

Faculty of Computer Science & Engineering

Operating Systems

Contact:
Trần Ngọc Anh Tú
51304672@hcmut.edu.vn



Lab 4 - Process

Objective

- ❖ Know how data are layout inside a process
- ❖ Understand dynamic memory allocation mechanism

Inside a process

- ❖ Each process has its own address space which is a set of memory addresses that a process can access.
- ❖ Address space is a large array of bytes starting at 0 and going up to some large number $2^{32} - 1$ or $2^{64} - 1$
- ❖ The address space is split into multiple parts, each part holds different parts of the process
 - ❖ Text
 - ❖ Data
 - ❖ Heap
 - ❖ Stack

Inside a process

- ❖ Exercise: Given following program, identify the memory segments which variables (or function) belongs to.

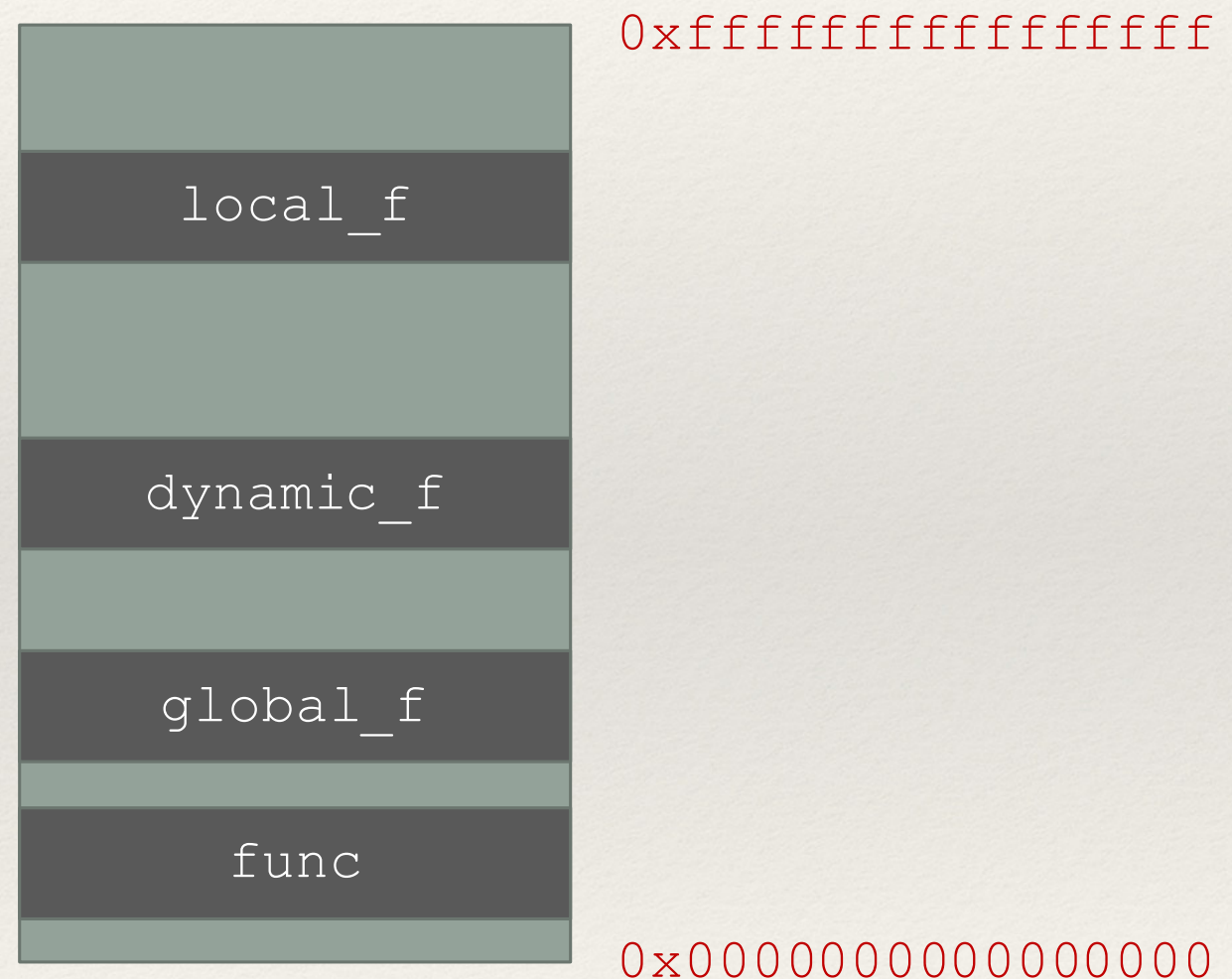
- global_f
- local_f
- func
- dynamic_f

```
#include <stdio.h>
#include <stdlib.h>
int global_f = 10;
int func() {
    return 0;
}
int main() {
    int local_f = 100;
    int *dynamic_f = (int*)malloc(sizeof(int));
    printf("global_f:    %p\n", &global_f);
    printf("local_f:      %p\n", &local_f);
    printf("func:           %p\n", &func);
    printf("dynamic_f: %p\n", dynamic_f);
}
```


