Lecture 22: Clustering, LLoyd's algorithm (k-means), spectral clustering

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Last time: Johnson-Lindenstrauss lemma

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- Let $0 < \epsilon < 1/2$, m > 4. Then, there exists a linear map $A : \mathbb{R}^d \to \mathbb{R}^n$ with $n = O(\epsilon^{-2} \log m)$ such that for all $x_i, x_j \in X$, $i, j \in [m], (1 \epsilon) \|x_i x_j\|^2 \leqslant \|Ax_i Ax_j\|^2 \leqslant (1 + \epsilon) \|x_i x_j\|^2$.

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- ▶ Informal: any set of points in high-dimensional space can be mapped to a lower-dimensional space while approximately preserving the distances between the points.

Proof

▶ Distortion by Gaussian random matrices: for any $x \in \mathbb{R}^d$, when the entries A_{ii} are iid standard Gaussian,

$$\mathbb{P}(n(1-\epsilon)\|x\|^2 \leqslant \|Ax\|^2 \leqslant n(1+\epsilon)\|x\|^2)$$

$$\geqslant 1 - 2\exp(-(\epsilon^2 - \epsilon^3)n/4).$$

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► Then, deterministic statement of J-L lemma follows from union bound over all *m*² pairs of points.

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- Chi-squared distribution: $\rho(x) = \frac{1}{2^{n/2}\Gamma(n/2)} x^{n/2-1} e^{-x/2}, x \geqslant 0.$
- Models sum of squares of n independent standard normal random variables.

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► (Candes,Romberg, Tao 2005) If *x* is *s*-sparse, then,

$$x = \operatorname{argmin}_{z \in \mathbb{R}^d} ||z||_1 \quad \text{s.t.} \quad Ax = Az.$$
 (2)



Convolutional Neural Networks (source: cs231n.stanford.edu)

- Suitable for image recognition. Won the 2012 ImageNet competition and subsequent ones.
- Three types of layers: convolutional, FC, pooling
- Convolutional layer: accepts a volume of size $W_1 \times H_1 \times D_1$ and outputs a volume of size $W_2 \times H_2 \times D_2$ where $W_2 = (W_1 F + 2P)/S + 1$ and $H_2 = (H_1 F + 2P)/S + 1$ and $D_2 = K$.
- K is number of filters, F is filter size, S is stride, P is padding.
- Pooling layer: downsamples along width and height, and optionally along depth.
- ► FC layer: computes class scores, resulting in volume of size 1 × 1 × K.



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- Closely related to dimensionality reduction.
- Definition of clustering depends on the definition of distance between points.
- ► Center-based clustering: k centers $\mu_1, \ldots, \mu_k \in \mathbb{R}^d$.

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▶ Given clusters C_1, \ldots, C_k , update centers $\mu_1, \ldots, \mu_k \in \mathbb{R}^d$ as

$$\mu_j = \frac{1}{|C_j|} \sum_{x_i \in C_j} x_i.$$



k-means algorithm (Lloyd's algorithm)

Lloyd's algorithm is an approximate method to solve the ERM problem:

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- Since $\mu_j^{(t)} = \frac{1}{|C_j^{(t)}|} \sum_{x_i \in C_j^{(t)}} x_i = \operatorname{argmin}_{\mu \in \mathbb{R}^d} \sum_{x_i \in C_j^{(t)}} ||x_i \mu||^2$,

$$\sum_{x_i \in \mathcal{C}_j^{(t)}} \|x_i - \mu_j^{(t)}\|^2 \leqslant \sum_{x_i \in \mathcal{C}_j^{(t)}} \|x_i - \mu_j^{(t-1)}\|^2, \quad \forall j \in [k].$$

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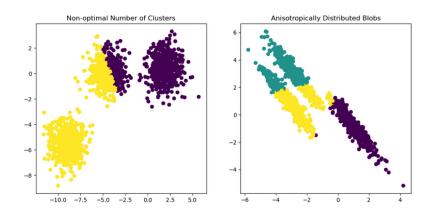
Thus, the ERM objective decreases at each iteration.

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Lloyd's algorithm properties

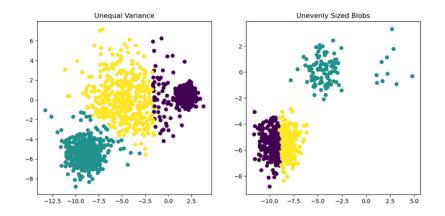
- k-means algorithm is sensitive to initialization of the centers.
- Complexity: O(mdk) per iteration, where m is the number of points, d is the dimension, and k is the number of clusters.

k-means failure modes



Source: sklearn's toy examples

k-means failure modes contd



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- ► ERM problem: $\min_{C_1,...,C_k} \sum_{j=1}^k \sum_{x_i \in C_j} \sum_{x_l \notin C_j} w_{il}$. Graph min-cut problem.

RatioCut problem: spectral clustering solution

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- H has orthonormal columns.

► Choose weighting, such as, $w_{ij} = \exp(-\|x_i - x_j\|^2/2\sigma^2)$. As $\sigma \to 0$, $w_{ij} \to \mathbb{1}_{i=j}$. The $m \times m$ matrix W is the adjacency matrix of a graph.

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- ▶ Let *D* be the diagonal matrix with $D_{ii} = \sum_{j=1}^{m} w_{ij}$.
- ▶ Graph laplacian: L = D W.
- Detects local structure / clusters in data.

Lemma proof: RatioCut objective and graph laplacian connection

▶ RatioCut objective(C_1, \dots, C_k)

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- L is positive semi-definite.

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- ► Kernel PCA with $K = L^{\dagger}$ is equivalent to Laplacian eigenmaps.