#29: Abstract Interpretation- Guided Synthesis

Sankha Narayan Guria

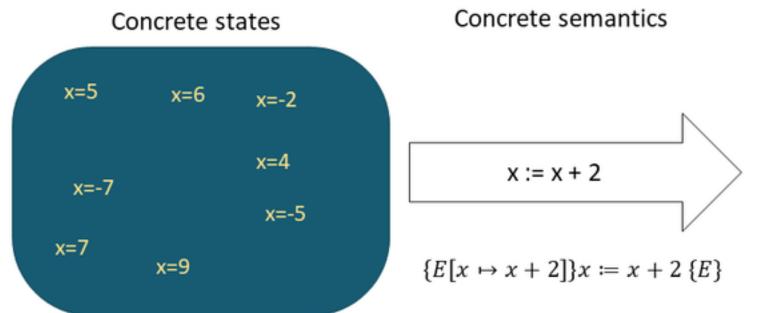
EECS 700: Introduction to Program Synthesis



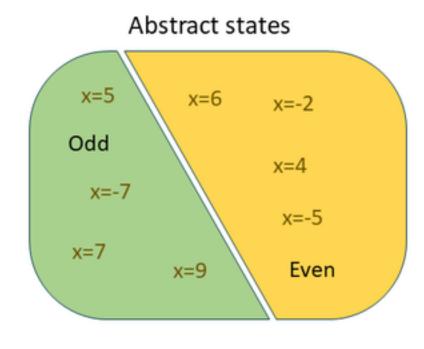
Today

- Synthesizing data-structure manipulation from storyboards
 - Rishabh Singh, Armando Solar-Lezama
- Absynthe: Abstract Interpretation-Guided Synthesis
 - Sankha Narayan Guria, Jeff Foster, David Van Horn