Inside a process

❖ Exercise: Given following program, identify the memory segments which variables (or function) belongs to.

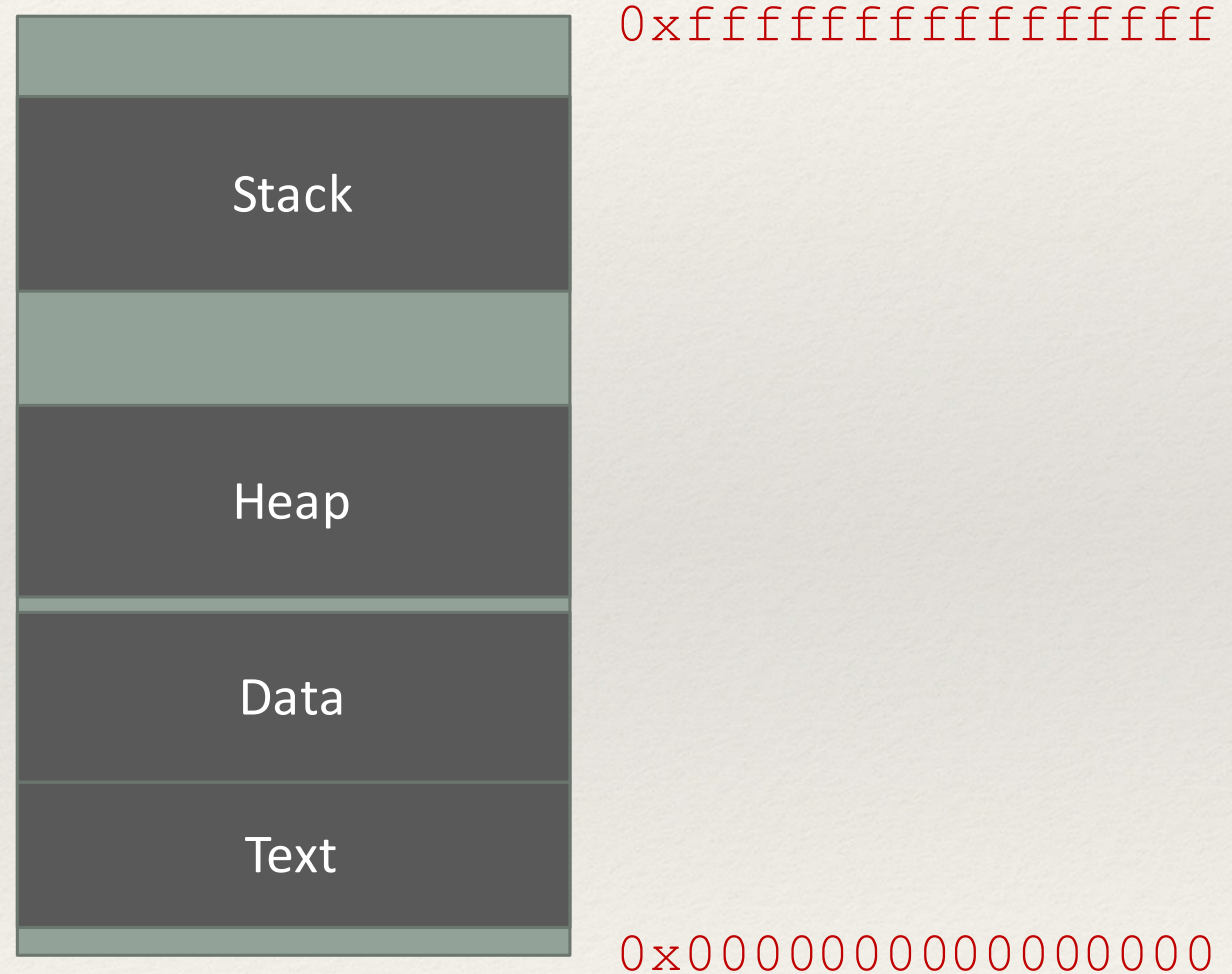
- global_f
- local_f
