

#2: Syntax-Guided Synthesis

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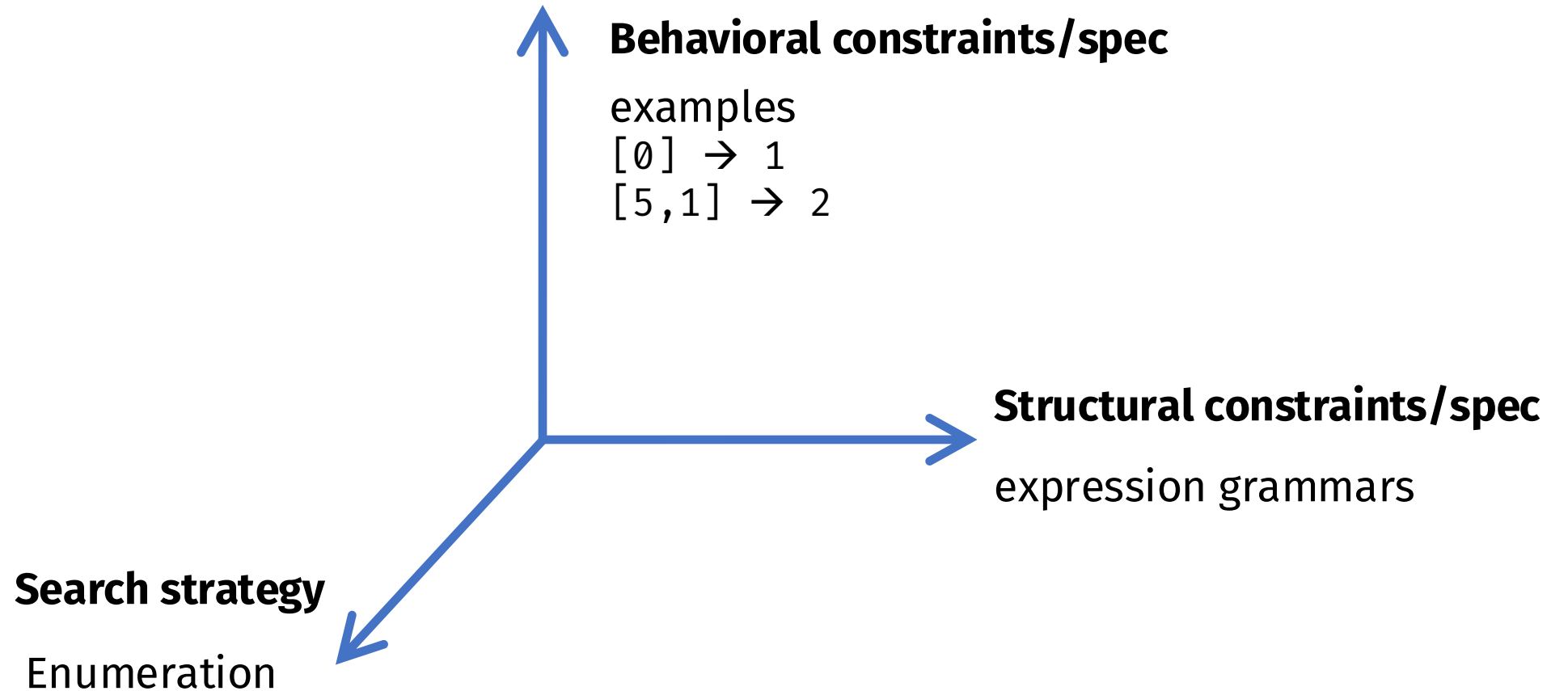
EECS 700: Introduction to Program Synthesis



Logistics

- **Office hours time changed**
 - Monday: 4pm – 5pm
- **Other questions?**

This week



Today

- **Synthesis from examples**
- **Syntax-guided synthesis**
 - expression grammars as structural constraints
 - the SyGuS project

Synthesis from examples

Synthesis from Examples

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Programming by Example

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Inductive Programming /
Inductive Learning

