#4: Equivalence Reduction

Sankha Narayan Guria

EECS 700: Introduction to Program Synthesis

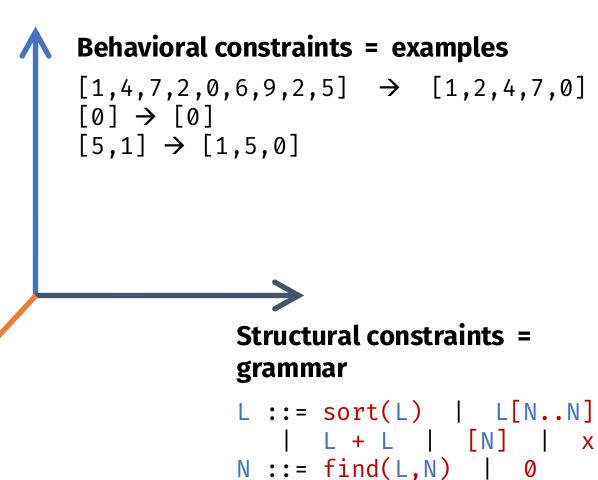


Logistics

• Paper 2 assignment released!

Demo code available at: https://github.com/ku-progsys/eecs700-demos

The problem statement



Search strategy?

Enumerative search

Explicit / Exhaustive Search

Idea: Enumerate programs from the grammar one by one and test them on the examples

```
L ::= sort(L)
                              L[N..N]
                              bottom-up
                                                      top-down
                       N ::= find(L,N)
                              0
X
   0
        x[0..0] x + x
                                                   L[N..N] L + L
sort(x)
                        [0]
                                        x sort(L)
                                                                  [N]
find(x,0)
sort(sort(x))
              sort(x[0..0])
                                                sort(sort(L))
                                        sort(x)
                                                               sort([N])
             sort([0])
                                        sort(L[N..N]) sort(L + L)
sort(x + x)
x[0..find(x,0)]
                                        x[N..N] (sort L)[N..N]
```

Bottom-up vs top-down

Top-down

Bottom-up

Smaller to larger depth

 Has to explore between 3*10⁹ and 10²³ programs to find sort(x[0..find(x, 0)]) + [0] (depth 6)

Candidates are whole but might not be complete

- Cannot always run on inputs
- Can always relate to outputs (?)

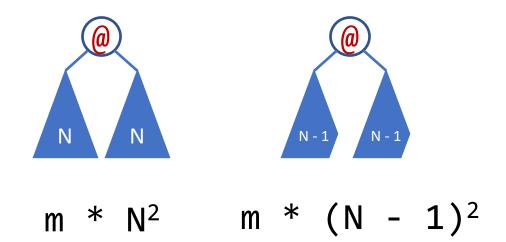
Candidates are complete but might not be whole

- Can always run on inputs
- Cannot always relate to outputs

How to make it scale

Prune

Discard useless subprograms



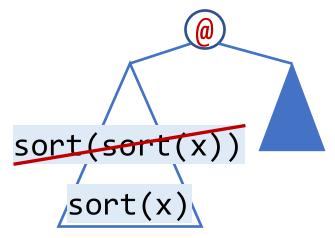
Prioritize

Explore promising candidates first

```
P = { [0][N..N] x[N..N] , dequeue this first ... }
```

When can we discard a subprogram?

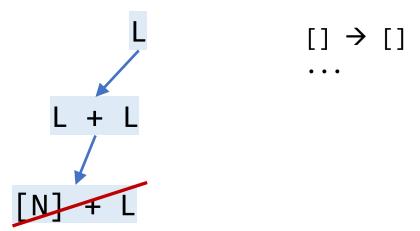
redundant



Equivalence reduction

(also: symmetry breaking)

infeasible



Top-down propagation

Equivalent programs

```
X
                                       0
                                   sort(x) x[0..0] x + x [0] find(x,0)
L ::= sort(L)
      L[N..N]
                                   sort(sort(x)) sort(x + x) sort(x[0..0])
                      bottom up
      L + L
                                   sort([0]) x[0..find(x,0)] x[find(x,0)..0]
      \lceil N \rceil
                                   x[find(x,0)..find(x,0)] sort(x)[0..0]
N ::= find(L,N)
                                   x[0..0][0..0] (x + x)[0..0] [0][0..0]
      0
                                   x + (x + x) x + [0] sort(x) + x x[0..0] + x
                                   (x + x) + x [0] + x x + x[0..0] x + sort(x)
```

