

# Lecture 22: Clustering, LLoyd's algorithm (k-means), spectral clustering

Nisha Chandramoorthy

November 14, 2023

## Last time: Johnson-Lindenstrauss lemma

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- ▶ Let  $0 < \epsilon < 1/2$ ,  $m > 4$ . Then, there exists a linear map  $A : \mathbb{R}^d \rightarrow \mathbb{R}^n$  with  $n = O(\epsilon^{-2} \log m)$  such that for all  $x_i, x_j \in X$ ,  $i, j \in [m]$ ,  $(1 - \epsilon)\|x_i - x_j\|^2 \leq \|Ax_i - Ax_j\|^2 \leq (1 + \epsilon)\|x_i - x_j\|^2$ .

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- ▶ Informal: any set of points in high-dimensional space can be mapped to a lower-dimensional space while approximately preserving the distances between the points.

# Proof

- Distortion by Gaussian random matrices: for any  $x \in \mathbb{R}^d$ , when the entries  $A_{ij}$  are iid standard Gaussian,

$$\begin{aligned}\mathbb{P}(n(1 - \epsilon)\|x\|^2 \leq \|Ax\|^2 \leq n(1 + \epsilon)\|x\|^2) \\ \geq 1 - 2 \exp(-(\epsilon^2 - \epsilon^3)n/4).\end{aligned}$$

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- ▶ Then, deterministic statement of J-L lemma follows from union bound over all  $m^2$  pairs of points.

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- ▶ Let  $A$  be a  $n \times d$  matrix with iid standard Gaussian entries. Then,  $E[(Ax)_j] = 0$  and  $\text{Var}((Ax)_j) = \|x\|^2$ , for all  $j \leq n$ .

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- ▶ Chi-squared distribution:  
$$\rho(x) = \frac{1}{2^{n/2} \Gamma(n/2)} x^{n/2-1} e^{-x/2}, x \geq 0.$$
- ▶ Models sum of squares of  $n$  independent standard normal random variables.

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- ▶ (Candes, Romberg, Tao 2005) If  $x$  is  $s$ -sparse, then,

$$x = \operatorname{argmin}_{z \in \mathbb{R}^d} \|z\|_1 \quad \text{s.t.} \quad Ax = Az. \quad (2)$$

# Convolutional Neural Networks (source: cs231n.stanford.edu)

- ▶ Suitable for image recognition. Won the 2012 ImageNet competition and subsequent ones.
- ▶ Three types of layers: convolutional, FC, pooling
- ▶ Convolutional layer: accepts a volume of size  $W_1 \times H_1 \times D_1$  and outputs a volume of size  $W_2 \times H_2 \times D_2$  where  $W_2 = (W_1 - F + 2P)/S + 1$  and  $H_2 = (H_1 - F + 2P)/S + 1$  and  $D_2 = K$ .
- ▶  $K$  is number of filters,  $F$  is filter size,  $S$  is stride,  $P$  is padding.
- ▶ Pooling layer: downsamples along width and height, and optionally along depth.
- ▶ FC layer: computes class scores, resulting in volume of size  $1 \times 1 \times K$ .

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- ▶ Closely related to dimensionality reduction.
- ▶ Definition of clustering depends on the definition of distance between points.
- ▶ Center-based clustering:  $k$  centers  $\mu_1, \dots, \mu_k \in \mathbb{R}^d$ .

# Lloyd's algorithm

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- ▶ Given clusters  $C_1, \dots, C_k$ , update centers  $\mu_1, \dots, \mu_k \in \mathbb{R}^d$  as

$$\mu_j = \frac{1}{|C_j|} \sum_{x_i \in C_j} x_i.$$

# k-means algorithm (Lloyd's algorithm)

- ▶ Lloyd's algorithm is an approximate method to solve the ERM problem:

$$\min_{C_1, \dots, C_k} \sum_{j=1}^k \sum_{x_i \in C_j} \|x_i - \mu(C_j)\|^2.$$

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- ▶ Lloyd's algorithm is a heuristic. It is not guaranteed to converge to the global optimum or even a local minimum.

