# CSE 301 – DATABASE

Lecture 10

# Chapter 5: Relational Database Design

#### **Functional Dependencies**

- A functional dependency (FD) is a relationship or association between two attributes (typically between the Primary Key and other non-key attributes within a table).
- For any relation R, attribute Y is functionally dependent on attribute X (usually the Primary Key), if for every valid instance of X, that value of X uniquely determines the value of Y.
- $\square$  This relationship is indicated by  $X \to Y$
- ☐ The left side of the above FD diagram is called the determinant, and the right side is the dependent.

## **Examples of Functional Dependencies**

R

A	В
1	6
2	7
3	8
4	9

 $\square$  A  $\rightarrow$  B, if for every valid instance of A, that value of A uniquely determines the value of B.

$$\square$$
 {1 \to 6}, {2 \to 7}, {3 \to 8}, {4 \to 9}

 $\square$  *B* has the same value for the same value as *A*.

R

A	В	
1	6	
2	7	
3	8	
2	9	

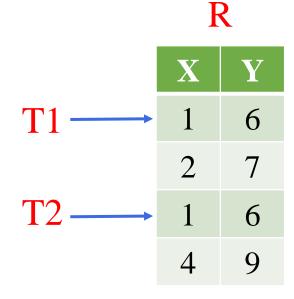
 $\square$  A  $\not\rightarrow$  B, if for every valid instance of A, that value of A not uniquely determines the value of B.

$$\square$$
 {1  $\rightarrow$  6}, {2  $\rightarrow$  7}, {3  $\rightarrow$  8}, {2  $\rightarrow$  9}

 $\square$  B has not the same value for the same value as A.

#### **Functional Dependencies**

- $\square$  We can mathematically represent  $X \to Y$ ,
  - ightharpoonup When  $X \subseteq R$  and  $Y \subseteq R$
  - ightharpoonup If T<sub>1</sub>[X] = T<sub>2</sub>[X], then T<sub>1</sub>[Y] = T<sub>2</sub>[Y]



### **Examples of Functional Dependencies**

■ Which of the following FD is not valid?

$$a) A \rightarrow B$$

$$(b) B \rightarrow C$$

c) 
$$BC \rightarrow A$$

$$d)AC \rightarrow B$$

R

A	В	C	
1	2	3	
4	2	3	
5	3	3	

- $\Box$   $A \to B$ ,  $\{1 \to 2\}$ ,  $\{4 \to 2\}$ ,  $\{5 \to 3\}$  *Valid*
- $\square$   $B \to C$ ,  $\{2 \to 3\}$ ,  $\{2 \to 3\}$ ,  $\{3 \to 3\}$  *Valid*
- $\square$  BC  $\rightarrow$  A,  $\{2,3 \rightarrow 1\}$ ,  $\{2,3 \rightarrow 4\}$ ,  $\{3,3 \rightarrow 5\}$  Not Valid
  - > Right hand side has not the same value for the same value as left hand side.
- $\square$   $AC \to B, \{1,3 \to 2\}, \{4,3 \to 2\}, \{5,3 \to 3\}$  Valid

#### **Examples of Functional Dependencies**

☐ Which of the following FD is valid?

a) 
$$XY \rightarrow Z$$
,  $Z \rightarrow Y$ 

$$b) XZ \rightarrow X, Y \rightarrow Z$$

c) 
$$YZ \rightarrow X$$
,  $Z \rightarrow X$ 

$$d) XZ \rightarrow Y, Y \rightarrow Z$$

R

c) 
$$YZ \rightarrow X$$
,  $Z \rightarrow X$ 

Option c is not valid

a)  $XY \rightarrow Z$ ,  $Z \rightarrow Y$ 

$$\square$$
  $XY \rightarrow Z$ 

$$\square$$
  $Z \rightarrow Y$ 

$$\begin{cases} 3 \to 4 \}, \{3 \to 5 \}, \\ \{3 \to 6 \}, \{2 \to 2 \} \end{cases}$$

