

Please write down all people in your team.

1.

2.

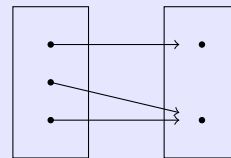
3.

4.

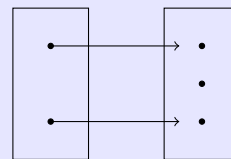
## Review: Functions

### Properties of Functions

- **Onto:** A function is **onto** if every element of the codomain has at least one element in the domain pointing to it. (Every output is attainable via at least one input.)
- **One-to-one:** A function is **one-to-one** if none of the elements in the codomain is the output from two *different* inputs from the domain.



Onto but not one-to-one

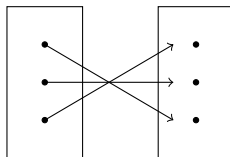


One-to-one but not onto

### Question 1

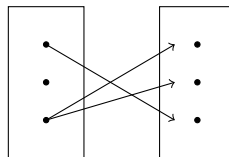
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Identify the properties for the following graphs.



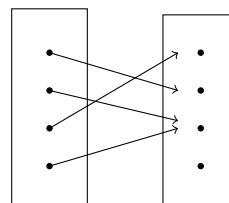
Onto?

One-to-one?



Onto?

One-to-one?



Onto?

One-to-one?

## 7.3 Isomorphism and Planarity

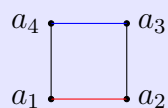
### 7.3.1 Isomorphism

#### Isomorphism

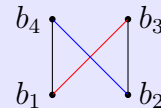
Simple graphs  $G$  and  $H$  are called **isomorphic** if there is a one-to-one and onto function  $f$  from the nodes of  $G$  to the nodes of  $H$  such that  $\{v, w\}$  is an edge of  $G$  if and only if  $\{f(v), f(w)\}$  is an edge of  $H$ . The function  $f$  is called an isomorphism. Hence, an isomorphism is simply a **rule** associating nodes that preserves the edges joining the nodes. <sup>a</sup>

In other words, two graphs are isomorphic if they're essentially the same graph, even if the vertices are in different positions.

**Example:**



$G$



$H$

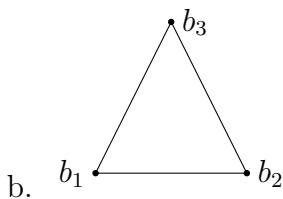
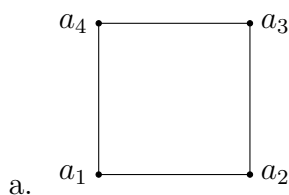
These are isomorphic... imagine taking  $a_2$  and  $a_3$  from graph  $G$  and physically flipping them with the edges still connected. In this case, our mapping is...  $a_1 \rightarrow b_1$        $a_2 \rightarrow b_3$        $a_3 \rightarrow b_2$        $a_4 \rightarrow b_4$ .

<sup>a</sup>Discrete Mathematics, Ensley and Crawley

#### Question 2

\_\_\_\_\_ / 5

Redraw the following graphs by moving the vertices around, but keeping the edges connected.



**Properties of isomorphic graphs**

Two graphs that are isomorphic to one another must have...:

- The same number of nodes
- The same number of edges
- The same number of nodes of any given degree.
- The same number of cycles.
- The same number of cycles of any given size.

<sup>a</sup>

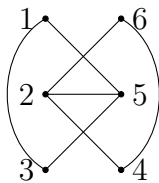
<sup>a</sup>Discrete Mathematics, Ensley and Crawley

**Question 3**

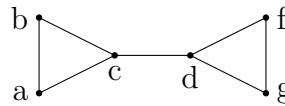
\_\_\_\_\_ / 4

Given the following two graphs...

G



H



- a. Write out all edges for both graphs.

$G: \{2, 5\}$

$H: \{d, c\}$

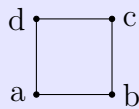
- b. For each edge from  $G$ , write out what edge in  $H$  corresponds to it.

