

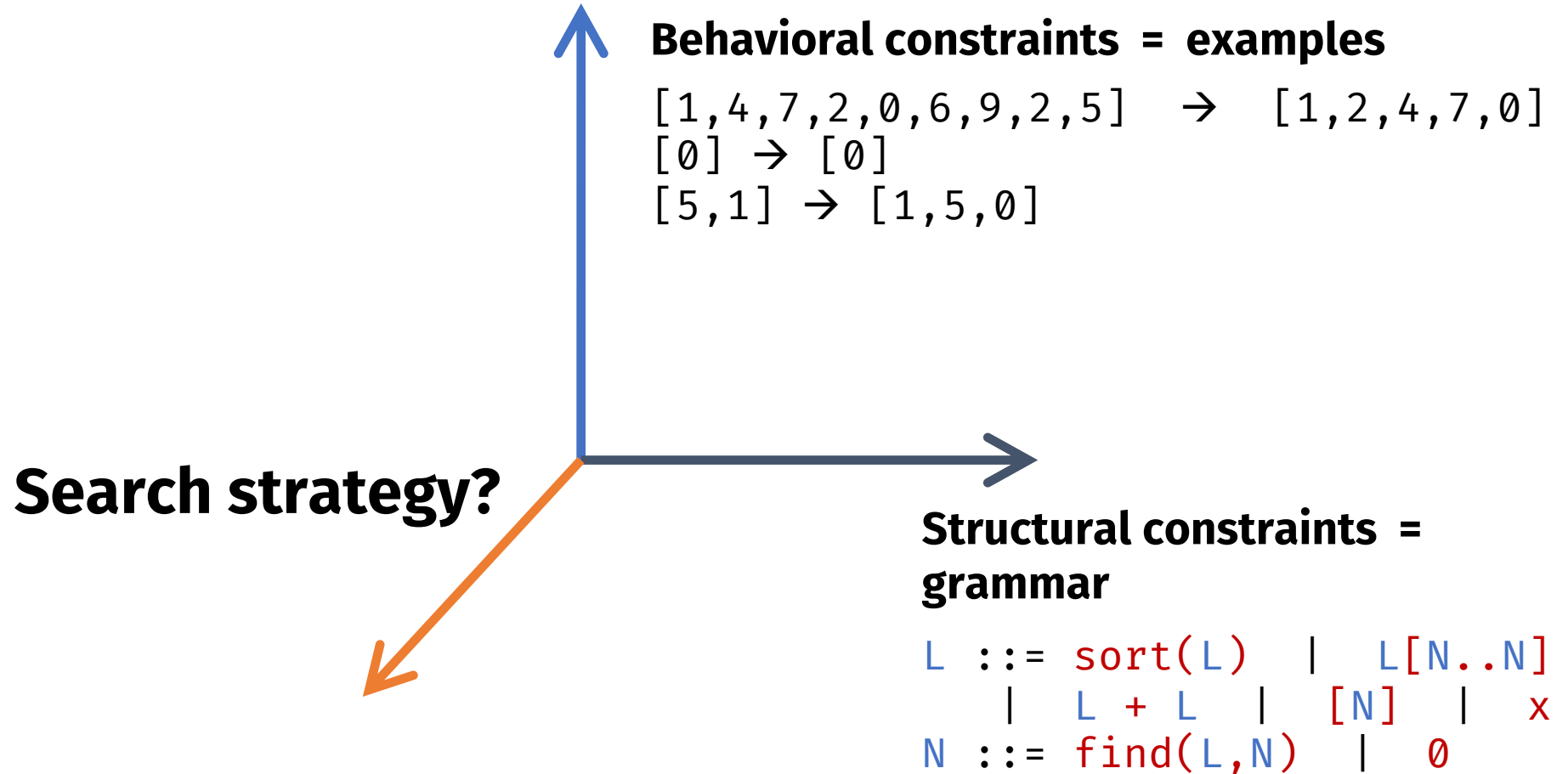
# #5: Top-down Propagation

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EECS 700: Introduction to Program Synthesis



