QUY NHON UNIVERSITY
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Optimization

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Roadmap

(1) Optimization Using Gradient Descent

(2) Constrained Optimization and Lagrange Multipliers

(3) Convex Sets and Functions

(4) Convex Optimization

(5) Convex Conjugate

Summary

- Training machine learning models = finding a good set of parameters
- A good set of parameters = Solution (or close to solution) to some optimization problem
- Directions: Unconstrained optimization, Constrained optimization, Convex optimization
- High-school math: A necessary condition for the optimal point: $f'(x) = 0$ (stationary point)
    - Gradient will play an important role

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Unconstrained Optimization and Gradient Algorithms

- Goal
$$\min f(\boldsymbol{x}), \quad f(\boldsymbol{x}) : \mathbb{R}^n \mapsto \mathbb{R}, \quad f \in C^1$$

- Graident-type algorithms
$$\boldsymbol{x}_{k+1} = \boldsymbol{x}_k + \gamma_k \boldsymbol{d}_k, \quad k = 0, 1, 2, \ldots$$

- Lemma. Any direction $\boldsymbol{d} \in \mathbb{R}^{n\times 1}$ that satisfies $\nabla f(\boldsymbol{x}) \cdot \boldsymbol{d} < 0$ is a descent direction of f at $\boldsymbol{x}$. That is, if we let $\boldsymbol{x}_\alpha = \boldsymbol{x} + \alpha\boldsymbol{d}$, $\exists \bar{\alpha} > 0$, such that for all $\alpha \in (0, \bar{\alpha}]$, $f(\boldsymbol{x}_\alpha) < f(\boldsymbol{x})$.
- Steepest gradient descent[1]. $\boldsymbol{d}_k = -\nabla f(\boldsymbol{x}_k)^{\mathsf{T}}$.
- Finding a local optimum $f(\boldsymbol{x}_\star)$, if the step-size γ_k is suitably chosen.
- Question. How do we choose $\boldsymbol{d}_k$ for a constrained optimization?

[1] In some cases, just gradient descent often means this steepest gradient descent.

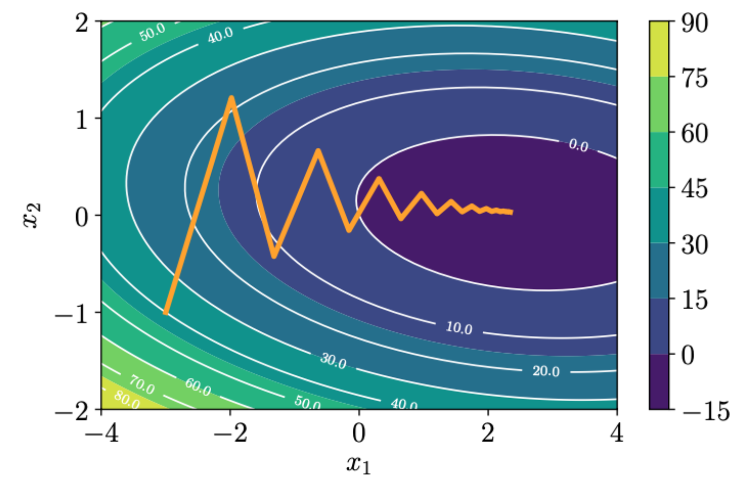
Example

- A quadratic function $f : \mathbb{R}^2 \mapsto \mathbb{R}$.

$$f\left(\begin{pmatrix} x_1 \\ x_2 \end{pmatrix}\right) = \frac{1}{2}\begin{pmatrix} x_1 \\ x_2 \end{pmatrix}^{\mathsf{T}} \begin{pmatrix} 2 & 1 \\ 1 & 20 \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \end{pmatrix} - \begin{pmatrix} 5 \\ 3 \end{pmatrix}^{\mathsf{T}} \begin{pmatrix} x_1 \\ x_2 \end{pmatrix},$$

whose gradient is $\begin{pmatrix} x_1 \\ x_2 \end{pmatrix}^{\mathsf{T}} \begin{pmatrix} 2 & 1 \\ 1 & 20 \end{pmatrix} - \begin{pmatrix} 5 \\ 3 \end{pmatrix}^{\mathsf{T}}$

- $\boldsymbol{x}_0 = (-3\ -1)^{\mathsf{T}}$
- constant step size $\alpha = 0.085$
- Zigzag pattern



Taxonomy

- Goal: $\min L(\boldsymbol{\theta})$ for n training data
- Based on the amount of training data used for each iteration
  - Batch gradient descent (the entire n)
  - Mini-batch gradient descent($k < n$ data)
  - Stochastic gradient descent (one sampled data)
- Based on the adaptive method of update
  - Momentum, NAG, Adagrad, RMSprop, Adam, etc
- `https://ruder.io/optimizing-gradient-descent/`

Stochastic Gradient Descent (SGD)

