

Smart Slicing with Q.Blender

ROADMAP

- Part I: Q.Blender
- Part II: Smart Slicing in Q.Blender

Part I: Q.Blender



DUAL SIDED BUSINESS MODEL



ARCHITECTURE

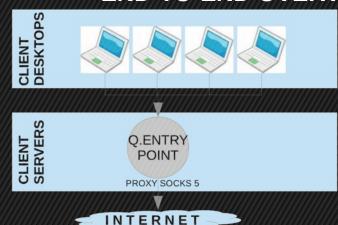
END-TO-END OVERVIEW

HPC CLIENT INFRASTRUCTURE

QARNOT INFRASTRUCTURE

DATA CENTER

Q.RADS HOSTS



Q.PKI
Q.MONITOR
Q.BILLING
Q.REPOSITORY

INTERNET

SITE #1

Q.BOX

Q.BOX

Q.RADS

Q.RADS



CLIENT APPLICATIONS

CUSTOM API CLIENT API

Q.API

QARNOT ENTRY POINT

Q.NETWORK

Q.NETWORK
Q.LINUX
(WINDOWS) CHROOT
CLIENT LIBS

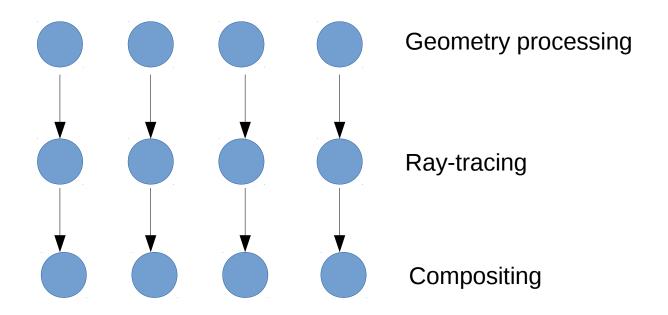
Python SDK for distributed rendering

```
import <u>qapy</u>
import <u>sys</u>
print("Loading config...")
api = gapy.QApy('garnot.conf')
print("Creating task...")
with api.create_task("My Blender Job", "blender", 1) as task:
     print("Sync resources from 'input' directory")
     task.resources.sync directory("input/")
     print("Setting constants...")
     task.constants['BLEND FILE'] = "model.blend"
     task.constants['BLEND_FORMAT'] = "PNG"
     task.constants['BLEND ENGINE'] = "CYCLES"
     task.constants['BLEND CYCLES SAMPLES'] = 1000
     print("Submitting task...")
     task.submit()
     print("Waiting for task completion...")
     task.wait()
     print("Retrieving results in 'output' directory")
     if not os.path.exists("output/"):
             os.makedirs("output/")
     task.download results("output/")
```



Part II: The smart slicing problem

Q.Blender parallel workflow



In Q.Blender we use a sort-first parallelization: the geometry processing task is duplicated between the nodes. In the ray-tracing phase, each node works on a sub-region of the image and the compositing task consists of joining the sub-images

Why sort-first?

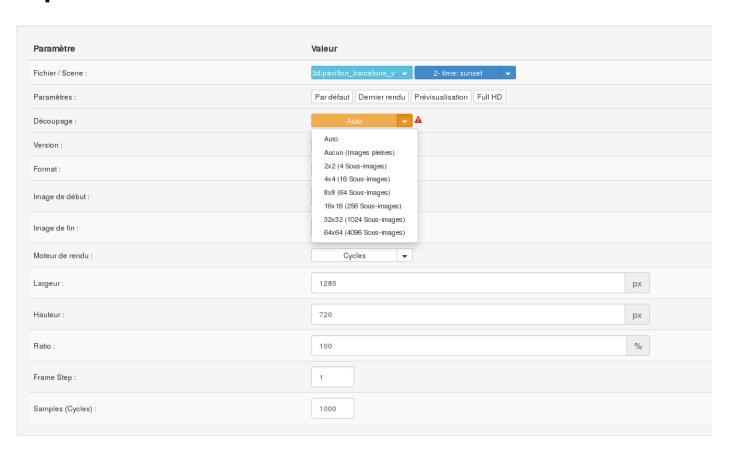
- Bad point: duplicated work in geometry processing
- Good point: We avoid expensive communication costs in a geo-distributed context

Creation of parallelism

- The decomposition in region is based on rectangle tiling
- Q.Blender supports various generic decompositions (2*2, 4*4)