Option a is not valid

b)  $XZ \rightarrow X, Y \rightarrow Z$ 

$$\square$$
  $XZ \rightarrow X$ 

$$\begin{cases} 1,3 \to 1 \}, & \{1,3 \to 1 \}, \\ \{4,3 \to 4 \}, & \{3,2 \to 3 \} \end{cases}$$

$$\square$$
  $Y \rightarrow Z$ 

$$\begin{cases} 4 \to 3 \}, \{5 \to 3 \}, \\ \{6 \to 3 \}, \{2 \to 2 \} \end{cases}$$

Option b is valid

 $d) XZ \rightarrow Y, Y \rightarrow Z$ 

$$\square$$
  $XZ \rightarrow X$ 

$$\begin{cases} 1,3 \to 4 \}, & \{1,3 \to 5 \}, \\ \{4,3 \to 6 \}, & \{3,2 \to 2 \} \end{cases}$$

$$\square$$
  $Y \rightarrow Z$ 

$$\{4 \to 3\}, \{5 \to 3\}, \{6 \to 3\}, \{2 \to 2\}$$

Option d is not valid Introduction: 1-7

#### **Exercise on Functional Dependencies**

■ Which of the following FD is correct?

$$a) A \rightarrow BC$$

b) 
$$DE \rightarrow C$$

c) 
$$C \rightarrow DE$$

$$d)$$
  $BC \rightarrow A$ 

A	В	C	D	E
A	2	3	4	5
2	A	3	4	5
A	2	3	6	5
a	2	3	6	6

R

☐ Which of the following FD is not correct? R

a) 
$$XY \rightarrow Z$$
,  $Z \rightarrow Y$ 

b) 
$$YZ \rightarrow X$$
,  $Y \rightarrow Z$ 

c) 
$$XZ \rightarrow X$$
,  $Z \rightarrow X$ 

$$d) XZ \rightarrow Y, Y \rightarrow Z$$

#### **Classification of Functional Dependencies**

- ☐ Trivial functional dependency
  - $\checkmark$  A  $\rightarrow$  B has trivial functional dependency if B is a subset of A (B  $\subseteq$  A).
  - $\checkmark$  Examples:  $A \rightarrow A$ ,  $AB \rightarrow B$ ,  $\{Employee\_id, Employee\_Name\} \rightarrow Employee\_Id$
- Non-trivial functional dependency
  - $\checkmark$  A  $\rightarrow$  B has a non-trivial functional dependency if B is not a subset of A.
  - ✓ If there is at least one attribute in right hand side that is not present in the left hand side.
  - ✓ When A intersection B is NULL, then  $A \rightarrow B$  is called as complete non-trivial.
  - $\checkmark$  Examples:  $AB \rightarrow BC$ ,  $\{ID\} \rightarrow \{ID, DOB\}$ ,  $\{Roll, Name\} \rightarrow \{Roll, Name, Phone\}$

#### Classification of Functional Dependencies

- ☐ Fully functional dependency
  - ✓ Given R and A → B, then B is fully functional dependent on A if there is no Z where Z is a proper subset of A ( $Z \subset A$ ) such that  $Z \to B$ .
  - $\checkmark$  Examples:  $\{AB \rightarrow C, A \rightarrow D\}$  is a fully FD,  $\{AB \rightarrow C, A \rightarrow C\}$  is not a fully FD
- ☐ Partial functional dependency
  - ✓ Given a relation R with FD F defined on the attributes of R and K as a candidate key, if A is a proper subset of K and if and only if  $X \to A$ , then A said to be partially dependent on K.
  - $\checkmark$  Examples: R(ABCD),  $C.K. \to AB$ ,  $FD: \{A \to C\}$ . C is partially dependent on A.