# The problem statement



# Enumerative search

=

Explicit / Exhaustive Search

Idea: Enumerate programs from the grammar one by one and test them on the examples

bottom-up

$L ::= \text{sort}(L)$   
 $L[N..N]$   
 $L + L$   
 $[N]$   
 $x$   
 $N ::= \text{find}(L, N)$   
 $\emptyset$

top-down

$x$     $\emptyset$

$\text{sort}(x)$     $x[0..0]$     $x + x$     $[0]$

$\text{find}(x, \emptyset)$

$\text{sort}(\text{sort}(x))$     $\text{sort}(x[0..0])$

$\text{sort}(x + x)$     $\text{sort}([0])$

$x[0..\text{find}(x, \emptyset)]$    ...

$L$

$x$     $\text{sort}(L)$     $L[N..N]$     $L + L$     $[N]$

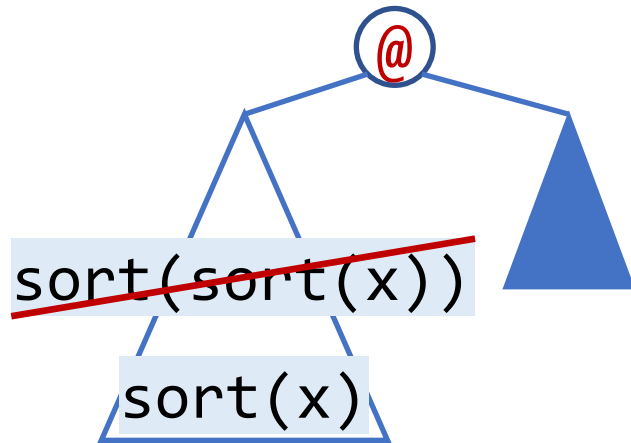
$\text{sort}(x)$     $\text{sort}(\text{sort}(L))$     $\text{sort}([N])$

$\text{sort}(L[N..N])$     $\text{sort}(L + L)$

$x[N..N]$     $(\text{sort } L)[N..N]$    ...

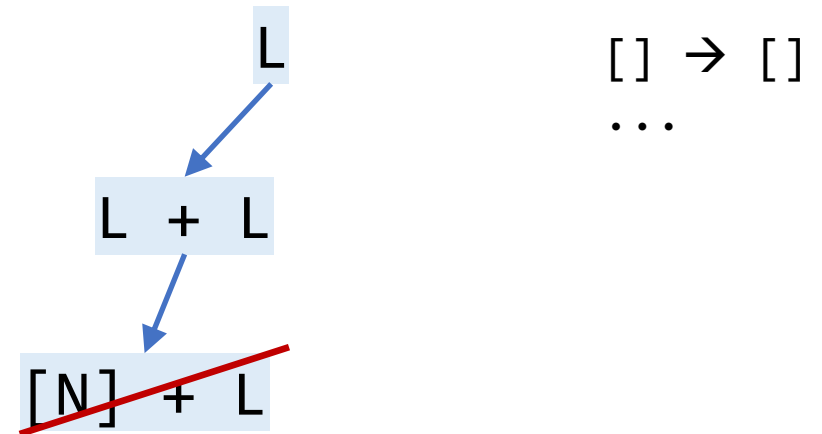
# When can we discard a subprogram?

redundant



**Equivalence reduction**  
(also: symmetry breaking)

infeasible




**Top-down propagation**

# Top-down search: reminder

generates a lot of incomplete terms  
only discards complete terms


iter 0: L

iter 1:  x L[N..N]

iter 2: L[N..N]


iter 3: x[N..N] L[N..N][N..N]

iter 4: x[0..N] L[N..N][N..N]

iter 5:  x[0..0] x[0.. find(L,N)] x[find(L,N)..N] ...

iter 6: x[0.. find(L,N)] x[find(L,N)..N] ...

