

PURDUE UNIVERSITY

CS 699

SPRING 2016

Research Thesis

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March 1, 2016

Chapter 1

Fourier Basics

1.1 Vector Space of Functions on Boolean Hyper-cube

Definition 1.1 (Inner Product). Consider the 2^n -dimensional vector space of all functions $f : \{0, 1\}^n \rightarrow \mathbb{R}$. We define an inner product on this space by

$$\langle f, g \rangle := \mathbb{E}[f \cdot g] = \frac{1}{2^n} \sum_{x \in \{0, 1\}^n} f(x)g(x)$$

.

1.2 Characteristic Functions

Definition 1.2 (Characteristic function). For each $S \subseteq [n] = \{1, 2, \dots, n\}$, we define the characteristic function of S as

$$\chi_S(x) = (-1)^{S \cdot x}, \text{ where } S \cdot x = \sum_{i=1}^n S_i \cdot x_i = \sum_{i \in S} x_i$$

.

Lemma 1.3. For every $S \subseteq [n]$,

$$\sum_{x \in \{0, 1\}^n} \chi_S(x) = \begin{cases} 2^n & \text{if } S = \emptyset \\ 0 & \text{if } S \neq \emptyset \end{cases}$$

Proof. If $S = \emptyset$, then $S \cdot x = 0$. So $\sum_{x \in \{0, 1\}^n} \chi_S(x) = \sum_{x \in \{0, 1\}^n} 1 = 2^n$.

If $S \neq \emptyset$, then there exists k such that $S_k \neq 0$. Hence,

$$\begin{aligned}
\sum_{x \in \{0,1\}^n} \chi_S(x) &= \sum_{x \in \{0,1\}^n} (-1)^{\sum_{i \in S} x_i} \\
&= \sum_{x \in \{0,1\}^n} [(-1)^{x_k} \cdot (-1)^{\sum_{i \in S \setminus \{k\}} x_i}] \\
&= \sum_{x_k \in \{0,1\}} (-1)^{x_k} \cdot \sum_{x \setminus x_k \in \{0,1\}^{n-1}} (-1)^{\sum_{i \in S \setminus \{k\}} x_i} \\
&= [(-1)^0 + (-1)^1] \sum_{x \setminus x_k \in \{0,1\}^{n-1}} (-1)^{\sum_{i \in S \setminus \{k\}} x_i} \\
&= 0
\end{aligned}$$

□

Theorem 1.4. For every $S, T \subseteq [n]$,

$$\langle \chi_S, \chi_T \rangle = \begin{cases} 1 & \text{if } S = T \\ 0 & \text{if } S \neq T \end{cases}$$

Proof.

$$\langle \chi_S, \chi_T \rangle = \frac{1}{2^n} \sum_{x \in \{0,1\}^n} (-1)^{S \cdot x + T \cdot x} = \frac{1}{2^n} \sum_{x \in \{0,1\}^n} (-1)^{(S \Delta T) \cdot x}$$

where Δ is the symmetric different between two sets S and T .

$S \Delta T = \emptyset$ if and only if $S = T$. Hence, our goal follows immediately from Lemma 1.3. □

1.3 Fourier Basis

Theorem 1.5. The set of all χ_S defines an orthonormal basis for the space of all real-valued function on $\{0,1\}^n$

Proof. From Theorem 1.4, the set of all χ_S is an orthonormal set. Also, there are 2^n different χ_S . Hence, the set of all χ_S must be an orthonormal basis for the space of all real-valued functions on $\{0,1\}^n$. □

The set of all χ_S is called the *Fourier basis*.

1.4 Fourier Transform

Definition 1.6. For each $S \subseteq [n]$, we define the Fourier transform of f at S as following:

$$\hat{f}(S) := \mathbb{E}[f \cdot \chi_S] = \langle f, \chi_S \rangle$$

Definition 1.7 (Fourier transform). The mapping $\mathcal{F} : f \mapsto \widehat{f}$ is called the Fourier transform.

If we view functions f, \widehat{f} as N -dimensional vectors, then we can write the Fourier transform as the product of f and some matrix F as following:

$$\begin{aligned}\mathcal{F}(f) &= f \cdot F \\ &= \frac{1}{N} (f(0), f(1), \dots, f(N-1)) \begin{pmatrix} \chi_0(0) & \chi_1(0) & \cdots & \chi_{N-1}(0) \\ \chi_0(1) & \chi_1(1) & \cdots & \chi_{N-1}(1) \\ \vdots & \vdots & \ddots & \vdots \\ \chi_0(N-1) & \chi_1(N-1) & \cdots & \chi_{N-1}(N-1) \end{pmatrix} \\ &= (\widehat{f}(0), \widehat{f}(1), \dots, \widehat{f}(N-1)) \\ &= \widehat{f}\end{aligned}$$

where $F_{ij} = \frac{\chi_i(j)}{N}$. Since $\chi_i(j) = \chi_j(i)$, F is symmetric.

