## Week1 friday

**Review**: Determine whether each statement below about regular expressions over the alphabet  $\{a, b, c\}$  is true or false:

True or False:  $a \in L((a \cup b) \cup c)$ 

True or False:  $ab \in L((a \cup b)^*)$ 

True or False:  $ba \in L(a^*b^*)$ 

True or False:  $\varepsilon \in L(a \cup b \cup c)$ 

True or False:  $\varepsilon \in L((a \cup b)^*)$ 

True or False:  $\varepsilon \in L(a^*b^*)$ 

From the pre-class reading, pages 34-36: A deterministic finite automaton (DFA) is specified by  $M = (Q, \Sigma, \delta, q_0, F)$ . This 5-tuple is called the **formal definition** of the DFA. The DFA can also be represented by its state diagram: with nodes for the state, labelled edges specifying the transition function, and decorations on nodes denoting the start and accept states.

Finite set of states Q can be labelled by any collection of distinct names. Often we use default state labels  $q0, q1, \ldots$ 

The alphabet  $\Sigma$  determines the possible inputs to the automaton. Each input to the automaton is a string over  $\Sigma$ , and the automaton "processes" the input one symbol (or character) at a time.

The transition function  $\delta$  gives the next state of the DFA based on the current state of the machine and on the next input symbol.

The start state  $q_0$  is an element of Q. Each computation of the machine starts at the start state.

The accept (final) states F form a subset of the states of the DFA,  $F \subseteq Q$ . These states are used to flag if the machine accepts or rejects an input string.

The computation of a machine on an input string is a sequence of states in the machine, starting with the start state, determined by transitions of the machine as it reads successive input symbols.

The DFA M accepts the given input string exactly when the computation of M on the input string ends in an accept state. M rejects the given input string exactly when the computation of M on the input string ends in a nonaccept state, that is, a state that is not in F.

The language of M, L(M), is defined as the set of all strings that are each accepted by the machine M. Each string that is rejected by M is not in L(M). The language of M is also called the language recognized by M.

What is **finite** about all deterministic finite automata? (Select all that apply)

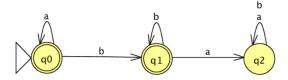
- ☐ The size of the machine (number of states, number of arrows)
- $\square$  The number of strings that are accepted by the machine
- $\square$  The length of each computation of the machine



The formal definition of this DFA is

Classify each string  $a, aa, ab, ba, bb, \varepsilon$  as accepted by the DFA or rejected by the DFA.

Why are these the only two options?



The language recognized by this DFA is



The language recognized by this DFA is