

# Contiguous Memory Allocation

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# Project Overview

This project will involve managing a contiguous region of memory of size MAX where addresses may range from  $0 \dots \text{MAX}-1$ . Program must respond to four different requests:

1. Request for a contiguous block of memory
2. Release of a contiguous block of memory
3. Compact unused holes of memory into one single block
4. Report the regions of free and allocated memory



# 01

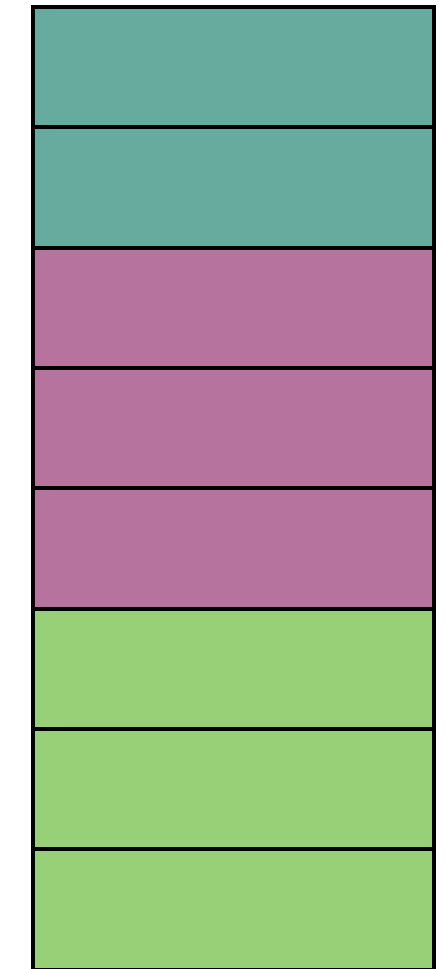
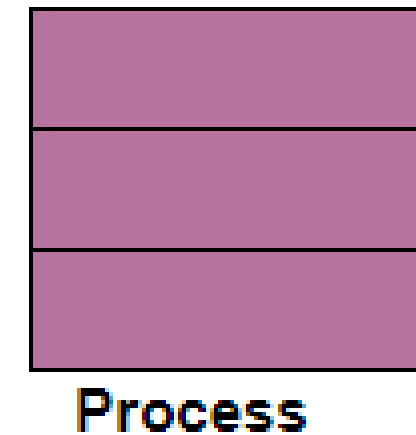
**Theoretical  
Foundation**

**Project**

**Advantages and  
Disadvantages**

# Definition:

- Contiguous Memory Allocation is a technique where the operating system allocates a single, continuous block of memory to a process.
- Advantages:
  - Simplifies memory management and access.
  - Ideal for systems with limited memory size where fast memory access is critical.



