## Maximum Matching in Bipartite Graph

```
const int MAXN1 = 50000; const int MAXN2 = 50000; const int MAXM = 150000;
int n1, n2, edges, last[MAXN1], prev[MAXM], head[MAXM];
int matching[MAXN2], dist[MAXN1], Q[MAXN1];
bool used[MAXN1], vis[MAXN1];
void init(int _n1, int _n2) {
       n1 = _{n1};

n2 = _{n2};
        edges = 0;
        fill(last, last + n1, -1);
void addEdge(int u, int v) { //use this in main function to build the graph
        head[edges] = v;
        prev[edges] = last[u];
        last[u] = edges++;
void bfs() {
        fill(dist, dist + n1, -1);
        int sizeQ = 0;
        for (int u = 0; u < n1; ++u) {
                if (!used[u]) {
                        Q[sizeQ++] = u;
                        dist[u] = 0;
                }
        for (int i = 0; i < sizeQ; i++) {
                int u1 = Q[i];
                for (int e = last[u1]; e \ge 0; e = prev[e]) {
                        int u2 = matching[head[e]];
                        if (u2 >= 0 \&\& dist[u2] < 0) {
                                dist[u2] = dist[u1] + 1;
                                Q[sizeQ++] = u2;
                        }
                }
        }
bool dfs(int u1) {
        vis[u1] = true;
        for (int e = last[u1]; e \ge 0; e = prev[e]) {
                int v = head[e];
                int u2 = matching[v];
                if (u2 < 0 \mid \mid !vis[u2] \&\& dist[u2] == dist[u1] + 1 \&\& dfs(u2)) {
                        matching[v] = u1;
                        used[u1] = true;
                        return true;
                }
        return false;
int maxMatching() {
        fill(used, used + n1, false);
        fill (matching, matching + n2, -1);
        for (int res = 0;;) {
                bfs();
                fill(vis, vis + n1, false);
                int f = 0;
                for (int u = 0; u < n1; ++u)
                        if (!used[u] && dfs(u))
                                ++f:
                if (!f)
                        return res;
                res += f;
        }
```

Dijkstra (fast)
typedef pair<int, int> PII;

```
int main() {
  int N, s, t;
  scanf ("%d%d%d", &N, &s, &t);
  vector<vector<PII> > edges(N);
  for (int i = 0; i < N; i++) {</pre>
    int M;
    scanf ("%d", &M);
    for (int j = 0; j < M; j++) {</pre>
      int vertex, dist;
      scanf ("%d%d", &vertex, &dist);
      edges[i].push back (make pair (dist, vertex)); // note order of arguments here
    }
  }
  // use priority queue in which top element has the "smallest" priority
  priority queue<PII, vector<PII>, greater<PII> > Q;
  vector<int> dist(N, INF), dad(N, -1);
  Q.push (make_pair (0, s));
  dist[s] = 0;
  while (!Q.empty()) {
    PII p = Q.top();
    if (p.second == t) break;
    Q.pop();
    int here = p.second;
    for (vector<PII>::iterator it=edges[here].begin(); it!=edges[here].end(); it++){
      if (dist[here] + it->first < dist[it->second]) {
        dist[it->second] = dist[here] + it->first;
        dad[it->second] = here;
        Q.push (make pair (dist[it->second], it->second));
      }
    }
  printf ("%d\n", dist[t]);
  if (dist[t] < INF)</pre>
    for (int i=t;i!=-1;i=dad[i])
      printf ("%d%c", i, (i==s?'\n':' '));
  return 0:
}
Miller Rabin Primality Test
const int maxIter = 10;
bool isPrime (unsigned long long N)
        if(N < 2) return false;</pre>
        if( N % 2 == 0) return N == 2;
        if(N % 3 == 0) return N == 3;
        if(N % 5 == 0) return N == 5;
        if ( N % 7 == 0) return N == 7;
        int d = 0;
        long long odd = N - 1;
        while ( (odd \& 1) == 0)
         {
                d++;
                odd>>= 1;
         }
        for (int i = 0; i < maxIter; i++)
         {
                long long a = rand() % (N - 1) + 1;
                                                               // a is random number from
1 to N -1
                long long mod = modulo( a, odd, N);
                                                                //(a^odd)%N
                bool passes = ( mod == 1 || mod == N -1 );
```

```
for (int r = 1; r < d \&\& !passes; <math>r ++)
                         mod = mulmod( mod, mod, N);
                                                         //(a*b)%N
                         passes = passes \mid \mid \mod == N - 1;
                if(!passes)
                        return false;
        return true;
}
Modular Multiplicative Inverse
int modmulinverse(int a,int m)
    int x = 0, y = 1, u = 1, v = 0;
    int e = m, f = a;
    int c,d,q,r;
    while (f != 1)
        q = e/f;
        r = e%f;
        c = x-q*u;
                           d = y-q*v;
        x = u;
                       y = v;
        u = c;
                       v = d;
                       f = r;
        e = f;
    u = (u+m) %m;
    return u;
Binary Index Tree
int read(int idx){
                                          void update(int idx ,int val){
        int sum = 0;
                                                   while (idx <= MaxVal) {</pre>
        while (idx > 0) {
                                                           tree[idx] += val;
```

```
int read(int idx) {
    int sum = 0;
    while (idx > 0) {
        sum += tree[idx];
        idx -= (idx & -idx);
    }
    return sum;
}
void update(int idx ,int val) {
        while (idx <= MaxVal) {
            tree[idx] += val;
            idx += (idx & -idx);
        }
    }
</pre>