Goal

Improve the « auto» mode



Computational Problem





$$t0 = 1516s$$
 $t1 = 1947s$
 $t2 = 408s$ $t3 = 713s$

T0 = 1947s

CPus: 4

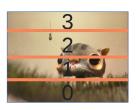
Rendering time of a region: ti



T1 = 1721s

Total rendering time:

$$Ti = max\{ti\}$$



$$t0 = 1698s$$
 $t1 = 1765s$
 $t2 = 788s$ $t3 = 333s$

T2 = 1765s

Question: How to tile an image for parallel processing such as to minimize Ti? Two sub-problems:

- What is the processing time on the different «regions» of an image ?
- How do we tile the image given the runtime of each region?

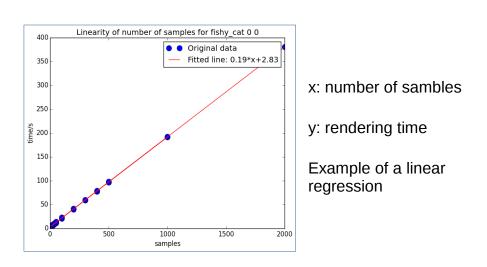
What impacts the load?

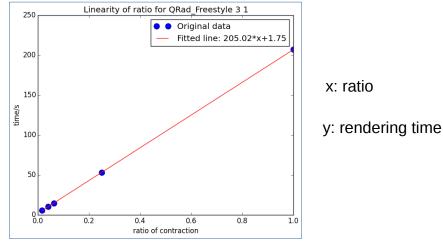
- Assumptions:
 - #Sample
 - Ratio
 - # cells (octree, polygons ..)
 - #frames
 - Pixel intensity ...

Load Estimation

Approach: Process the image with a small number of samples and ratio and project the measured runtime (we assume an initial fine and squared tiling in regions)

Main result: Linear regression of the runtime on the #number of samples and on the ratio



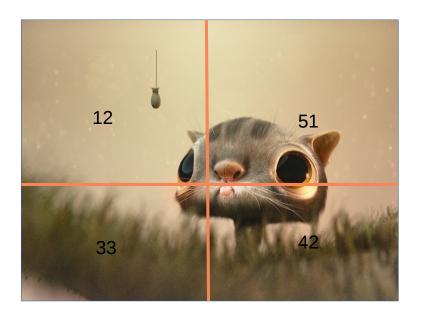


A regression was found for more than 96 images. The constant term is very small

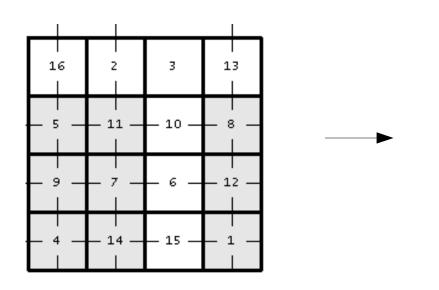
We considered more than 80 images.

Load estimation (2)

• With few samples, we can have a *good* estimation of the runtime difference between regions of the image



So What is the problem?

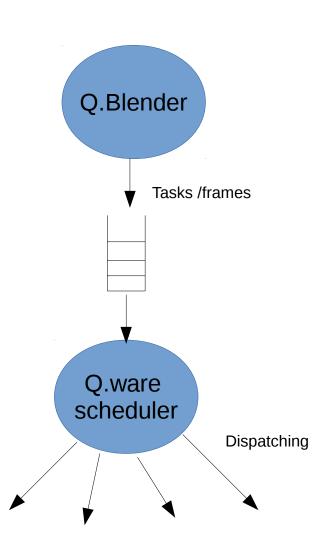


 Assuming a grid with g*g « unit » frames of different load, how to share this load between M processors?

A proposition

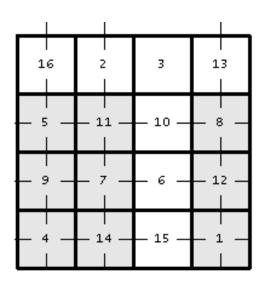
- Assign the « unit » frames such as to reduce imbalance
- M-partition problem, NP-hard but there are several good heuristics and semiexact algorithms

But...