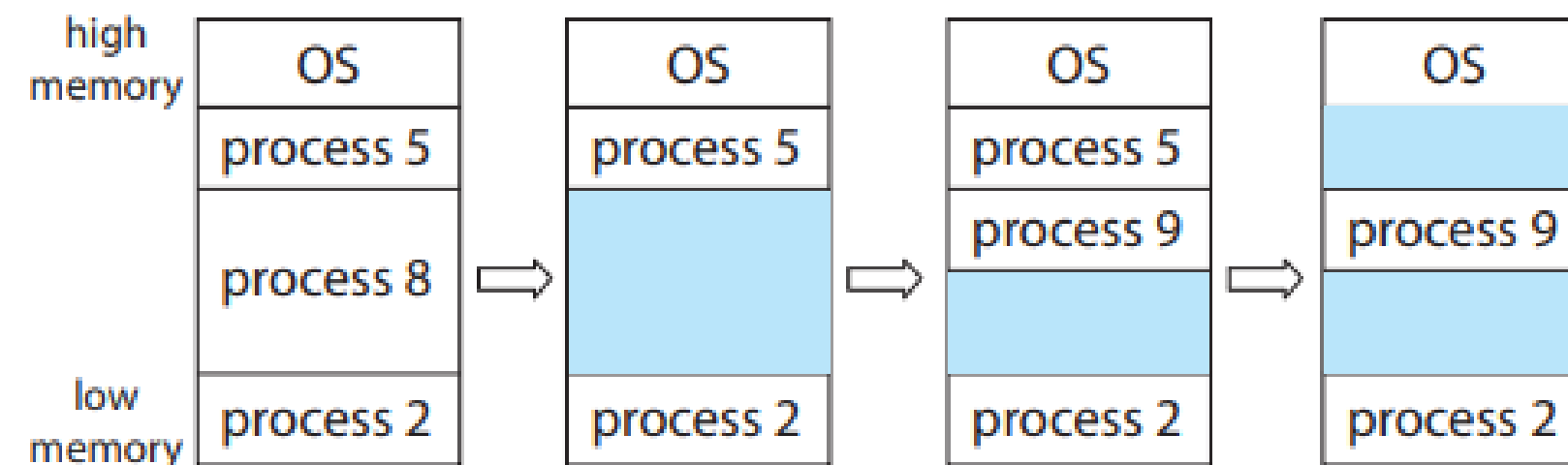
**Memory Blocks**

# Memory Allocation

- A method of assigning processes to variably sized memory partitions, where each partition holds exactly one process.
- **Management:**
  - The operating system maintains a table to track available and occupied memory.
  - Memory starts as one large block ("hole") and becomes fragmented as processes are added and removed.

# Partition Management

- **Initially:** Memory is a single large block ("hole").
- As processes arrive and leave: Holes of varying sizes are created.



# Memory Allocation and Release

- **Allocation:**
  - The OS considers process requirements and available memory to allocate space.
- **Release:**
  - Upon completion, memory is freed and reused.
  - Insufficient memory:
    - Reject the process with an error message.
    - Place the process in a wait queue for later allocation.