iter 7: x[0.. find(x,N)] x[0.. find(L[N..N],N)] ...

iter 8:  x[0.. find(x,0)] x[0.. find(x,find(L,N))] ...

iter 9:

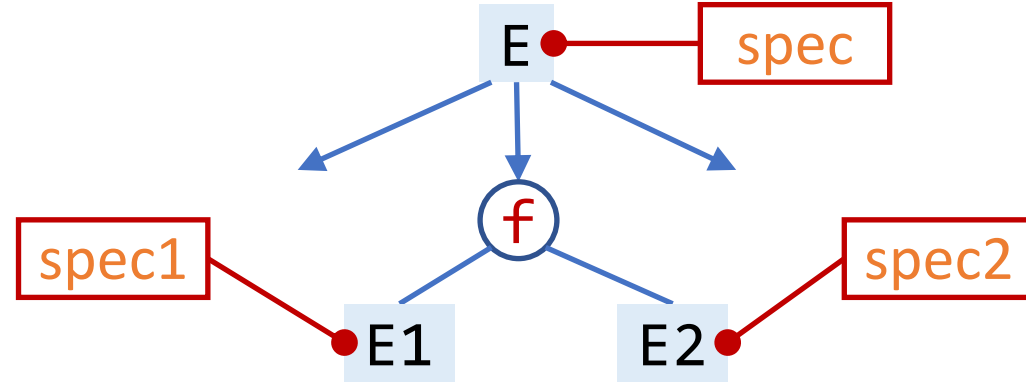
need to reject hopeless programs early!

```
L ::= L[N..N] |  
      x  
N ::= find(L,N) |  
      0
```

`[[1,4,0,6]]`  $\rightarrow$  `[1,4]`

# Top-down propagation

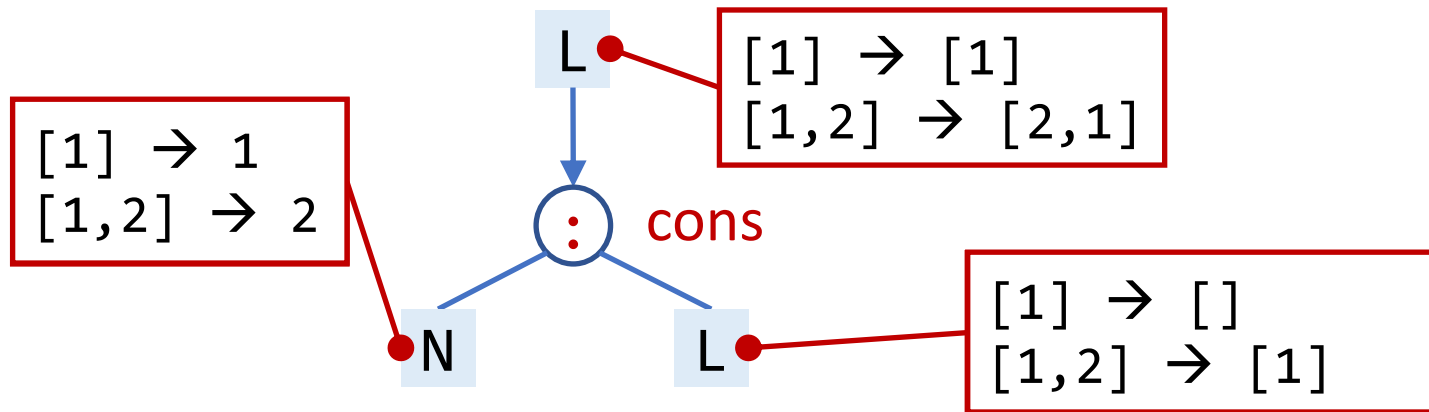
- **Idea:** once we pick the production, infer specs for subprograms



- If  $\text{spec1} = \perp$  or  $\text{spec2} = \perp$  discard  $f(E1, E2)$ !
- For now:  $\text{spec} = \text{examples}$