```

## Stable Marriage

```
int m[501][501];
int w[501][501];
int m explored[501];
int w engaged[501];
queue<int> unmarried;
int main() {
        int T;
        T=read int();
        while(T) {
                int n=read int();
                int i=1;
                while(i<=n) {
                        w engaged[i]=0;
                        int j=1;
                        read int();
                        while (j \le n) {
                                w[i][read int()]=j;
                                 j++;
                        i++;
                i=1;
                while (i \le n) {
                        unmarried.push(i);
```

```
m explored[i]=0;
                            int j=1;
                            read int();
                            while (j \le n) {
                                     m[i][j]=read_int();
                                     j++;
                            }
                            i++;
                  while(!unmarried.empty()) {
                           int man=unmarried.front();
                           unmarried.pop();
                           while(1) {
                                     m explored[man]++;
                                     int next_w=m[man][m_explored[man]];
                                     if(w_engaged[next_w]) {
                                              if(w[next_w][w_engaged[next_w]]<w[next_w][man]) {</pre>
                                                       continue;
                                              } else {
                                                       unmarried.push(w engaged[next w]);
                                                       w engaged[next w]=man;
                                                       break;
                                     } else {
                                              w engaged[next w]=man;
                                              break;
                  i=1;
                  while(i<=n) {
                           printf("%d %d\n", w_engaged[i], i);
                            i++;
                  T--;
         return 0;
}
Maximum flow
#define MAX 200
int edge[MAX][MAX];
int parent[MAX];
bool search_path(int N,int s, int t)
{
bool visited[N];
int i;
for(i=0; i < N; i++)
visited[i] = false;
queue<int> myqueue;
myqueue.push(s);
parent[s] = -1;
visited[s] = true;
while(!(myqueue.empty()))
{
int f = myqueue.front();
myqueue.pop();
for(int i = 0; i < N; i++)
if(edge[f][i] > 0 \&\& visited[i] == false)
myqueue.push(i);
```

```
parent[i] = f;
visited[i] = true;
}
return visited[t];
}
int ford_fulkerson(int N, int flow,int s, int t)
int min_wt;
while(search_path(N,s,t))
  min_wt = 99999;
int tt = t;
while(tt != s)
int temp = parent[tt];
if(min_wt > edge[temp][tt])
min_wt = edge[temp][tt];
tt = temp;
tt = t;
while(tt != s)
int temp = parent[tt];
edge[temp][tt] -= min_wt;
edge[tt][temp] += min_wt ;
tt = temp;
flow = flow + min wt;
return flow;
ConvexHull.cc 7/27
struct Point {
    long long x;
     long long y;
     double angle;
} p [100 + 5], hull [100 + 5];
double angle (long long x, long long y)
{
     double theta = atan2 (fabs (y), fabs (x));
     if (x \ge 0) return theta;
     return pi - theta;
}
bool cmp (Point a, Point b)
{
     if (fabs (a.angle - b.angle) < 1e-6 )
         return (a.x * a.x + a.y * a.y) < (b.x * b.x + b.y * b.y);
     return a.angle < b.angle;</pre>
}
bool isInRight (Point a, Point b, Point c)
{
     if (c.x * (a.y - b.y) + c.y * (b.x - a.x) + (a.x * b.y - a.y * b.x) < 0)
         return true;
     return false;
}
double square (double a)
```

```
return a * a;
}
int main ()
{
    int testCase;
    scanf ("%d", &testCase);
    while ( testCase-- ) {
        long long initialLength=0;
        int totalPoints;
        cin >>totalPoints;
        for ( int i = 0; i < totalPoints; i++ )
            cin >> p [i].x >> p [i].y;
        // special case if only 1 totalPoints
        if ( totalPoints == 1 ) {
            printf ("%.5lf\n", (double) initialLength);
            continue;
        int id = 0;
        for ( int i = 1; i < totalPoints; i++ )</pre>
            if (p[i].y < p[id].y) id = i;
        swap (p [id], p [0]);
        // scaling down respective to id
        for ( int i = totalPoints - 1; i >= 0; i-- ) {
            p [i].x -= p [0].x;
            p [i].y -= p [0].y;
        // measure angle according to id
        for ( int i = 1; i < totalPoints; i++ )</pre>
            p [i].angle = angle (p [i].x, p [i].y);
        sort (p + 1, p + totalPoints, cmp);
        hull [0] = p [0];
        hull [1] = p [1];
        int top = 1;
        for ( int i = 2; i < totalPoints; i++ ) {
            while (isInRight (hull [top - 1], hull [top], p [i]))
                top--;
            hull [++top] = p [i];
        double finalLength = 0;
        // perimeter of convex hull
        for ( int i = 1; i <= top; i++ )
            finalLength += sqrt (square (hull [i - 1].x - hull [i].x) + square (hull [i
- 1].y - hull [i].y));
        finalLength += sqrt (square (hull [top].x - hull [0].x) + square (hull [top].y
- hull [0].y));
        if ( finalLength < initialLength ) finalLength = initialLength;</pre>
        printf ("%.21f\n", finalLength);
        for(int i=0;i<totalPoints;i++)</pre>
                       cout<<hull[i].x<<" "<<hull[i].y<<" ";
    }
```

```
return 0;
Geometry.cc 8/27
