

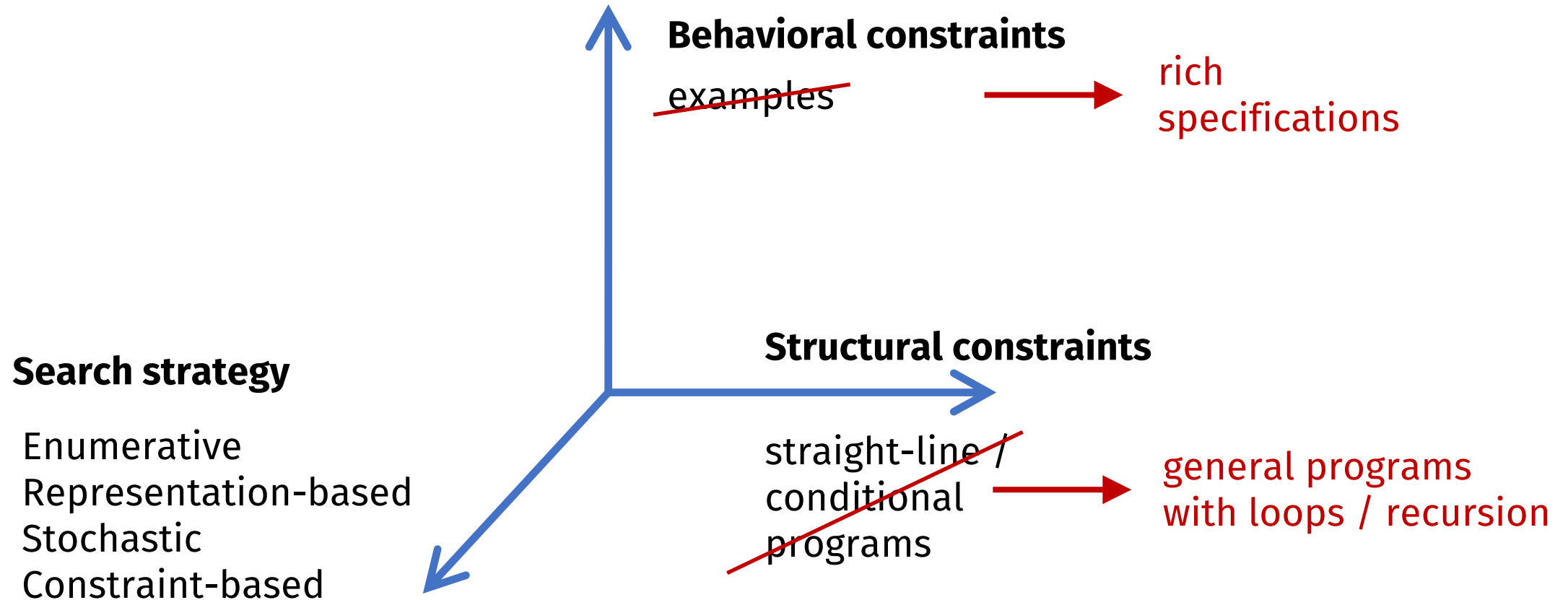
# #18: Specifications

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EECS 700: Introduction to Program Synthesis



# 3-axes of synthesis



# Examples of rich specifications

- Reference implementation
- Assertions
- Pre- and post-condition
- Fancy types

# Reference Implementation

Easy to compute the result, but hard to compute it efficiently or under structural constraints

```
bit[W] AES_round (bit[W] in, bit[W] rkey)
{
    ... // Transcribe NIST standard
}

bit[W] AES_round_sk (bit[W] in, bit[W] rkey) implements AES_round {
    ... // Sketch for table lookup
}
```

# Assertions

Hard to compute the result, but easy to check its properties

```
split_seconds (int totsec) {  
    int h := ?;  
    int m := ?;  
    int s := ?;  
    assert totsec == h*3600 + m*60 + s;  
    assert 0 <= h && 0 <= m < 60 && 0 <= s < 60;  
}
```