# When is TDP possible?

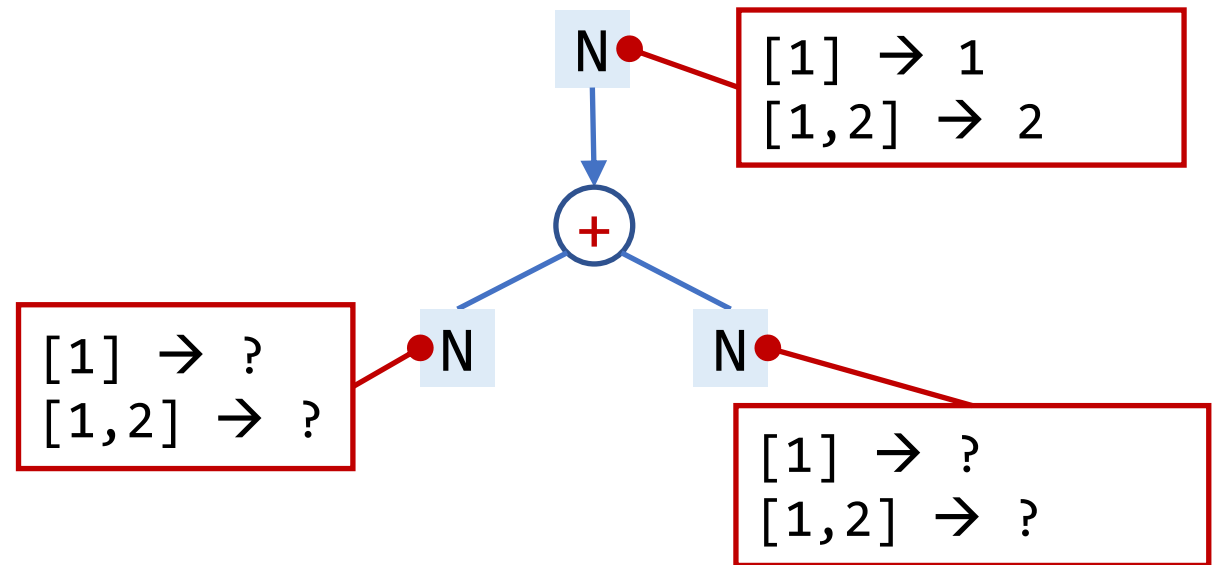
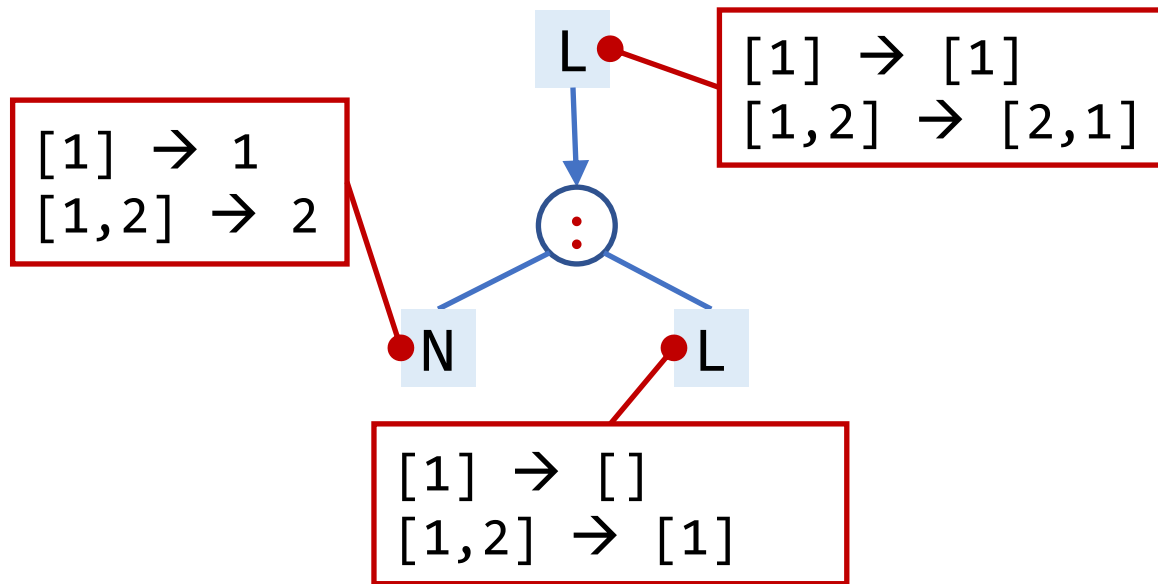
Depends on **f**!