double INF = 1e100;
double EPS = 1e-12;
struct PT {
  double x, y;
  PT() {}
  PT (double x, double y) : x(x), y(y) {}
  PT(const PT &p) : x(p.x), y(p.y) {}
  PT operator + (const PT &p) const { return PT(x+p.x, y+p.y); }
  PT operator - (const PT &p) const { return PT(x-p.x, y-p.y); }
 PT operator * (double c) const { return PT(x*c, y*c ); } PT operator / (double c) const { return PT(x/c, y/c ); }
  PT operator * (double c)
};
double dot(PT p, PT q)
                            { return p.x*q.x+p.y*q.y; }
double dist2(PT p, PT q)
                            { return dot(p-q,p-q); }
                          { return p.x*q.y-p.y*q.x; }
double cross(PT p, PT q)
ostream &operator<<(ostream &os, const PT &p) {
  os << "(" << p.x << "," << p.y << ")"; }
// rotate a point CCW or CW around the origin
PT RotateCCW90 (PT p) { return PT(-p.y,p.x); }
                       { return PT(p.y,-p.x); }
PT RotateCW90 (PT p)
PT RotateCCW(PT p, double t) {
  return PT (p.x*cos(t)-p.y*sin(t), p.x*sin(t)+p.y*cos(t));
// project point c onto line through a and b
// assuming a != b
PT ProjectPointLine (PT a, PT b, PT c) {
 return a + (b-a) *dot(c-a, b-a) /dot(b-a, b-a);
// project point c onto line segment through a and b
PT ProjectPointSegment(PT a, PT b, PT c) {
 double r = dot(b-a,b-a);
  if (fabs(r) < EPS) return a;</pre>
 r = dot(c-a, b-a)/r;
  if (r < 0) return a;</pre>
  if (r > 1) return b;
  return a + (b-a) *r;
}
// compute distance from c to segment between a and b
double DistancePointSegment(PT a, PT b, PT c) {
 return sqrt(dist2(c, ProjectPointSegment(a, b, c)));
// compute distance between point (x,y,z) and plane ax+by+cz=d
double DistancePointPlane (double x, double y, double z,
                           double a, double b, double c, double d)
  return fabs(a*x+b*y+c*z-d)/sqrt(a*a+b*b+c*c);
}
// determine if lines from a to b and c to d are parallel or collinear
bool LinesParallel (PT a, PT b, PT c, PT d) {
  return fabs(cross(b-a, c-d)) < EPS;</pre>
bool LinesCollinear(PT a, PT b, PT c, PT d) {
  return LinesParallel(a, b, c, d)
      && fabs(cross(a-b, a-c)) < EPS
      && fabs(cross(c-d, c-a)) < EPS;
}
```

```
// determine if line segment from a to b intersects with
// line segment from c to d
bool SegmentsIntersect(PT a, PT b, PT c, PT d) {
  if (LinesCollinear(a, b, c, d)) {
    if (dist2(a, c) < EPS || dist2(a, d) < EPS ||</pre>
      dist2(b, c) < EPS || dist2(b, d) < EPS) return true;
    if (dot(c-a, c-b) > 0 \&\& dot(d-a, d-b) > 0 \&\& dot(c-b, d-b) > 0)
      return false;
    return true;
  }
  if (cross(d-a, b-a) * cross(c-a, b-a) > 0) return false;
  if (cross(a-c, d-c) * cross(b-c, d-c) > 0) return false;
  return true:
}
// compute intersection of line passing through a and b with line passing through c and
d, assuming unique intersection exists; for segment intersection, check if segments
intersect first
PT ComputeLineIntersection(PT a, PT b, PT c, PT d) {
 b=b-a; d=c-d; c=c-a;
  assert(dot(b, b) > EPS && dot(d, d) > EPS);
  return a + b*cross(c, d)/cross(b, d);}
PT ComputeCircleCenter(PT a, PT b, PT c) {
 b = (a+b)/2;
  c = (a+c)/2;
  return ComputeLineIntersection(b, b+RotateCW90(a-b), c, c+RotateCW90(a-c));
// determine if point is in a possibly non-convex polygon (by William
// Randolph Franklin); returns 1 for strictly interior points, 0 for
// strictly exterior points, and 0 or 1 for the remaining points.
// Note that it is possible to convert this into an *exact* test using
// integer arithmetic by taking care of the division appropriately
// (making sure to deal with signs properly) and then by writing exact
// tests for checking point on polygon boundary
bool PointInPolygon(const vector<PT> &p, PT q) {
 bool c = 0;
  for (int i = 0; i < p.size(); i++) {</pre>
    int j = (i+1)%p.size();
    if ((p[i].y <= q.y && q.y < p[j].y ||</pre>
     p[j].y \le q.y \&\& q.y < p[i].y) \&\&
      q.x < p[i].x + (p[j].x - p[i].x) * (q.y - p[i].y) / (p[j].y - p[i].y))
      c = !c;
  return c;
// determine if point is on the boundary of a polygon
bool PointOnPolygon(const vector<PT> &p, PT q) {
  for (int i = 0; i < p.size(); i++)</pre>
    if (dist2(ProjectPointSegment(p[i], p[(i+1)%p.size()], q), q) < EPS)</pre>
      return true;
    return false;
// compute intersection of line through points a and b with circle centered at c with
rad r > 0
vector<PT> CircleLineIntersection(PT a, PT b, PT c, double r) {
  vector<PT> ret;
 b = b-a;
  a = a-c;
  double A = dot(b, b);
  double B = dot(a, b);
  double C = dot(a, a) - r*r;
  double D = B*B - A*C;
  if (D < -EPS) return ret;</pre>
  ret.push back(c+a+b*(-B+sqrt(D+EPS))/A);
```

```
if (D > EPS)
    ret.push back(c+a+b*(-B-sqrt(D))/A);
  return ret;
}
// compute intersection of circle centered at a with rad r with circle centered at b
with rad R
vector<PT> CircleCircleIntersection(PT a, PT b, double r, double R) {
  vector<PT> ret;
  double d = sqrt(dist2(a, b));
  if (d > r+R \mid | d+min(r, R) < max(r, R)) return ret;
  double x = (d*d-R*R+r*r)/(2*d);
  double y = sqrt(r*r-x*x);
  PT v = (b-a)/d;
  ret.push back(a+v*x + RotateCCW90(v)*y);
  if (y > 0)
    ret.push_back(a+v*x - RotateCCW90(v)*y);
  return ret;
}
// This code computes the area or centroid of a (possibly nonconvex)polygon, assuming
that the coordinates are listed in a clockwise or counterclockwise fashion.
the centroid is often known as the "center of gravity" or "center of mass".