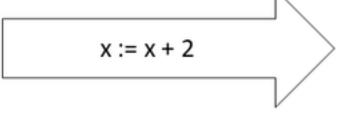
Key idea 1: Abstract domain



Key idea 1: Abstract domain



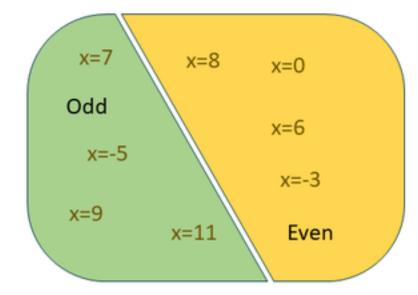
Abstract semantics



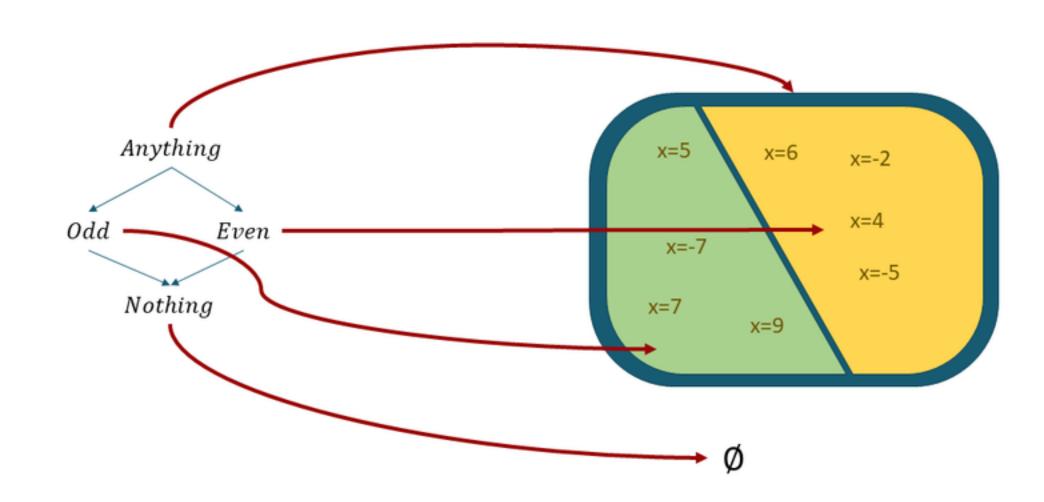
$$Odd + Odd = Even$$

 $Odd + Even = Odd$
 $Even + Even = Even$
 $Even + Odd = Odd$

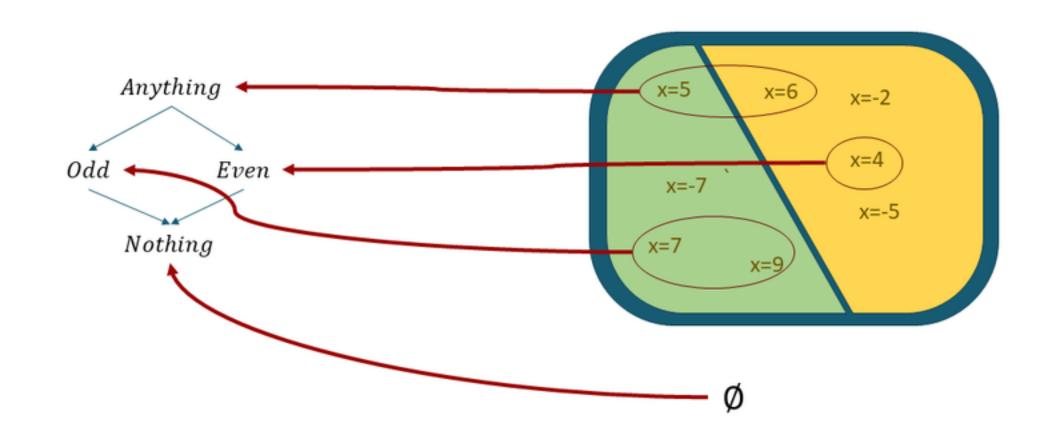
New abstract states



Concretization



Abstraction



Key idea 2: Abstract Interpretation

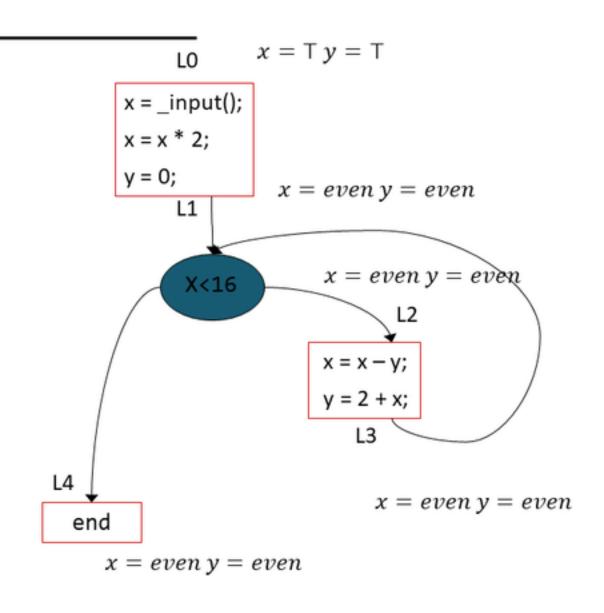
Compute an abstract value for every program point

Abstraction of the set of states possible at that point

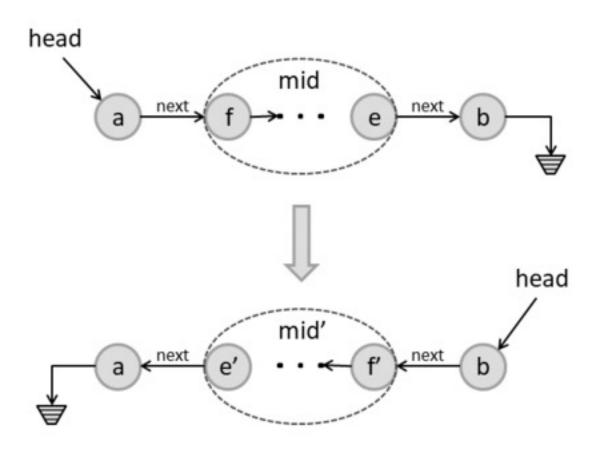
Iterate until computation converges

```
x = T y = T
                                                      LO
                                                x = _input();
                                                x = x * 2;
LO
                                                y = 0;
        x = _input();
                                                      L1
        x = x * 2;
        y = 0;
                                                      X<16
L1
        while(x < 16){
L2
         x = x - y;
                                                                     x = x - y;
         y = 2 + x;
L3
                                                                     y = 2 + x;
                                                                        L3
L4
                                         L4
                                           end
```

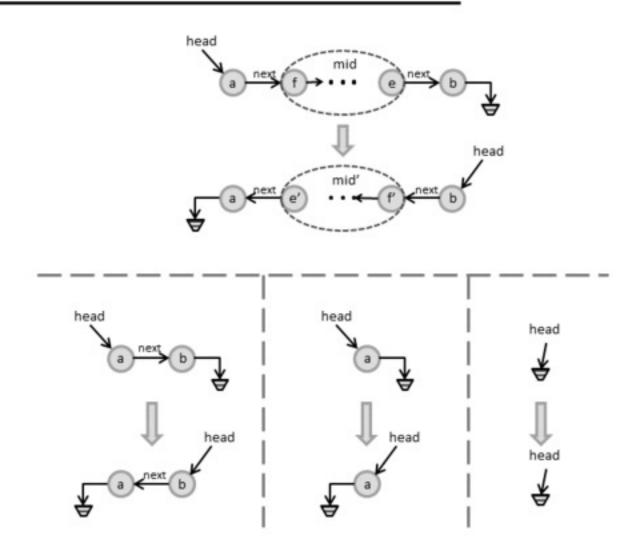
```
LO
        x = _input();
        x = x * 2;
        y = 0;
L1
        while(x < 16){
L2
          x = x - y;
          y = 2 + x;
L3
L4
```