Inductive learning: History

MIT/LCS/TR-76

LEARNING STRUCTURAL DESCRIPTIONS FROM EXAMPLES

Patrick H. Winston

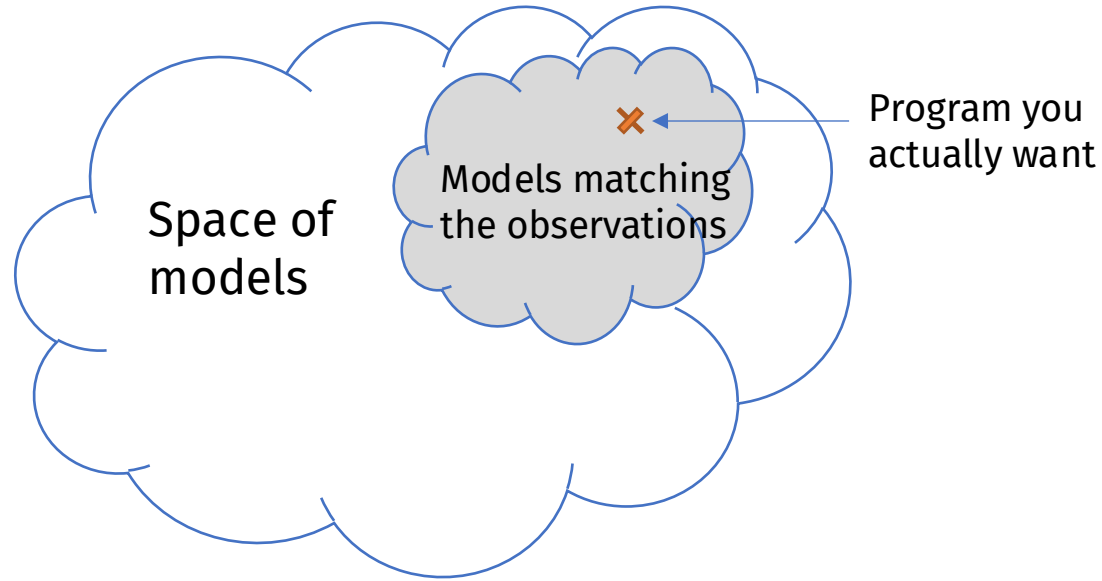
September 1970



Patrick
Winston

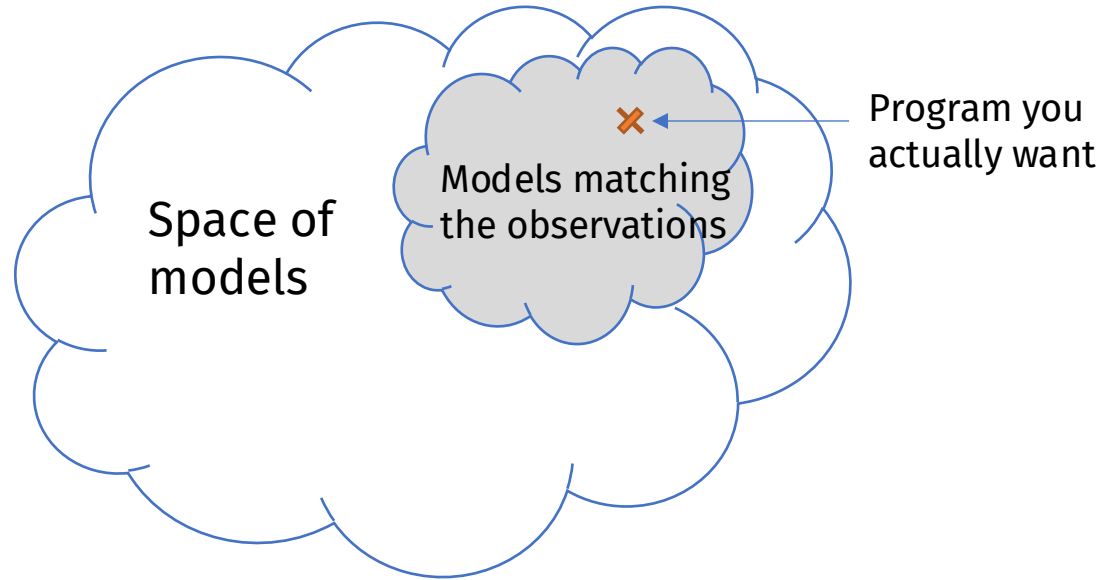
- Explored the question of generalizing from a set of observations
- Became the foundation of machine learning

Key issues in inductive learning



- (1) How do you find a model that matches the observations?
- (2) How do you know it is the model you are looking for?

