

2. a) $\int_0^{\infty} \frac{\alpha x_0^\alpha}{x^{\alpha+1}} dx = \alpha x_0^\alpha \int_{x_0}^{\infty} x^{-\alpha-1} dx = \alpha x_0^\alpha \left(\frac{1}{-\alpha} x^{-\alpha} \right) \Big|_{x_0}^{\infty}$
 $= \alpha x_0^\alpha \left(\frac{1}{-\alpha} - \frac{1}{x_0^\alpha} \right)$
 $= \frac{\alpha x_0^\alpha}{-\alpha} \left(-\frac{1}{x_0^\alpha} \right) = -x_0^\alpha \left(-\frac{1}{x_0^\alpha} \right)$
 $= \boxed{1} \Rightarrow \text{valid PDF}$

$F(X \geq x) = \int_x^{\infty} \frac{\alpha x_0^\alpha}{x^{\alpha+1}} dx$
 $= \alpha x_0^\alpha \left(\frac{1}{-\alpha} x^{-\alpha} \right) \Big|_x^{\infty}$
 $= \frac{\alpha x_0^\alpha}{-\alpha} \left(\frac{1}{\infty} - \frac{1}{x^\alpha} \right) = -x_0^\alpha \left(-\frac{1}{x^\alpha} \right) = \left(\frac{x_0}{x} \right)^\alpha$
 $\boxed{F(X \geq x) = \left(\frac{x_0}{x} \right)^\alpha}$

c) $l_n(\alpha) = \log \left(\prod_{i=1}^n \frac{\alpha x_0^\alpha}{x_i^{\alpha+1}} \right) = \sum_{i=1}^n \log \left(\frac{\alpha x_0^\alpha}{x_i^{\alpha+1}} \right) = \sum_{i=1}^n (\log(\alpha x_0^\alpha) - \log(x_i^{\alpha+1}))$
 $l_n(\alpha) = \sum_{i=1}^n (\log \alpha + \alpha \log x_0 - (\alpha+1) \log x_i)$

$\frac{dl_n(\alpha)}{d\alpha} = \sum_{i=1}^n \left(\frac{1}{\alpha} + \log x_0 - \log x_i \right)$

$0 = \frac{n}{\alpha} + n \log x_0 - \sum_{i=1}^n \log x_i$
 $\frac{n}{\alpha} = \sum_{i=1}^n \log x_i - n \log x_0 \Rightarrow \hat{\alpha} = \frac{n}{\sum_{i=1}^n \log x_i - n \log x_0}$

3. a) $cl(G) = \frac{3 \cdot \# \text{ triangles}}{\# \text{ connected triples}}$; $d_i = 2k \forall i$

• # triangles per node, clockwise: choose 2 of the node's clockwise nbs within k hops
 \Rightarrow avoid double-counting

$$= \binom{k}{2} = \frac{k!}{2!(k-2)!} = \frac{k(k-1)}{2}$$

• # connected triplets per node: $\binom{2k}{2} = \frac{2k(2k-1)}{2}$

$$\Rightarrow cl(G) = \frac{3nk(k-1)}{2nk(2k-1)} = \frac{3(k-1)}{2(2k-1)} = \boxed{\frac{3k-3}{4k-2}}$$

b) • # triangles;
 - still $\frac{n k(k-1)}{2}$ original

new:

• can be made by connecting nodes between $k+1$ and $2k$ hops

away; prob that they're connected is $\frac{1}{2} n 2k \cdot p / \frac{n(n-1)}{2} \approx \frac{2kp}{n}$

• # triples;
 - still $\frac{2nk(2k-1)}{2}$ original

- each new edge creates triples w/all of the original $2k$ nodes @ either of its ends ($= (2k)^2$ triples)

• n possible, w/prob. p

$$\Rightarrow (2k)^2 np$$

however, as $n \rightarrow \infty$ this term goes to 0, so it's negligible.

$$\Rightarrow \# \Delta \approx \frac{nk(k-1)}{2}$$

- new edges also create pairs w/other new edges:

- l new edges $\rightarrow \binom{l}{2} = \frac{l(l-1)}{2}$ new triples

- expected # of new edges to a node is $\frac{2|E|p}{n} = \frac{2(\frac{2nk}{2}p)}{n} = 2kp$

\Rightarrow exp. new triples is $n \frac{2kp(2kp-1)}{2} \approx (2k)^2 p^2 n$

$$\Rightarrow cl(G) = \frac{\frac{1}{2}nk(k-1) \cdot 3}{\frac{1}{2}2nk(2k-1) + (2k)^2 np + \underbrace{\frac{1}{2}n(2k^2)p^2}_{\text{ignore}}} = \frac{\frac{3}{2}nk(k-1)}{nk(2k-1) + (2k)^2 np} = \frac{\frac{3}{2}nk(k-1)}{2nk^2 - nk + 4k^2 np}$$

$$= \frac{3k(k-1)}{2(2k^2 - k + 4k^2 p)}$$

$$= \frac{3k-3}{4k-2+8kp}$$

4. a) $\overset{i > m}{k_i(i) = M}$ (each node connects to m others @ birth; node i is born @ time i)

b) $\boxed{\frac{dk_i(t)}{dt} \approx \frac{M}{t}}$ (a ratio of $\frac{M}{t}$ nodes have new connections added from time $t \rightarrow t+1$)

$$k_i(t) = \int \frac{M}{t} dt = M \log t + c \rightarrow k_i(i) = m = m \log i + c$$

$$\Rightarrow \underline{c = M - m \log i}$$

c) $k_i(t) = m + m \log\left(\frac{t}{i}\right) \geq d \Rightarrow k_i(t) = m \log t + m - m \log i$

$$m \log\left(\frac{t}{i}\right) \geq d - m$$

$$\frac{t}{i} \geq e^{\frac{d-m}{m}}$$

$$i \leq \frac{t}{e^{\frac{d-m}{m}}} \rightarrow \text{t nodes by time t, so } \boxed{\bar{F}(d) = e^{\frac{(m-d)}{d}}}$$

$$= m + m \log\left(\frac{t}{i}\right) \quad \square$$

$$p(d) = -\frac{d \bar{F}(d)}{d d} = -\frac{d}{d d} e^{\frac{m-d}{d}} = -\frac{d}{d d} e^{\frac{m}{d} - 1} = (e^{\frac{m}{d} - 1}) \cdot -m d^{-2}$$

$$p(d) = \boxed{\frac{1}{m d^2} \cdot e^{\frac{m-d}{d}}}$$