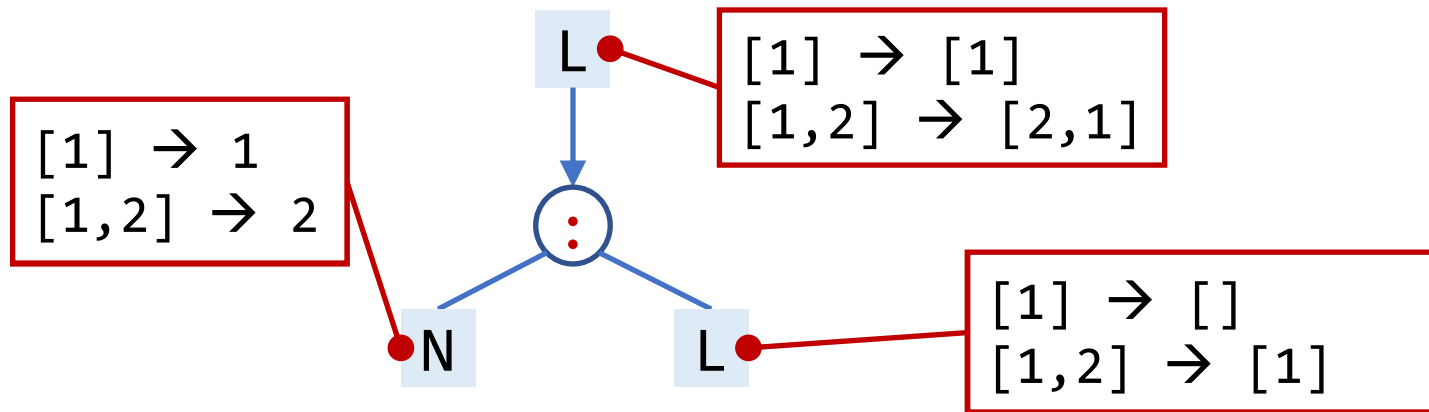
# When is TDP possible?

Depends on **f**!