Key issues in inductive learning

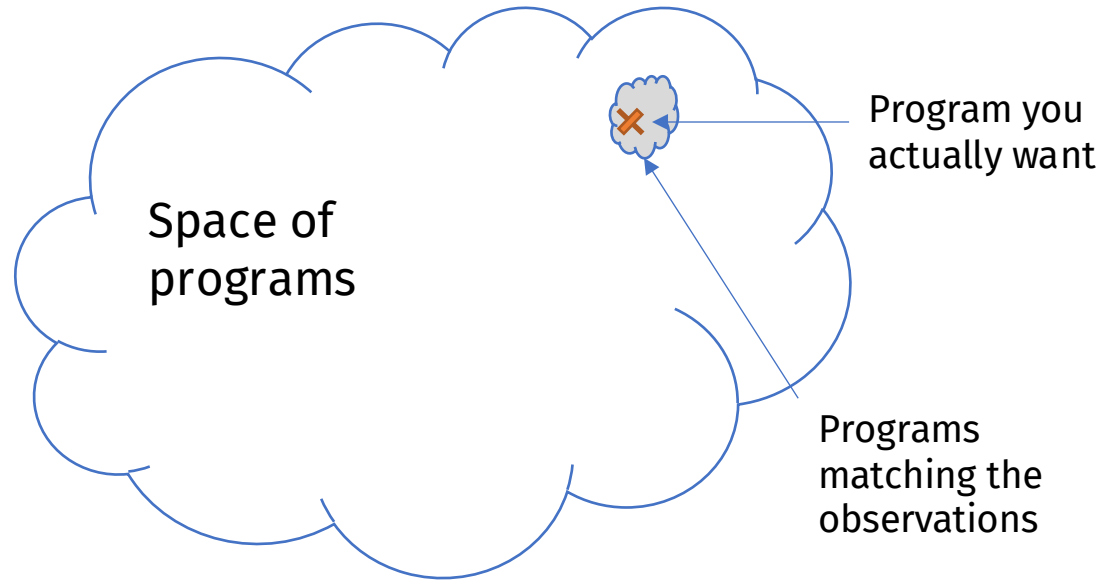


Traditional ML:

- Fix the space so that (1) is easy
- (2) becomes the main challenge

- (1) How do you find a model that matches the observations?
- (2) How do you know it is the model you are looking for?