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- ▶ Since

$$\mu_j^{(t)} = \frac{1}{|C_j^{(t)}|} \sum_{x_i \in C_j^{(t)}} x_i = \operatorname{argmin}_{\mu \in \mathbb{R}^d} \sum_{x_i \in C_j^{(t)}} \|x_i - \mu\|^2,$$

$$\sum_{x_i \in C_j^{(t)}} \|x_i - \mu_j^{(t)}\|^2 \leq \sum_{x_i \in C_j^{(t)}} \|x_i - \mu_j^{(t-1)}\|^2, \quad \forall j \in [k].$$



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- Proof (contd.): by definition of  $C_j^{(t)}$ ,

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- ▶ Summing over  $j \in [k]$ ,

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- Thus, the ERM objective decreases at each iteration.

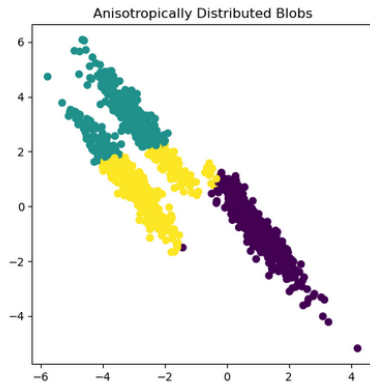
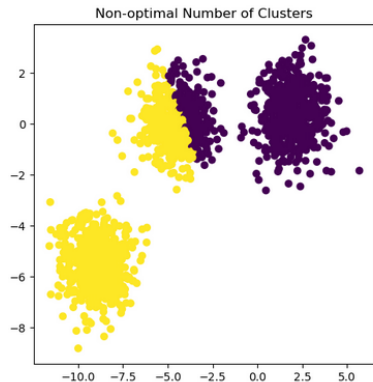
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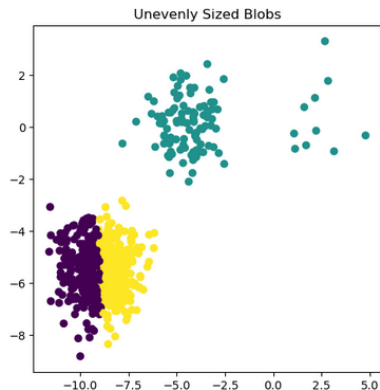
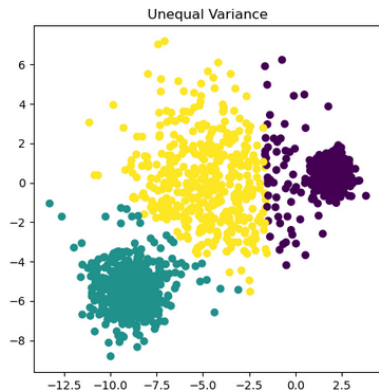
- ▶ k-means algorithm is sensitive to initialization of the centers.
- ▶ Complexity:  $O(mdk)$  per iteration, where  $m$  is the number of points,  $d$  is the dimension, and  $k$  is the number of clusters.

# k-means failure modes



Source: [sklearn's toy examples](#)

# k-means failure modes contd



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# Spectral clustering

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- ▶ ERM problem:  $\min_{C_1, \dots, C_k} \sum_{j=1}^k \sum_{x_i \in C_j} \sum_{x_l \notin C_j} w_{il}$ . Graph min-cut problem.

# RatioCut problem: spectral clustering solution

► RatioCut problem:  $\min_{C_1, \dots, C_k} \sum_{j=1}^k \frac{\sum_{x_i \in C_j} \sum_{x_l \notin C_j} w_{il}}{|C_j|}.$

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- ▶  $h_i$  ( $i$ th column of  $H$ ) is nonzero at row  $j$  if  $x_j$  is in cluster  $i$ .
- ▶  $H$  has orthonormal columns.