double ComputeSignedArea(const vector<PT> &p) {
  double area = 0;
  for(int i = 0; i < p.size(); i++) {</pre>
    int j = (i+1) % p.size();
    area += p[i].x*p[j].y - p[j].x*p[i].y;
  return area / 2.0;
double ComputeArea(const vector<PT> &p) {
  return fabs(ComputeSignedArea(p));
PT ComputeCentroid(const vector<PT> &p) {
  PT c(0,0);
  double scale = 6.0 * ComputeSignedArea(p);
  for (int i = 0; i < p.size(); i++) {</pre>
    int j = (i+1) % p.size();
    c = c + (p[i]+p[j])*(p[i].x*p[j].y - p[j].x*p[i].y);
  return c / scale;
}
// tests whether or not a given polygon (in CW or CCW order) is simple
bool IsSimple(const vector<PT> &p) {
  for (int i = 0; i < p.size(); i++) {</pre>
    for (int k = i+1; k < p.size(); k++) {</pre>
      int j = (i+1) % p.size();
      int l = (k+1) % p.size();
      if (i == 1 || j == k) continue;
      if (SegmentsIntersect(p[i], p[j], p[k], p[l]))
        return false;
  }
  return true;
int main() {
    cerr << RotateCCW90(PT(2,5)) << endl; // expected: (-5,2)</pre>
  cerr \ll RotateCW90(PT(2,5)) \ll endl; // expected: (5,-2)
  cerr << RotateCCW(PT(2,5),M PI/2) << endl;// expected: (-5,2)</pre>
  cerr << ProjectPointLine(PT(-5,-2), PT(10,4), PT(3,7)) << endl; // expected: (5,2)
  cerr << ProjectPointSegment(PT(-5,-2), PT(10,4), PT(3,7)) << ""
       << ProjectPointSegment(PT(7.5,3), PT(10,4), PT(3,7)) << " "
       << ProjectPointSegment(PT(-5,-2), PT(2.5,1), PT(3,7)) << endl;
  // expected: 6.78903
  cerr << DistancePointPlane(4,-4,3,2,-2,5,-8) << endl;</pre>
```

```
// expected: 1 0 1
cerr << LinesParallel(PT(1,1), PT(3,5), PT(2,1), PT(4,5)) << " "
     << LinesParallel(PT(1,1), PT(3,5), PT(2,0), PT(4,5)) << ""
     << LinesParallel(PT(1,1), PT(3,5), PT(5,9), PT(7,13)) << endl;
// expected: 0 0 1
cerr << LinesCollinear(PT(1,1), PT(3,5), PT(2,1), PT(4,5)) << " "
     << LinesCollinear(PT(1,1), PT(3,5), PT(2,0), PT(4,5)) << ""
     << LinesCollinear(PT(1,1), PT(3,5), PT(5,9), PT(7,13)) << endl;
// expected: 1 1 1 0
cerr << SegmentsIntersect(PT(0,0), PT(2,4), PT(3,1), PT(-1,3)) << " "
     << SegmentsIntersect(PT(0,0), PT(2,4), PT(4,3), PT(0,5)) << " "
     << SegmentsIntersect(PT(0,0), PT(2,4), PT(2,-1), PT(-2,1)) << " "
     << SegmentsIntersect(PT(0,0), PT(2,4), PT(5,5), PT(1,7)) << endl;
// expected: (1,2)
cerr << ComputeLineIntersection(PT(0,0), PT(2,4), PT(3,1), PT(-1,3)) << endl;
// expected: (1,1)
cerr << ComputeCircleCenter(PT(-3,4), PT(6,1), PT(4,5)) << endl;
vector<PT> v;
v.push back(PT(0,0));
v.push back(PT(5,0));
v.push back(PT(5,5));
v.push back(PT(0,5));
// expected: 1 1 1 0 0
cerr << PointInPolygon(v, PT(2,2)) << " "</pre>
     << PointInPolygon(v, PT(2,0)) << " "
     << PointInPolygon(v, PT(0,2)) << " "
     << PointInPolygon(v, PT(5,2)) << " "
     << PointInPolygon(v, PT(2,5)) << endl;
// expected: 0 1 1 1 1
cerr << PointOnPolygon(v, PT(2,2)) << " "</pre>
     << PointOnPolygon(v, PT(2,0)) << " "
     << PointOnPolygon(v, PT(0,2)) << " "
     << PointOnPolygon(v, PT(5,2)) << " "
     << PointOnPolygon(v, PT(2,5)) << endl;
// expected: (1,6)
//
             (5,4) (4,5)
//
            blank line
             (4,5) (5,4)
            blank line
             (4,5) (5,4)
vector < PT > u = CircleLineIntersection(PT(0,6), PT(2,6), PT(1,1), 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;</pre>
u = CircleLineIntersection(PT(0,9), PT(9,0), PT(1,1), 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;</pre>
u = CircleCircleIntersection(PT(1,1), PT(10,10), 5, 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;</pre>
u = CircleCircleIntersection(PT(1,1), PT(8,8), 5, 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;</pre>
u = CircleCircleIntersection(PT(1,1), PT(4.5,4.5), 10, sqrt(2.0)/2.0);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;</pre>
u = CircleCircleIntersection(PT(1,1), PT(4.5,4.5), 5, sqrt(2.0)/2.0);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;</pre>
// area should be 5.0
// centroid should be (1.1666666, 1.166666)
PT pa[] = { PT(0,0), PT(5,0), PT(1,1), PT(0,5) };
vector<PT> p(pa, pa+4);
PT c = ComputeCentroid(p);
cerr << "Area: " << ComputeArea(p) << endl;</pre>
```

```
cerr << "Centroid: " << c << endl;
  return 0;
}
Euclid.cc 12/27
// This is a collection of useful code for solving problems that
// involve modular linear equations. Note that all of the// returns d = qcd(a,b);
finds x, y such that d = ax + by
int extended euclid(int a, int b, int &x, int &y) {
  int xx = y = 0;
 int yy = x = 1;
 while (b) {
   int q = a/b;
   int t = b; b = a%b; a = t;
   t = xx; xx = x-q*xx; x = t;
    t = yy; yy = y-q*yy; y = t;
  }
  return a;
}
// finds all solutions to ax = b (mod n)
VI modular linear equation solver(int a, int b, int n) {
  int x, y;
  VI solutions;
  int d = extended euclid(a, n, x, y);
  if (!(b%d)) {
   x = mod (x*(b/d), n);
    for (int i = 0; i < d; i++)
      solutions.push back (mod(x + i*(n/d), n));
  return solutions;
// Chinese remainder theorem (special case): find z such that
// z % x = a, z % y = b. Here, z is unique modulo M = lcm(x,y).