The synthesis approach

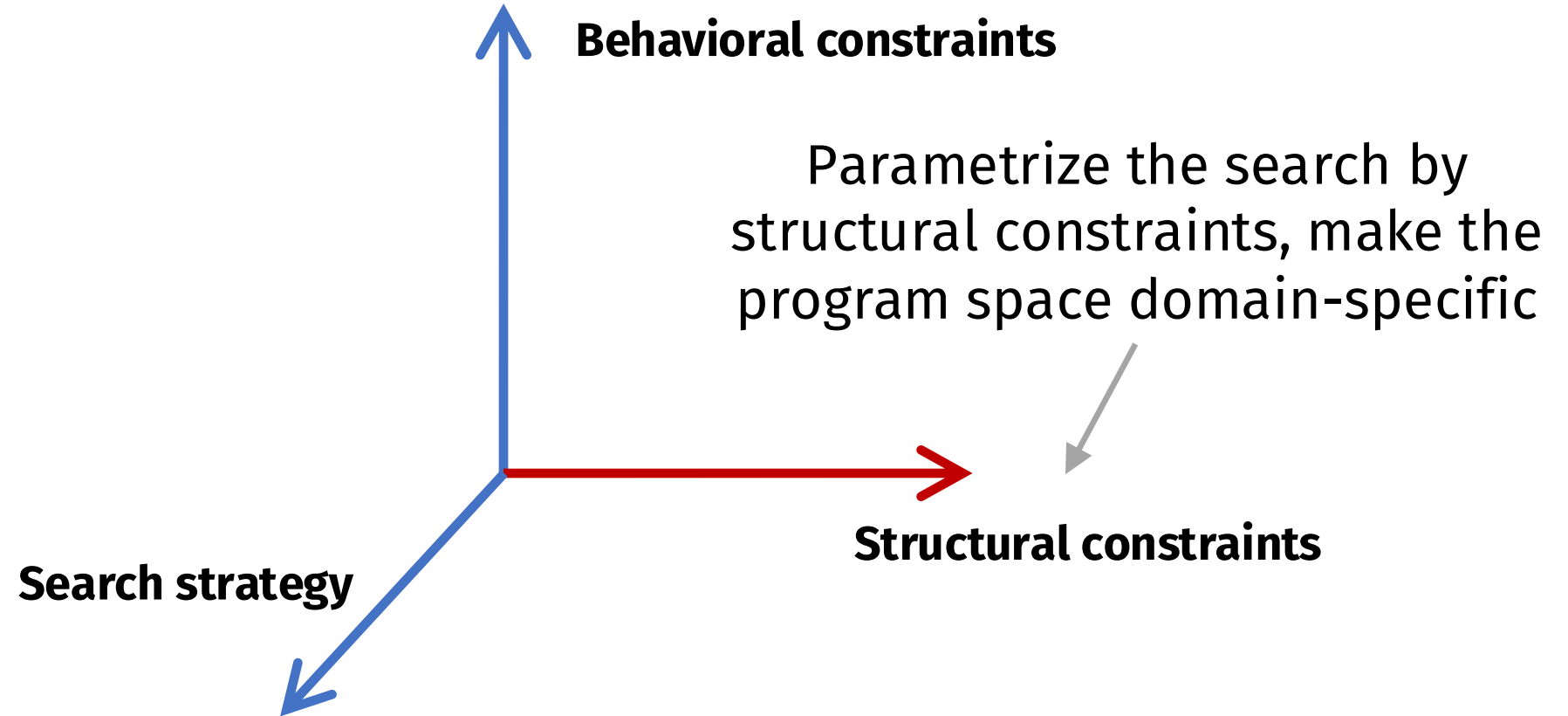


Program synthesis:

- Customize the space so that (2) becomes easier
- (1) is now the main challenge

- (1) How do you find a program that matches the observations?
- (2) How do you know it is the program you are looking for?

