Google Bigtable

(Bigtable: A Distributed Storage System for Structured Data)

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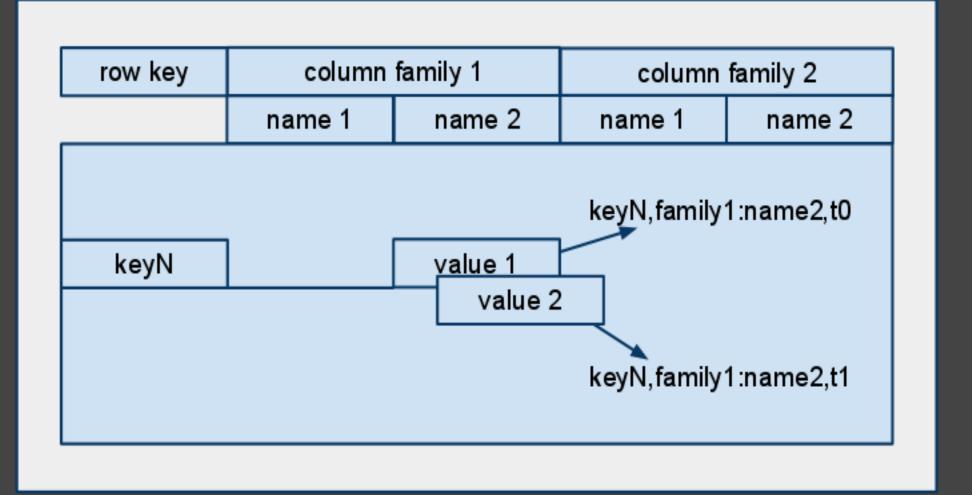
Presentation overview

- introduction
- design
- basic implementation
- GFS HDFS introduction
- MapReduce introduction
- implementation
- HBase Apache Bigtable solution
- performances and usage case
- some thoughts for discussion

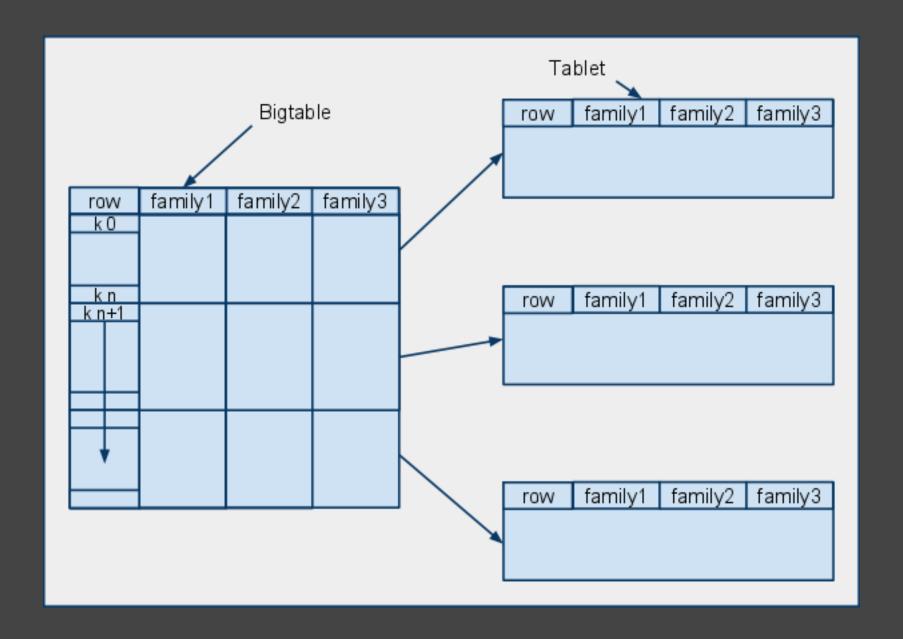
Introduction

- what is bigtable?
- why need for something more than SQL?
- machines number increase scaling
- fault tolerance
- row, column, cell
- versioning
- storage on distributed file system
- read, write operations on row or set of rows
- scanner iterator
- no need for indexing