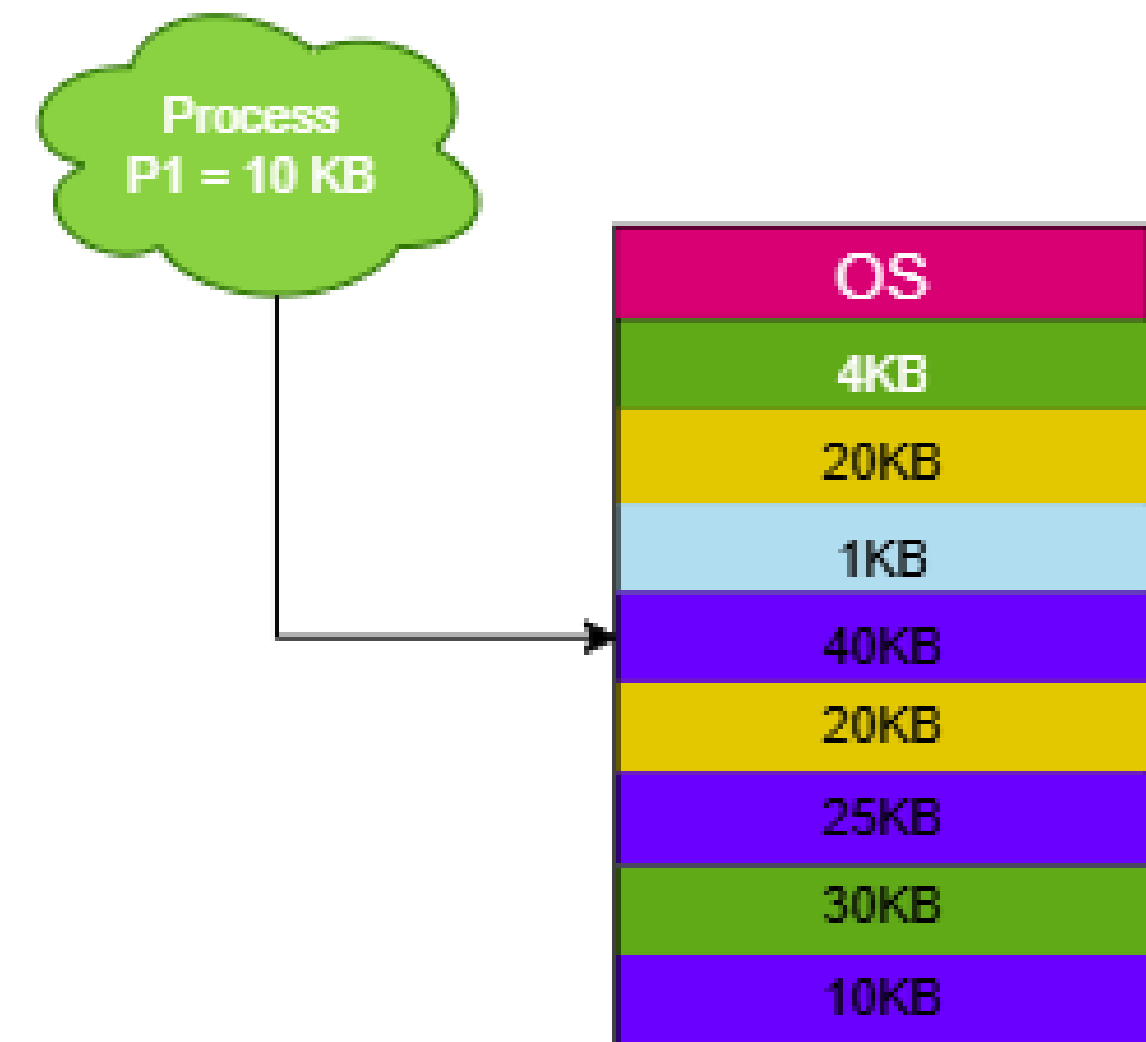
# Free Space Management

- **Fragmentation:**
  - Memory holes of varying sizes are scattered throughout memory.
- **Handling Large Holes:**
  - If a hole is too large:
    - Split into allocated and free space.
  - Adjacent free holes:
    - Merged into a larger block.

# Allocation Strategies:

- **Strategies to select memory holes:**
  - **First-fit:** Choose the first hole large enough.
  - **Best-fit:** Choose the smallest hole large enough.
  - **Worst-fit:** Choose the largest hole.
- **Comparison:**
  - First-fit and Best-fit are generally faster and more efficient than Worst-fit.