- func
- dynamic_f



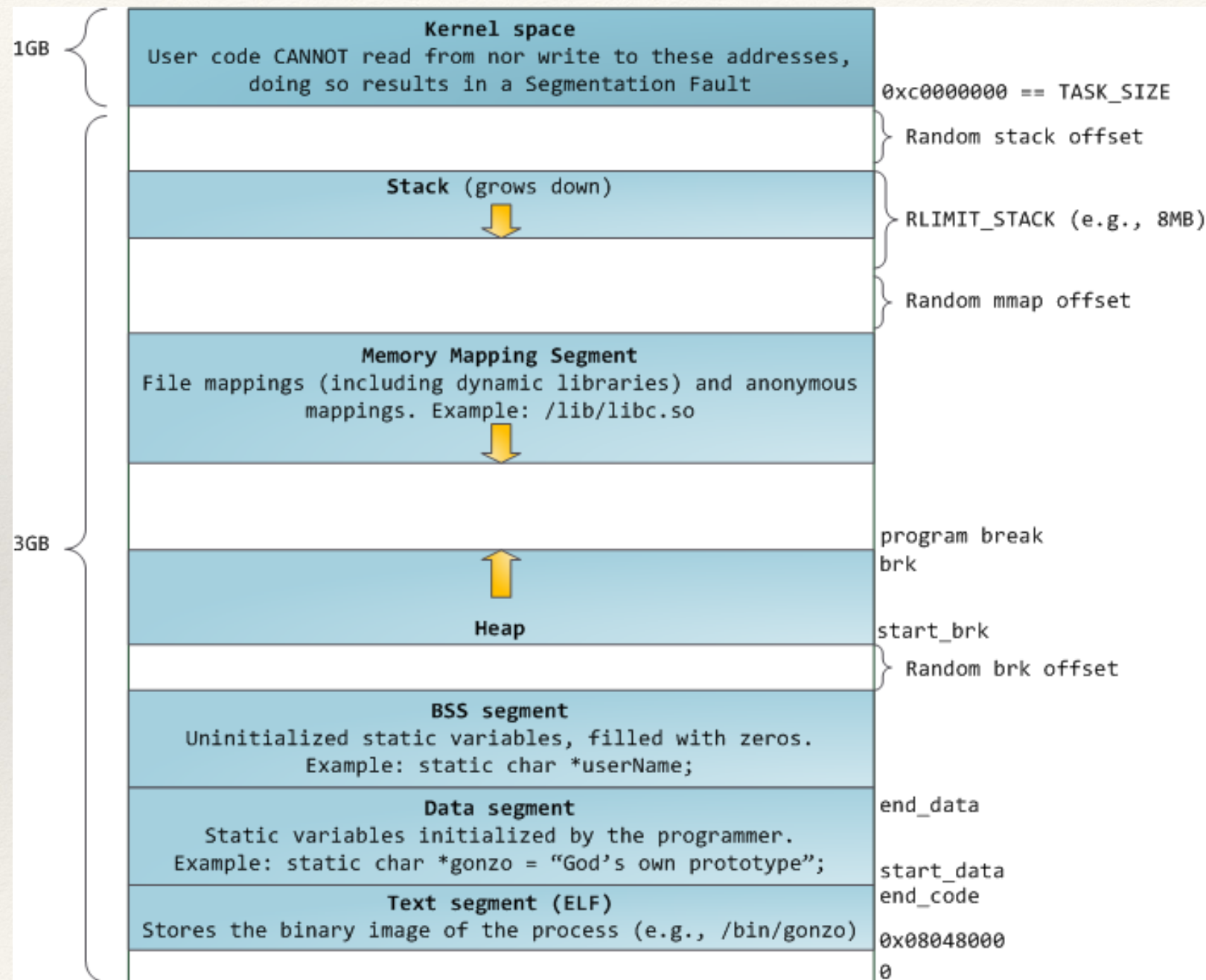
Inside a process

❖ Exercise: Given following program, identify the memory segments which variables (or function) belongs to.