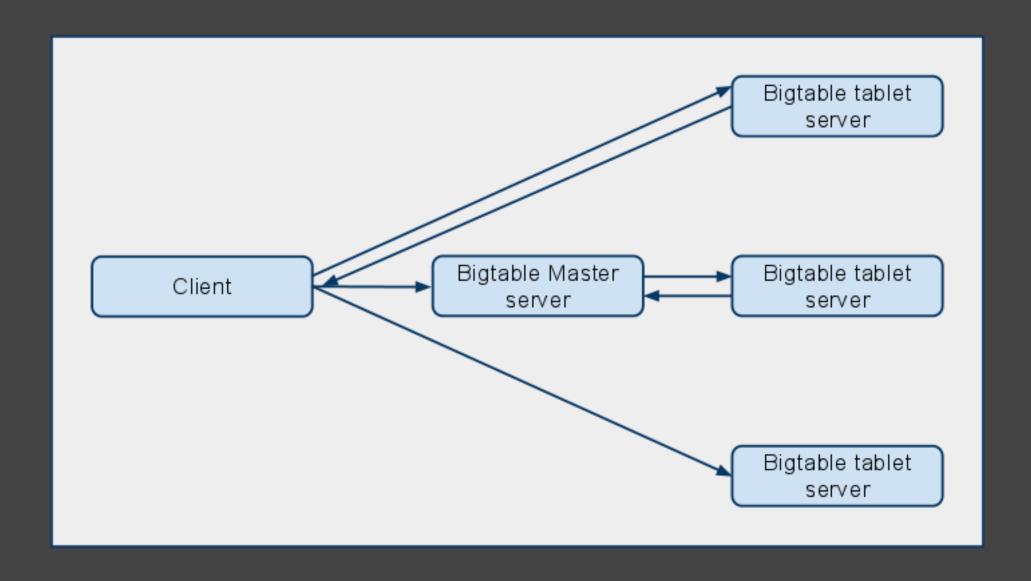
- column family basic unit of access control
- column family must be created before data insertion
- column labels number is unlimited
- row key is arbitrary string, max. size 64Kb
- every cell has copies in time configurable



- table is divided in tablets
- tablets are distributed on machines
- each tablet contains N rows
- tablet size 100 200 Mb
- tablet server splits tablet when it grows, only operation initiated by tablet server
- tablets are stored on GFS



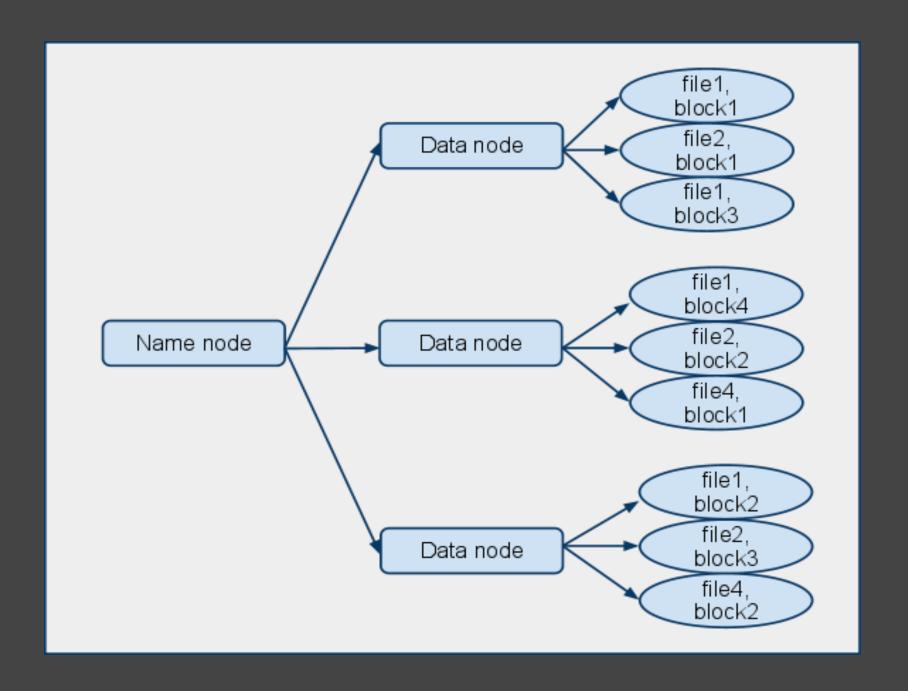
Basic implementation



GFS - HDFS introduction

- distributed file system
- files are divided in relatively large blocks, 64Mb
- blocks of one file are stored on multiple machines
- every block is replicated in configurable number of copies
- fault tolerance guaranteed
- Master Slaves architecture

