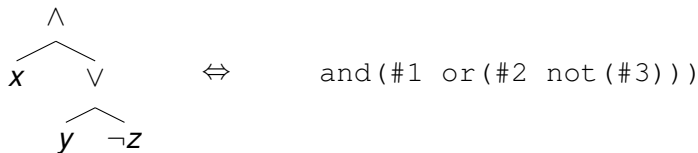
Combo is:

- Programmatic language of MOSES

Reduct Operates on Combo

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- Programmatic language of MOSES
- Syntactic tree in prefix notation



Reduct Operates on Combo

Combo is:

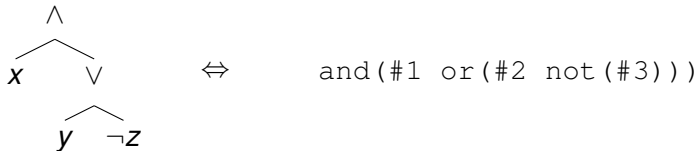
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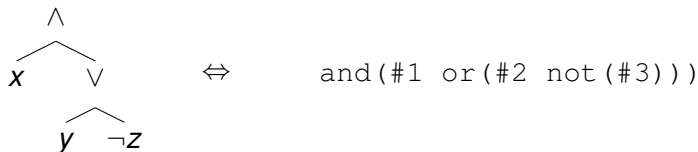


- Boolean: and, or, not, true, false

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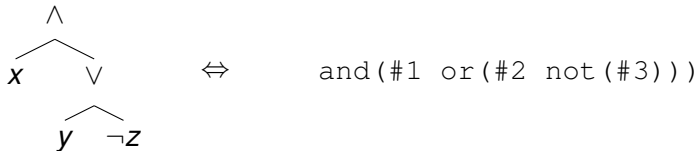


- **Boolean:** `and`, `or`, `not`, `true`, `false`
- **Continuous:** `+`, `*`, `/`, `exp`, `log`, `sin`

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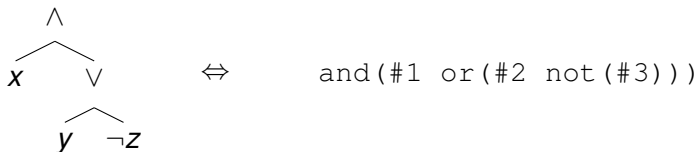


- **Boolean:** `and`, `or`, `not`, `true`, `false`
- **Continuous:** `+`, `*`, `/`, `exp`, `log`, `sin`
- **Numerico-Boolean:** `0<`, `impulse`, `contin_if`

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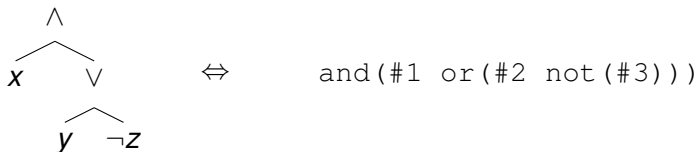


- **Boolean:** and, or, not, true, false
- **Continuous:** +, *, /, exp, log, sin
- **Numerico-Boolean:** 0<, impulse, contin_if
- **Action:** and_seq, action_if, action_while, ...

Reduct Operates on Combo

Combo is:

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- Boolean: and, or, not, true, false
- Continuous: +, *, /, exp, log, sin
- Numerico-Boolean: 0<, impulse, contin_if
- Action: and_seq, action_if, action_while, ...
- Higher order: *procedure_name*(arity) := *body*

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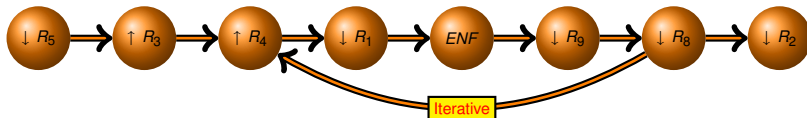
2 The Reduct Toolkit

- Recall of Combo
- **How it works**
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Reduction

How it works

Apply series of reduction rules according to a given strategy.