- Assume $L(\boldsymbol{\theta}) = \sum_{i=1}^{n} L_n(\boldsymbol{\theta})$ (which happens in many cases in machine learning, e.g., negative log-likelihood in regression)
- Gradient update

$$\boldsymbol{\theta}_{k+1} = \boldsymbol{\theta}_k - \gamma_k \nabla L(\boldsymbol{\theta}_k)^{\mathsf{T}} = \boldsymbol{\theta}_k - \gamma_k \sum_{n=1}^{N} \nabla L_n(\boldsymbol{\theta}_k)^{\mathsf{T}}$$

  - Batch gradient: $\sum_{n=1}^{N} \nabla L_n(\boldsymbol{\theta}_k)^{\mathsf{T}}$
  - Mini-batch gradient: $\sum_{n \in \mathcal{K}} \nabla L_n(\boldsymbol{\theta}_k)^{\mathsf{T}}$ for a suitable choice of $\mathcal{K}, |\mathcal{K}| < n$
  - Stochastic gradient: $\nabla L_n(\boldsymbol{\theta}_i)^{\mathsf{T}}$ for some (randomly chosen) i. Noisy approximation to the real gradient.
- Tradeoff: computation burden vs. exactness

Adaptivity for Better Convergence: Momemtum

- Step size.
  - Too small: slow update, Too big: overshoot, zig-zag, often fail to converge
- Adaptive update: smooth out the erratic behavior and dampens oscillations
- Gradient descent with momentum

$$
\begin{aligned}
\boldsymbol{x}_{k+1} &= \boldsymbol{x}_k - \gamma_i \nabla f(\boldsymbol{x}_k)^{\mathsf{T}} + \alpha \Delta \boldsymbol{x}_k, \quad \alpha \in [0, 1] \\
\Delta \boldsymbol{x}_k &= \boldsymbol{x}_k - \boldsymbol{x}_{k-1}
\end{aligned}
$$

  - Memory term: $\alpha \Delta \boldsymbol{x}_k$, where α is the degree of how much we remember the past
  - Next update $=$ a linear combination of current and previous updates

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Standard Constrained Optimization Problem

- An optimization problem in standard form:

$$
\begin{aligned}
&\text{minimize } f(\boldsymbol{x}) \\
&\text{subject to } g_i(\boldsymbol{x}) \leq 0, \quad i = 1, 2, \ldots, m \quad \textit{(Inequality constraints)} \\
&\qquad\qquad\quad h_j(\boldsymbol{x}) = 0, \quad j = 1, 2, \ldots, p \quad \textit{(Equality constraints)}
\end{aligned}
$$

- Variables: $\boldsymbol{x} \in \mathbb{R}^n$. Assume nonempty feasible set

- Optimal value: p^*. Optimizer: $\boldsymbol{x}^*$

Problem Solving via Langrange Multipliers

- Duality Mentality
  - Bound or solve an optimization problem via a different optimization problem!
  - We'll develop the basic Lagrange duality theory for a general optimization problem, then specialize for convex optimization
- Idea: augment the objective with a weighted sum of constraints
  - Lagrangian:

$$\mathcal{L}(\boldsymbol{x}, \boldsymbol{\lambda}, \boldsymbol{\nu}) = f(\boldsymbol{x}) + \sum_{i=1}^{m} \lambda_i g_i(\boldsymbol{x}) + \sum_{i=1}^{p} \nu_i h_i(\boldsymbol{x})$$

  - Lagrange multipliers (dual variables): $\boldsymbol{\lambda} = (\lambda_i : i = 1, \cdots, m) \succeq 0$, $\boldsymbol{\nu} = (\nu_1, \cdots, \nu_p)$
  - Lagrange dual function:

$$\mathcal{D}(\boldsymbol{\lambda}, \boldsymbol{\nu}) = \inf_{\boldsymbol{x}} \mathcal{L}(\boldsymbol{x}, \boldsymbol{\lambda}, \boldsymbol{\nu})$$

Lower Bound on the Optimal Value

- The dual function $\mathcal{D}(\boldsymbol{\lambda}, \boldsymbol{\nu})$ is the lower bound on the optimal value p^*.
- Theorem. $\mathcal{D}(\boldsymbol{\lambda}, \boldsymbol{\nu}) \leq p^*, \ \ \forall \boldsymbol{\lambda} \succeq 0, \ \boldsymbol{\nu}$
- Proof. Consider feasible $\tilde{\boldsymbol{x}}$. Then,

$$\mathcal{L}(\tilde{\boldsymbol{x}}, \boldsymbol{\lambda}, \boldsymbol{\nu}) = f(\tilde{\boldsymbol{x}}) + \sum_{i=1}^{m} \lambda_i g_i(\tilde{\boldsymbol{x}}) + \sum_{i=1}^{p} \nu_i h_i(\tilde{\boldsymbol{x}}) \leq f(\tilde{\boldsymbol{x}})$$

since $f_i(\tilde{\boldsymbol{x}}) \leq 0$ and $\lambda_i \geq 0$.
Hence, $\mathcal{D}(\boldsymbol{\lambda}, \boldsymbol{\nu}) \leq \mathcal{L}(\tilde{\boldsymbol{x}}, \boldsymbol{\lambda}, \boldsymbol{\nu}) \leq f(\tilde{\boldsymbol{x}})$ for all feasible $\tilde{\boldsymbol{x}}$. Therefore, $\mathcal{D}(\boldsymbol{\lambda}, \boldsymbol{\nu}) \leq p^*$.