#### **Classification of Functional Dependencies**

- ☐ Transitive functional dependency
  - ✓ If  $A \rightarrow B$  and  $B \rightarrow C$ , then C is transitively functional dependent on A such that  $A \rightarrow C$ .
  - $\checkmark$  Examples:  $\{AB \rightarrow C, C \rightarrow D\}, So, \{AB \rightarrow D\}$

#### **Armstrong's axioms / Inference Rules**

- ☐ Reflexive Rule (IR1)
  - ✓ If Y is a subset of X (X  $\supseteq$  Y), then X determines Y (X  $\rightarrow$  Y).
  - $\checkmark$  Examples:  $\{AB \rightarrow A\}$ ,  $\{Employee\_id, Employee\_Name\} \rightarrow \{Employee\_Name\}$
- ☐ Augmentation Rule (IR2)
  - $\checkmark$  If X determines Y (X  $\rightarrow$  Y), then XZ determines YZ (XZ  $\rightarrow$  YZ). for any Z.
  - $\checkmark$  Examples: R(ABCD), if  $A \rightarrow B$  then  $AC \rightarrow BC$
- ☐ Transitive Rule (IR3)
  - ✓ If X determines Y (X → Y) and Y determine Z (Y → Z), then X must also determine Z (X  $\rightarrow$  Z).

#### **Armstrong's axioms / Inference Rules**

- Union Rule (IR1)
  - ✓ If X determines Y (X → Y) and X determines Z (X → Z), then X must also determine Y and Z (X → YZ).
  - $\checkmark$  Examples: R(ABCD), if  $A \rightarrow B$  and  $A \rightarrow C$  then  $A \rightarrow BC$ .
- ☐ Decomposition Rule (IR2)
  - ✓ If X determines Y and Z (X  $\rightarrow$  YZ), then X determines Y (X  $\rightarrow$  Y) and X determines Z (Y  $\rightarrow$  Z) separately.
  - $\checkmark$  Examples: R(ABCD), if  $A \rightarrow BC$ , then  $A \rightarrow B$  and  $A \rightarrow C$
- Pseudo Transitive Rule (IR3)
  - ✓ If X determines Y (X → Y) and YZ determines W (YZ → W), then XZ determines W (XZ → W).

#### **Closure Set of Attributes**

 $\square$  X<sup>+</sup> is the set of all attributes that can be determine using the given set X(attributes).

Or

- The Closure Of Functional Dependency means the complete set of all possible attributes that can be functionally derived from given functional dependency using the inference rules known as Armstrong's Rules.
- $\square$  If "F" is a functional dependency then closure of functional dependency can be denoted using " $\{F\}$ +".

#### **Closure Set of Attributes**

- ☐ There are three steps to calculate closure of functional dependency. These are:
  - ✓ Step-1 : Add the attributes which are present on Left Hand Side in the original functional dependency.
  - ✓ Step-2 : Now, add the attributes present on the Right Hand Side of the functional dependency.
  - $\checkmark$  Step-3:
    - With the help of attributes present on Right Hand Side, check the other attributes that can be derived from the other given functional dependencies.
    - Repeat this process until all the possible attributes which can be derived are added in the closure.

• Consider the table Student\_details having (Roll\_no, Name, Marks, Location) as the attributes and having two functional dependencies.

```
FD1 : Roll_no → Name, Marks
FD2 : Name → Marks, Location
```

- Now, we will calculate the closure of all the attributes present in the relation using the three steps mentioned below. Find closure set of attributes of {Roll\_no}+
- Step-1: add attributes present on the LHS of the first functional dependency to the closure.
  {Roll\_no}+= {Roll\_no}
- Step-2: add attributes present on the RHS of the original functional dependency to the closure.

```
{Roll_no} + = {Roll_no, Name, Marks}
```

- Step-3: Add the other possible attributes which can be derived using attributes present on the RHS of the closure.
  - So, Roll\_No attribute cannot functionally determine any attribute, but Name attribute can determine other attributes such as Marks and Location using 2<sup>nd</sup> Functional Dependency.
  - Therefore, complete closure of Roll\_No will be:

```
{Roll_no}+= {Roll_No, Marks, Name, Location}
```