Storyboard Programming

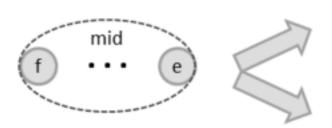


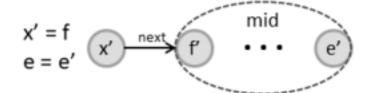
Scenarios for LL-reversal



Inductive insights with fold/unfold

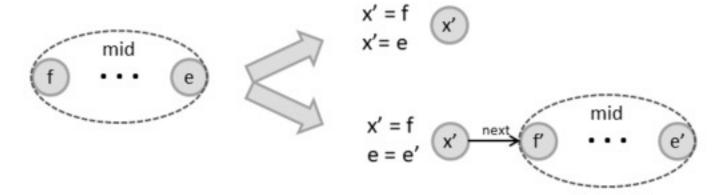
Unfold:



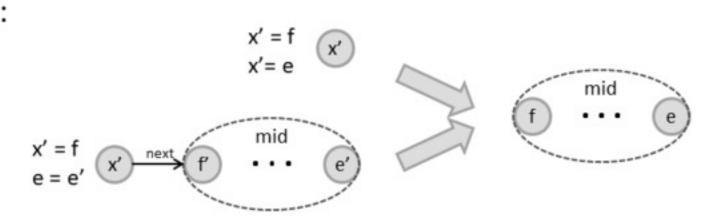


Inductive insights with fold/unfold

Unfold:



Fold:



Concrete Domain

Memory locations: $\mathcal{L}^{\#}$

Variables: $v_0, v_1, \dots v_k$

Variable predicates: v_i : $\mathcal{L}^\# \to \text{Bool } v_i(l)$ indicates that variable v_i points to loc l

Fields: sel_0 , sel_1 , ... sel_k

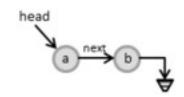
Field predicates : sel_0 : $\mathcal{L}^\# \times \mathcal{L}^\# \to Bool$

 $sel_i(l_1,l_2)$ indicates that there is a field sel_i from object l_1 to object l_2

$$\mathcal{L}^{\#} = \{a, b\}$$

head(a) = truehead(b) = false

Next	а	b
a	false	true
\boldsymbol{b}	false	false



Abstract Domain

Abstract memory locations: \mathcal{L}

represents a set of concrete locations

Summary location indicator: $sm: \mathcal{L} \rightarrow Tree\ Valued\ Logic\ (TVL)$

indicates if a location represents more than one concrete loc

Attachment Points: $A: \mathcal{L} \to \{\mathcal{L}\}$

maps a summary node to a set of locations that serve as attachment points

Variable predicates: $v_i : \mathcal{L} \to \text{TVL } v_i(l)$ indicates that variable v_i points to loc l

Field predicates : $sel_0: \mathcal{L} \times \mathcal{L} \rightarrow TVL$

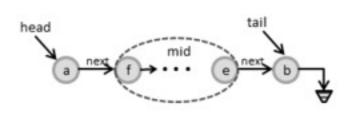
 $sel_i(l_1,l_2)$ indicates that there is a field sel_i from object l_1 to object l_2

 $\mathcal{L} = \{a, f, e, mid, b\}$

sm		\mathcal{A}	
а	false		
f	false		
e	false		
mid	true	mid	{f,e}
b	false		

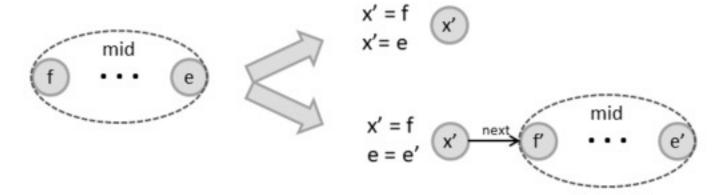
head	
а	true
f	false
e	false
mid	false
b	false

tail		next	а	f	e	mie	b
а	false	а	F	T	1	1	F
f	false	f	F	F	F	1	F
е	false	e	F	F	F	F	T
mid	false	mid	F	F	/	F	/
b	true	b	F	F	F	F	F