Lagrangian Dual Problem

- Lower bound from Lagrange dual function depends on $(\boldsymbol{\lambda}, \boldsymbol{\nu})$.
- Question. What's the best lower bound?

$$\textbf{Langrangian dual problem} \quad \begin{array}{ll} \text{maximize} & \mathcal{D}(\boldsymbol{\lambda}, \boldsymbol{\nu}) \\ \text{subject to} & \boldsymbol{\lambda} \succeq 0 \end{array}$$

- Dual variables: $(\boldsymbol{\lambda}, \boldsymbol{\nu})$
- Always a convex optimization, because $\mathcal{D}(\boldsymbol{\lambda}, \boldsymbol{\nu})$ is always concave over $\boldsymbol{\lambda}, \boldsymbol{\nu}$.
  - Infimum over $\boldsymbol{x}$ of a family of affine functions in $(\boldsymbol{\lambda}, \boldsymbol{\nu})$ (we will see this later)
- Denote the optimal value of Lagrange dual problem by d^*.

Weak Duality

- What's the relationship between d^* and p^*?

Weak Duality

$$d^* \le p^*$$

- Weak duality always hold (even if the primal problem is not convex):
- Optimal duality gap: $p^* - d^*$
- Efficient generation of the lower bounds through the dual problem

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Convex Optimization

- Convex optimization problem

$$
\begin{aligned}
\text{minimize} \quad & f(\boldsymbol{x}) \\
\text{subject to} \quad & \boldsymbol{x} \in \mathcal{X},
\end{aligned}
$$

  where $f(\boldsymbol{x}) : \mathbb{R}^n \mapsto \mathbb{R}$ is a convex function, and $\mathcal{X}$ is a convex set.

- The watershed between easily solvable problem and intractable ones is not ‘linearity’, but ‘convexity’
- Let’s overview the background of convex functions, convex sets, and their basic properties.

Convex Set

- Set $\mathcal{C}$ is a convex set if the line segment between any two points in $\mathcal{C}$ lies in $\mathcal{C}$, i.e., if for any $x_1, x_2 \in \mathcal{C}$ and any $\theta \in [0, 1]$, we have $\theta x_1 + (1 - \theta)x_2 \in \mathcal{C}$
- Convex hull of $\mathcal{C}$ is the set of all convex combinations of points in $\mathcal{C}$:

$$\left\{ \sum_{i=1}^{k} \theta_i x_i \mid x_i \in \mathcal{C}, \theta_i \geq 0, i = 1, 2, \ldots, k, \sum_{i=1}^{k} \theta_i = 1 \right\}$$

- What is k? For all k? For some k?

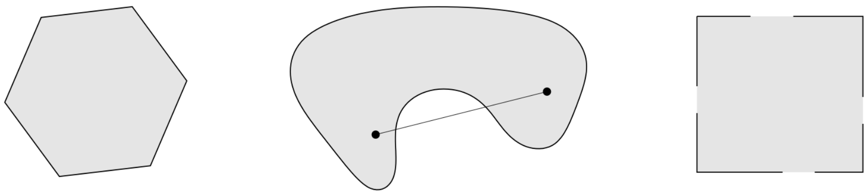
- Generalize to infinite sums and integrals:

$$\sum_{i=1}^{\infty} \theta_i x_i \in \mathcal{C}, \quad \int_{\mathcal{C}} p(x) x dx \in \mathcal{C},$$

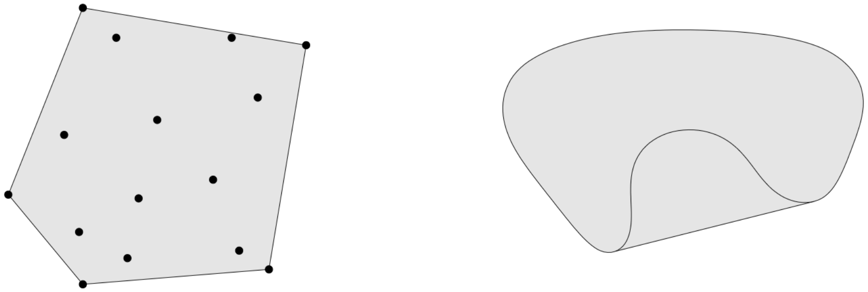
where $\sum_{i=1}^{\infty} \theta_i = 1$ and $p(x)$ is a pdf of some random variable.