- Example 1: R (ABCDEFG)
  - $\{A \rightarrow B, BC \rightarrow DE, AEG \rightarrow G\}$
  - Find  $(AC)^+ = ?$
  - $(AC)^+ = AC$
  - $= ABC (A \rightarrow B)$
- = ABCDE (BC  $\rightarrow$  DE)
- = ABCDE

- Example 2: R (ABCDE)
  - $\{A \rightarrow BC, CD \rightarrow E, B \rightarrow D, E \rightarrow A\}$
  - Find  $(B)^+ = ?$
- $(B)^+ = B$
- $= BD (B \rightarrow D)$
- = BD

- Example 3: R (ABCDEF)
  - {AB  $\rightarrow$  C, CD  $\rightarrow$  E, DE  $\rightarrow$  B}
  - Find  $(AB)^+ = ?$
- $(AB)^+ = AB$
- =ABC (AB $\rightarrow$ C)

- Example 1: R (ABCDEFGH)
  - {A  $\rightarrow$  BC, CD  $\rightarrow$  E, E  $\rightarrow$  C, D  $\rightarrow$  AEH, ABH  $\rightarrow$  BD, DH  $\rightarrow$  BC}
  - Find  $(BCD)^+ \rightarrow H$ ?
  - $(BCD)^+ = BCD$
  - = BCDE (CD  $\longrightarrow$  E)
  - = ABCDEH (D  $\rightarrow$  AEH)
  - So,  $(BCD)^+ \rightarrow H$  is valid

- Exercise 1: R (ABCDEF)
  - $\{A \rightarrow BC, CD \rightarrow EF, B \rightarrow D, E \rightarrow A\}$
  - Find  $(AE)^+ = ?$
- Exercise 2: R (ABCDEF)
  - $\blacksquare \{AB \to C, CD \to E, DE \to B\}$
  - Find  $(AB)^+ = ?$
- Exercise 3: Student (Roll, Name, DoB, Phone, Course)
  - Roll, Name} → Phone, {Course, DoB} → Roll, Course
     → Name, Roll → Name}
  - Find  $(Course)^+ = ?$

#### **Equivalence Of Functional Dependencies**

R (ACDEH)

• F: 
$$\{A \rightarrow C, AC \rightarrow D, E \rightarrow AD, E \rightarrow H\}$$

• G: 
$$\{A \rightarrow CD, E \rightarrow AH\}$$

$$\bullet (A)^+ = ACD$$

$$\bullet (AC)^+ = ACD$$

• 
$$(E)^+ = EAH = EAHCD$$

• So, 
$$F \subseteq G$$

So, 
$$G \subseteq F$$

• 
$$F \subset G$$
 and  $G \subset F$ , so,  $F = G$ 

Find the correct option

a) 
$$F \subseteq G$$

b) 
$$G \subset F$$

c) 
$$F = G$$

d) 
$$F \neq G$$

#### **Equivalence Of Functional Dependencies**

- R (ACDEH)
- F:  $\{A \rightarrow C, AC \rightarrow D, E \rightarrow AD, E \rightarrow H\}$
- G:  $\{A \rightarrow CD, E \rightarrow AH\}$
- Check both FD's are equivalent or not.

- R (ABC)
- F:  $\{A \rightarrow B, B \rightarrow C, C \rightarrow A\}$
- G:  $\{A \rightarrow BC, B \rightarrow A, C \rightarrow A\}$
- Find the correct option

a) 
$$F \subseteq G$$

b) 
$$G \subseteq F$$

c) 
$$F = G$$

d) 
$$F \neq G$$

- R (VWXYZ)
- F:  $\{W \to X, WX \to Y, Z \to WY, Z \to V\}$ 
  - G:  $\{W \rightarrow XY, Z \rightarrow WX\}$
  - Find the correct option

a) 
$$F \subseteq G$$

b) 
$$G \subseteq F$$

$$c) F = G$$

d) 
$$F \neq G$$