

Week4 wednesday

Language	$s \in L$	$s \notin L$	Is the language regular or nonregular?
$\{a^n b^n \mid 0 \leq n \leq 5\}$			
$\{b^n a^n \mid n \geq 2\}$			
$\{a^i b^j \mid 0 \leq i \leq j\}$			
$\{a^i b^j \mid i \geq j + 3, j \geq 0\}$			
$\{b^i a^j \mid i \geq 1, j \geq 3\}$			
$\{w \in \{a, b\}^* \mid w = w^{\mathcal{R}}\}$			
$\{ww^{\mathcal{R}} \mid w \in \{a, b\}^*\}$			

Regular sets are not the end of the story

- Many nice / simple / important sets are not regular
- Limitation of the finite-state automaton model: Can't "count", Can only remember finitely far into the past, Can't backtrack, Must make decisions in "real-time"
- We know actual computers are more powerful than this model...

The **next** model of computation. Idea: allow some memory of unbounded size. How?

- To generalize regular expressions: **context-free grammars**
- To generalize NFA: **Pushdown automata**, which is like an NFA with access to a stack: Number of states is fixed, number of entries in stack is unbounded. At each step (1) Transition to new state based on current state, letter read, and top letter of stack, then (2) (Possibly) push or pop a letter to (or from) top of stack. Accept a string iff there is some sequence of states and some sequence of stack contents which helps the PDA processes the entire input string and ends in an accepting state.



Trace the computation of this PDA on the input string 01.

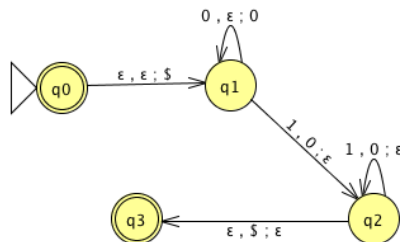
Trace the computation of this PDA on the input string 011.

Definition A **pushdown automaton** (PDA) is specified by a 6-tuple $(Q, \Sigma, \Gamma, \delta, q_0, F)$ where Q is the finite set of states, Σ is the input alphabet, Γ is the stack alphabet,

$$\delta : Q \times \Sigma_{\epsilon} \times \Gamma_{\epsilon} \rightarrow \mathcal{P}(Q \times \Gamma_{\epsilon})$$

is the transition function, $q_0 \in Q$ is the start state, $F \subseteq Q$ is the set of accept states.

Formal definition



Draw the state diagram of a PDA with $\Sigma = \Gamma$.

Draw the state diagram of a PDA with $\Sigma \cap \Gamma = \emptyset$.

A PDA recognizing the set $\{ \quad \}$ can be informally described as:

Read symbols from the input. As each 0 is read, push it onto the stack. As soon as 1s are seen, pop a 0 off the stack for each 1 read. If the stack becomes empty and there is exactly one 1 left to read, read that 1 and accept the input. If the stack becomes empty and there are either zero or more than one 1s left to read, or if the 1s are finished while the stack still contains 0s, or if any 0s appear in the input following 1s, reject the input.

State diagram for this PDA:

$$\Sigma = \{a, b, c\}. \quad L = \{a^i b^j c^k \mid i = j \text{ or } i = k, \ i \geq 0, j \geq 0, k \geq 0\}$$

Week4 friday

Consider the state diagram of a PDA with input alphabet Σ and stack alphabet Γ .

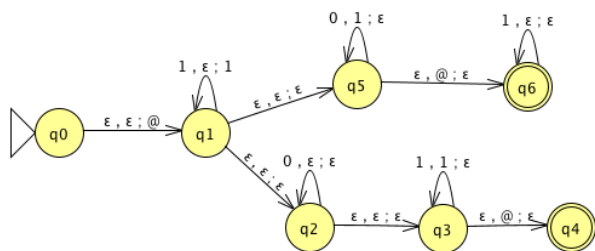
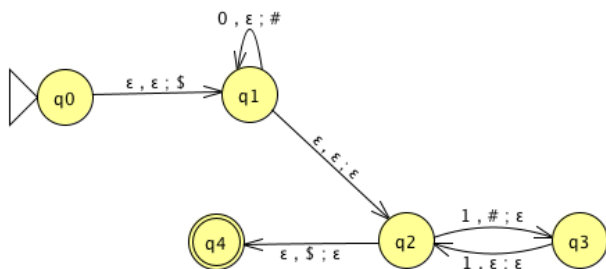
Label	means
$a, b \rightarrow c$ when $a \in \Sigma, b \in \Gamma, c \in \Gamma$	
$a, \varepsilon \rightarrow c$ when $a \in \Sigma, c \in \Gamma$	
$a, b \rightarrow \varepsilon$ when $a \in \Sigma, b \in \Gamma$	
$a, \varepsilon \rightarrow \varepsilon$ when $a \in \Sigma$	

How does the meaning change if a is replaced by ε ?

For the PDA state diagrams below, $\Sigma = \{0, 1\}$.

Mathematical description of language

State diagram of PDA recognizing language



$$\{0^i 1^j 0^k \mid i, j, k \geq 0\}$$

Assume $a \in \Sigma$ and L is a language over Σ . Recall that

$$aL = \{aw \mid w \in L\}$$

If $M = (Q, \Sigma, \Gamma, \delta, q_0, F)$ is a PDA with $L(M) = L$, a PDA M_1 that recognizes aL is

$$M_1 = (\quad, \Sigma, \Gamma, \delta_1, \quad, \quad)$$

with