Examples

- Convex and Non-convex sets



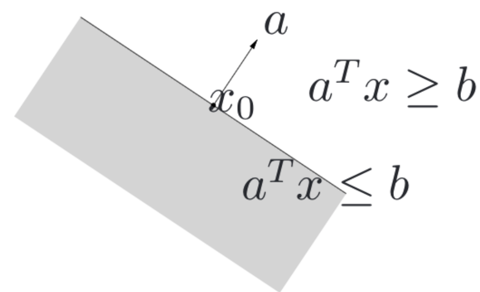
- Convex hulls



Examples of Convex Sets

- Hyperplane in $\mathbb{R}^n$ is a set: $\{x \mid a^\top x = b\}$ where $a \in \mathbb{R}^n, a \neq 0, b \in \mathbb{R}$
  In other words, $\{x \mid a^\top (x - x_0) = 0\}$, where x_0 is any point in the hyperplane, i.e., $a^\top x_0 = b$.

- Divides $\mathbb{R}^n$ into two halfspaces: $\{x | a^\top x \leq b\}$ and $\{x | a^\top x > b\}$

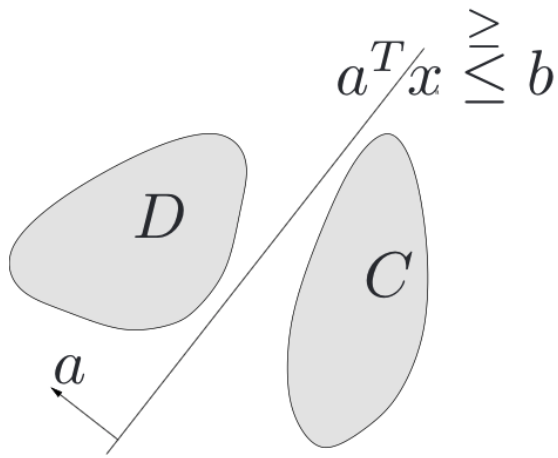


- Polyhedron is the solution set of a finite number of linear equalities and inequalities (intersection of finite number of halfspaces and hyperplanes)

$$\mathcal{P} = \{x \mid a_j^\top x \leq b_j, j = 1, \ldots, m, c_j^\top x = d_j, j = 1, \ldots, p\} = \{x \mid Ax \leq b, C$$

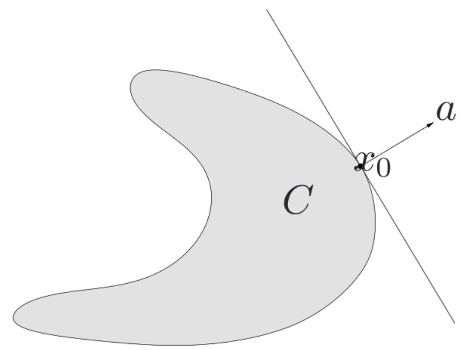
- Polytope: a bounded polyhedron

Separating Hyperplane Theorem



- $\mathcal{C}$ and $\mathcal{D}$: non-intersecting convex sets, i.e., $\mathcal{C} \bigcap \mathcal{D} = \phi$.
- Then, there exist $a \neq 0$ and b such that $a^\mathsf{T} x \leq b$ for all $x \in \mathcal{C}$ and $a^\mathsf{T} x \geq b$ for all $x \in \mathcal{D}$.

Supporting Hyperplane Theorem



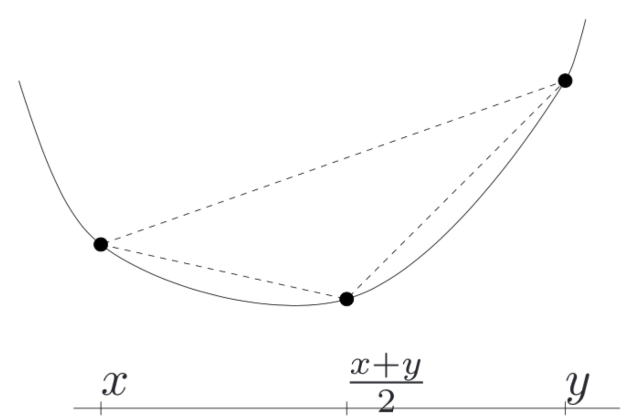
- Given a set $\mathcal{C} \in \mathbb{R}^n$ and a point x_0 on its boundary, if $a \neq 0$ satisfies $a^\mathsf{T} x \leq a^\mathsf{T} x_0$ for all $x \in \mathcal{C}$, then $\{x | a^\mathsf{T} x = a^\mathsf{T} x_0\}$ is called a supporting hyperplane to $\mathcal{C}$ at x_0
- For any nonempty convex set $\mathcal{C}$ and any x_0 on boundary of $\mathcal{C}$, there exists a supporting hyperplane to $\mathcal{C}$ at x_0
- What happens if $\mathcal{C}$ is non-convex?