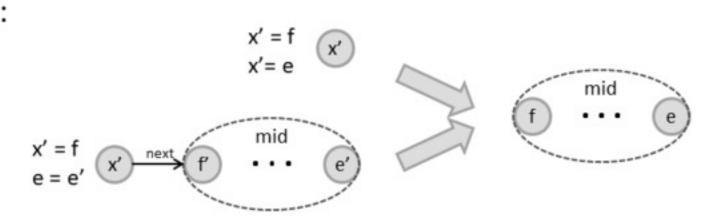


Inductive insights with fold/unfold

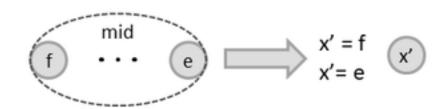
Unfold:



Fold:



Unfold



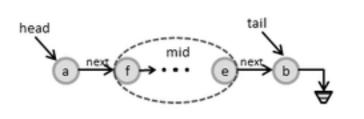
sm	
а	false
f	false
е	false
mid	true
b	false

\mathcal{A}	
mid	{f,e}

head	
а	true
f	false
e	false
mid	false
b	false

tail	
а	false
f	false
e	false
mid	false
b	true

next	а	f	e	mie	b
а	F	T	/	/	F
f	F	F	F	/	F
e	F	F	F	F	Т
mid	F	F	/	F	/
b	F	F	F	F	F





Unfold(head.next)



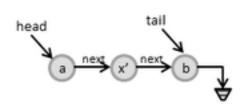
sm	
а	false
x'	false
b	false

\mathcal{A}	

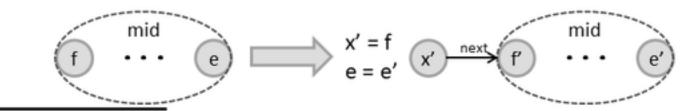
head	
а	true
x'	false
b	false

tail	
а	false
x'	false
b	true

ı	next	а	x'	b
	а	F	T	F
	x'	F	F	Т
	b	F	F	F



Unfold



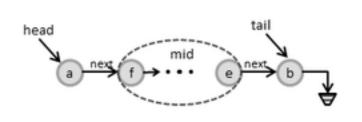
sm	
а	false
f	false
e	false
mid	true
b	false

\mathcal{A}	
mid	$\{f,e\}$

head	
а	true
f	false
е	false
mid	false
b	false

tail	
а	false
f	false
e	false
mid	false
b	true

next	а	f	e	mie	b
а	F	Т	/	/	F
f	F	F	F	/	F
e	F	F	F	F	Т
mid	F	F	/	F	/
b	F	F	F	F	F





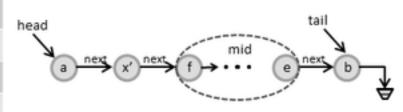
Unfold(head.next)



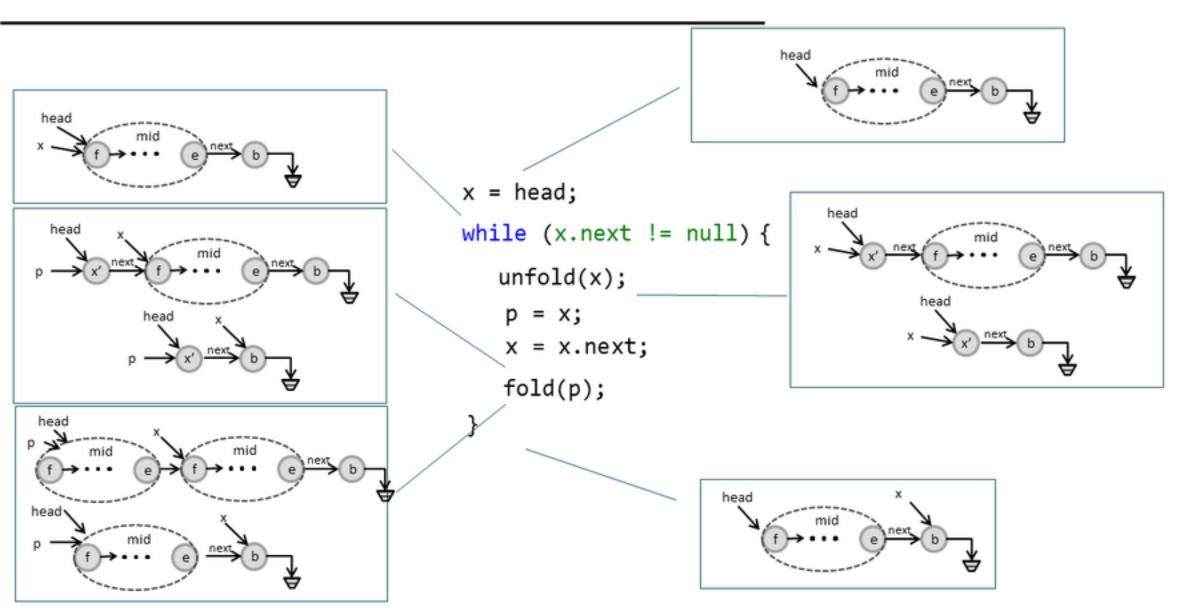
		4	
sm		\mathcal{A}	
а	false		
f	false		
e	false		
x'	false		
mid	true	mid	{f,e}
b	false		