Lemma 1.8. F is invertible

Proof. We have

$$(F \cdot F)_{ij} = \frac{1}{N^2} \sum_{k=0}^{N-1} \langle \chi_i(k), \chi_j(k) \rangle$$

Based on Theorem 1.4, it is easy to see that

$$(F \cdot F)_{ij} = \begin{cases} \frac{1}{N^2} & \text{if } i = j \\ 0 & \text{otherwise} \end{cases}$$

So $F \cdot F = N^2 \cdot I$, which implies that F is invertible □

Theorem 1.9. The mapping \mathcal{F} is linear.

Proof. This follows from the properties of inner product. For any S ,

$$\widehat{af + bg}(S) = \langle af + bg, \chi_S \rangle = \langle af, \chi_S \rangle + \langle bg, \chi_S \rangle = a \langle f, \chi_S \rangle + b \langle g, \chi_S \rangle = a \widehat{f}(S) + b \widehat{g}(S)$$

Thus, $\widehat{af + bg} = a \widehat{f} + b \widehat{g}$, which means \mathcal{F} is linear.

Here is another way to show that.

$$\widehat{af + bg} = \mathcal{F}(af + bg) = (af + bg)F = afF + bgF = a(fF) + b(gF) = a \widehat{f} + b \widehat{g}$$

□

Theorem 1.10. The linear map \mathcal{F} is a bijection.

Proof. It suffices to show that F is invertible, which follows immediately from Lemma 1.8. □

1.5 Parseval's Identity

Since the set of χ_S forms an orthonormal basis,

$$f = \sum_S \widehat{f}(S) \chi_S. \quad (1.1)$$

Hence,

$$\langle f, g \rangle = \sum_S \widehat{f}(S) \widehat{g}(S) \quad (1.2)$$

In particular, when $f = g$ we get Parseval's identity:

$$\|f\|_2^2 = \sum_S \widehat{f}(S)^2 \quad (1.3)$$

This also implies:

$$\|f - g\|_2^2 = \sum_S (\widehat{f}(S) - \widehat{g}(S))^2 \quad (1.4)$$

1.6 Convolution

Definition 1.11. Given any two function f and $g : \{0, 1\}^n \rightarrow \mathbb{R}$, the convolution of $f * g : \{0, 1\}^n \rightarrow \mathbb{R}$ is defined as

$$(f * g)(x) := \frac{1}{2^n} \sum_{y \in \{0, 1\}^n} f(x \oplus y) g(y)$$

Theorem 1.12. If X and Y are n -bits random independent variables with probability distributions f and g , respectively, then $2^n(f * g)$ is the distribution of the random variable $Z = X \oplus Y$.

Proof.

$$\begin{aligned} \Pr[Z = z] &= \Pr[X = z \oplus Y] \\ &= \sum_{y \in \{0, 1\}^n} \Pr[X = z \oplus y | Y = y] \\ &= \sum_{y \in \{0, 1\}^n} \Pr[X = z \oplus y] \cdot \Pr[Y = y] \\ &= \sum_{y \in \{0, 1\}^n} f((z \oplus y)) \cdot g(y) \\ &= 2^n(f * g)(z) \end{aligned}$$

□

Theorem 1.13. For every $S \subseteq [n]$,

$$\widehat{f * g}(S) = \widehat{f}(S) \cdot \widehat{g}(S)$$

Proof.

$$\begin{aligned}
\widehat{f * g}(S) &= \frac{1}{2^n} \sum_x (f * g)(x) \chi_S(x) \\
&= \frac{1}{2^n} \sum_x \left(\frac{1}{2^n} \sum_y f(x \oplus y) g(y) \right) \chi_S(x) \\
&= \frac{1}{2^{2n}} \sum_x \sum_y f(x \oplus y) g(y) \chi_S(x \oplus y) \chi_S(y) \\
&= \frac{1}{2^n} \sum_x f(x \oplus y) \chi_S(x \oplus y) \left(\frac{1}{2^n} \sum_y g(y) \chi_S(y) \right) \\
&= \widehat{f}(S) \cdot \widehat{g}(S)
\end{aligned}$$

□

Intuitively, the convolution $f * g$ is the product of the Fourier transforms of f and g .