Convex Functions

- $f : \mathbb{R}^n \mapsto \mathbb{R}$ is a convex function if dom f is a convex set and for all $x, y \in$ dom f and $\theta \in [0, 1]$, we have

$$f(\theta x + (1-\theta)y) \leq \theta f(x) + (1-\theta) f(y)$$

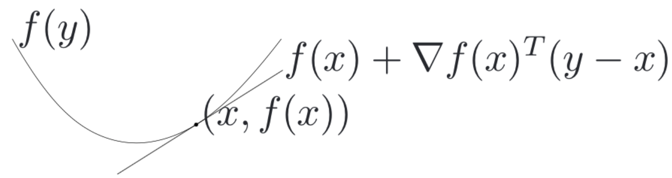
- f is strictly convex if the strict inequality in the above holds for all $x \neq y$ and $0 < \theta < 1$.
  - f is concave if $-f$ is convex
  - Affine functions are convex and concave
  - Jensen's inequality. For a rv X, $f(\mathbb{E}[X]) \leq \mathbb{E}[f(X)]$.



Conditions of Convex Functions (1)

- First-order condition. For differentiable functions, f is convex iff

$$f(y) - f(x) \geq \nabla f(x)^{\mathsf{T}}(y - x), \quad \forall x, y \in \text{dom } f, \text{and dom } f \text{ is convex}$$



- Example. $f(y) = y^2$.
- $f(y) \geq \tilde{f}_x(y)$ where $\tilde{f}_x(y)$ is the first order Taylor expansion of $f(y)$ at x.
- Local information (first order Taylor approximation) about a convex function provides global information (global underestimator).
- If $\nabla f(x) = 0$, then $f(y) \geq f(x)$, $\forall y$. Thus, x is a global minimizer of f

Conditions of Convex Functions (2)

- Second-order condition. For twice differentiable functions, f is convex iff

$$\nabla^2 f(x) \succeq 0$$

  for all $x \in \operatorname{dom} f$ (upward slope) and $\operatorname{dom} f$ is convex

- Example: $f(x) = x^2$.
- Meaning: The graph of the function have positive (upward) curvature at x.

Examples of Convex or Concave Functions

- e^{ax} is convex on $\mathbb{R}$, for any $a \in \mathbb{R}$
- x^a is convex on $\mathbb{R}_{++}$ when $a \geq 1$ or $a \leq 0$, and concave for $0 \leq a \leq 1$
- $|x|^p$ is convex on $\mathbb{R}$ for $p \geq 1$
- $\log x$ is concave on $\mathbb{R}_{++}$
- $x \log x$ is strictly convex on $\mathbb{R}_{++}$
- Every norm on $\mathbb{R}^n$ is convex
- $f(x) = \max\{x_1, \ldots, x_n\}$ is convex on $\mathbb{R}^n$
- $f(x) = \log \sum_{i=1}^{n} e^{x_i}$ is convex on $\mathbb{R}^n$
- $f(x) = \left(\prod_{i=1}^{n} x_i\right)^{\frac{1}{n}}$ is concave on $\mathbb{R}^n_{++}$

Convexity-Preserving Operations

- $f = \sum_{i=1}^{n} w_i f_i$ convex if f_i are all convex and $w_i \geq 0$
- $g(x) = f(Ax + b)$ is convex iff $f(x)$ is convex
- $f(x) = \max\{f_1(x), f_2(x)\}$ convex if f_i convex, e.g., sum of r largest components is convex
- $f(x) = h(g(x))$, where $h : \mathbb{R}^k \mapsto \mathbb{R}$ and $g : \mathbb{R}^n \mapsto \mathbb{R}^k$.

  If $k = 1$: $f''(x) = h''(g(x))g'(x)^2 + h'(g(x))g''(x)$. So, f is convex if h is convex and nondecreasing and g is convex, or if h is convex and nonincreasing and g is concave ...
- $g(x) = \inf_{y \in \mathcal{C}} f(x, y)$ is convex if f is convex in (x, y) and $\mathcal{C}$ is convex

Point-wise Supremum

- If $f(x, y)$ is convex in x for each $y \in \mathcal{A}$, then

$$g(x) = \sup_{y \in \mathcal{A}} f(x, y)$$

  is convex. Similarly, if $f(x, y)$ is concave in x for each $y \in \mathcal{A}$, then

$$g(x) = \inf_{y \in \mathcal{A}} f(x, y)$$

  is concave.

- Example. distance to farthest point in a set $\mathcal{C}$: $f(x) = \sup_{y \in \mathcal{C}} ||x - y||$ is convex.

- Example. Lagrange dual function

$$\mathcal{D}(\boldsymbol{\lambda}, \boldsymbol{\nu}) = \inf_{\boldsymbol{x}} \mathcal{L}(\boldsymbol{x}, \boldsymbol{\lambda}, \boldsymbol{\nu})$$

  is concave.