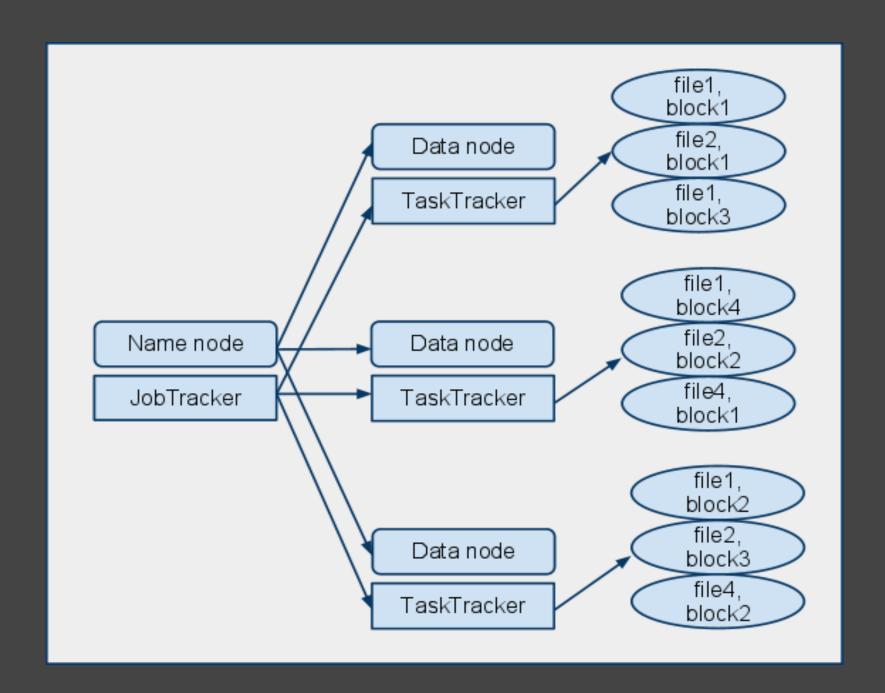
GFS - HDFS introduction: overview



MapReduce introduction

- framework build on top of (GFS) HDFS
- used for distributed data processing
- Master Slaves architecture
- between data elements there is no dependency
- everything is executed in parallel
- two parts: Mapper and Reducer
- Mapper reads input data and emits (key, value pairs) to Reducer
- Data on Reducer is grouped by key
- Reducer produces output data
- Combiner and Partitioner who are they, and what they want?