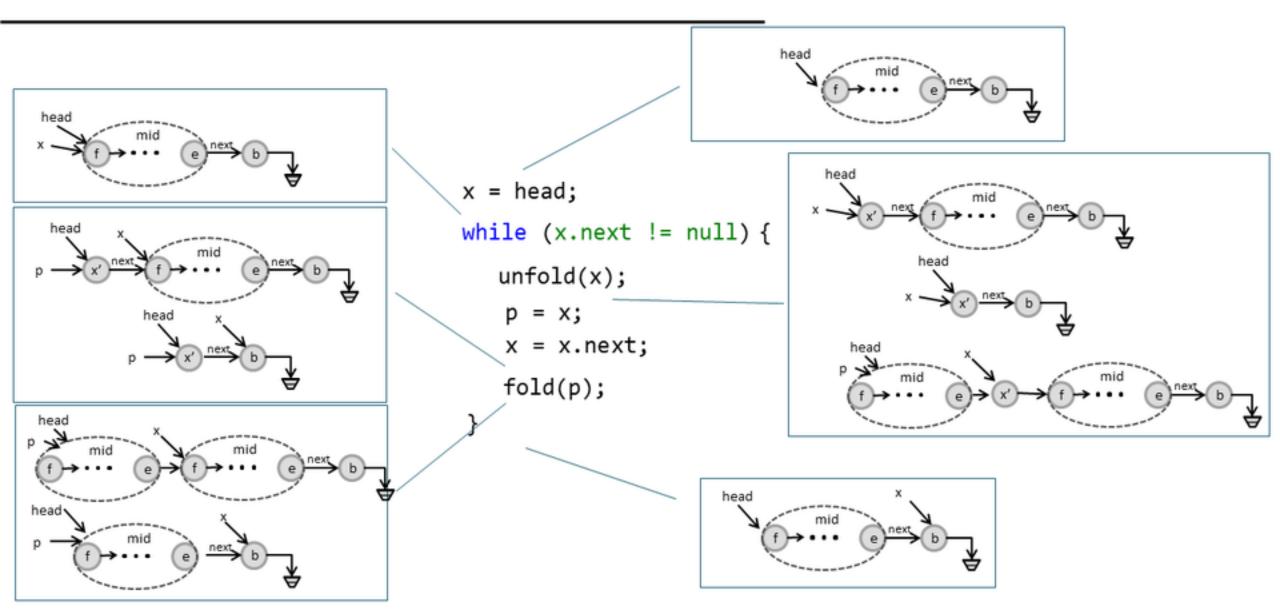
head	
а	true
f	false
e	false
x'	false
mid	false
b	false

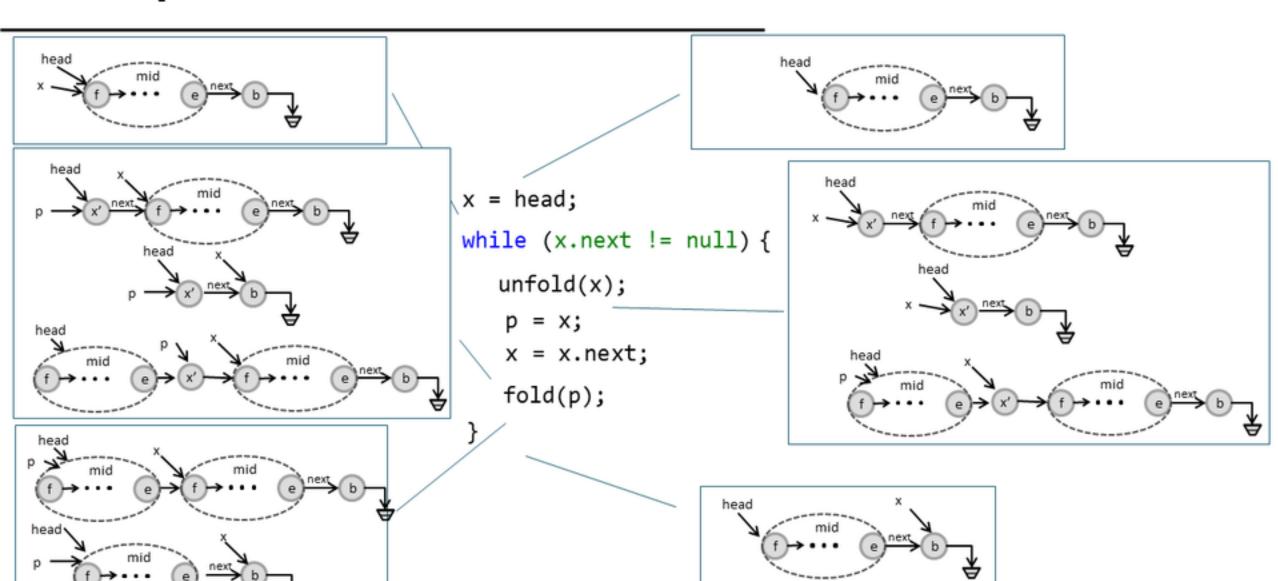
nex	а	f	e	x'	mid	b
а	F	F	F	Т	F	F
f	F	F	F	F	/	F
e	F	F	F	F	F	T
x'	F	Т	/	F	/	F
mid	F	F	/	F	F	/
b	F	F	F	F	F	F



```
head
x = head;
while (x.next != null) {
  unfold(x);
   x = x.next;
   fold(x);
```