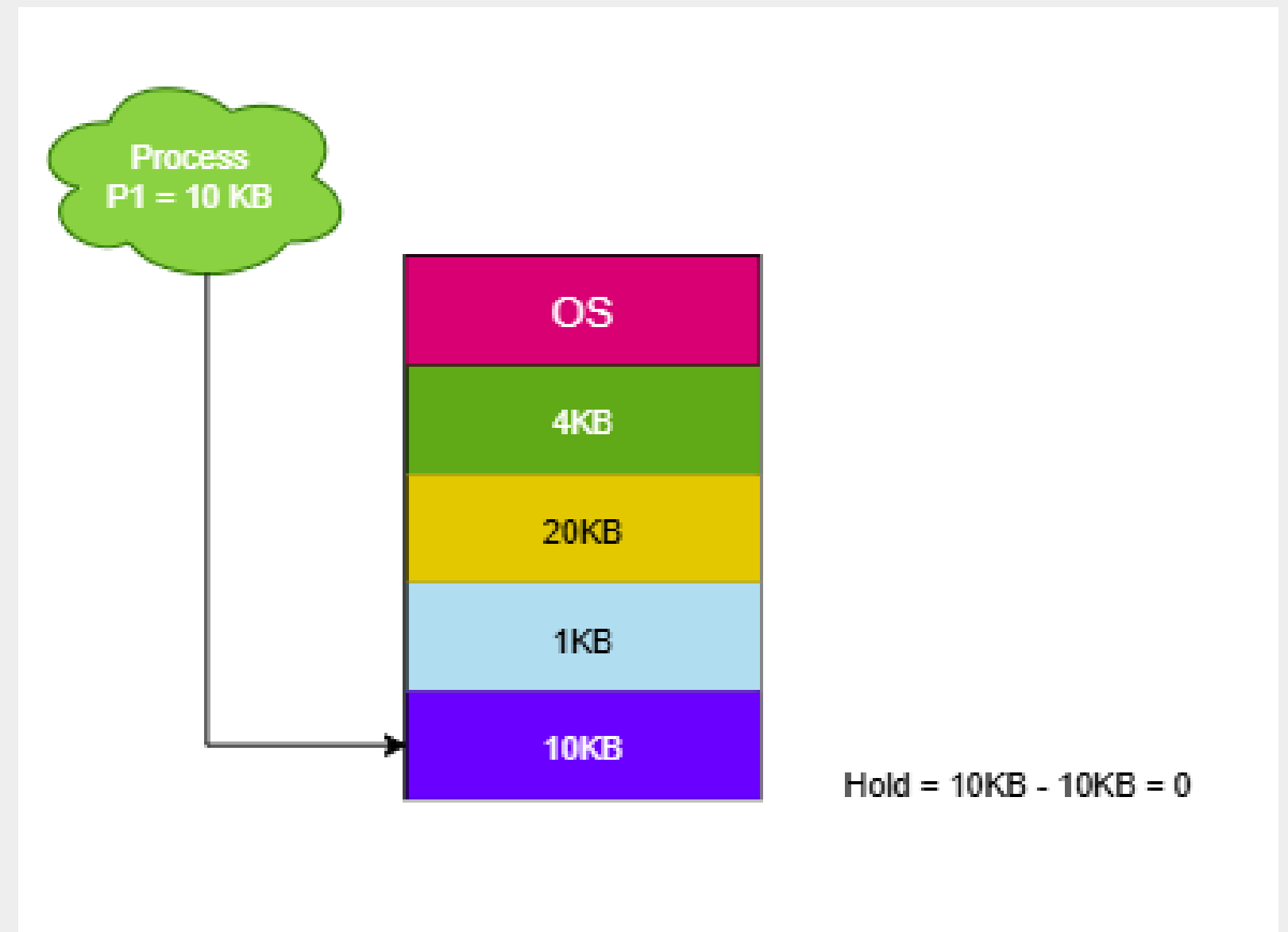
## FIRST-FIT



# Allocation Strategies:

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## BEST-FIT

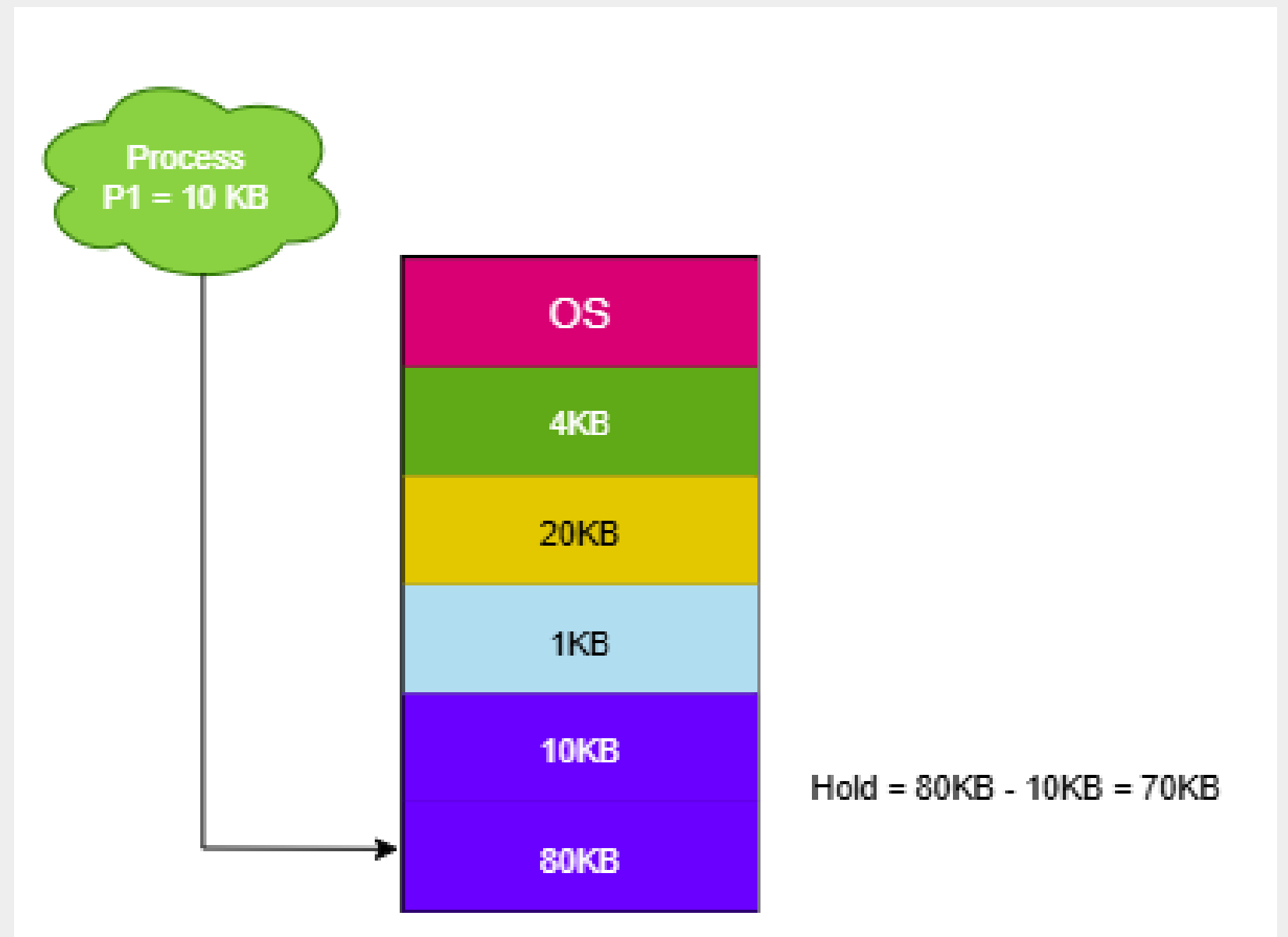




# Allocation Strategies:

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## WORST-FIT



**Theoretical  
Foundation**

# 02

**Project**

**Advantages and  
Disadvantages**

# Contiguous Memory Allocation

- **Purpose:**

- Simulate memory allocation and management processes.
- Handle requests, releases, and compaction of memory dynamically.

- **Initialization Command:**

*/allocator 1048576*

-> Allocates 1 MB (1,048,576 bytes) of memory.

- **User Interaction Prompt:**

*allocator>*

# Commands

- **RQ (Request):**
  - Request memory for a process.
  - Syntax: *RQ* <Process> <Bytes> <Strategy>
    - Example: *RQ P0 40000 W* (Worst-fit)
- **RL (Release):**
  - Release memory allocated to a process.
  - Syntax: *RL* <Process>
    - Example: *RL P0*
- **STAT (Status Report):**
  - Display allocated and unused memory regions.
  - Syntax: *STAT*
- **X (Exit):**
  - Terminate the program.

```
allocator>RQ P0 20 F
allocator>RQ P1 45 W
allocator>STAT
Addresses [0:19] Process P0
Addresses [20:64] Process P1
Addresses [65:199] Unused
allocator>RL P0
allocator>STAT
Addresses [0:19] Unused
Addresses [20:64] Process P1
Addresses [65:199] Unused
allocator>C
allocator>STAT
Addresses [0:44] Process P1
Addresses [45:199] Unused
allocator>X
Bye
```