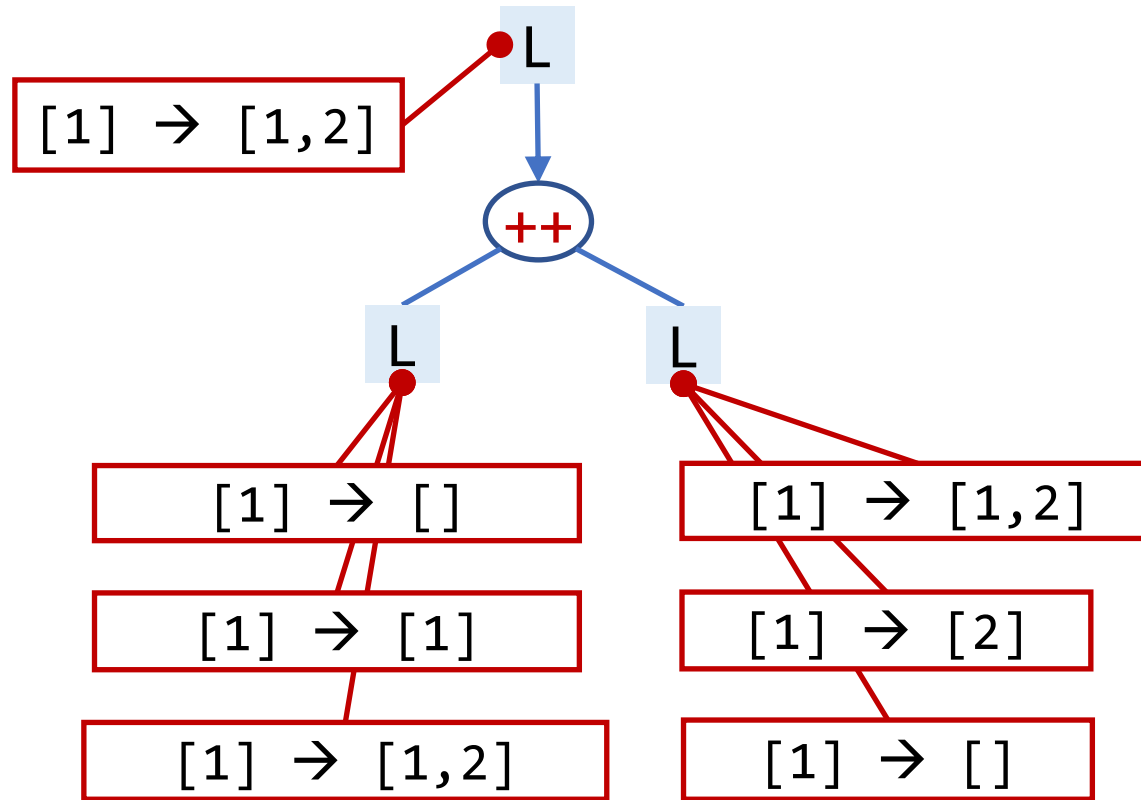
# When is TDP possible?

- Depends on **f**!



- Works when the function is injective!
- **Q:** when would we infer  $\perp$ ? **A:** If at least one of the outputs is  $[]$ !

# Something in between?



Works when the function has a  
“small inverse”

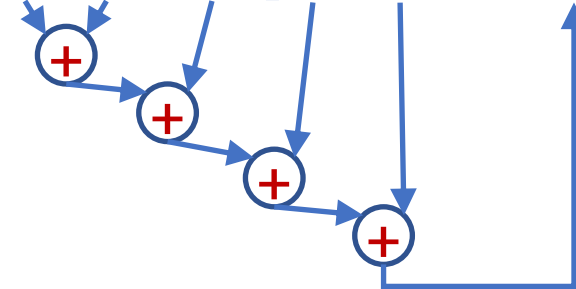
- or just the output examples have a small inverse

# $\lambda^2$ : TDP for list combinators

• map  $f$   $x$       map  $(\backslash y . y + 1)$   $[1, -3, 1, 7] \rightarrow [2, -2, 2, 8]$

• filter  $f$   $x$       filter  $(\backslash y . y > 0)$   $[1, -3, 1, 7] \rightarrow [1, 1, 7]$