// Return (z,M). On failure, M=-1.
PII chinese remainder theorem(int x, int a, int y, int b) {
  int s, t;
  int d = extended euclid(x, y, s, t);
 if (a%d != b%d) return make pair(0, -1);
 return make pair (mod (s*b*x+t*a*y, x*y) /d, x*y/d);
// Chinese remainder theorem: find z such that
// z % x[i] = a[i] for all i. Note that the solution is
// unique modulo M = lcm \ i \ (x[i]). Return (z,M). On
// failure, M = -1. Note that we do not require the a[i]'s
// to be relatively prime.
PII chinese_remainder_theorem(const VI &x, const VI &a) {
 PII ret = make_pair(a[0], x[0]);
  for (int i = 1; i < x.size(); i++) {</pre>
   ret = chinese remainder theorem(ret.first, ret.second, x[i], a[i]);
    if (ret.second == -1) break;
  }
  return ret;
// computes x and y such that ax + by = c; on failure, x = y = -1
void linear diophantine(int a, int b, int c, int &x, int &y) {
 int d = gcd(a,b);
 if (c%d) {
   x = y = -1;
  } else {
   x = c/d * mod inverse(a/d, b/d);
    y = (c-a*x)/b;
  }
}
```

```
// Uses:
    (1) solving systems of linear equations (AX=B)
     (2) inverting matrices (AX=I)
//
     (3) computing determinants of square matrices
//
// Running time: O(n^3)
//
// INPUT:
             a[][] = an nxn matrix
//
             b[][] = an nxm matrix
//
// OUTPUT:
                    = an nxm matrix (stored in b[][])
             X
//
             A^{-1} = an nxn matrix (stored in a[][])
//
             returns determinant of a[][]
const double EPS = 1e-10;
typedef vector<int> VI;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
T GaussJordan (VVT &a, VVT &b) {
  const int n = a.size();
  const int m = b[0].size();
  VI irow(n), icol(n), ipiv(n);
  T \det = 1;
  for (int i = 0; i < n; i++) {</pre>
    int pj = -1, pk = -1;
    for (int j = 0; j < n; j++) if (!ipiv[j])</pre>
      for (int k = 0; k < n; k++) if (!ipiv[k])</pre>
        if (pj == -1 \mid | fabs(a[j][k]) > fabs(a[pj][pk])) { pj = j; pk = k; }
    if (fabs(a[pj][pk]) < EPS) { cerr << "Matrix is singular." << endl; exit(0); }</pre>
    ipiv[pk]++;
    swap(a[pj], a[pk]);
    swap(b[pj], b[pk]);
    if (pj != pk) det *= -1;
    irow[i] = pj;
    icol[i] = pk;
    T c = 1.0 / a[pk][pk];
    det *= a[pk][pk];
    a[pk][pk] = 1.0;
    for (int p = 0; p < n; p++) a[pk][p] *= c;</pre>
    for (int p = 0; p < m; p++) b[pk][p] *= c;</pre>
    for (int p = 0; p < n; p++) if (p != pk) {</pre>
      c = a[p][pk];
      a[p][pk] = 0;
      for (int q = 0; q < n; q++) a[p][q] -= a[pk][q] * c;</pre>
      for (int q = 0; q < m; q++) b[p][q] -= b[pk][q] * c;</pre>
    }
  }
  for (int p = n-1; p >= 0; p--) if (irow[p] != icol[p]) {
    for (int k = 0; k < n; k++) swap(a[k][irow[p]], a[k][icol[p]]);
  return det;
}
int main() {
  const int n = 4;
  const int m = 2;
  double A[n][n] = { {1,2,3,4}, {1,0,1,0}, {5,3,2,4}, {6,1,4,6} };
  double B[n][m] = \{ \{1,2\}, \{4,3\}, \{5,6\}, \{8,7\} \};
  VVT a(n), b(n);
  for (int i = 0; i < n; i++) {</pre>
```

```
a[i] = VT(A[i], A[i] + n);
    b[i] = VT(B[i], B[i] + m);
  double det = GaussJordan(a, b);
  // expected: 60
  cout << "Determinant: " << det << endl;</pre>
  cout << "Inverse: " << endl;</pre>
  for (int i = 0; i < n; i++) {</pre>
    for (int j = 0; j < n; j++)</pre>
     cout << a[i][j] << ' ';
    cout << endl;</pre>
  }
  // expected: 1.63333 1.3
                -0.166667 0.5
  //
                2.36667 1.7
  //
                -1.85 -1.35
  cout << "Solution: " << endl;</pre>
  for (int i = 0; i < n; i++) {</pre>
    for (int j = 0; j < m; j++)</pre>
      cout << b[i][j] << '
    cout << endl;</pre>
  }
SCC.cc 18/27
#include<memory.h>
struct edge{int e, nxt;};
int V, E;
edge e[MAXE], er[MAXE];
int sp[MAXV], spr[MAXV];
int group cnt, group num[MAXV];
bool v[MAXV];
int stk[MAXV];
void fill forward(int x)
  int i;
  v[x]=true;
  for(i=sp[x];i;i=e[i].nxt) if(!v[e[i].e]) fill_forward(e[i].e);
  stk[++stk[0]]=x;
void fill_backward(int x)
  int i;
  v[x] = false;
  group_num[x]=group_cnt;
  for(i=spr[x];i;i=er[i].nxt) if(v[er[i].e]) fill backward(er[i].e);
void add edge(int v1, int v2) //add edge v1->v2
  e [++E].e=v2; e [E].nxt=sp [v1]; sp [v1]=E;
  er[ E].e=v1; er[E].nxt=spr[v2]; spr[v2]=E;
void SCC()
  int i;
  stk[0]=0;
  memset(v, false, sizeof(v));
  for(i=1;i<=V;i++) if(!v[i]) fill_forward(i);</pre>
  group cnt=0;
  for(i=stk[0];i>=1;i--) if(v[stk[i]]) {group_cnt++; fill backward(stk[i]);}
```

```
struct SuffixArray {
  const int L;
  string s;
  vector<vector<int> > P;