Reduct Rules

Example of reduction rules in the Boolean domain:

- $R_4: \text{and}(X) \rightarrow X$

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- $R_8: \text{or}(X \text{ true}) \rightarrow \text{true}$

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- $R_{13}: X/Z + Y/Z \rightarrow (X+Y)/Z$

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In the action domain:

- $R_{43}: \text{action_if}(B \ X \ Y) \rightarrow X$ if B is reducible to `true`

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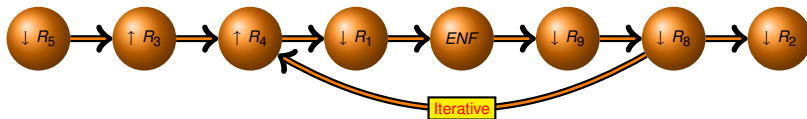
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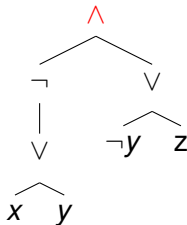
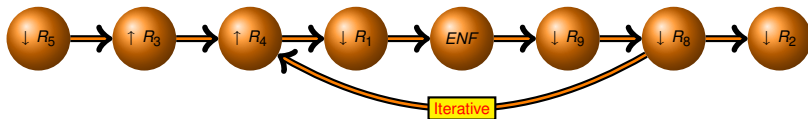
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about 85 rules in total...

Combining the reduction rules

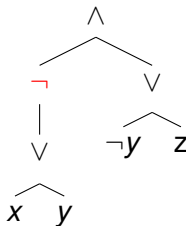
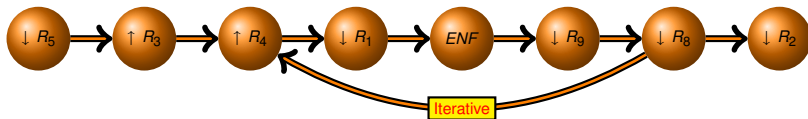


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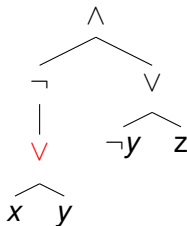
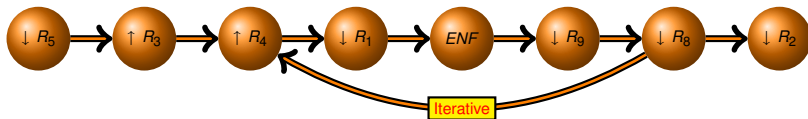
Apply R_i from the root down to the leaves.