- `global_f`
- `local_f`
- `func`
- `dynamic_f`

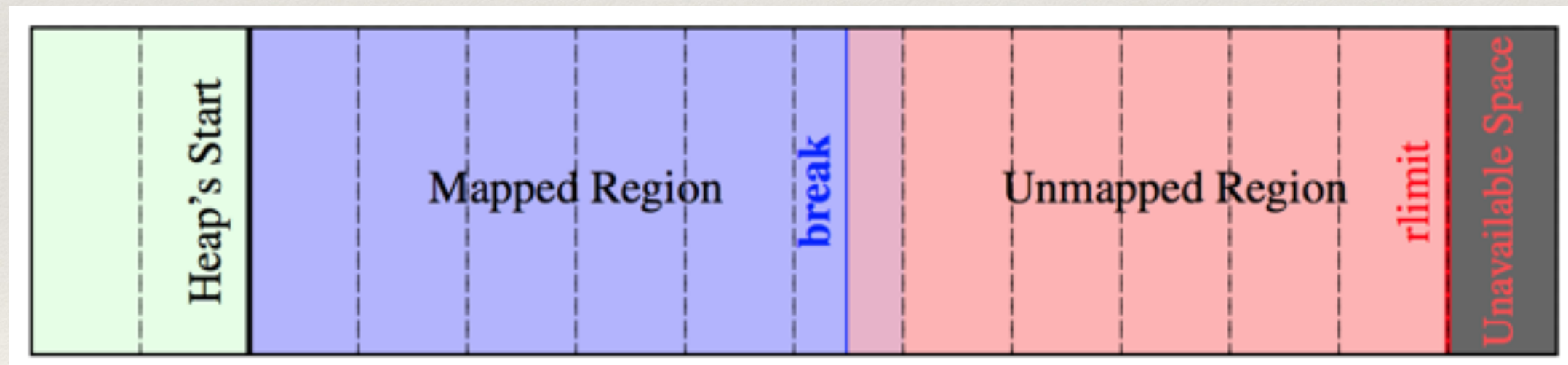


Process address space in Linux



Heap region

- ❖ The heap is placed right after data region in address space.
- ❖ break (program break) separates used regions and unused regions.
- ❖ Heap increases upward but the break cannot go through rlimit



Heap region

- ❖ We could increase or decrease the size of heap region by using two system calls:
 - ❖ `brk` -> changes the value of break pointer.
 - ❖ `sbrk` -> increases the value of break by a given number of bytes.
- ❖ We could use those system calls to implement `malloc` and `free`.

Heap region

- ❖ A simple implementation of malloc

```
int *ptr = (int *) malloc(10 * sizeof (int));  
if (ptr == NULL) {  
} else {  
    /* Allocation succeeded.  Do something.  */  
    free(ptr);  
    ptr = NULL;  
}
```

