Week8 monday

Theorem: A	4_{TM} is	s not	Turing-o	decidable.
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Proof: Suppose towards a contradiction that there is a Turing machine that decides A_{TM} . We call this presumed machine M_{ATM} .

By assumption, for every Turing machine M and every string w

- If $w \in L(M)$, then the computation of M_{ATM} on $\langle M, w \rangle$ ______
- If $w \notin L(M)$, then the computation of M_{ATM} on $\langle M, w \rangle$ ______

Define a **new** Turing machine using the high-level description:

D = "On input $\langle M \rangle$, where M is a Turing machine:

- 1. Run M_{ATM} on $\langle M, \langle M \rangle \rangle$.
- 2. If M_{ATM} accepts, reject; if M_{ATM} rejects, accept."

Is D a Turing machine?

Is D a decider?

What is the result of the computation of D on $\langle D \rangle$?

