

# Data Structures and Algorithms

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# What is Algorithm?

As per Donald Knuth

## Algorithm

A definite, effective and finite process that receives input and produces an output

**Definite** : steps are clear, concise and unambiguous

**Effective** : you can perform each operation precisely

**Finite** : finite number of steps

## Analysis

When two programs solve the same problem, Analysis is finding answer to the question which one is **better**?

# Better on?

- Readability :

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# Better on?

- Readability : changes with programming language
- Number of Lines : changes with programming language
- Amount of computing resources : changes with programming language
- Run time : changes with processor speed, compiler and programming language

# An example : Checking the run time

our first example

# Big-O Notation

## Requirement

To characterize an algorithm's efficiency in terms of execution time, independent of any particular program or computer

## Solution

To quantify the algorithm in terms of number of operations or steps

$$T(n)$$

$T(n)$  is a function that indicates the time an algorithm takes to solve a problem of size  $n$

# Example 1

```
1  def sum_of_n(n):  
2      total = 0  
3      for i in range(n):  
4          total+=i  
5      return total  
6  
7
```

- For **sum\_of\_n**, we can take the basic compute step as the assignment operations
- In **sum\_of\_n** following are the assignment operations
  - **sum = 0**
  - **sum+ = n**
- $T(n) = n + 1$
- We are only interested in the dominant term in  $T(n)$ , because as  $n$  increases faster compared to other terms, i.e it overpowers the rest

## Big-O

The dominant term in  $T(n)$ , which can be termed as **order of magnitude** function. Big-O  $\implies$  Biggest Order



frametitleCommon Big-O functions

# Quiz 1

What is the Big-O for the program given below :

```
1  a=5
2  b=6
3  c = 10
4  for i in range(n):
5      for j in range(n):
6          x=i*i
7          y=j*j
8          z=i*j
9  for k in range(n):
10     w = a * k + 45
11     v=b*b
12 d = 33
13
```

# Quiz 1

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```

$$T(n) = 3n^2 + 2n + 4 \implies O(n^2)$$

# Anagram

A

string is an anagram of another if second is simply a rearrangement of the first. For example **python** and **typhon**

# Solution 1: Checking off

```
1  def anagram_sol1(word1, word2):
2      word2_list = list(word2)
3      index1=0
4      is_anagram = True
5      while index1 < len(word1) and is_anagram:
6          index2=0
7          is_continue = True
8          while index2 < len(word2_list) and
is_continue:
9              if word1[index1] == word2_list[index2]:
10                 word2_list[index2] = None
11                 is_anagram = True
12                 is_continue = False
13             else:
14                 is_anagram = False
15                 is_continue = True
16                 index2+=1
17             index1+=1
18         return is_anagram
19
```

# Solution 1: Big-O

## Example

Each letter in word1 has to iterate a maximum of  $n$  locations to find a match. That is

$$\sum_{i=1}^n i = \frac{n(n+1)}{2} = \frac{n^2 + n}{2}$$

## Solution 2: Sort and Compare

```
20 def anagram_sol2(word1, word2):
21     word1_list = list(word1)
22     word2_list = list(word2)
23     word1_list.sort()
24     word2_list.sort()
25     index=0
26     while index < len(word1_list):
27         if word1_list[index] != word2_list[index]:
28             return False
29         index+=1
30     return True
31
```

## Solution 2: Big-O

There is a single iteration of  $n$  if there is a match but the sort operation takes the precedence due to it's  $O(n^2)$  or  $O(n \log n)$  complexity



## Solution 3: Brute Force

This tries to exhaust all possibilities. Here we generate all possible anagrams of **word1** and matches this with **word2**. For a word of length  $n$ , there are  $n!$

## Solution 4: Count and Compare

```
1  def anagram_sol4(word1, word2):
2      counter1 = [0]*26
3      counter2 = [0]*26
4      offset = ord('a')
5      for letter in word1:
6          counter1[ord(letter)-offset]+=1
7      for letter in word2:
8          counter2[ord(letter)-offset]+=1
9      for index in range(26):
10         if counter1[index] != counter2[index]:
11             return False
12     return True
13
```

## Solution 4: Count and Compare

```
1  def anagram_sol4(word1, word2):
2      counter1 = [0]*26
3      counter2 = [0]*26
4      offset = ord('a')
5      for letter in word1:
6          counter1[ord(letter)-offset]+=1
7      for letter in word2:
8          counter2[ord(letter)-offset]+=1
9      for index in range(26):
10         if counter1[index] != counter2[index]:
11             return False
12     return True
13
```

$$T(n) = 2n + 26 \implies O(n)$$

The solution above can run in linear time but required more space requirements than the other solutions

## What is Linear Data Structure?

- Data structures in which each element stays its position relative to the elements before and after.
- Examples are Stacks, Queues, Deques and Lists
- The difference are in the way items are added or removed

## What is a Stack?

A **stack** is an ordered collection of items where the addition of new item and the removal of existing items always takes place at the same end

- also known as **Last-In-First-Out(LIFO)**
- Newer items are in the top and older items are at the bottom
- Browser **Back** button is an example of stack

# Stack-ADT: Python Implementation

- **Stack()** : creates new stack. Needs no parameters rather returns an empty stack
- **push(item)** : adds a new item to top of the stack. It needs the item and returns nothing
- **pop()** : removes the top item from the stack. It needs no parameter and return the item. stack is modified
- **peek()** : returns top item from the stack. No item parameter is required and did not modify the stack
- **is\_empty()** : tests to see whether stack is empty. It needs no parameter and returns a boolean value
- **size()** : returns the number of items in the stack. It needs no parameter and returns an integer value

# Stack ADT: Python Implementation

```
1 class Stack:
2     def __init__(self) -> None:
3         self.items = []
4
5     def is_empty(self):
6         return self.items == []
7
8     def push(self, item):
9         return self.items.append(item)
10
11    def pop(self):
12        return self.items.pop()
13
14    def peek(self):
15        return self.items[-1]
16
17    def size(self):
18        return len(self.items)
19
```

# Stack Example: Paranthesis Checker

```
1 def para_check(symbol_str):
2     # creating symbol templates and the key - value
    pairs
3     open_symbol = {'(':0, '[':1, '{':2}
4     close_symbol = {')':0, ']':1, '}':2}
5     # initialising the stack
6     para_stack = Stack()
7     #to check the unbalance condition during symbol
    iteration and exit the loop is exists
8     is_check = True
9     index = 0
10    while index < len(symbol_str) and is_check:
11        symbol = symbol_str[index]
12        #checking for opening symbols
13        if symbol in open_symbol.keys():
14            para_stack.push(symbol)
```



```

1         #checking for closing symbols
2         elif symbol in close_symbol.keys():
3             # check is stack is empty and still a close
symbol exists resulting in unbalance
4             if para_stack.is_empty():
5                 # exit the loop if exists
6                 is_check = False
7                 # check if the closing symbol is equal to
opening symbol , if not, exit the loop
8                 elif close_symbol[symbol] != open_symbol[
para_stack.peek()] :
9                     is_check = False
10                else:
11                    # if both of the above condition fails,
then there is a balance exists and pop out the symbol
12                    para_stack.pop()
13

```

```
1         else:
2             pass
3             index+=1
4             # to check if the stack is empty and check the
balance flag
5             if para_stack.is_empty() and is_check:
6                 return f"balanced"
7             else:
8                 return f"unbalanced"
```