Roadmap

Standard Convex Optimization

- A standard convex optimization problem with variables $\boldsymbol{x}$:

$$\begin{array}{ll} \text{minimize} & f(\boldsymbol{x}) \\ \text{subject to} & g_i(\boldsymbol{x}) \leq 0, \quad i = 1, 2, \ldots, m \\ & {a_i}^\mathsf{T} \boldsymbol{x} = b_i, \quad i = 1, 2, \ldots, p \end{array}$$

  where $f, f_1, \ldots, f_m$ are convex functions.

- Minimize convex objective function (or maximize concave objective function)
- Upper bound inequality constraints on convex functions ($\Rightarrow$ Constraint set is convex)
- Equality constraints must be affine (Only affine functions leads to a convex set for the equality constraints)

Properties for Optimality

- Local optimality implies global optimality.
  - Given x is locally optimal for a convex optimization problem, i.e., $\boldsymbol{x}$ is feasible and for some $R > 0$,

$$f(\boldsymbol{x}) = \inf\{f(\boldsymbol{z}) \mid \boldsymbol{z} \text{ is feasible }, \|\boldsymbol{z} - \boldsymbol{x}\|_2 \leq R\}$$

  - Theorem. if $\boldsymbol{x}$ is locally optimal in convex program, then globally optimal.
- Optimal condition for differentiable f
  - $\boldsymbol{x}$ is optimal for a convex optimization problem iff $\boldsymbol{x}$ is feasible and for all feasible $\boldsymbol{y}$:

$$\nabla f(\boldsymbol{x})^{\mathsf{T}}(\boldsymbol{y} - \boldsymbol{x}) \geq 0$$

  - $-\nabla f(\boldsymbol{x})$ defines a supporting hyperplane to the feasible set $(\{\boldsymbol{y} \mid -\nabla f(\boldsymbol{x})^{\mathsf{T}}\boldsymbol{y} \leq -\nabla f(\boldsymbol{x})^{\mathsf{T}}\boldsymbol{x}\})$.
  - (Note) Unconstrained convex optimization: condition reduces to: $\nabla f(\boldsymbol{x}) = 0$

Strong Duality

- Strong duality (zero optimal duality gap):

$$d^* = p^*$$

- If strong duality holds, solving dual is 'equivalent' to solving primal. But strong duality does not always hold
- Convexity and constraint qualifications $\implies$ Strong duality
- A simple constraint qualification: Slater's condition (there exists strictly feasible primal variables $f_i(\boldsymbol{x}) < 0$ for non-affine f_i) (see Section 5.3.2 of Boyd's book).
- Another reason why convex optimization is 'easy'

Complementary Slackness

- Assume strong duality holds. Then,

$$f(\boldsymbol{x}^*) = \mathcal{D}(\boldsymbol{\lambda}^*, \boldsymbol{\nu}^*) = \inf_{\boldsymbol{x}} \left(f(\boldsymbol{x}) + \sum_{i=1}^{m} \lambda_i^* f_i(\boldsymbol{x}) + \sum_{i=1}^{p} \nu_i^* h_i(\boldsymbol{x}) \right)$$

$$\leq f(\boldsymbol{x}^*) + \sum_{i=1}^{m} \lambda_i^* f_i(\boldsymbol{x}^*) + \sum_{i=1}^{p} \nu_i^* h_i(\boldsymbol{x}^*) \leq f(\boldsymbol{x}^*)$$

- Thus, the two inequalities must hold with equality, implying: $\lambda_i^* f_i(\boldsymbol{x}^*) = 0, \ \forall i$
- Complementary slackness condition:

$$\lambda_i^* > 0 \implies f_i(\boldsymbol{x}^*) = 0$$
$$f_i(\boldsymbol{x}^*) < 0 \implies \lambda_i^* = 0$$

- i-th optimal Lagrange multiplier is zero unless the ith constraint is active at the optimum.

KKT Condition

- Since $\boldsymbol{x}^*$ minimizes $\mathcal{L}(\boldsymbol{x}, \boldsymbol{\lambda}^*, \boldsymbol{\nu}^*)$ over $\boldsymbol{x}$,

$$\nabla f(\boldsymbol{x}^*) + \sum_{i=1}^{m} \lambda_i^* \nabla f_i(\boldsymbol{x}^*) + \sum_{i=1}^{p} \nu_i^* \nabla h_i(\boldsymbol{x}^*) = 0$$

Karush-Kuhn-Tucker optimality condition

$$f_i(\boldsymbol{x}^*) \leq 0, \quad h_i(\boldsymbol{x}^*) = 0, \quad \lambda_i^* \succeq 0$$
$$\lambda_i^* f_i(\boldsymbol{x}^*) = 0$$
$$\nabla f(\boldsymbol{x}^*) + \sum_{i=1}^{m} \lambda_i^* \nabla f_i(\boldsymbol{x}^*) + \sum_{i=1}^{p} \nu_i^* \nabla h_i(\boldsymbol{x}^*) = 0$$

- Any optimization with strong duality, KKT condition is necessary for primal-dual optimality
- Convex optimization with Slater's condition, KKT is also sufficient for primal-dual optimality.