Theorem 1.14. *Let V be a subspace of dimension k of $\{0, 1\}^n$ and let V^\perp be the dual of V . Define*

$$f(x) = \begin{cases} \frac{1}{2^k} & \text{if } x \in V \\ 0 & \text{otherwise.} \end{cases}$$

Then

$$\widehat{f}(S) = \begin{cases} \frac{1}{N} & \text{if } S \in V^\perp \\ 0 & \text{otherwise.} \end{cases}$$

Proof. Suppose v_1, v_2, \dots, v_k is a basis of V .

Claim 1.15. $V = \langle v_1 \rangle \oplus \langle v_2 \rangle \oplus \dots \oplus \langle v_k \rangle$

Claim 1.16. $\langle v_1 \rangle^\perp \cap \dots \cap \langle v_k \rangle^\perp = \langle v_1, \dots, v_k \rangle^\perp = V^\perp$

Let $f_i = \begin{cases} \frac{1}{2} & \text{if } S \in \langle v_i \rangle \\ 0 & \text{otherwise} \end{cases}$, then $\widehat{f}_i = \begin{cases} \frac{1}{N} & \text{if } S \in \langle v_i \rangle^\perp \\ 0 & \text{otherwise.} \end{cases}$

From above claims, we immediately obtain following result.

Claim 1.17. $f = f_1 \oplus f_2 \oplus \dots \oplus f_k$

Hence,

$$\widehat{f}(S) = N^{k-1} \widehat{f}_1(S) \dots \widehat{f}_k(S)$$

If $S \in V^\perp$, then $S \in \langle v_i \rangle^\perp$ for every i , so $\widehat{f}(S) = N^{k-1} \cdot \left(\frac{1}{N}\right)^k = \frac{1}{N}$.

If $S \notin V^\perp$, then there exists some i such that $S \notin \langle v_i \rangle^\perp$, which implies $\widehat{f}_i(S) = 0$. Hence, $\widehat{f}(S) = 0$ □

Chapter 2

Min Entropy

Let $X = (x_0, x_1, \dots, x_{N-1})$ be a distribution function of a random variable over $\{0, 1\}^n$, where $N := 2^n$.

Definition 2.1 (Min Entropy). We define the min entropy of X as follow.

$$H_\infty(X) := \min_i (-\log x_i)$$

This implies that if $H_\infty(X) \geq k$ then $x_i \leq \frac{1}{2^k}$ for every $0 \leq i \leq N - 1$

Theorem 2.2 (Collision Probability). *If we sample X twice, then the probability we get the same result, denoted $Col(X)$, is $\sum_{i=0}^{N-1} x_i^2 = N \cdot \|X\|_2^2$.*

Definition 2.3 (Flat Distribution). A probability distribution function $f: \{0, 1\}^n \rightarrow (0, 1)$ is a T -flat if there $\exists S \subseteq \{0, 1\}^n$ such that $|S| = T$ and $f(x) = \begin{cases} \frac{1}{T} & \text{if } x \in S \\ 0 & \text{otherwise} \end{cases}$

Lemma 2.4. For every $\alpha \geq \beta$, every α -flat distribution can be written as the sum of β -flat distributions.

Theorem 2.5. *For every integer $k \geq 0$, if $H_\infty(X) \geq k$, then $X = \sum \alpha_i X_i$, where each X_i is a 2^k -flat, $\alpha_i \in [0, 1]$ for every i , and $\sum_i \alpha_i = 1$.*

Proof. Let S be the set of all the probability distributions X with $H_\infty(X) \geq k$, then S is a compact (closed and bounded) convex set in \mathbb{R}^N .

Claim 2.6. The set of all 2^k -flat distributions is the set of all extreme points of S .

First, every 2^k -flat distribution X is an extreme point of S since if $X = \alpha Y + (1 - \alpha)Z$ for some $Y, Z \in S$ and $\alpha \in [0, 1]$, then $X = Y = Z$. Next, we want to show that for any X , which is not a 2^k -flat distribution, X is not an extreme point of S . Since X is not a 2^k -flat distribution, there exist $i < j$ such that $0 < x_i, x_j < \frac{1}{2^k}$. So we can choose $\delta > 0$ such that $x_i + \delta, x_j + \delta \leq \frac{1}{2^k}$ and $x_i - \delta, x_j - \delta \geq 0$. Now let $y_k = z_k = x_k$ for every $k \neq i, k \neq j$, $y_i = x_i + \delta, y_j = x_j - \delta, z_i = x_i - \delta$, and $z_j = x_j + \delta$. Then $Y, Z \in S$ and $X = \frac{1}{2}Y + \frac{1}{2}Z$, which implies that X is not an extreme point.

