Please write down all people in your team.

1.

2.

3.

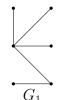
4.

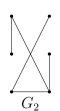
#### 7.2 **Proofs about Graphs and Trees**

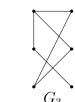
Although this section is named "Proofs", we are actually going to focus on Trees for this section.

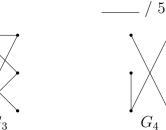
#### 7.2.1Introduction to Trees

Question 1









a. How many vertices does each graph have?

 $G_2$  \_\_\_\_\_

 $G_3$  \_\_\_\_  $G_4$  \_\_

b. How many edges does each graph have?

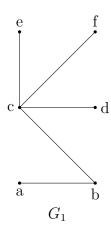
 $G_2$  \_\_\_\_  $G_3$  \_\_\_  $G_4$  \_\_

- c. Which graph is NOT a connected graph?
- d. Which of the graphs has at least one cycle?
- e. Which of the graphs is a tree?

A simple connected graph with no cycles is a **tree**.

<sup>&</sup>lt;sup>1</sup>From Jim Van Horn's POGIL Activity 16

### ${\rm Question}\,\,{\bf 2}$ \_\_\_\_\_/2



a. What is the degree of each of the vertices in  $G_1$ ?

$$deg(a)$$
 \_\_\_\_  $deg(c)$  \_\_\_\_

$$deg(b)$$
 \_\_\_\_\_

$$deg(c)$$
 \_\_\_\_\_

$$deg(d)$$
 \_\_\_\_\_

$$deg(e)$$
 \_\_\_\_\_

$$deg(f)$$
 \_\_\_\_\_

b. List the leaves for  $G_1$ .

Vertices of degree 1 in a tree are called **leaves** of the tree.

 $<sup>^2{\</sup>rm From~Jim~Van~Horn's~POGIL~Activity~16}$ 

A **Tree** is a connected simple graph that has no cycles. Vertices of degree 1 in a tree are called **Leaves** of the tree. a

<sup>a</sup>Discrete Mathematics, Ensley and Crawley

# Question 3 \_\_\_\_\_ / 4

Given these 6 vertices, draw a tree other than  $G_1$  or  $G_2$ .

e f

c • d

å b

- a. How many edges are in your new tree?
- b. How many leaves on your new tree?
- c. If you removed one edge, would the graph still be connected?

A tree with n vertices will have n-1 edges. In other words, it is a connected graph and if you remove an edge then it will become a disconnected graph.

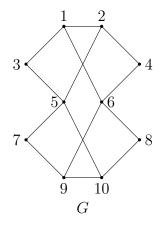
<sup>&</sup>lt;sup>3</sup>From Jim Van Horn's POGIL Activity 16

## 7.2.2 Subgraphs and Trees

A graph H is a **subgraph** of a graph G if all nodes and edges in H are also nodes and edges in G.

<sup>a</sup>Discrete Mathematics, Ensley and Crawley

 $\operatorname{Question}_4 4$ 



 $G_1$   $G_2$ 

a. Draw a graph  $G_1$  above using vertices and edges from G... Vertices: 1, 2, 5, Edges:  $\{1, 2\}$  and  $\{2, 5\}$ .

Is this a **subgraph**?

Are all the vertices of  $G_1$  also nodes of G?

Are all the edges of  $G_1$  also edges of G?

b. Draw a graph  $G_2$  above using vertices and edges from  $G_{...}$  Vertices: 1, 3, 4, Edges:  $\{1, 3\}$  and  $\{3, 4\}$ 

Is this a **subgraph**?

Are all the vertices of  $G_2$  also nodes of G?

Are all the edges of  $G_2$  also edges of G?

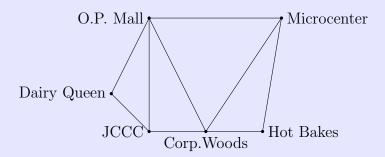
\_ / 4

<sup>&</sup>lt;sup>4</sup>From Jim Van Horn's POGIL Activity 16

### 7.2.3 Spanning Trees

Let G be a simple connected graph. The subgraph T is a **spanning tree** of G if T is a tree and every node in G is a node in T.

### Example: <sup>b</sup>

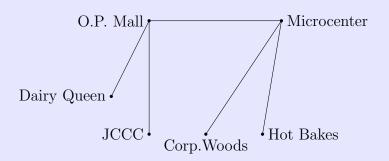


To get from Corporate Woods to JCCC, there are three paths leading into JCCC:

- (1) Directly from  $CW \to JCCC$ ,
- (2) CW  $\rightarrow$  OP Mall  $\rightarrow$  JCCC, and
- (3) CW  $\rightarrow$  OP Mall  $\rightarrow$  Dairy Queen  $\rightarrow$  JCCC

We want to make a graph that connects all locations with the fewest paths. One way to do this is to remove edges of a cycle until no additional edges can be removed without getting a disconnected graph.

One example result is this:



<sup>&</sup>lt;sup>a</sup>Discrete Mathematics, Ensley and Crawley

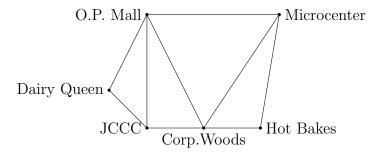
<sup>&</sup>lt;sup>b</sup>From Jim Van Horn's POGIL Activity 16

#### Spanning Tree algorithm <sup>a</sup>

- 1. Begin with a simple connected graph  $G_0$ .
- 2. For each  $i \geq 1$ , as long as there is a cycle in  $G_{i-1}$ ...
  - (a) Choose an edge e in any cycle of  $G_{i-1}$ , and form the subgraph  $G_i$  of  $G_{i-1}$  by deleting e from  $G_{i-1}$
- 3. The final result  $G_k$  will be a spanning tree of  $G_0$ . This is a spanning tree.

Question 5 \_\_\_\_\_ / 2

Follow the algorithm to create a Spanning Tree from this map. "x" out edges that you choose to delete as you go. Draw your spanning tree below.



 $<sup>^</sup>a\mathrm{Discrete}$  Mathematics, Ensley and Crawley

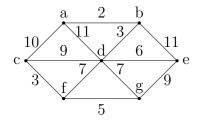
### 7.2.4 Minimal Spanning Trees

#### Prim's Minimal Spanning Tree algorithm <sup>a</sup>

- 1. Given a connected simple graph G with n+1 nodes.
- 2. Let  $v_0$  be any node in G, and let  $T_0 = \{v_0\}$  be a tree with one node and no edges.
- 3. For each k from  $\{1, 2, ..., n\}$ ...
  - (a) Let  $E_k = \{e \text{ an edge in } G : e \text{ has one endpoint in } T_{k-1} \text{ and the other endpoint not in } T_{k-1}\}.$
  - (b) Let  $e_k$  be the edge in  $E_k$  with the smallest weight. (In case of a tie, choose any edge of the smallest weight.)
  - (c) Let  $T_k$  be the tree obtained by adding edge  $e_k$  (along with its node not already in  $T_{k-1}$  to  $T_{k-1}$ .
- 4.  $T_n$  is the tree returned by the algorithm.

<sup>a</sup>Discrete Mathematics, Ensley and Crawley

Question 6 \_\_\_\_\_ / 2



Use Prim's algorithm to find a minimal spanning tree for the graph.