  vector<pair<int,int>,int> > M;
  SuffixArray(const string &s): L(s.length()), s(s), P(1, vector<int>(L, 0)), M(L) {
    for (int i = 0; i < L; i++) P[0][i] = int(s[i]);</pre>
    for (int skip = 1, level = 1; skip < L; skip *= 2, level++) {</pre>
      P.push back(vector<int>(L, 0));
      for (int i = 0; i < L; i++)</pre>
        M[i] = make pair(make pair(P[level-1][i], i + skip < L ? P[level-1][i + skip] :
-1000), i);
      sort(M.begin(), M.end());
      for (int i = 0; i < L; i++)</pre>
        P[level][M[i].second] = (i > 0 && M[i].first == M[i-1].first) ? P[level][M[i-1].first)
1].second] : i;
    }
  }
  vector<int> GetSuffixArray() { return P.back(); }
  // returns the length of the longest common prefix of s[i...L-1] and s[j...L-1]
  int LongestCommonPrefix(int i, int j) {
    int len = 0;
    if (i == j) return L - i;
    for (int k = P.size() - 1; k >= 0 && i < L && j < L; k--) {</pre>
      if (P[k][i] == P[k][j]) {
        i += 1 << k;
        j += 1 << k;
        len += 1 << k;
    return len;
};
UnionFind.cc 21/27
//union-find set: the vector/array contains the parent of each node
int find(vector \langle int \rangle_{C} C, int x){return (C[x]==x) ? x : C[x]=find(C, C[x]);} //C++
int find(int x) {return (C[x]==x)?x:C[x]=find(C[x]);} //C
KMP
function build failure function(pattern[])
 // let m be the length of the pattern
 F[0] = F[1] = 0; // always true
 for(i = 2; i \le m; i++) {
   // j is the index of the largest next partial match
   // (the largest suffix/prefix) of the string under
   // index i - 1
   j = F[i - 1];
   for(;;) {
     // check to see if the last character of string i -
     // - pattern[i - 1] "expands" the current "candidate"
     // best partial match - the prefix under index j
```

if(pattern[i] == pattern[i - 1])

```
F[i] = j + 1; break;
       // if we cannot "expand" even the empty string
       if(j == 0) \{ F[i] = 0; break; \}
      // else go to the next best "candidate" partial match
      i = F[i];
 }
function Knuth_Morris_Pratt(text[], pattern[])
 // let n be the size of the text, m the
 // size of the pattern, and F[] - the
 // "failure function"
 build_failure_function(pattern[]);
 i = 0; // the initial state of the automaton is
 // the empty string
 i = 0; // the first character of the text
 for(;;) {
    if(j == n) break; // we reached the end of the text
    // if the current character of the text "expands" the
    // current match
    if(text[i] == pattern[i]) {
       i++; // change the state of the automaton
      j++; // get the next character from the text
      if(i == m) // match found
    }
    // if the current state is not zero (we have not
    // reached the empty string yet) we try to
    // "expand" the next best (largest) match
    else if(i > 0) i = F[i];
    // if we reached the empty string and failed to
    // "expand" even it; we go to the next
    // character from the text, the state of the
    // automaton remains zero
    else j++;
 }
}
KRUSKAL
using namespace std;
vector< pair<int, edge > > Graph, MST;
int n,e,parent[MAX];
void initialize(int n)
     for(int i=1; i<=n; i++)
           parent[i] = i;
     MST.clear();
     Graph.clear();
}
```

```
//finding parent with path compression
int findSet(int x, int *parent)
     if(x!=parent[x])
          parent[x] = findSet(parent[x], parent);
     return parent[x];
}
void kruskal()
     sort(Graph.begin(),Graph.end());
     int total = 0;
     for(int i=0; i<e; i++)
          int pu = findSet(Graph[i].second.first,parent);
          int pv = findSet(Graph[i].second.second,parent);
          if(pu!=pv)
                total+=Graph[i].first;
                MST.push back(Graph[i]);
                parent[pu] = parent[pv];
     cout<<total<<"\n";
int main()
     cin>>n>>e;
     initialize(n);
     int u, v, w;
     for(int i=0; i<e; i++)
          cin>>u>>v>>w;
          Graph.push_back(pair<int, edge>(w,edge(u,v)));
     kruskal();
     return 0;
Baby step Gaint Step Algorithm:
To calculate value of x such that
Output: A value x satisfying
       1.m \leftarrow \text{Ceiling}(\sqrt{n})
       2.For all j where 0 \le j < m:
                1. Compute \alpha j and store the pair (j, \alpha j) in a table. (See section "In practice")
       3.Compute \alpha–m.
       4.\gamma \leftarrow \beta. (set \gamma = \beta)
       5. For i = 0 to (m - 1):
                1. Check to see if \gamma is the second component (\alpha j) of any pair in the table.
               2.If so, return im + j.