Look Sketch

```
void llReverse(Node head)
{
          ?? /*1*/
          while (?? /*p*/)
          {
                ?? /*2*/
           }
          ?? /*3*/
     }
```

Look Sketch

```
void llReverse(Node head)
    {
        cstmt* /*1*/
        while (cond /*p*/)
        {
            cstmt* /*2*/
        }
        cstmt* /*3*/
    }
```

Conditional Statements

var(.ptr?) op var(.ptr?) | null



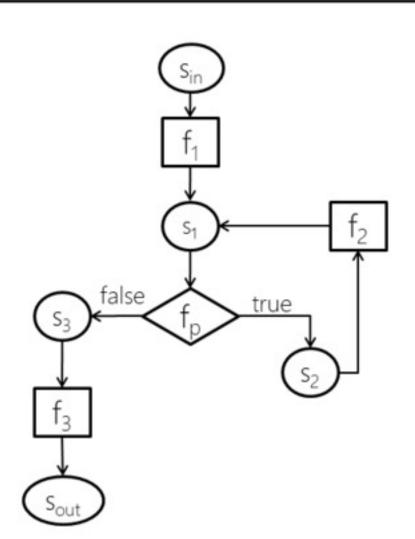
cstmt: if(COND) then STMT



var(.ptr?) = var(.ptr?)

unfold/fold var

Data flow equations



$$s_1 = f_1(s_{in}) \cup f_2(s_2)$$

$$s_2 = f_p(s_1)$$

$$s_3 = \overline{f_p}(s_1)$$

$$s_{out} = f_3(s_3)$$

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arg0

	id	valueA
0	255	1141
1	91	1130
2	347	830
:	:	:
8	225	638
9	257	616

arg1

	id	valueB	
0	255	1231	
1	91	1170	
2	5247	954	
:	:	:	
12	211	575	
13	25	530	

	id	valueA	valueB
0	255	1141	1231
1	91	1130	1170
2	347	830	870
5	159	715	734
8	225	638	644

arg2

"valueA ≠ valueB"

Types and column labels are a potential good abstraction

{"id", "valueA", "valueB"} x DataFrame

Types Abstract Interpreter

class PyTypeInterp

Parameter to Absynthe for a class of problems

Pandas data frame merge

left.merge(right, opts)
df1.merge(df2, on = ['id'])

Types Abstract Interpreter

```
class PyTypeInterp
  def self.pd_merge(left, right, opt)
   if left ⊆ DataFrame &&
      right ⊆ DataFrame &&
      opt ⊆ { on: Array<String>}
      DataFrame
  end
  end
end
```

Pandas data frame query

```
# df.query(pred)
df.query('valueA > 10')
```

Columns Abstract Interpreter

Pandas data frame merge

df1.merge(df2, on = ['id'])

Final data frame is union of both

end

Columns Abstract Interpreter

```
class ColNameInterp

def self.pd_merge(left, right, opt)
   left U right
end
```

Pandas data frame query

df.query('valueA > 10')

