

# Mathworks Candidate Presentation

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# Agenda

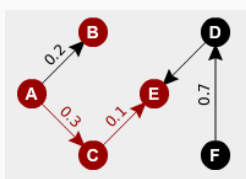
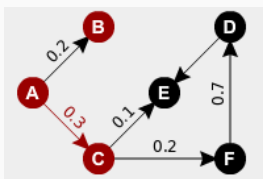
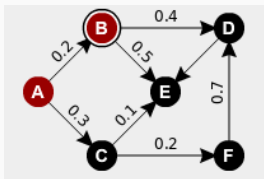
- Introduction
- C++ Project
- Another Project

# Introduction

- Education: PhD in Computer Science (December 2019)
- Experience:
  - Engineer (onsite contractor), Nortel (Avaya), Ireland. Verification and performance testing, telecom-software integrated system.
  - E-commercial PHP SWE, Lazada (Alibaba), Vietnam.
  - Intern SWE, Splice Machine, USA. SQL database over Hadoop and Spark.
  - Teaching/Research Assistant.
- Interests: Concurrency, distributed computing, compiler, functional programming.
- Matlab experience: basic engineering computing (course work).

# C++ Project: Graph Influence Problem

- Diffusion process over a graph.  $G(V, E)$ , edge  $(u, v) = w$ : some probability.
- $S$ : seed set
- Activation of neighbors by the edge probability, until no more new nodes. Size of the final set:  $I(S)$ , the *Influence* of  $S$ .



# Proposed Problem

- Influence Maximization with probabilistic guarantee.

Motivation: Expectation is limited, for example, low variance is also desired.

$$M_{\delta}(I(S)) = \max\{a \mid \Pr[I(S) \geq a] \geq \delta\}$$

Optimization problem MAXPROBINF:  $|S| = k$ , maximize  $M_{\delta}(I(S))$ .

- Multi-criteria Approximation with theoretical bound.
- Implementation: C++ with OpenMP.

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**Algorithm 1**  $k$ -Influence set by MULTICRITMDelta

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```
1: function MULTICRITMDelta( $G(V, E), k, \delta, n, N$ )
2:    $l \leftarrow 1$ 
3:    $h \leftarrow n$ 
4:   for  $step \leftarrow [1 \dots \log n]$  do
5:      $S \leftarrow \emptyset$ 
6:      $\lambda \leftarrow (l + h)/2$ 
7:      $C \leftarrow \delta\lambda$ , and  $F(S) \leftarrow F_\lambda(S)$ 
8:     while  $F(S) < C$  do
9:       Pick node  $j$  that minimizes  $\frac{w_j}{F(S \cup \{j\}) - F(S)}$ 
10:       $S \leftarrow S \cup \{j\}$ 
11:      feasible if  $|S| \leq k \ln(N\lambda)$  :  $l \leftarrow \lambda$ 
12:      infeasible if  $|S| > k \ln(N\lambda)$ :  $h \leftarrow \lambda$ 
return  $S$ 
```

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- Data structures + functions, organized by namespaces.
- Simple interface, no requirement for API reuse  $\implies$  No need for OO.
- Smart pointers and modern C++ container operations.
- Simple parallel computing: Each thread performs sampling, accumulates the measurements, update the shared total accumulation.

```

namespace datatypes {
    using LInt = long long int;
    using Vertex = std::tuple<LInt, LInt>;
    using Edge = std::tuple<LInt, LInt, double>;
    struct GraphByEdges {
        std::unique_ptr<std::vector<Vertex>> vertexes;
        std::unique_ptr<std::vector<Edge>> edges;
        GraphByEdges(
            std::unique_ptr<std::vector<Vertex>> &&pV,
            std::unique_ptr<std::vector<Edge>> &&pE)
            : vertexes(std::move(pV)), edges(std::move(pE)) {}
        GraphByEdges (const GraphByEdges&) = delete;
        GraphByEdges& operator= (const GraphByEdges&) = delete;
    };
    struct NodeIndexedCover {
        LInt cc_id;
        LInt cc_size;
    };
}

```



```

namespace graph {
    std::unique_ptr<datatypes::GraphByEdges> read_edges(...);
    std::unique_ptr<std::vector<datatypes::Node>>
        get_graph_nodes(...);
    std::unique_ptr<datatypes::Graph> sample_graph(..., const
        std::unique_ptr<util::Dice> &dice);
    std::unique_ptr<std::vector<std::unique_ptr<datatypes::
        ConnectedComp>>> connected_component(
        std::unique_ptr<datatypes::Graph>& graph);
    std::unique_ptr<std::vector<datatypes::NodeIndexedCover>>
        get_cover(
        std::unique_ptr<datatypes::Graph>& graph);
}

```

```

auto node_measure =
    make_unique<vector<NodeMeasure>>(nodes->size());
#pragma omp parallel for
for (int i = 0; i < num_threads; i++) {
    auto node_indexed_measure =
        make_unique<vector<LInt>>(nodes->size(), 0);
    for (int j = 0; j < batch_size; j++) {
        auto g =
            graph::sample_graph(graph_edges, nodes, dices[i]);
        auto temp = graph::connected_component(g);
        unique_ptr<vector<NodeIndexedCover>> nics =
            graph::get_cover(g);
        _update_node_measure(
            node_indexed_measure, nics, kset_ids);
    }
#pragma omp critical
    _gather_node_measure(node_indexed_measure, node_measure);
}

```

```

class Task {
public:
    CyclicBarrier* cb;
    std::mutex* update_lock;
    Task(CyclicBarrier* cb, std::mutex* update_lock)
        : cb(cb), update_lock(update_lock) {}
    void operator()(...) {...}
}

for (int j = 0; j < nthreads; j++) {
    Task task(cb, update_lock);
    threads.push_back(std::thread(task, ...));
}
for (auto& th : threads) { th.join(); }

```

```

class CyclicBarrier {
public:
    int const count;
    std::atomic<int> seats;
    std::atomic<int> gen;
    explicit CyclicBarrier(int count)
        : count(count), seats(count), gen(0) {}

    void wait() {
        int const cur = gen;
        if (--seats == 0) {
            seats = count; ++gen;
        } else {
            while (gen == cur) {
                std::this_thread::yield();
            }
        }
    }
};

```

- Actor: message-passing concurrent entities.
- Messages: Multiple writers, single consumer.
- Networking: Asynchronous TCP, (almost) transparent.
- Replace the low-level threads and event loops.
- Concurrency and distribution by actors and messages.

# Why not Actor in C++

Requirement: Using only Standard Library.

- C++: No thread pools.
- C++: Simulate *pattern matching* is difficult (via metaprogramming).
- C++: Serialization and asynchronous networking are difficult.

Other difficulties:

- Type safety: difficult to enforce message types for actors.
- Shared messages life cycle (one message sent to many actors).

A nice implementation: CAF: <https://actor-framework.org/>

**Thank you!**

**Q & A**