Assignment 3

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Part One: The social media features

From the ER-diagram, there is the entity Subscriber with the primary keys <u>username</u> and <u>number</u>. Thus, we will have a table named Subscriber and we will make (<u>username</u>, <u>number</u>) to be the primary key pair. There are also attributes email, hash and salt in the entity Subscriber. We will use large strings CLOB for <u>username</u>. For <u>number</u>, INT is used for the domain of attribute. Since both <u>username</u> and <u>number</u> are used as primary keys, two different subscribers can either have the same username or the same number but not both (exclusive or). We would use VARCHAR(320) to store email. Since hash and salt values require 64-byte each, we would use BINARY(64) for each of them. Since all the above attributes are required to create a subscriber, all of them need to be not null (except for <u>number</u> since it is already automatically incremented). The relational schema for Subscriber table is

Subscriber(<u>username</u>: CLOB,

number: INT,

email: VARCHAR(320), hash: BINARY(64), salt: BINARY(64)).

The ER-diagram has a many-to-many relationship Friend_Of between Subscriber and Subscriber. Thus, we will have a table name Friend_Of. The primary keys of the Subscriber will be the primary keys of the Friend_Of. Since there are two subscribers, we will have a 4-tuple primary key (funame, funum, tuname, tunum). funame is short for from username, which is the username of the subscriber who is following and funum is short for from user number, which is the number of the subscriber who is following. Similarly, tuname and tunum is the username and the user number of the subscriber who is being followed, respectively. Thus, funame and tuname would reference username (column) of the Subscriber table while funum and tunum would reference number (column) of the Subscriber table. Since the domain of the attribute username in the Subscriber table is CLOB, we need to use CLOB to be the domain of the attributes (keys) funame and tuname. Similarly, we need to use INT for funum and tunum. All of the attributes would be not null as well. We also need to

check that a subscriber is not following themselves, that is, we need to check either $\underline{\text{funame}} <> \underline{\text{tunum}} <> \underline{\text{tunum}}$ or both are different (inclusive or). The relational schema for the Friend_Of table is

Friend_Of(<u>funame</u>: CLOB, <u>funum</u>: INT, <u>tuname</u>: CLOB, tunum: INT).

We also have the entity Reaction, which will be the table Reaction in the relational schema. Notice that Reaction, ThreadR, and ReviewR has an ISA relationship, thus, we will use the ER method to create the relational schema for this entity-relationship model. Since Reaction has a strict-one-to-many relationship with Subscriber, we can add the primary keys of the table Subscriber to be the attributes of the table Reaction to represent this relationship without having another table. We name these attributes to be uname and unum reference to the <u>username</u> and <u>number</u> in the table Subscriber with the domain of CLOB and INT, respectively. The primary key for Reaction is <u>id</u> with INT as the type. We will automatically increment the <u>id</u> of the Reaction. Furthermore, Reaction has title and content as its attributes. Since the title would be shorter than the content of a reaction, we will use VARCHAR(100) for title and CLOB for content. The relational schema for Reaction is

Reaction(id: INT, title: VARCHAR(100), content: CLOB, uname: CLOB, unum: INT

Notice that ThreadR has strict-one-to-many and ISA relationships with Reaction. As mentioned before, we use the ER method for the ISA relationship. Thus, we will have threadfrom as the primary key of the table ThreadR to reference the id of the current reaction. We also have an attribute threadto to reference the id of the reaction that the current reaction is reacting to. Since both threadfrom and threadto reference the id of the reactions, both have the type INT and must be not null. To make sure a reaction is not reacting to itself, we need to check that threadfrom <> threadto (CHECK(threadfrom <> threadto)).