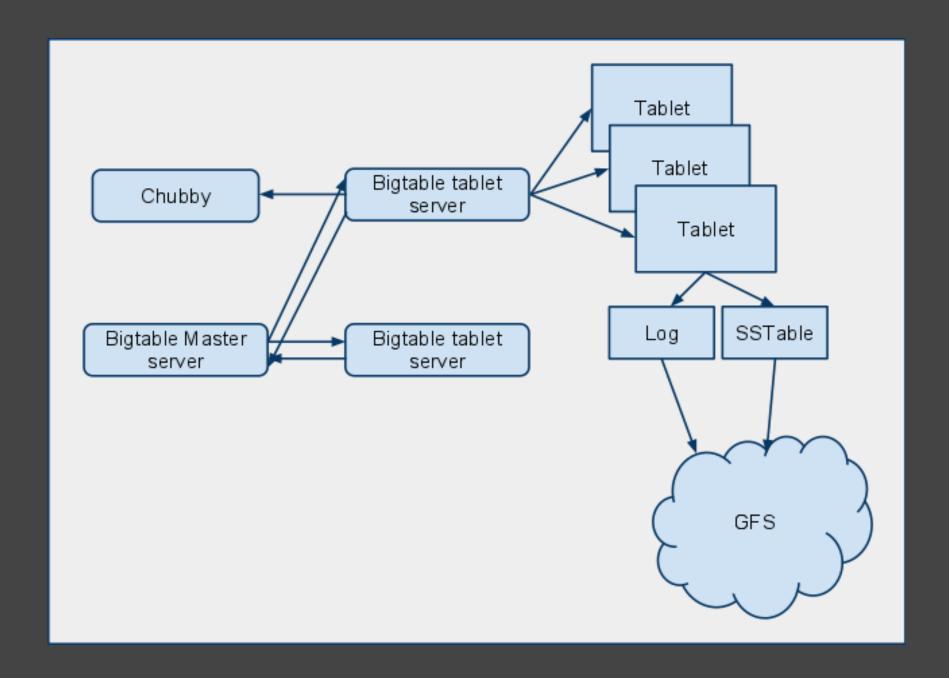
MapReduce introduction: overview



Implementation

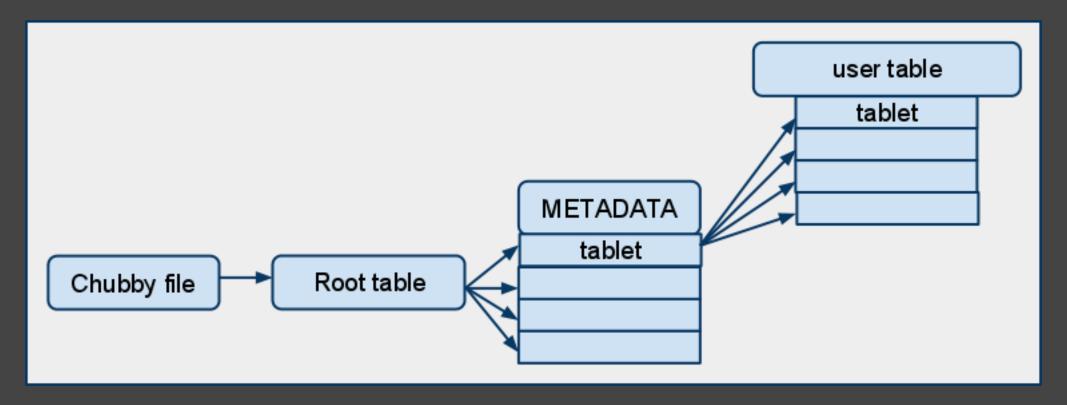
- Chubby distributed lock service, synchronization point, start point for master, tablet server and even clients
- Master server, only one, guaranteed by Chubby
- in case of Master tragical death tablet servers continue to serve client
- Master decides who serves each tablet (including ROOT and METATABLE)
- tablets are saved on GFS
- all trafic between clients and Bigtable is done directly within client and tablet server (Master is not required)
- client caches METATABLE, on first fault recache is done

Implementation: overview



Implementation: oragnization

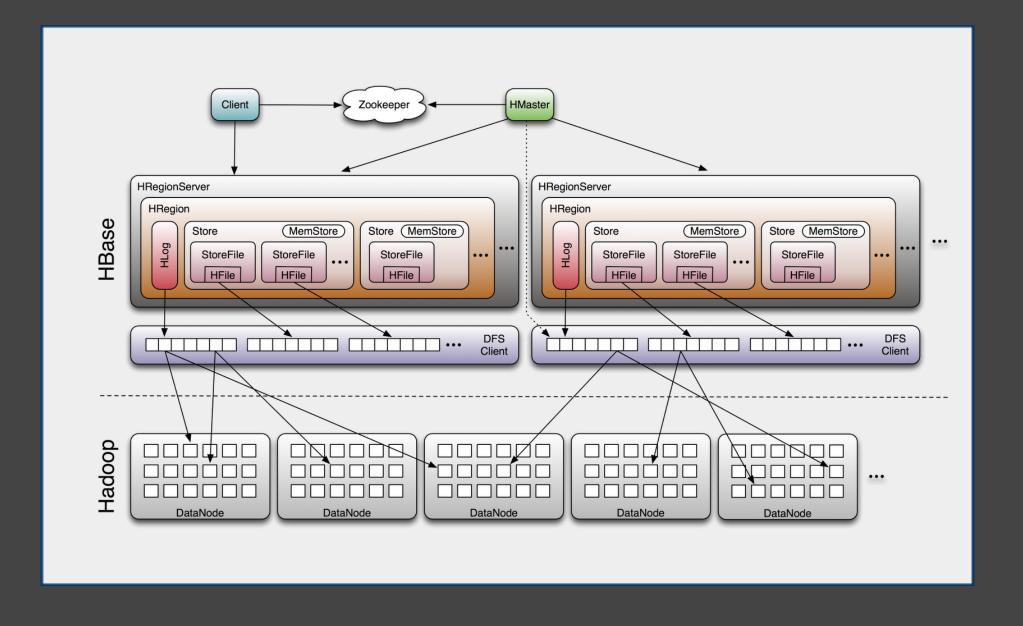
- two special tables: ROOT and METADATA
- ROOT never gets splitted
- location of ROOT is in Chubby file
- estimated ROOT size 128Mb => 2^34 tablets
- METADATA contains all user tablets, row key is an encoding of the tablet's table identifier and its end row