Back to the problem, since S is a compact convex set, the set of all convex combinations of its vertices is identical to S . Hence, every distribution X with $H_\infty(X) \geq k$ can be written as a convex combination of k -flat distributions. \square

Theorem 2.7. *If $H_\infty(X) \geq k$, then $Col(X) \leq \frac{1}{2^k}$.*

Proof. By theorem 2.5, we can write X as $X = \sum_i \alpha_i X_i$, where each X_i is a 2^k -flat, $\sum \alpha_i = 1$, and $\alpha_i \in [0, 1]$ for every i . It is obvious that $Col(X_i) = \|X_i\|_2^2 = \frac{1}{2^k}$. Collision functions are convex, so by Jensen's inequality,

$$Col(X) = Col\left(\sum_i \alpha_i X_i\right) \leq \sum_i \alpha_i \cdot Col(X_i) = \sum_i \alpha_i \frac{1}{2^k} = \frac{1}{2^k} \sum_i \alpha_i = \frac{1}{2^k}$$

\square

Theorem 2.8. *If $H_\infty(X) \geq k$, then $\sum_S \widehat{X}(S)^2 \leq \frac{1}{N \cdot 2^k}$.*

This follows immediately from the Parseval's identity.

Definition 2.9 (Small Bias Distribution). Let \mathcal{D} be a probability distribution function over $\{0, 1\}^n$. We say that \mathcal{D} is α -bias if $\widehat{\mathcal{D}}(S) \leq \frac{\alpha}{N}$.

Definition 2.10. *Statistical Different* between two distributions A and B is defined as follow:

$$SD(A, B) = \frac{1}{2} \sum_i |a_i - b_i|$$

Theorem 2.11. *Let \mathcal{D} be a small bias distribution with $\widehat{\mathcal{D}}(S) \leq \frac{\alpha}{N}$ for all S , let \mathcal{M} be a min entropy source such that $H_\infty(\mathcal{M}) \geq k$, and let \mathcal{U} be the uniform distribution over n -bits string. Then*

$$SD(\mathcal{D} \oplus \mathcal{M}, \mathcal{U}) \leq \frac{\alpha \sqrt{N}}{2^{1+k/2}}$$

Proof.

$$\begin{aligned} SD(\mathcal{D} \oplus \mathcal{M}, \mathcal{U}) &= \frac{1}{2} \sum_i |(\mathcal{D} \oplus \mathcal{M})(i) - \mathcal{U}(i)| \\ &\leq \frac{1}{2} \sqrt{N \sum_i [(\mathcal{D} \oplus \mathcal{M})(i) - \mathcal{U}(i)]^2} \\ &= \frac{1}{2} \sqrt{N^2 \cdot \|(\mathcal{D} \oplus \mathcal{M}) - \mathcal{U}\|_2^2} \\ &= \frac{N}{2} \sqrt{\sum_S [\widehat{\mathcal{D} \oplus \mathcal{M}}(S) - \mathcal{U}(S)]^2} \\ &= \frac{N}{2} \sqrt{\sum_{S \neq \emptyset} \widehat{\mathcal{D} \oplus \mathcal{M}}(S)^2} \end{aligned}$$

By convolution,

$$\begin{aligned}
\sum_{S \neq \emptyset} \widehat{\mathcal{D} \oplus \mathcal{M}}(S)^2 &= \sum_{S \neq \emptyset} N^2 \cdot \widehat{\mathcal{D} * \mathcal{M}}(S)^2 \\
&= N^2 \sum_{S \neq \emptyset} \widehat{\mathcal{D}}(S)^2 \cdot \widehat{\mathcal{M}}(S)^2 \\
&\leq N^2 \cdot \sum_{S \neq \emptyset} \left(\frac{\alpha}{N}\right)^2 \cdot \widehat{\mathcal{M}}(S)^2 \\
&= \alpha^2 \cdot \sum_{S \neq \emptyset} \widehat{\mathcal{M}}(S)^2 \\
&\leq \frac{\alpha^2}{N \cdot 2^k}
\end{aligned}$$

Hence,

$$SD(\mathcal{D} \oplus \mathcal{M}, \mathcal{U}) \leq \frac{\alpha \sqrt{N}}{2^{1+k/2}}$$

