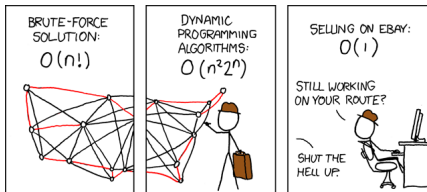


CS319: Scientific Computing (with MATLAB)

Sorting, Complexity, and structs

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Week 7: **9am and 4pm**, 22 Feb 2023



<http://xkcd.com/399>

Important: you should read:

- ▶ Learning MATLAB, Section 6.7 (Structures).
<https://doi-org.nuigalway.idm.oclc.org/10.1137/1.9780898717662>
- ▶ The MATLAB Guide, Chapters 18 and 19:
<https://doi-org.nuigalway.idm.oclc.org/10.1137/1.9781611974669>

In-Class test (updated)

We'll have an in-class test **10am** on Wednesday of next week (Week 8). There will be no 4pm lecture that day.

- ▶ Sample question at <https://www.niallmadden.ie/2223-CS319/CS319-SampleTest.pdf>
- ▶ Sample solutions are available on Blackboard.

This week, CS319 will be concerned with...

- 1 1: A note on complexity
- 2 2: Merge Sort
 - Why is Merge Sort is fast
- 3 3: Comparing in practice
- 4 4: The Password Problem
 - Algorithm (high-level)
 - Implementation
- 5 5: Structures
 - Example: pchip

1: A note on complexity

In Lab 4 (right after this class) you'll be provided with code for a sorting algorithm, Bubble Sort. You'll then be challenged to write a “better” sorting function, using **Merge Sort**.

But what does “better” mean?

There are many ways that one algorithm could be considered superior to another, for example:

- ▶ takes less time to run;
- ▶ takes less memory to run;
- ▶ takes less time to program;
- ▶ is more accurate;
- ▶ is more reliable;
- ▶ ...?

1: A note on complexity

Focusing on efficiency, we now need a way of discussing how the time taken by an algorithm depends on the problem size.

The usual way to discuss this is in terms of “Big \mathcal{O} ” notation, which classifies how, e.g., algorithms’ run-times grow as the input size grows.

For example, if we say an algorithm for a problem of size n has complexity $\mathcal{O}(n^2)$, then we mean there is some constant, C such that the run-time is at most Cn^2 .

Often, we don’t really care too much about what C is. For example, if Algorithm 1 had complexity $0.1n^2$, and Algorithm 2 had complexity $100n$, then...

1: A note on complexity

The best to worst, some common complexities are

- ▶ $\mathcal{O}(1)$
- ▶ $\mathcal{O}(\log n)$
- ▶ $\mathcal{O}(n)$
- ▶ $\mathcal{O}(n \log n)$
- ▶ $\mathcal{O}(n^2)$
- ▶ $\mathcal{O}(n^3)$
- ▶ $\mathcal{O}(2^n)$
- ▶ $\mathcal{O}(n!)$

2: Merge Sort

The **Bubble Sort algorithm** too slow for project we will undertake: its worse-case complexity is $\mathcal{O}(N^2)$ for a list of length N .

Instead we'll implement the **Merge Sort** algorithm. It has complexity $\mathcal{O}(N \log N)$.

Merge Sort

- ▶ Split the list into two smaller lists,
- ▶ Split each of those into 2 smaller lists.
- ▶ Keep doing this until each list is of length 1.
- ▶ A list of length 1 is already sorted, so...
- ▶ Reassemble each of your sub-lists by merging these sorted list.

2: Merge Sort

It is useful to write this as a **recursive algorithm**:

Recursive Merge Sort Algorithm

```
procedure mergesort( $L = a_1, a_2, \dots, a_n$ )  
  if  $n > 1$  then  
     $m := \text{floor}(n/2)$   
     $L_1 := (a_1, a_2, \dots, a_m)$   
     $L_2 := (a_{m+1}, a_{m+2}, \dots, a_n)$   
     $L := \text{merge}(\text{mergesort}(L_1), \text{mergesort}(L_2)).$   
  end if
```

So we need two functions:

- (i) A `Merge()` function to merge two sorted list
- (ii) A `MergeSort()` function that
 - ▶ splits the list in two,
 - ▶ calls `MergeSort()` for each half
 - ▶ calls the `Merge()` function

2: Merge Sort

Example (Merge Sort)

Show how **Merge Sort** would sort the list

9 5 1 2 6 3 4 9 4

3: Comparing in practice

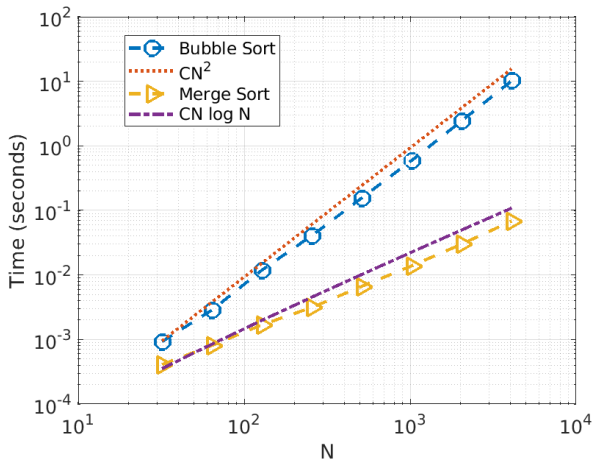
In Lab 3 you should find that...