If not, \gamma \leftarrow \gamma \bullet \alpha - m
Pollard Factorization
def pollardRho(N):
        if N%2==0:
```

```
return 2
   x = random.randint(1, N-1)
   y = x
   c = random.randint(1, N-1)
   g = 1
   while g==1:
      x = ((x*x)\%N+c)\%N
      y = ((y*y)\%N+c)\%N
      y = ((y*y)\%N+c)\%N
      g = gcd(abs(x-y),N)
   return g
Method for Finding the Lowest Common Ancestor of 2 nodes in a tree.
void process(int N) //N nodes
      memset(Root, -1, sizeof(Root)); //Root[N][log N]
      for(int i=1; i \le N; i++) Root[i][0]=parent[i]; //stores the 2i
      for (int i=1; (1<<i) <= N; i++)
      for (int j=1; j<=N; j++)</pre>
      if(Root[j][i-1]!=-1)
      Root[j][i]=Root[Root[j][i-1]][i-1];
int lca(int p,int q) //Fnds the LCA of 2 nodes
      int temp;
      if(depth[p]>depth[q])
             int steps=store[depth[q]];
      for(int i=steps;i>=0;i--) //putting p & q on same level
             swap(p,q);
      if(depth[q]-(1<< i) >= depth[p])
      q=Root[q][i];
      if(p==q) return p;
      for(int i=steps;i>=0;i--)
             if(Root[p][i]!=Root[q][i])
                    p=Root[p][i],q=Root[q][i];
      return parent[p];
void init() //store[i] = int( log2i )
      store[0]=0; store[1]=0; store[2]=1;
      int cmp=4;
      for (int i=3; i<1008; i++)
```

```
if(cmp>i) store[i]=store[i-1];
               else
                       store[i]=store[i-1]+1;
                       cmp<<=1;
//Maximum cost weighted matching
//Lifted from topcoder, tested on SPOJ:GREED
# define MAXN 500
# define MAXD 1000
# define INF 1000000000000011
int N,R;
int matched;
long long adjmat[MAXN][MAXN];
long long l1[MAXN],l2[MAXN],sl[MAXN];
int m1[MAXN],m2[MAXN],bfsq[MAXN],par[MAXN],sx[MAXN];
char S[MAXN],T[MAXN];
void updatetree(int v,int prev)
 S[v]=true;
 par[v]=prev;
 for(int i=0;i< N;i++)
   if(11[v]+12[i]-adjmat[v][i] < sl[i])
      sl[i]=l1[v]+l2[i]-adjmat[v][i];
      sx[i]=v;
   }
void updatelabels()
 long long delta=111<<62;
 for(int i=0;i<N;i++)
   if((!T[i])\&\&(delta>sl[i]))
      delta=sl[i];
 for(int i=0;i<N;i++)
   if(S[i])l1[i]-=delta;
   if(T[i])12[i]+=delta;
   else sl[i]-=delta;
void augment()
 if(matched==N)return;
 int qmax=0,qpos=0,x,y,root;
 memset(S,0,N);
 memset(T,0,N);
 memset(par,-1,N<<2);
 for(x=0;x< N;x++)
   if(m1[x]==-1)
```

{

}

{

} }

```
bfsq[qmax++]=root=x;
     S[x]=1;
     par[x]=-2;
     break;
for(y=0;y< N;y++)
  sl[y]=11[root]+12[y]-adjmat[root][y];
  sx[y]=root;
}
while(1)
  while(qpos<qmax)
     x=bfsq[qpos++];
     for(y=0;y< N;y++)
     {
        if((adjmat[x][y]==11[x]+12[y])\&\&(!T[y]))
          if(m2[y]==-1)
             break;
          T[y]=1;
          bfsq[qmax++]=m2[y];
          updatetree(m2[y],x);
     }
     if(y<N)break;
  if(y<N)break;
  updatelabels();
  qmax=qpos=0;
  for(y=0;y<\!N;y++)
  {
     if((!T[y])&&(sl[y]==0))
     {
        if(m2[y]==-1)
          x=sx[y];
          break;
        }
        else
          T[y]=1;
          if(!S[m2[y]]) \\
             bfsq[qmax++]=m2[y];
             updatetree(m2[y],sx[y]);\\
        }
  if(y < N)
     break;
if(y < N)
  for(int cx=x,cy=y,ty;cx!=-2;cx=par[cx],cy=ty)
     ty = m1[cx];
```

```
m2[cy] = cx;
      m1[cx] = cy;
    augment();
//Way shorter but slower Hungarian Algorithm Use if Number of vertices is small
//Warning: Vertices to be 1-indexed Tested on SPOJ: GREED
const int INF=(1 << 30)-1;
int a[505][505];
int Hungarian(int n,int m){
 vector<int> u (n+1), v (m+1), p (m+1), way (m+1);
 for (int i=1; i <= n; ++i) {
    p[0] = i;
    int j0 = 0;
    vector<int> minv (m+1, INF);
    vector<char> used (m+1, false);
    do {
      used[i0] = true;
      int i0 = p[j0], delta = INF, j1;
      for (int j=1; j <= m; ++j)
         if (!used[j]) {
            int cur = a[i0][j]-u[i0]-v[j];
           if (cur < minv[j])
              minv[j] = cur, way[j] = j0;
            if (minv[j] < delta)
              delta = minv[j], j1 = j;
         }
      for (int j=0; j<=m; ++j)
         if (used[j])
            u[p[j]] += delta, v[j] -= delta;
         else
           minv[j] -= delta;
      j0 = j1;
    \} while (p[j0] != 0);
    do {
      int i1 = way[i0];
      p[j0] = p[j1];
      j0 = j1;
    } while (j0);
 return -v[0];
}
Closest pair
typedef pair<long long, long long> pairll;
int n;
pairll pnts [100000];
set<pairll> box;
double best;
int compx(pairll a, pairll b) { return a.px<b.px; }</pre>
int main () {
        scanf("%d", &n);
        for (int i=0;i<n;++i) scanf("%lld %lld", &pnts[i].px, &pnts[i].py);</pre>
        sort(pnts, pnts+n, compx);
```

```
best = 1500000000; // INF
      box.insert(pnts[0]);
      int left = 0;
      for (int i=1;i<n;++i) {
            while (left<i && pnts[i].px-pnts[left].px > best) box.erase(pnts[left++]);
            for (typeof(box.begin()) it=box.lower bound(make pair(pnts[i].py-best,
pnts[i].px-best));
            it!=box.end() && pnts[i].py+best>=it->py; it++)
            best = min(best, sqrt(pow(pnts[i].py - it->py, 2.0)+pow(pnts[i].px - it-
>px, 2.0)));
     box.insert(pnts[i]);
     printf("%.2f\n", best);
      return 0;
Articulation Point
#define NIL -1
class Graph
int V; // No. of vertices
list<int> *adj; // A dynamic array of adjacency lists
void APUtil(int v, bool visited[], int disc[], int low[],
int parent[], bool ap[]);
public:
Graph(int V); // Constructor
void addEdge(int v, int w); // function to add an edge to graph
void AP(); // prints articulation points
};
Graph::Graph(int V)
this->V = V;
 adj = new list<int>[V];
}
void Graph::addEdge(int v, int w)
adj[v].push_back(w);
adj[w].push back(v); // Note: the graph is undirected
}
// u --> The vertex to be visited next
```

```
// disc[] --> Stores discovery times of visited vertices
// ap[] --> Store articulation points
void Graph::APUtil(int u, bool visited[], int disc[], int low[], int parent[], bool
ap[])
 // A static variable is used for simplicity, we can avoid use of static // variable by
passing a pointer.