**Theoretical  
Foundation**

**Project**

**Advantages and  
Disadvantages**





# Advantages:

- **Simplicity**
- **Efficiency**
- **Low Fragmentation**

# Disadvantages:

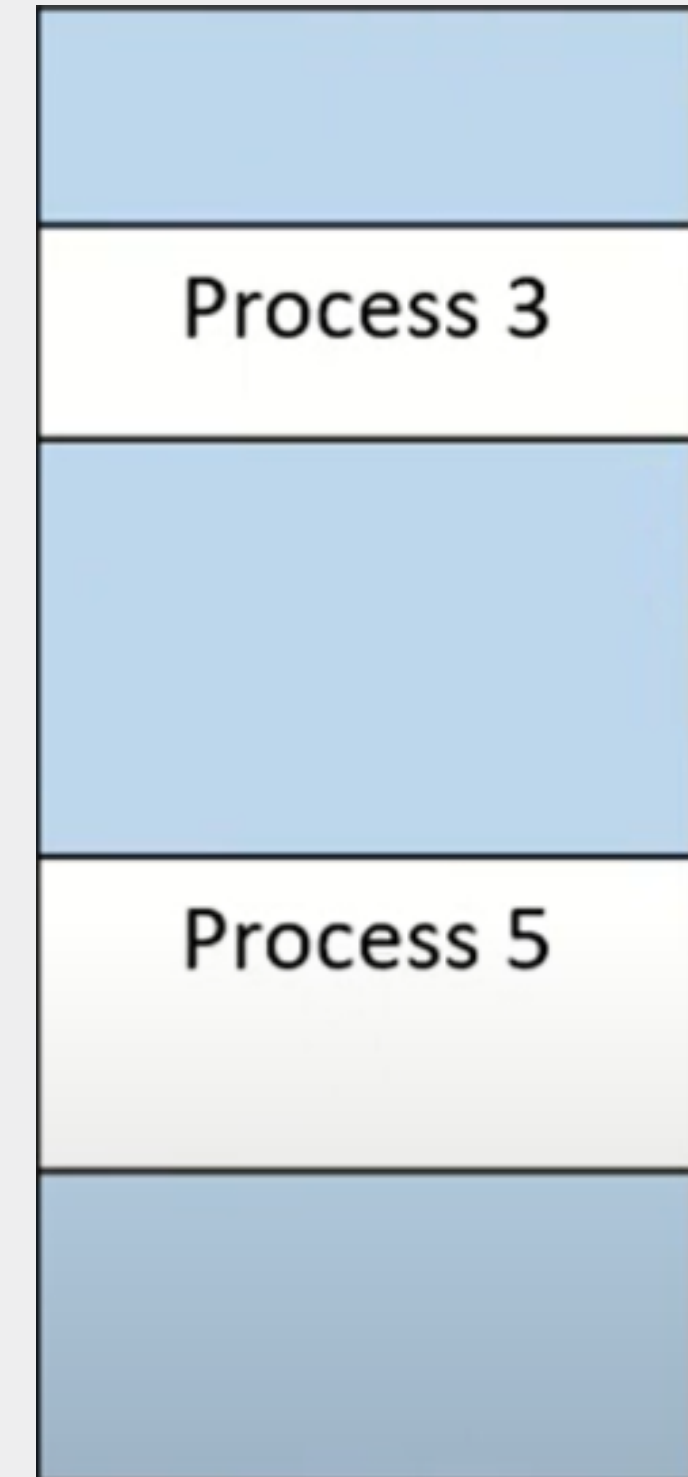
- **Limited Flexibility**
- **Memory Wastage**
- **Difficulty Managing Larger Memory Sizes**





# External Fragmentation:

- Occurs with both first-fit and best-fit strategies.
- Leads to fragmented, non-contiguous free regions.
- The "50-percent rule": 1/3 of memory can remain unused.



The background features several light gray squares of varying sizes scattered across the slide. There are two 2x2 grids of squares in the upper corners, a horizontal pair in the lower left, and a single square in the lower right.

Q&A

Thank You