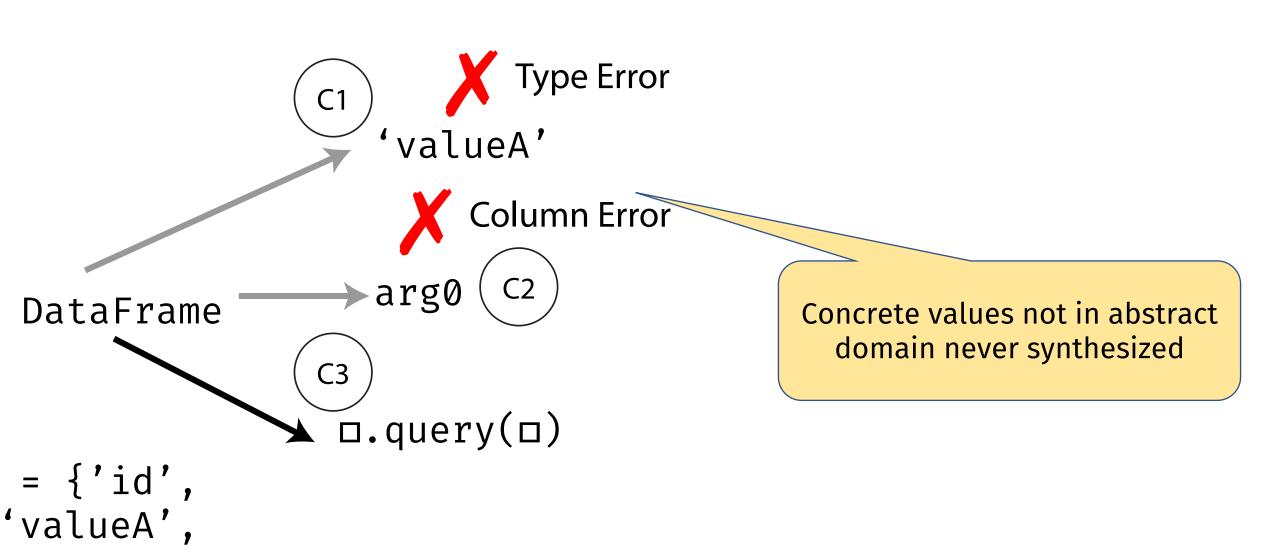
Final data frame has same columns

end

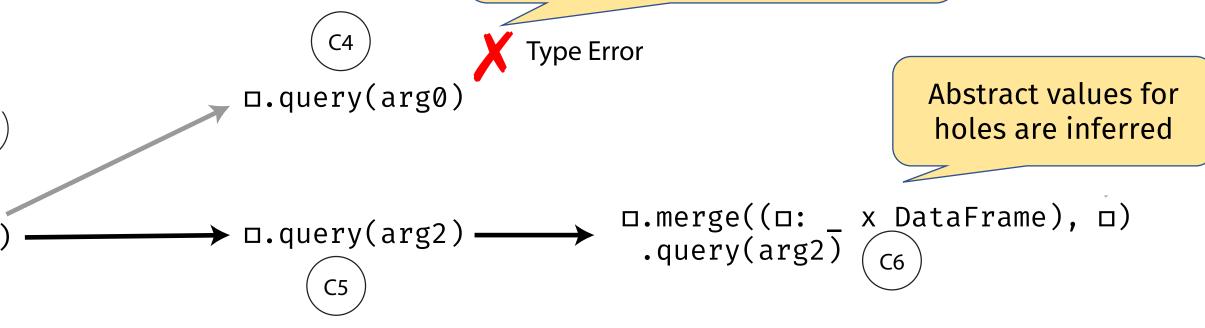
Starting candidate derived from the synthesis goal

 $\left(\mathsf{C0}\right)$

□: Col x DataFrame



Partial programs are evaluated through the abstract interpreter

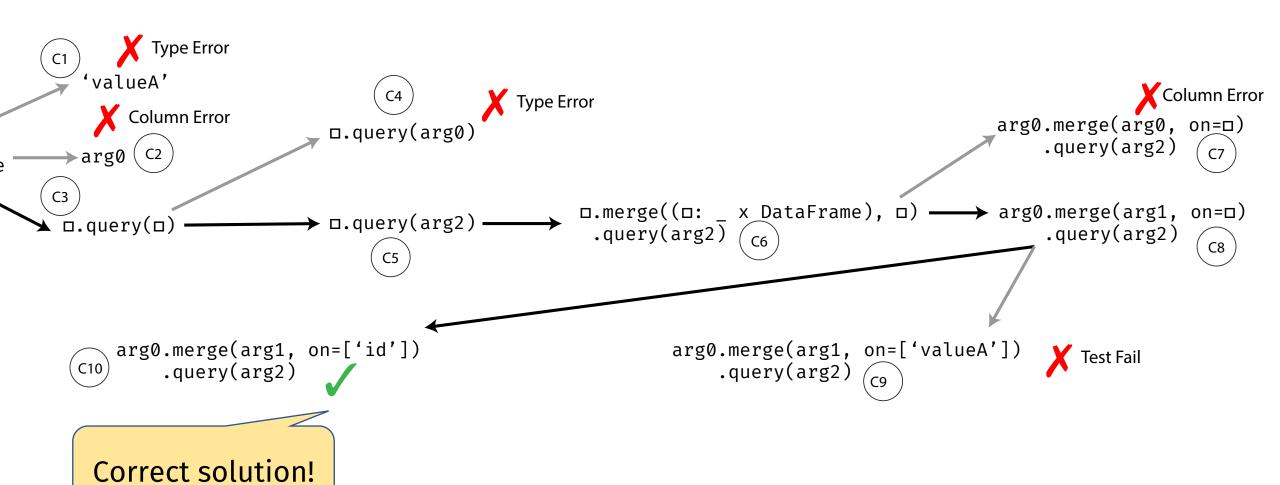


```
Column Error

arg0.merge(arg0, on=□)
.query(arg2)

□.merge(□: _ x DataFrame), □) → arg0.merge(arg1, on=□)
.query(arg2)

(C8)
```