 static int time = 0;
 // Count of children in DFS Tree
 int children = 0;
 // Mark the current node as visited
 visited[u] = true;
 disc[u] = low[u] = ++time;
 // Go through all vertices aadjacent to this
 list<int>::iterator i;
 for (i = adj[u].begin(); i != adj[u].end(); ++i)
 int v = *i; // v is current adjacent of u // If v is not visited yet, then make it a
child of u // in DFS tree and recur for it
 if (!visited[v])
 children++;
 parent[v] = u;
APUtil(v, visited, disc, low, parent, ap);
 // Check if the subtree rooted with v has a connection to // one of the ancestors of u
low[u] = min(low[u], low[v]);
 // u is an articulation point in following cases // (1) u is root of DFS tree and has
two or more chilren.
if (parent[u] == NIL && children > 1)
 ap[u] = true;
 // (2) If u is not root and low value of one of its child is more // than discovery
value of u.
if (parent[u] != NIL && low[v] >= disc[u])
 ap[u] = true;
 // Update low value of u for parent function calls.
 else if (v != parent[u])
 low[u] = min(low[u], disc[v]);
```

```
// The function to do DFS traversal. It uses recursive function APUtil()
void Graph::AP()
bool *visited = new bool[V];
 int *disc = new int[V];
 int *low = new int[V];
 int *parent = new int[V];
bool *ap = new bool[V]; // To store articulation points
 // Initialize parent and visited, and ap(articulation point) arrays
 for (int i = 0; i < V; i++)
parent[i] = NIL;
 visited[i] = false;
 ap[i] = false;
 // Call the recursive helper function to find articulation points // in DFS tree
rooted with vertex 'i'
for (int i = 0; i < V; i++)
 if (visited[i] == false)
APUtil(i, visited, disc, low, parent, ap); // Now ap[] contains articulation points,
print them
for (int i = 0; i < V; i++)
if (ap[i] == true)
cout << i << " ";
}
// Driver program to test above function
int main()
 // Create graphs given in above diagrams
 cout << "\nArticulation points in first graph \n";Graph g1(5);</pre>
 g1.addEdge(1, 0);g1.addEdge(0, 2);g1.addEdge(2, 1);g1.addEdge(0, 3);g1.addEdge(3,
4);q1.AP();
 cout << "\nArticulation points in second graph \n";</pre>
 Graph g2(4);g2.addEdge(0, 1);g2.addEdge(1, 2);g2.addEdge(2, 3);g2.AP();
```

}

```
cout << "\nArticulation points in third graph \n";Graph g3(7);</pre>
 g3.addEdge(0, 1);g3.addEdge(1, 2);g3.addEdge(2, 0);g3.addEdge(1, 3);g3.addEdge(1, 4);
g3.addEdge(1, 6);g3.addEdge(3, 5);g3.addEdge(4, 5);g3.AP();
return 0;}
HIST2
int main(){
for (i=0 ; i< (1<<N) ; i++)
for (j=0 ; j < N ; j++)
      count[i][j]=fact[i][j] = 0;
      for (i = 0 ; i < N ; i++)
      {
            count[(1 << i)][i] = 1;
            fact[(1 << i)][i] = 2*A[i]+2;
      for (i = 3 ; i < (1 << N) ; i++)
      for(j = 0 ; j < N ; j + +)
      {
            if(i&(1<<j))
                  long long int temp max = 2*A[j]+2;
                  long long int c = 0;
                  for (k = 0 ; k < N ; k++)
                         if(((i - (1 << j)) & (1 << k)) != 0)
                         {//do something}
            fact[i][j] = temp max; }}}
JAVA Fast IO
BufferedReader fin = new BufferedReader (new InputStreamReader ( System . in ) ) ;
StringTokenizer toker;
String line ;
while ( ( line = fin . readLine ( ) . trim ( ) ) != null )
      toker = new StringTokenizer ( line ) ;
      int m = Integer . parseInt ( toker . nextToken ( ) );
      int n = Integer . parseInt ( toker . NextToken ( ) );
Longest Increasing Subsequence
int lis( int arr[], int n, int *max ref)
{
```

```
if(n == 1)
       return 1;
    int res, max ending here = 1; // length of LIS ending with arr[n-1]
    /* Recursively get all LIS ending with arr[0], arr[1] ... ar[n-2]. If
       arr[i-1] is smaller than arr[n-1], and max ending with arr[n-1] needs
       to be updated, then update it */
    for(int i = 1; i < n; i++)
           res = lis(arr, i, max ref);
        if (arr[i-1] < arr[n-1] \&\& res + 1 > max ending here)
           max ending here = res + 1;
    }
    // Compare max ending here with the overall max. And update the // overall max
if needed
    if (*max ref < max ending here)</pre>
       *max_ref = max_ending_here;
    // Return length of LIS ending with arr[n-1]
    return max ending here;}
// The wrapper function for lis()
int lis(int arr[], int n)
    // The max variable holds the result
    int max = 1;
    // The function lis() stores its result in max
   lis( arr, n, &max );
   return max;
```