Useful Tips

- Minimization problem (min-min-max rule)
  - Problem: $\min f(x) \quad \text{s.t.} \quad f_i(x) \leq 0, \quad g_i(x) = 0, \quad x$
  - $f(x)$: convex, $f_i(x)$: convex, $g_i(x)$: affine
  - $L(x, \lambda, \mu) = f(x) + \sum_i \lambda_i f_i(x) + \sum_i \mu_i g_i(x)$
  - $\inf_x L(x, \lambda, \mu) = \mathcal{D}(\lambda, \mu)$
  - $\max_{\lambda \geq 0} \mathcal{D}(\lambda, \mu)$
- Maximization problem (max-max-min rule)
  - Problem: $\max f(x) \quad \text{s.t.} \quad f_i(x) \geq 0, \quad g_i(x) = 0, \quad x$
  - $f(x)$: concave, $f_i(x)$: concave, $g_i(x)$: affine
  - $L(x, \lambda, \mu) = f(x) + \sum_i \lambda_i f_i(x) + \sum_i \mu_i g_i(x)$
  - $\sup_x L(x, \lambda, \mu) = \mathcal{D}(\lambda, \mu)$
  - $\min_{\lambda \geq 0} \mathcal{D}(\lambda, \mu)$

Linear Programming

- Primal problem

$$\begin{array}{ll} \min_{\boldsymbol{x}\in\mathbb{R}^d} & \boldsymbol{c}^\top \boldsymbol{x} \\ \text{subject to} & \boldsymbol{A}\boldsymbol{x} \preceq \boldsymbol{b}, \end{array}$$

where $\boldsymbol{A} \in \mathbb{R}^{m\times d}$ and $\boldsymbol{b} \in \mathbb{R}^m$.

- Dual problem

$$\begin{array}{ll} \max_{\boldsymbol{\lambda}\in\mathbb{R}^m} & -\boldsymbol{b}^\top \boldsymbol{\lambda} \\ \text{subject to} & \boldsymbol{c} + \boldsymbol{A}^\top\boldsymbol{\lambda} = 0,\ \boldsymbol{\lambda} \succeq 0, \end{array}$$

where $\boldsymbol{\lambda} \in \mathbb{R}^m$.

- The Lagrangian: $\mathcal{L}(\boldsymbol{x},\boldsymbol{\lambda}) = (\boldsymbol{c} + \boldsymbol{A}^\top\boldsymbol{\lambda})^\top \boldsymbol{x} - \boldsymbol{\lambda}^\top \boldsymbol{b}$, whose derivative w.r.t. $\boldsymbol{x}$ becomes zero, when $\boldsymbol{c} + \boldsymbol{A}^\top\boldsymbol{\lambda} = 0$.
- The dual function: $\mathcal{D}(\boldsymbol{\lambda}) = -\boldsymbol{\lambda}^\top \boldsymbol{b}$

Quadratic Programming

- Primal problem

$$\begin{aligned} \min_{\boldsymbol{x}\in\mathbb{R}^d} \quad & \frac{1}{2}\boldsymbol{x}^\top \boldsymbol{Q}\boldsymbol{x} + c^\top \boldsymbol{x} \\ \text{subject to} \quad & \boldsymbol{A}\boldsymbol{x} \preceq \boldsymbol{b}, \end{aligned}$$

where $\boldsymbol{A} \in \mathbb{R}^{m\times d}$, $\boldsymbol{b} \in \mathbb{R}^m$, $\boldsymbol{c} \in \mathbb{R}^d$, the square matrix $\boldsymbol{Q}$ is symmetric, positive definite.

- Dual problem

$$\begin{aligned} \max_{\boldsymbol{\lambda}\in\mathbb{R}^m} \quad & \left(-\frac{1}{2}(\boldsymbol{c} + \boldsymbol{A}^\top\boldsymbol{\lambda})^\top \boldsymbol{A}\boldsymbol{Q}^{-1}(\boldsymbol{c} + \boldsymbol{A}^\top\boldsymbol{\lambda}) - \boldsymbol{\lambda}^\top \boldsymbol{b} \right) \\ \text{subject to} \quad & \boldsymbol{\lambda} \succeq 0, \end{aligned}$$

where $\boldsymbol{\lambda} \in \mathbb{R}^m$.