Equivalent programs

```
0
                                    sort(x) x[0..0] x + x [0] find(x,0)
L ::= sort(L)
      L[N..N]
                                    sort(sort(x)) sort(x + x) sort(x[0..0])
                      bottom up
      L + L
                                    sort([0]) \times [0..find(x,0)] \times [find(x,0)..0]
      \lceil N \rceil
                                    x[find(x,0)..find(x,0)] sort(x)[0..0]
N ::= find(L,N)
                                    x[0..0][0..0] (x + x)[0..0] [0][0..0]
      0
                                    x + (x + x) x + [0] sort(x) + x x[0..0] + x
                                    (x + x) + x [0] + x x + x[0..0] x + sort(x)
```

Equivalent programs

```
0
                                            x[0..0] x + x [0] find(x,0)
                                  sort(x)
L ::= sort(L)
     L[N..N]
                                               sort(x + x)
                     bottom_up
     L + L
                                           x[0..find(x,0)]
      [N]
N ::= find(L,N)
      0
                                 x + (x + x) x + [0] sort(x) + x
                                             [0] + x
                                                                 x + sort(x)
```

Bottom-up + equivalence reduction

```
bottom-up (\langle T, N, R, S \rangle, [i \rightarrow o]) {
  bank := [t | A ::= t in R]
  while (true)
    forall (p in bank)
      if (p([i]) = [o])
        return p;
    bank += grow(bank);
                                                  m * N^2 m * (N - 1)^2
grow (bank) {
  bank' := []
  forall (A ::= rhs in R)
    bank' += [rhs[B -> p] | p in bank, B \rightarrow^* p]
  return [p' in bank' | forall p in bank: !equiv(p, p')];
```

Bottom-up + equivalence reduction

```
bottom-up (\langle T, N, R, S \rangle, [i \rightarrow o]) {
  bank := [t | A ::= t in R]
  while (true)
    forall (p in bank)
      if (p([i]) = [o])
         return p;
    bank += grow(bank);
grow (bank) {
  bank' := []
  forall (A ::= rhs in R)
    bank' += [rhs[B -> p] | p in bank, B \rightarrow^* p]
  return [p' in bank' | forall p in bank: !equiv(p, p')];
```

- How do we implement equiv?
 - In general, undecidable
 - For SyGuS problems: expensive
 - Doing expensive checks on every candidate defeats the purpose of pruning the space!

```
bottom-up (<T, N, R, S>, [i → o])
{ ... }

equiv(p, p') {
    return p([i]) = p'([i])
}

sort(x) x[0..0] x + x [0] find(x,0)
```

- In PBE, all we care about is equivalence on the given inputs!
 - easy to check efficiently
 - even more programs are equivalent

```
x[0..find(x,0)]

x + (x + x) x + [0] sort(x) + x
```

x + sort(x)

sort(x + x)

[0] + x

$$x + (x + x) x + [0] sort(x) + x$$
 $[0] + x$
 $x + sort(x)$

```
bottom-up (<T, N, R, S>, [i → o])
{ ... }

equiv(p, p') {
   return p([i]) = p'([i])
}
x[0..0] x + x
```

$$x + (x + x)$$

- Proposed simultaneously in two papers:
 - Udupa, Raghavan, Deshmukh, Mador-Haim, Martin, Alur: <u>TRANSIT:</u> <u>specifying protocols with concolic snippets</u>. PLDI'13
 - Albarghouthi, Gulwani, Kincaid: Recursive Program Synthesis. CAV'13
- Variations used in most bottom-up PBE tools:
 - ESolver (baseline SyGuS enumerative solver)
 - **EUSolver** [Alur et al. TACAS'17]
 - Probe [Barke et al. OOPSLA'20]
 - TFCoder [Shi et al. TOPLAS'22]

User-specifies equations

[Smith, Albarghouthi: VMCAI'19]

sort(sort(x)) rule 1 applies, not in normal form

Built-in equivalences

 For a predefined set of operations, equivalence reduction can be hard-coded in the tool or built into the grammar

```
L ::= sort(L)

L[N..N]

L + L

[N]

X

N ::= find(L,N)

0
```

Built-in equivalences

- Used by:
 - **λ**² [Feser et al.'15]
 - Leon [Kneuss et al.'13]
- Leon's implementation using attribute grammars described in:
 - Koukoutos, Kneuss, Kuncak: An Update on Deductive Synthesis and Repair in the Leon tool [SYNT'16]

Equivalence reduction: comparison

- Observational
 - Very general, no user input required
 - Finds more equivalences
 - Can be costly (with many examples, large outputs)
 - If new examples are added, has to restart the search
- User-specified
 - Fast
 - Requires equations
- Built-in
 - Even faster
 - Restricted to built-in operators
 - Only certain symmetries can be eliminated by modifying the grammar
- Q1: Can any of them apply to top-down?
- Q2: Can any of them apply beyond PBE?