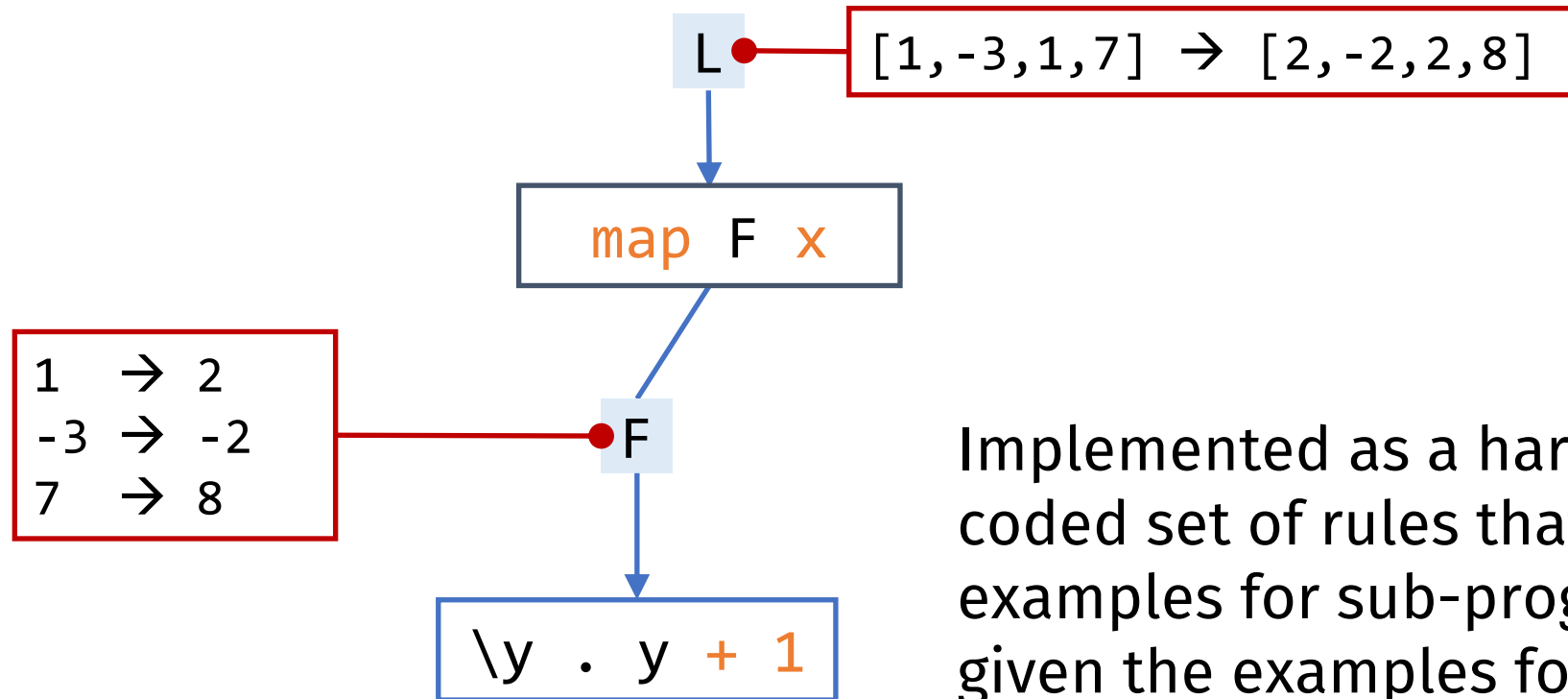
• fold  $f$  acc  $x$       fold  $(\backslash y z . y + z)$   $0$   $[1, -3, 1, 7] \rightarrow 6$