# Tests

Hard to specify results as inputs/outputs, but can assert postconditions about state of the program

```
post = Post.create(author: 'author', slug: 'hello-world', title: 'Hello World')
updated = update_post('author', 'hello-world', author: 'dummy',
                     title: 'Foo Bar', slug: 'foobar')

assert { updated.id == post.id }           # the returned post should be same
assert { updated.author == "author" }      # the author field should not be updated
assert { updated.title == "Foo Bar" }      # the title should be updated
assert { updated.slug == 'hello-world' }
```

# Pre-/post-conditions

Hard to compute the result; need correctness guarantees

sort (**int[]** in, **int** n) **returns** (**int[]** out)

**requires**  $n \geq 0$

**ensures**  $\forall i, j. 0 \leq i < j < n \Rightarrow out[i] \leq out[j]$

$\forall i. 0 \leq i < n \Rightarrow \exists j. 0 \leq j < n \wedge in[i] = out[j]$

{

**?**

}

# Refinement types

Same as pre-/post-conditions but logic goes inside the types

```
data RBT a where
  Empty :: RBT a
  Node  :: x: a ->
    black: Bool ->
    left:  { RBT {a | _v < x} | !black ==> isBlack _v } ->
    right: { RBT {a | x < _v} | (!black ==> isBlack _v) &&
      (blackHeight _v == blackHeight left) } ->
    RBT a

insert :: x: a -> t: RBT a -> {RBT a | elems _v == elems t + [x]}
insert = ?
```

binary search tree

red nodes have black children

same number of black nodes on every path to leaves



# Why go beyond examples?

- Might need too many
  - **Example:** Myth needs 12 for insert\_sorted, 24 for list\_n\_th
  - Examples contain *too little* information
  - Successful tools use domain-specific ranking
- Output difficult to construct
  - **Example:** AES cypher, RBT
  - Examples also contain *too much* information (concrete outputs)
- Need strong guarantees
  - **Example:** AES cypher
- Reasoning about non-functional properties
  - **Example:** security protocols

# Why is this hard?

```
gcd (int a, int b) returns (int c)
```

```
requires  $a > 0 \wedge b > 0$ 
```

```
ensures  $a \% c = 0 \wedge b \% c = 0$ 
```

```
 $\forall d . c < d \Rightarrow a \% d \neq 0 \vee b \% d \neq 0$ 
```

```
{
```

```
int x , y := a, b;
```

```
while (x != y) {
```

```
    if (x > y) x := ?;
```

```
    else y := ?;
```

```
}}
```

**infinitely many inputs**

cannot validate by testing

**infinitely many paths!**

hard to generate constraints

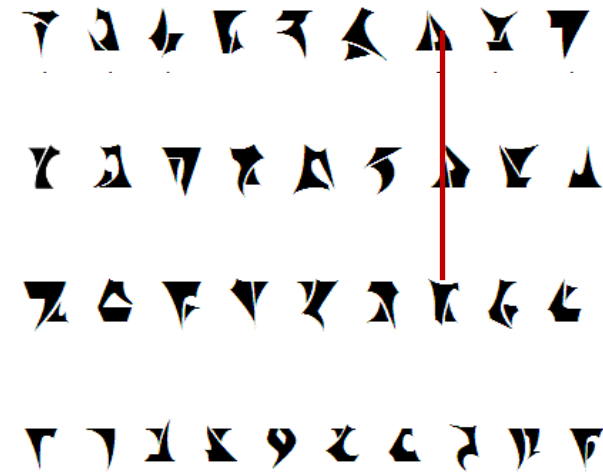
# Why is this hard?

Synthesis from examples



validation was easy!

Synthesis from specifications



SEE IF YOU CAN FIND ANY KLINGON FRUIT!

validation is hard!  
(and search is still hard)

# Upcoming lectures