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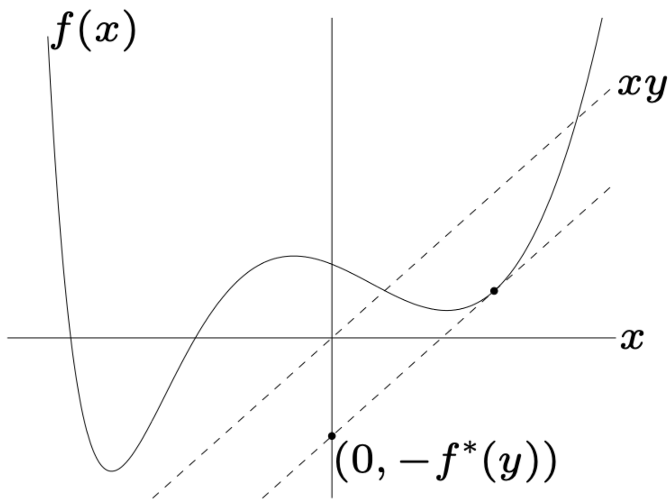
Conjugate Function: Definition and Meaning

- Given $f : \mathbb{R}^n \mapsto \mathbb{R}$, the conjugate function $f^* : \mathbb{R}^n \mapsto \mathbb{R}$ defined as:

$$f^*(\boldsymbol{y}) = \sup_{\boldsymbol{x} \in \text{dom } f} (\boldsymbol{y}^\top \boldsymbol{x} - f(\boldsymbol{x}))$$

with domain consisting of $\boldsymbol{y} \in \mathbb{R}^n$ for which the supremum is finite

- Assume $\mathbb{R}^1$.
- For a given slope of y, $yx - f(x)$ is the vertical distance between the line yx and $f(x)$.
- Thus, $f^*(y)$ is the maximum distance



Conjugate Function: Properties

- Given $f : \mathbb{R}^n \mapsto \mathbb{R}$, the conjugate function $f^* : \mathbb{R}^n \mapsto \mathbb{R}$ defined as:

$$f^*(\boldsymbol{y}) = \sup_{\boldsymbol{x} \in \text{dom } f} (\boldsymbol{y}^\top \boldsymbol{x} - f(\boldsymbol{x}))$$

  with domain consisting of $\boldsymbol{y} \in \mathbb{R}^n$ for which the supremum is finite

- $f^*(\boldsymbol{y})$: always convex (the pointwise supremum of a family of affine functions of $\boldsymbol{y}$)
- $f^{**} = f$ if f is convex and closed
- Fenchel's inequality: $f(\boldsymbol{x}) + f^*(\boldsymbol{y}) \geq \boldsymbol{x}^\top \boldsymbol{y}$ for all $\boldsymbol{x}, \boldsymbol{y}$ (by definition)
  - Example. $f(x) = |x|^2/2$. Then, $f^*(y) = |y|^2/2$. Thus, F-inequality tells us:

$$\frac{1}{2}(|x|^2 + |y|^2) \geq xy$$

Examples of Conjugate Functions

- $f(x) = ax + b,\ f^*(a) = -b$
- $f(x) = -\log x,\ f^*(y) = -\log(-y) - 1$ for $y < 0$
- $f(x) = e^x,\ f^*(y) = y \log y - y$
- $f(x) = x \log x,\ f^*(y) = e^{y-1}$
- $f(x) = \frac{1}{2}x^\top Q x,\ f^*(y) = \frac{1}{2}y^\top Q^{-1} y$ (Q is positive definite)
- $f(x) = \log \sum_{i=1}^n e^{x_i}$,

$$f^*(y) = \begin{cases} \sum_{i=1}^n y_i \log y_i & \text{if } y \succeq 0 \text{ and } \sum_{i=1}^n y_i = 1, \\ \infty & \text{otherwise} \end{cases}$$

Conjugate Function and Lagrangian Dual Function

- They are closely related. Consider the following problem:

$$\begin{array}{ll} \text{minimize} & f(\boldsymbol{x}) \\ \text{subject to} & \boldsymbol{A}\boldsymbol{x} \preceq \boldsymbol{b}, \\ & \boldsymbol{C}\boldsymbol{x} = \boldsymbol{d} \end{array}$$

- Using the conjugate of f, we can write the dual function as:

$$\begin{aligned} \mathcal{D}(\boldsymbol{\lambda}, \boldsymbol{\nu}) &= \inf_{\boldsymbol{x}} \left(f(\boldsymbol{x}) + \boldsymbol{\lambda}^\top (\boldsymbol{A}\boldsymbol{x} - \boldsymbol{b}) + \boldsymbol{\nu}^\top (\boldsymbol{C}\boldsymbol{x} - \boldsymbol{d}) \right) \\ &= -\boldsymbol{b}^\top \boldsymbol{\lambda} - \boldsymbol{d}^\top \boldsymbol{\nu} + \inf_{\boldsymbol{x}} \left(f(\boldsymbol{x}) + (\boldsymbol{A}^\top \boldsymbol{\lambda} + \boldsymbol{C}^\top \boldsymbol{\nu})^\top \boldsymbol{x} \right) \\ &= -\boldsymbol{b}^\top \boldsymbol{\lambda} - \boldsymbol{d}^\top \boldsymbol{\nu} - f^* \left(-\boldsymbol{A}^\top \boldsymbol{\lambda} - \boldsymbol{C}^\top \boldsymbol{\nu} \right) \end{aligned}$$

THANKS FOR YOUR ATTENTION

Discussions