fold  $(\backslash y z . y + z)$   $0$   $[] \rightarrow 0$

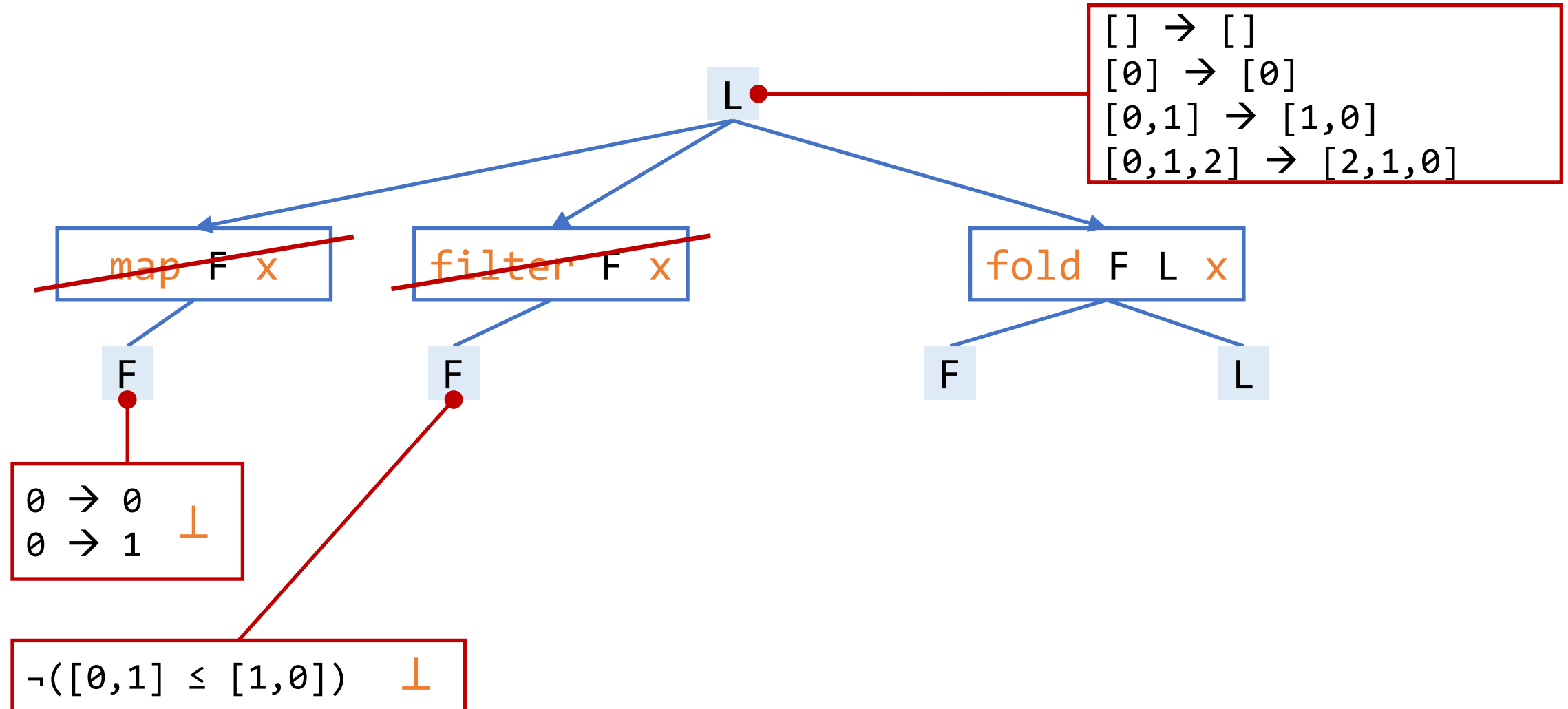


# $\lambda^2$ : TDP for list combinators



Implemented as a hard-coded set of rules that derive examples for sub-program(s) given the examples for the whole program and the combinator

# $\lambda^2$ : TDP for list combinators



# $\lambda^2$ : TDP for list combinators

fold F L []  $\rightarrow$  []

fold F [] [0]  $\rightarrow$  [0]

fold F [] [0,1]  $\rightarrow$  [1,0]

fold F [] [0,1,2]  $\rightarrow$  [2,1,0]

L

[ ]  $\rightarrow$  [ ]  
[0]  $\rightarrow$  [0]  
[0,1]  $\rightarrow$  [1,0]  
[0,1,2]  $\rightarrow$  [2,1,0]

fold F L x

F

L

<>  $\rightarrow$  [ ]

[ ]

<[ ], 0>  $\rightarrow$  [0]  
<[0], 1>  $\rightarrow$  [1,0]  
<[1,0], 2>  $\rightarrow$  [2,1,0]

\y z. z : y

# Condition abduction

- Smart way to synthesize conditionals
- Used in many tools (under different names):
  - **FlashFill** [Gulwani '11]
  - **Escher** [Albarghouthi et al. '13]
  - **Leon** [Kneuss et al. '13]
  - **Synquid** [Polikarpova et al. '16]
  - **EUSolver** [Alur et al. '17]
- In fact, an instance of TDP!



# Condition abduction