Example:  $\{2,5\} \rightarrow \{d, c\}$

### 7.3.2 Adjacency matrix

We can also use a matrix to list out which vertices are adjacent to which other vertices in order to help us figure out if two graphs are isomorphic.

**Example:**



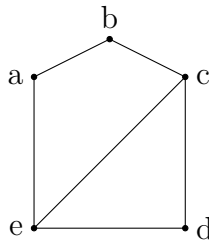
	a	b	c	d
a	0	1	0	1
b	1	0	1	0
c	0	1	0	1
d	1	0	1	0

$a$  is adjacent to  $b$  and  $d$ , so in the  $a$  row we have 1's under the  $b$  and  $d$  columns.

#### Question 4

\_\_\_\_\_ / 4

For the following graph...



a. Finish the adjacency matrix:

	a	b	c	d	e
a					
b					
c					
d					
e					

b. Fill out the degrees of each:

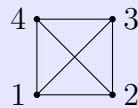
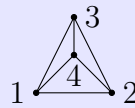
a	b	c	d	e

### 7.3.3 Planarity

1. A simple, connected graph is called **planar** if there is a way to draw it (on a plane) so that no edges cross (i.e., they can only meet at a node). We will call “drawing” of a graph on a plane surface with no edge-crossings an **embedding** of the graph in the plane.
2. A graph is called **bipartite** if its set of nodes can be partitioned into two disjoint sets  $S_1$  and  $S_2$  so that every edge in the graph has one endpoint in  $S_1$  and one endpoint in  $S_2$ .
3. The **complete graph** on  $n$  nodes, denoted by  $K_n$ , is the simple graph with nodes  $\{1, \dots, n\}$  and an edge between every pair of distinct nodes.
4. The **complete bipartite graph** on  $n, m$  nodes, denoted by  $K_{n,m}$ , is the simple bipartite graph with nodes  $S_1 = \{a_1, a_2, \dots, a_n\}$  and  $S_2 = \{b_1, b_2, \dots, b_m\}$  and with edges connecting each node in  $S_1$  to every node in  $S_2$ .

<sup>a</sup>

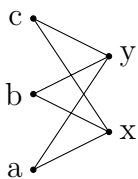
**Example:** Let's redraw the graph  $K_4$  so it has no overlapping edges.

 $K_4$  $K_4$  redrawn<sup>a</sup>Discrete Mathematics, Ensley and Crawley

#### Question 5

\_\_\_\_\_ / 2

Redraw the following graph,  $K_{3,2}$ , so that no edges are overlapping.



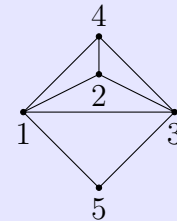
**Faces**

For a planar graph  $G$  embedded in the plane, a **face** of the graph is a region of the plane created by the drawing. Since the plane is an unbounded surface, every embedding of a finite planar graph will have exactly one **unbound** face.<sup>a</sup>

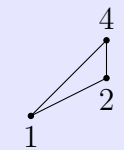
**Unbound (external) face:** Think of the external face as the “canvas” that all other faces are painted on to. Or, if you were viewing a silhouette of the drawing, you would only see the unbounded face - the sum of all the faces.

**Example:**

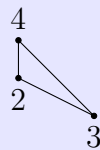
For the drawing, identify the faces by giving the cycle that creates each face, and highlight the unbounded face.



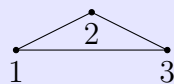
Faces:



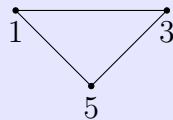
1, 2, 4, 1



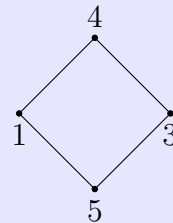
2, 3, 4, 2



1, 2, 3, 1



1, 3, 5, 1



1, 4, 3, 5, 1  
(unbounded)

<sup>a</sup>Discrete Mathematics, Ensley and Crawley

**Question 6**

\_\_\_\_\_ / 2

For both graphs, draw out each of its **faces**, then write out all the cycles bordering faces and identify the unbounded cycle.