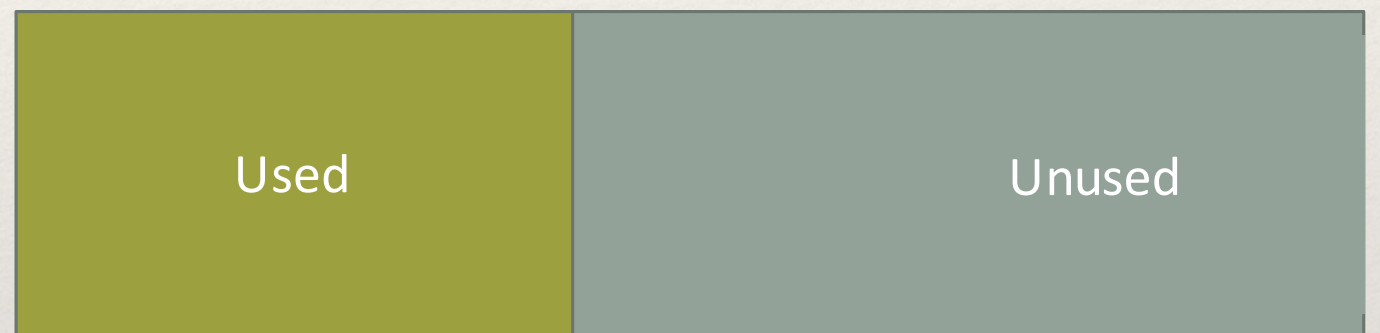
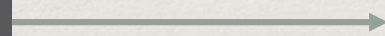
Dynamic allocation

- ❖ Actually, malloc does not allocate a new memory region whose size exactly equals to the size given by user. It allocates a few extra bytes to hold needed information (including the size of allocated region) so that free function could easily clean this region.

Dynamic allocation

- ❖ How to allocate and clean memory region on heap?

```
void *p;  
p = malloc(100);  
  
// do something  
  
free(p);
```



Dynamic allocation

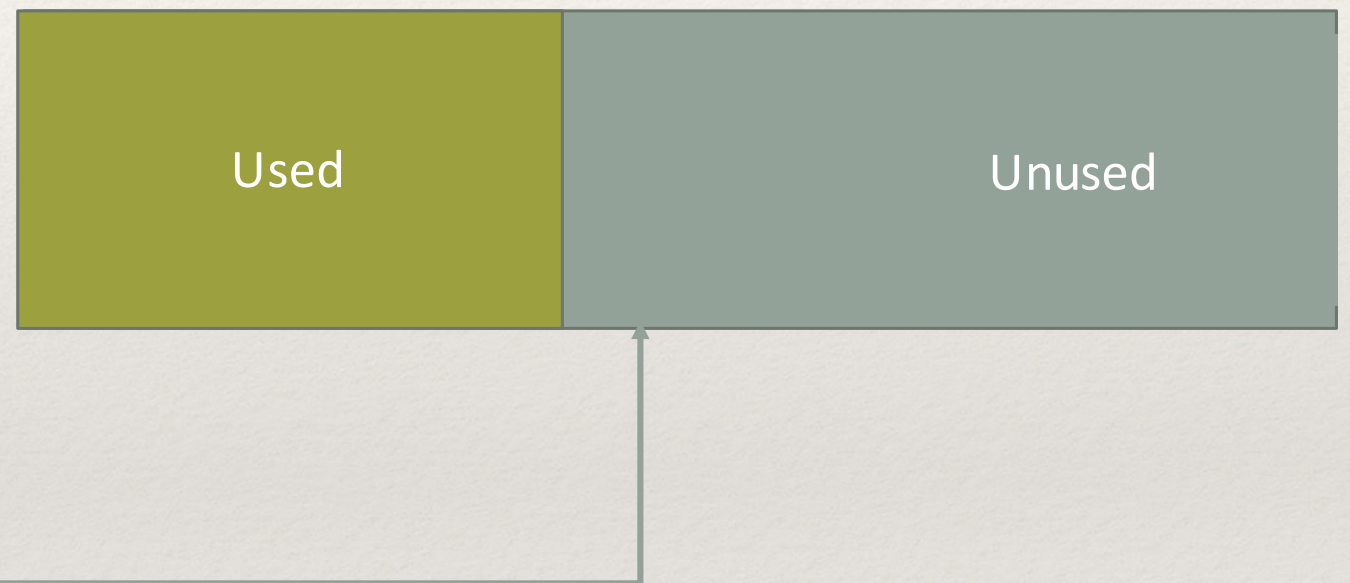
- ❖ How to allocate and clean memory region on heap?



Dynamic allocation

- ❖ How to allocate and clean memory region on heap?

```
void *p;  
p = malloc(100);  
  
// do something  
  
free(p);
```



Dynamic allocation

- ❖ Exercise: Implement your version of malloc and free using the technique described above.

End

Thanks!