Combining the reduction rules



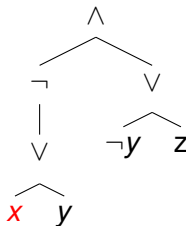
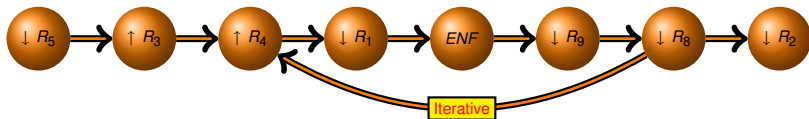
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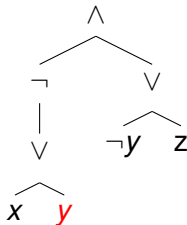
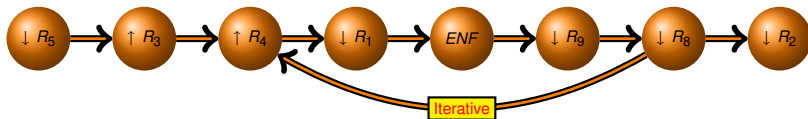
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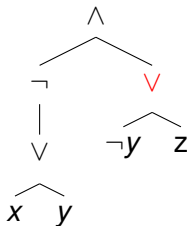
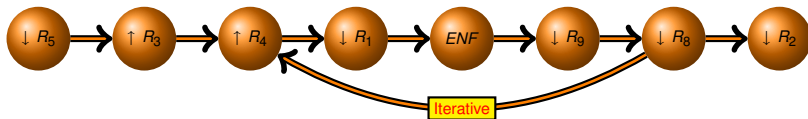
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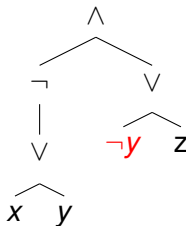
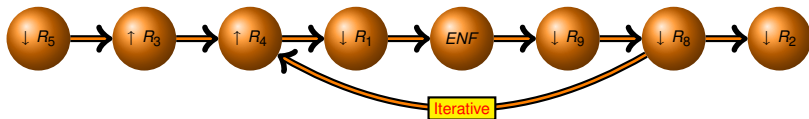
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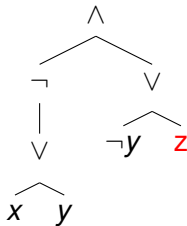
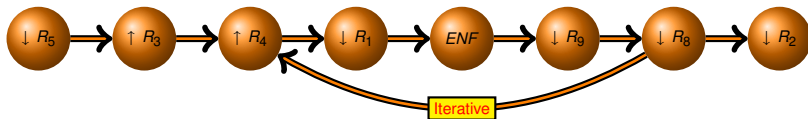
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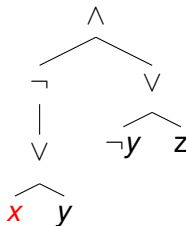
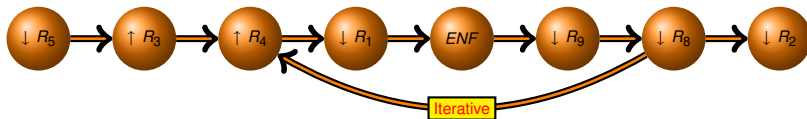
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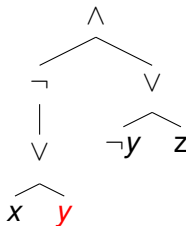
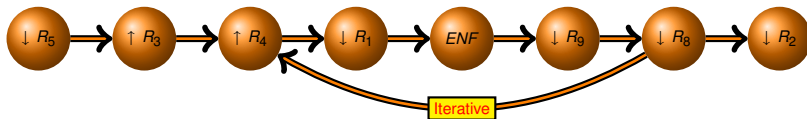
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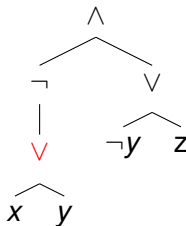
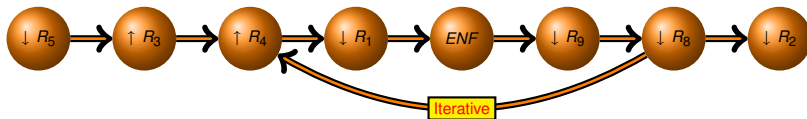
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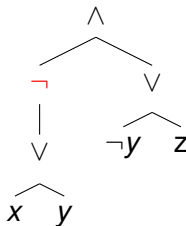
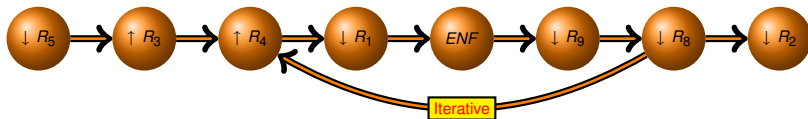
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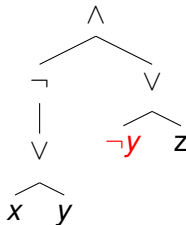
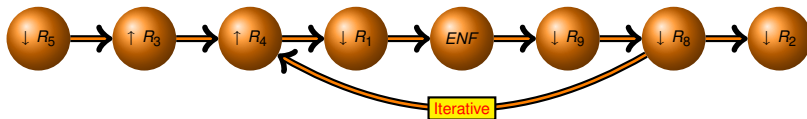
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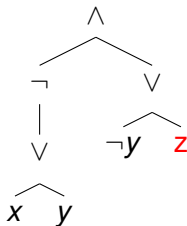
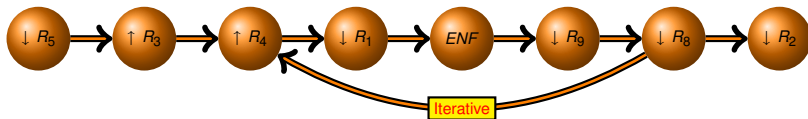
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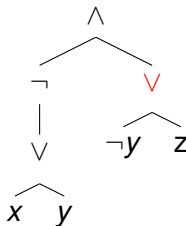
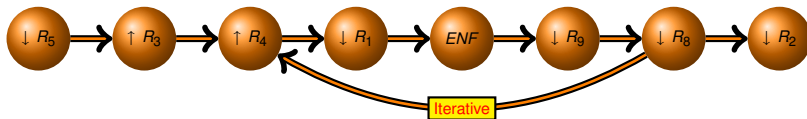
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Combining the reduction rules