Key idea



Syntax-Guided Synthesis

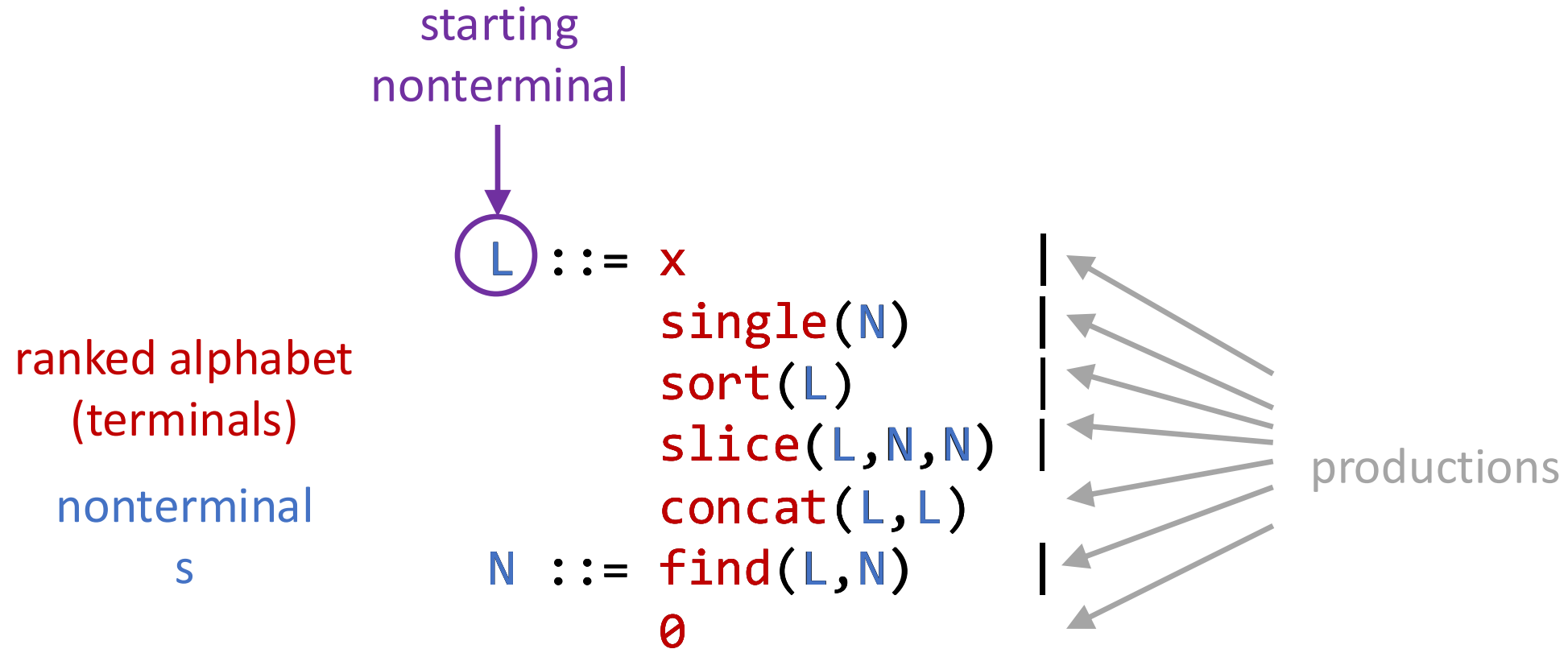
Example

`[1,4,7,2,0,6,9,2,5,0] → [1,2,4,7,0]`

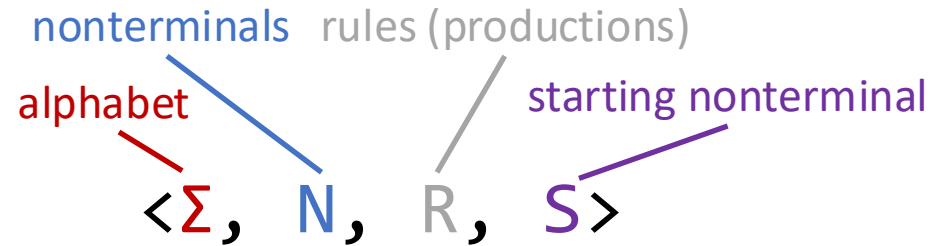
<code>L ::= x</code>		the input
<code>single(N)</code>		<code>single(1) = [1]</code>
<code>sort(L)</code>		<code>sort([6,9,2,5]) = [2,5,6,9]</code>
<code>slice(L,N,N)</code>		<code>slice([6,9,2,5],0,2) = [6,9]</code>
<code>concat(L,L)</code>		<code>concat([6,9],[2,5]) = [6,9,2,5]</code>
<code>N ::= find(L,N)</code>		<code>find([6,9],9) = 1</code>
<code>0</code>		<code>0</code>

`f(x) := concat(sort(slice(x,0,find(x,0))), single(0))`

Regular tree grammars (RTGs)



Regular tree grammars (RTGs)



- Trees: $\tau \in T_{\Sigma}(N)$ = all trees made from $N \cup \Sigma$
- Rules are of the form: $A \rightarrow \sigma(A_1, \dots, A_n)$
- Derives in one step: $\mathcal{C}[A] \rightarrow \mathcal{C}[t]$ if $(A \rightarrow t) \in R$
 A is the leftmost non-terminal in $\mathcal{C}[A]$
- Incomplete terms/programs: $\{\tau \in T_{\Sigma}(N) \mid A \rightarrow^* \tau\}$
- Complete terms/programs: $\{t \in T_{\Sigma} \mid A \rightarrow^* t\}$
 = programs without holes
- Whole programs: $\{t \in T_{\Sigma} \mid S \rightarrow^* t\}$
 = roughly, programs of the right type