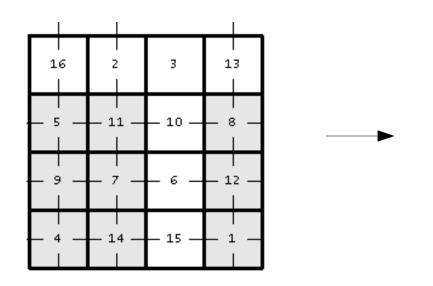
Resource oblivious SaaS logic: Q.Blender interacts with the general cloud resource manager: it cannot explicitely decide on the processor that will compute a given frame

So what can we do?



Make sure that the imbalance between the frames submitted for processing to the resource manager is not too high

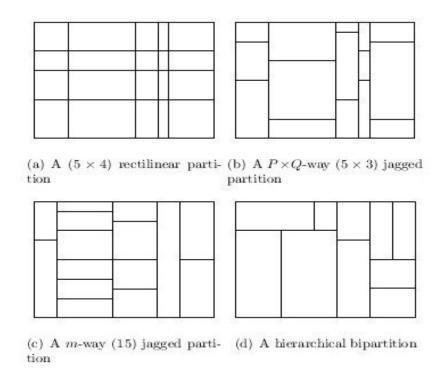
So What is finally the problem?



 Create M rectangles such as to minimize Lmax-Lmin. Here, Lmax is the load of the most heavy rectangle. Lmin is the load for the less heavy. M could be set as the number of processors.

Old problem

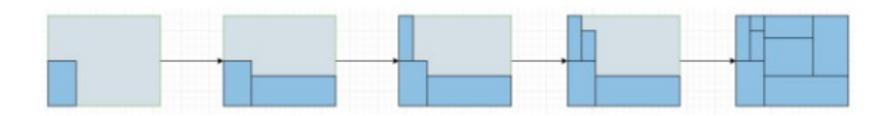
NP-hard. See https://arxiv.org/abs/1104.2566



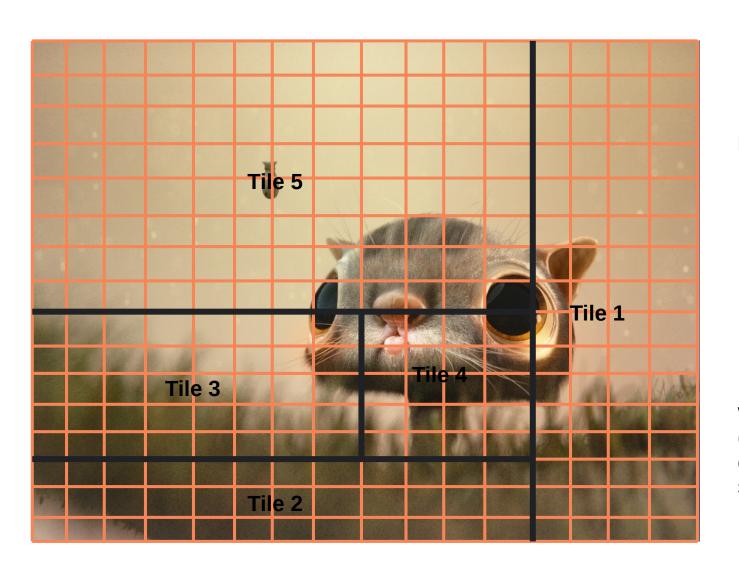
 Objective: a more efficient partitioning that only requires the number of rectangles as input.

Proposition

- A greedy «rectangle stacking » algorithm that iteratively chooses a rectangle whose load is close to the average of the area that was not already partitioned.
- Works like brick laying (layer by layer)
- No brick is placed in vaccum



Proposition(2)



Runtime

T1 = 1005s

T2 = 1011s

T3 = 996s

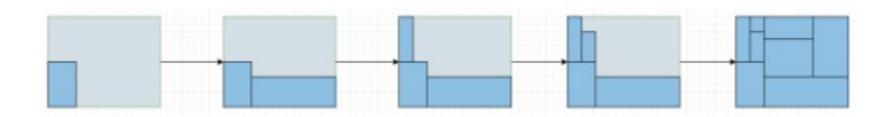
T4 = 952s

T5 = 971s

We gain more than 600 seconds over our traditional tiling solution.

Analysis

- The choice of the first tile is critical. It almost causes a 1D partitioning
- The rectangles become « more vertical » as we increase layers.
- The algorithm is not optimal



Open challenges

- Comparative analysis with related approaches
- Integration in the current «auto» mode
- Improvement of the rectangle stacking algorithm