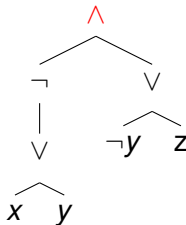
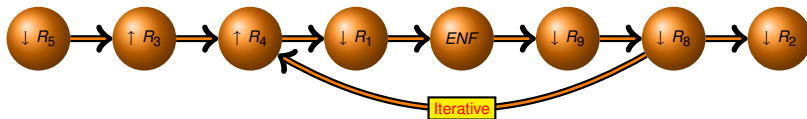
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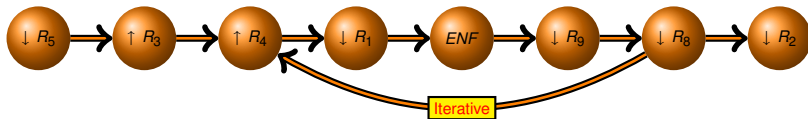
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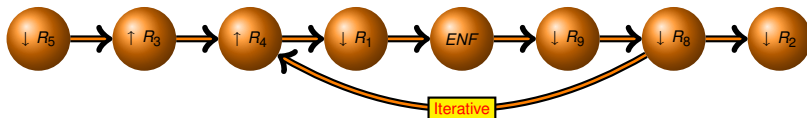
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Combining the reduction rules



Repeat the strategy
wrapped by
Iterative until the
combo tree no longer
changes.

Combining the reduction rules

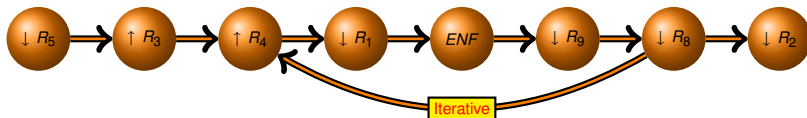


ENF: **E**legant **N**ormal **F**orm
(Holman, '90). A strategy onto itself.



- Efficient

Combining the reduction rules

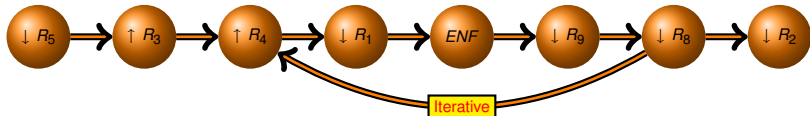


ENF: **Elegant Normal Form**
(Holman, '90). A strategy onto itself.



- Efficient
- Retain the structure (or even improve it)

Combining the reduction rules



```
emacs@a-lin-desktop
File Edit Options Buffers Tools C++ Help

namespace reduct {
  const rule& logical_reduction() {
    using namespace combo;
    static sequential r=
      sequential(downwards(reduce_not(),id::boolean_type),
        upwards(remove_dangling_junctors()),
        iterative(sequential(
          downwards(level()),
          downwards(insert_and(),id::boolean_type),
          subtree_to_enf(),
          downwards(reduce_and(),id::boolean_type),
          downwards(reduce_or(),
            id::boolean_type)),
          downwards(remove_unary_junctors(),id::boolean_type));
    return r;
  }
} //~namespace reduct

--:-- logical_reduction.cc Bot L39 Bzr-2712 (C++/1 Abbrev)-----
```

Outline

1 Introduction

- Program Space
- What are Programs in Normal Forms?
- Effects of Reducing Programs in Normal Forms

2 The Reduct Toolkit

- Recall of Combo
- How it works
- **Demo...**
- What Remains to Be Implemented

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- Boolean, continuous, action-perception: virtually **complete**
 - Numerico-Boolean: dealing with **non-linear systems** like for instance $0 < (* (+ (\#1 \ 1) \ + (\#1 \ 1)))$
 - Higher order, list, recursion, etc: **everything** remains to be done.
 - Factorizing the rules, operator properties rather than operators themselves, for instance
 - $+(X \ 0) \rightarrow X$
 - $\text{or}(X \ \text{false}) \rightarrow X$
- 2 instances of the *neutral_element* reduction rule.