Implementation

- SSTable file representation of tablet
 - consist of blocks
 - column family labels are stored side by side
 - index of blocks at end of file
- compaction
 - minor
 - major
- locality groups
- caching

HBase - Apache Bigtable solution



HBase - Bigtable synonyms

- Bigtable ~ HBase
- tablet ~ region
- Master server ~ HBase master
- Tablet server ~ HBase Region server
- Chubby ~ Zookeeper ???
- GFS ~ HDFS
- MapReduce ~ MapReduce :)
- SSTable file ~ MapFile file

HBase differences

- MapFiles index is stored in separate file instead at end of file as in SSTable
- Unlike Bigtable which identifies a row range by the table name and end-key, HBase identifies a row range by the table name and start-key
- when the HBase master dies, the cluster will shut down

Performance and usage case

- number of 1000-byte values read/written per second
- table shows the rate per tablet server

Experiment	<u>1 TS</u>	<u>50 TS</u>	<u>250 TS</u>	<u>500 TS</u>
random reads	1212	593	479	241
random reads mem	10811	8511	8000	6250
random writes	8850	3745	3425	2000
sequential reads	4425	2463	2625	2469
sequential writes	8547	3623	2451	1905
scans	15385	10526	9524	7843

Performances and usage case

Scaling:

- performance of random reads from memory increases by almost a factor of 300 as the number of tablet server increases by a factor of 500
- significant drop in per-server throughput when going from 1 to 50 tablet servers
- the random read benchmark shows the worst scaling (an increase in aggregate throughput by only a factor of 100 for a 500-fold increase in number of servers)

Performances and usage case

Real users:

- Google Analytics:
- It provides aggregate statistics, such as the number of unique visitors per day and the page views per URL per day
- two tables raw user clicks (~200Tb, 14% compression) and summary (~20Tb, 29% compression)
- Google Earth:
- Google operates a collection of services that provide users with access to high-resolution satellite imagery of the world's surface
- two tables (~70Tb, imagery, no compression), (~500Gb, cache, served on hundreds of tablet servers

Cooked by Vast: RecordStack

Introduction:

- long-term archival storage for vertical data to HDFS
- optimized for later analytical processing.
- stores:
 - listings data
 - transactional log data
 - user clicks
 - leads
 - other reporting stuff

Cooked by Vast: RecordStack

Goals:

- Efficient archival of records to HDFS conserving space
- Checksuming records and tracking change
- Optmization of HDFS stored structure by periodical data compacting
- Simple access from Hadoop tasks for analytical processing
- Simple access via client application Remote console
- generation of reports

Some thoughts for discussion

- doing Bigtable with some number of "classic" databases?
- expose of Chubby to client, is it required?
- why not open source Bigtable?
- run of tablet servers on data nodes closer to data?

Useful links

- Scalability: http://scalability.rs/
- NoSQL Summer: http://nosqlsummer.org/
- Google Bigtable paper: http://tinyurl.com/37rlevv
- HBase Architecture: http://tinyurl.com/3y3fhbk
- HBase related blog: http://www.larsgeorge.com/
- Hadoop: The Definitive guide: http://tinyurl.com/36ponz3

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Thanks !!!