□

Theorem 2.12. *Let M be a distribution with min entropy k over $\{0, 1\}^n$, let G_0, G be 1 by $\frac{n}{2}$ and $\frac{n}{2}$ by n matrices respectively, and let X be a distribution over $\{0, 1\}^{\frac{n}{2}}$. Then*

$$SD\{(XG_0, XG \oplus M, G_0, G), (U, XG \oplus M, G_0, G)\} \leq \frac{1}{2^k}$$

Proof. For simplicity, let $A = (XG_0, XG \oplus M|G_0, G)$ and $B = (U, XG \oplus M|G_0, G)$.

$$\begin{aligned}
&SD\{(XG_0, XG \oplus M, G_0, G), (U, XG \oplus M, G_0, G)\} \\
&= \mathbb{E}_{G_0, G} [SD(A, B)] \\
&= \mathbb{E}_{G_0, G} \left[\frac{1}{2} \sum_{i, j} |A(i) - B(i)| \right] \\
&\leq \frac{1}{2} \mathbb{E}_{G_0, G} \left[\sqrt{2N \cdot \sum_i ((A(i) - B(i))^2)} \right] \\
&\leq \frac{\sqrt{2N}}{2} \sqrt{\mathbb{E}_{G_0, G} \left[\sum_i (A(i) - B(i))^2 \right]} \\
&= \frac{\sqrt{2N}}{2} \sqrt{\sum_S \mathbb{E}_{G_0, G} [(\widehat{A}(S) - \widehat{B}(S))^2]}
\end{aligned}$$

□

Appendix A

A.1 Dual of a Vector Space

Definition A.1 (Dual space). Let V be a subspace of $\{0, 1\}^n$. We define the dual of V as $V^\perp = \{x \in \{0, 1\}^n \mid x \cdot v = 0 \ \forall v \in V\}$.

Theorem A.2. V^\perp is a subspace of $\{0, 1\}^n$.

Proof. For any $x, y \in V^\perp, a \in \{0, 1\}, (a \cdot x + y) \cdot v = a \cdot (x \cdot v) + y \cdot v = 0 + 0 = 0$. \square

Lemma A.3. $\sum_{i:\text{even}}^t \binom{n}{i} = \sum_{i:\text{odd}}^t \binom{n}{i} = 2^{t-1}$.

Theorem A.4. For any subspace V of dimension k of $\{0, 1\}^n$, there exists a unique dual space V^\perp of dimension $(n - k)$.

Proof. We will show that $|V^\perp| = 2^{n-k}$ by induction on k .

If $k = 0$, then $V = \{0\}$. Clearly, $V^\perp = \{0, 1\}^n$.

If $k = 1$, let $V = \{0, v\}$. Suppose the number of $v_i = 1$ is t , then the number of x such that $x \cdot v = 0$ is $\sum_{i: 2|t-i} \binom{n}{i} 2^{n-t} = 2^{t-1} \cdot 2^{n-t} = 2^{n-1}$ by Lemma A.3.

Suppose that there exists a unique orthogonal subspace V^\perp of dimension $(n - k + 1)$ for any subspace V of dimension $k - 1$ of $\{0, 1\}^n$, where $k \geq 2$.

Let $V = \langle v_1, v_2, \dots, v_k \rangle$, $S_1 = \langle v_1, v_2, \dots, v_{k-1} \rangle$, and $S_2 = \langle v_k \rangle$. Then, $V^\perp = S_1^\perp \cap S_2^\perp$.

Suppose $\dim(V^\perp) = t$. We want to show $t = n - k$.

By induction hypothesis, $\dim(S_1^\perp) = n - k + 1$ and $\dim(S_2^\perp) = n - 1$.

If $t \leq n - k - 1$, then we need $[(n - k + 1) - t]$ independent vectors to cover S_1^\perp from extending V^\perp , and we need $[(n - 1) - t]$ independent vectors to cover S_2^\perp from extending V^\perp . Since $S_1^\perp \cup S_2^\perp \subseteq \{0, 1\}^n$, we must have $[(n - k + 1) - t] + [(n - 1) - t] + t \leq n$, which is equivalent to $t \geq n - k$, contradiction.

If $t \geq n - k + 1$, then $S_1^\perp \subseteq S_2^\perp$, this is impossible since v_k is independent from v_1, v_2, \dots, v_{k-1} . Thus, $t = n - k$. So $|V^\perp| = 2^{n-k}$. \square