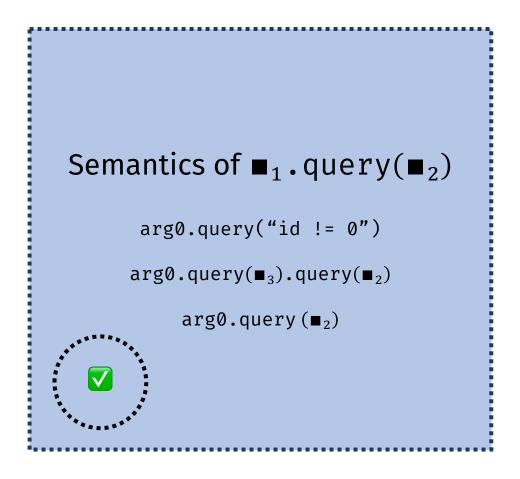


Searching for Programs

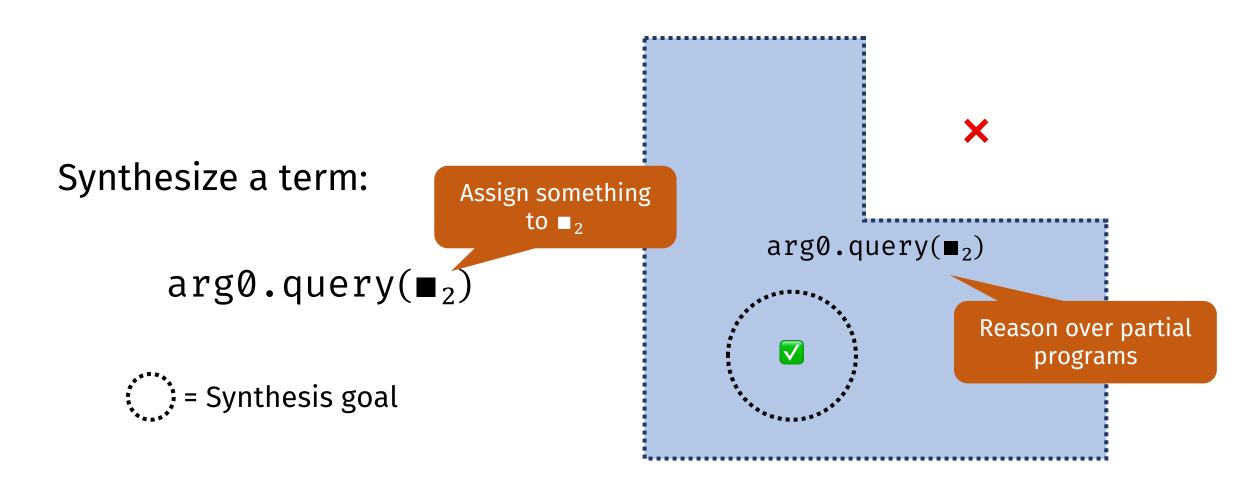
Synthesize a term:

$$\blacksquare_1$$
 •query(\blacksquare_2)

such that it satisfies a synthesis goal :



Searching for Programs



Inferring abstract values

Finite abstract domains:

Types: Int, Str, DataFrame

Infinite abstract dor

Enumerate through valid abstract values

Solver-aided:

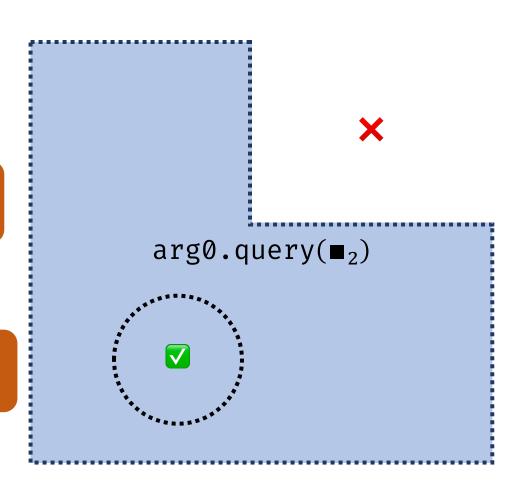
String Length: Linear integer arithmetic

Other:

Data frame columns

Keep 1 hole symbolic and solve for it

Fall back to term enumeration



Absynthe: Abstract Interpretation- Guided Synthesis

- Abstract domains are good at pruning search space
- Framework uses abstract interpreters as a parameter to guide search
- Abstractions for holes are inferred from abstract semantics
- Solves AutoPandas with simple abstract semantics without GPUs