$\text{concat}(L, \emptyset)$

$L \rightarrow \text{concat}(L, L)$

$\text{concat}(L, L) \rightarrow \text{concat}(x, L)$

$\text{find}(\text{concat}(L, L), N)$

$\text{find}(\text{concat}(x, x), \emptyset)$

$\text{sort}(\text{concat}(L, L))$

RTGs as structural constraints

Space of programs
= the *language* of an RTG $L(G)$
= all **complete**, **whole** programs

$\textcircled{L} ::= x$
 $\text{single}(N)$
 $\text{sort}(L)$
 $\text{slice}(L, N, N)$
 $\text{concat}(L, L)$
 $N ::= \text{find}(L, N)$
 \emptyset



x $\text{sort}(x)$ $\text{concat}(x, x)$ $\text{slice}(x, \emptyset, \emptyset)$
...
 $\text{slice}(x, \emptyset, \text{find}(x, \emptyset))$
...
 $\text{concat}(\text{sort}(\text{slice}(x, \emptyset, \text{find}(x, \emptyset))), \text{single}(\emptyset))$
...

How big is the space?

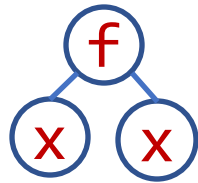
$$E ::= x \mid f(E, E)$$

depth ≤ 0



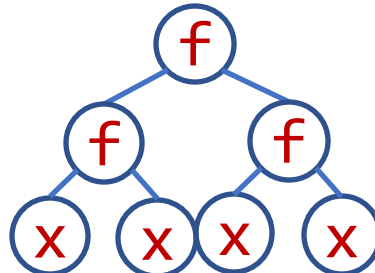
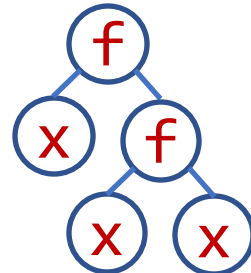
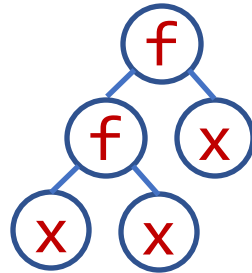
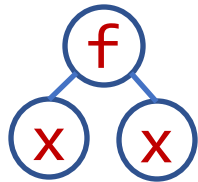
$$N(0) = 1$$

depth ≤ 1



$$N(1) = 2$$

depth ≤ 2



$$N(2) = 5$$

$$N(d) = 1 + N(d - 1)^2$$

How big is the space?

$E ::= x \mid f(E, E)$

$$N(d) = 1 + N(d-1)^2$$

$$N(d) \sim c^{2^d} \quad (c > 1)$$

$$N(1) = 1$$

$$N(2) = 2$$

$$N(3) = 5$$

$$N(4) = 26$$

$$N(5) = 677$$

$$N(6) = 458330$$

$$N(7) = 210066388901$$

$$N(8) = 44127887745906175987802$$

$$N(9) = 1947270476915296449559703445493848930452791205$$

$$N(10) = 3791862310265926082868235028027893277370233152247388584761734150717768254410341175325352026$$

How big is the space?

$$E ::= \begin{array}{c} x_1 \mid \dots \mid x_k \\ f_1(E, E) \mid \dots \mid f_m(E, E) \end{array}$$

$$N(\emptyset) = k$$

$$N(d) = k + m * N(d - 1)^2$$

$$N(1) = 3$$

$$N(2) = 30$$

$$N(3) = 2703$$

$$N(4) = 21918630$$

$$N(5) = 1441279023230703$$

$$N(6) = 6231855668414547953818685622630$$

$$N(7) = 116508075215851596766492219468227024724121520304443212304350703$$

$$k = m = 3$$

Syntactic sugar

Instead of this:

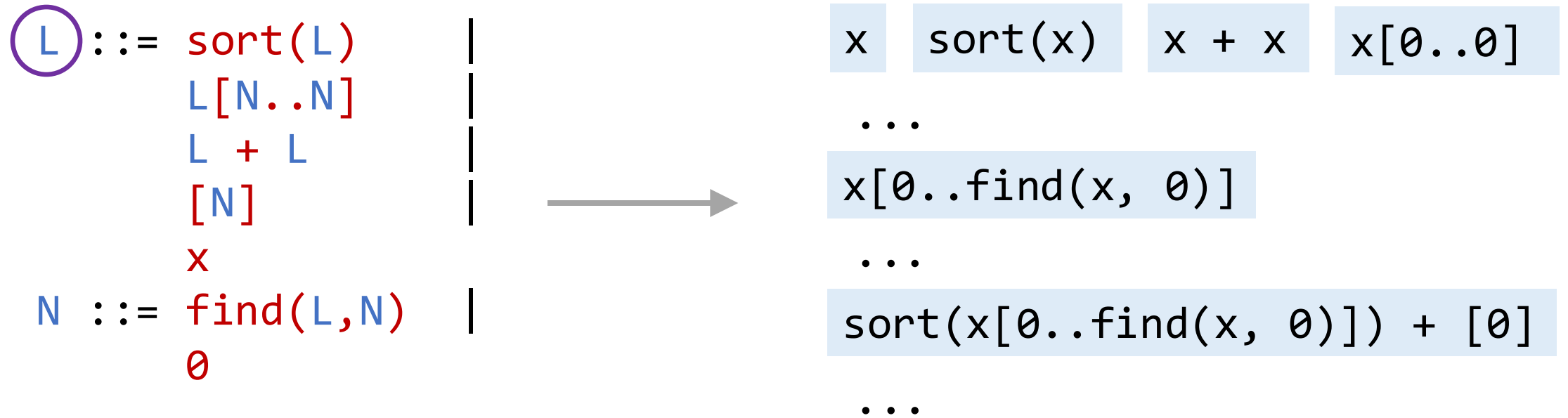
```
L ::= x |  
      single(N) |  
      sort(L) |  
      slice(L, N, N) |  
      concat(L, L)  
N ::= find(L, N) |  
      ∅
```

We will often write this:

```
L ::= x |  
      [N] |  
      sort(L) |  
      L[N..N] |  
      L + L  
N ::= find(L, N) |  
      ∅
```

- allow custom syntax for terminal symbols
- not an RTG strictly speaking, but you know what we mean...

Syntactic sugar



The SyGuS project

<https://sygus.org/>

- Goal: Unify different syntax-guided approaches
- Collection of synthesis benchmarks + yearly competition
 - 6 competitions since 2013
 - consider writing a SyGuS solver for your project!
- Common input format + supporting tools
 - parser, baseline synthesizers

[Alur et al. 2013]

SyGuS problems

- SyGuS problem = $\langle \text{theory, spec, grammar} \rangle$

A “library” of types and function symbols

Example: Linear Integer Arithmetic (LIA)

True, False

0, 1, 2, ...

\wedge , \vee , \neg , $+$, \leq , ite

RTG with terminals in the theory
(+ input variables)

Example: Conditional LIA expressions w/o sums

$E ::= x \mid 0 \mid \text{ite } C \ E \ E$

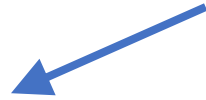
$C ::= E \leq E \mid C \wedge C \mid \neg C$

SyGuS problems

SyGuS problem = $\langle \text{theory, spec, grammar} \rangle$



A first-order logic formula over
the theory



Examples:

$$f(0, 1) = 1 \wedge$$

$$f(1, 0) = 1 \wedge$$

$$f(1, 1) = 1 \wedge$$

$$f(2, 0) = 2$$

SyGuS demo

SyGuS problems

- SyGuS problem = $\langle \text{theory, spec, grammar} \rangle$

A first-order logic formula over
the theory

Examples:

$$\begin{aligned} f(0, 1) &= 1 \wedge \\ f(1, 0) &= 1 \wedge \\ f(1, 1) &= 1 \wedge \\ f(2, 0) &= 2 \end{aligned}$$

Formula with free variables:

$$\begin{aligned} x &\leq f(x, y) \wedge \\ y &\leq f(x, y) \wedge \\ (f(x, y) &= x \vee f(x, y) = y) \end{aligned}$$

can inductive synthesis
handle these?