## Recall: graphical representation of $X$

- ▶ Choose weighting, such as,  $w_{ij} = \exp(-\|x_i - x_j\|^2/2\sigma^2)$ . As  $\sigma \rightarrow 0$ ,  $w_{ij} \rightarrow \mathbb{1}_{i=j}$ . The  $m \times m$  matrix  $W$  is the adjacency matrix of a graph.

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- ▶ Let  $D$  be the diagonal matrix with  $D_{ii} = \sum_{j=1}^m w_{ij}$ .
- ▶ Graph laplacian:  $L = D - W$ .
- ▶ Detects local structure / clusters in data.

# Lemma proof: RatioCut objective and graph laplacian connection

► RatioCut objective( $C_1, \dots, C_k$ )

$$:= \sum_{j=1}^k \frac{\sum_{x_i \in C_j} \sum_{x_l \notin C_j} w_{il}}{|C_j|}.$$



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- ▶ For any vector  $v$ ,  $v^\top L v = (1/2) \sum_{i,j=1}^m w_{ij} (v_i - v_j)^2$ .
- ▶  $L$  is positive semi-definite.

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- ▶ Kernel PCA with  $K = L^\dagger$  is equivalent to Laplacian eigenmaps.



# Final project deliverables

→ Proposal 10.1. (2<sup>nd</sup> Nov)

→ Presentation 25.1.

= 1 per team 5-6 minutes

3-4 minutes

3 sessions

1 overflow session

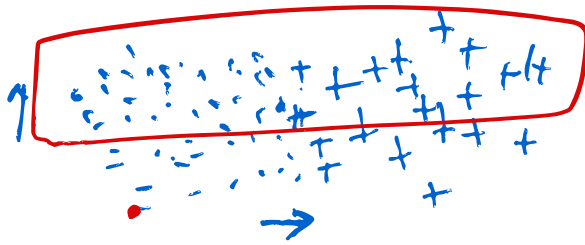
} Zoom  
4-8<sup>th</sup>

7-9 pm

Finalize presentation schedules by  
14<sup>th</sup> Nov

Rules for attendance ( % of points  
for asking questions)

→ Reports : 7<sup>th</sup> Dec  
(65.1.)

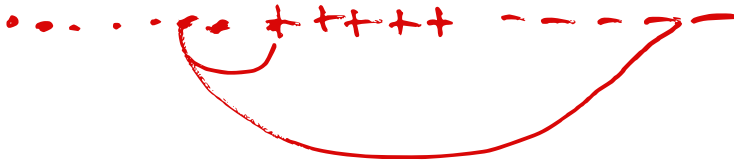


$k=2$

$\bullet \rightarrow 1$

$+$   $\rightarrow 2$

Distance b/w two points



$$\rightarrow w_{ij} = e^{-\frac{\|x_i - x_j\|^2}{2\sigma^2}}$$

$$L = D - W$$

$$v^T L v = \frac{1}{2} \sum_{i,j \in [m]} w_{ij} (v_i - v_j)^2$$

$L \in \mathbb{R}^{m \times m}$        $v \in \mathbb{R}^m$

Laplacian  
eigen maps  
objective

Rayleigh quotient

$$\min_{\|v\|=1} v^T L v = \lambda_{\min}(L)$$

$$\|v\|=1$$

$$\min_{V \in \mathbb{R}^{m \times k}} V^T L V = \min_{\substack{v_1, v_2, \dots, v_k \\ v_i \cdot v_j = 0}} \sum_{i=1}^k v_i^T L v_i$$

$\|v_i\|=1$

$v_i$ :  $i^{\text{th}}$  column of  $V$

$$\operatorname{argmin} \sum_{i=1}^k v_i^T L v_i = \left\{ \text{bottom } k \text{ eigenvectors of } L \right\}$$