- ▶ **Bubble Sort** has a worst-case complexity of $\mathcal{O}(N^2)$ for a list of length N .
- ▶ **Merge Sort** has a worst-case complexity of $\mathcal{O}(N \log N)$ for a list of length N .

This means that if we have a list of length N , then the expected time taken for the methods are $C_B N^2$ and $C_M N \log N$, for some constants C_B and C_M .

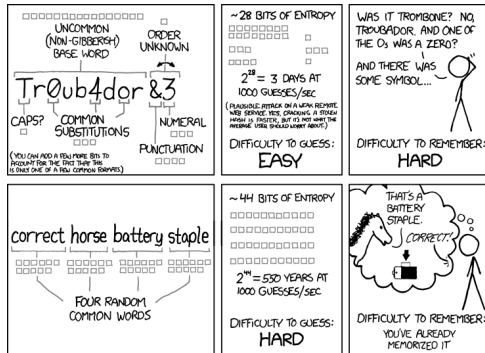
How to estimate these?

3: Comparing in practice



The data for this figure was collected using the `tic` and `toc` functions, which we saw before in Week 4.

4: The Password Problem



THROUGH 20 YEARS OF EFFORT, WE'VE SUCCESSFULLY TRAINED EVERYONE TO USE PASSWORDS THAT ARE HARD FOR HUMANS TO REMEMBER, BUT EASY FOR COMPUTERS TO GUESS.

<https://xkcd.com/936/>

4: The Password Problem

In Lab 3, we considered the problem of sorting a long list of “stolen” passwords. Actually, our goal is the to determine the most common.

The source of the data is the infamous **RockYou** password file, a list of over 30,000,000 unencrypted passwords [stolen from RockYou in 2009](#), and now widely available online.

The file contains one password per line, in no particular order. The first few are

```
password
mekster11
mekster11
progr4sm
khas8950
emilio1
holiday2
caitlin1
```

Given a list of 30,000,000 passwords, how shall we work out which 10 (say) occur most frequently?

Idea:

1. Load the list of passwords from a file.
2. Sort the list alphabetically, using MATLAB's `sort` function.
3. Create a list of words that contains no repetitions, using `unique`.
4. `[U, ai] = unique(P)` also returns a vector, `ai`, such that $U = P(ai)$. Since P is sorted, this can tell us the frequency. Example:

5. Sort the word frequency list, in descending order, storing the “key”.
6. Use the key to output the top 10.

The first step is to load array `Passwords`, from the file, count the number of entries, and sort it.

`PasswordFrequency.m`

```
6 %% Load the Passwords array
  load UserAccount-1e6;
8 NumberOfPwDs = length>Passwords);
  fprintf("Array Passwords has %d entries\n", NumberOfPwDs);
10 %% Sort the passwords
  Passwords = sort>Passwords);
```

Make a list of the unique words, and their frequency.

PasswordFrequency.m

```
16 %% Find the most frequently occurring word.  
%   - create a new list of unique words  
18 %   - a corresponding count of the number on instances.  
[UniqueWords, ai] = unique>Passwords);  
20 WordFreq      = diff([ai; NumberOfPWDs+1]);
```

Sort the frequency count, in descending order, and use the key to to order the *UniqueWords* array.

PasswordFrequency.m

```
22 %% Sort by Frequency
% Again use the "sort" function, but keep the "key"
24 [WordFreq,key]=sort(WordFreq,'descend');
UniqueWords = UniqueWords(key);
```

Output the top 10:

PasswordFrequency.m

```
%% Output top 10
28 fprintf("The 10 most common words (and freqs) are:\n");
   for i=1:10
30     fprintf("\t%10s (%3d)\n", UniqueWords(i), WordFreq(i));
   end
```

```
1 The 10 most common words (and freqs) are:
   password (215)
3   iloveyou (175)
   123456 (126)
5   password1 (117)
   abc123 (114)
7   iloveyou1 ( 93)
   princess ( 93)
9   love ( 77)
   princess1 ( 73)
```

5: Structures

So far, for all the arrays we have studied, we access specific elements using index/number. E.g, `x=[3 1 4 1 5 9]`, and access the 3rd element as `x(3)`.

A *structure* is a type of array where the entries have names. But more than that is true:

- ▶ Elements of the structure can be of different types;
- ▶ Elements can scalars, arrays, or even other structures.

5: Structures

In this simple example, we'll create a structure for a module.

ModuleStructure.m

```
1 %% Using a simple struture:
2 Module.code='2223-CS319';
3 Module.name='Scientific Computing';
4 Module.Students=[20123456, 19876543, 21212121];
5 Module.Graded=["A", "C", "B"];
6 disp(Module)
```

The variable *Module* is now of type **struct**.

```
>> Module
2 Module =
    struct with fields:
4     code: 'CS319'
     name: 'Scientific Computing'
6     Students: [20123456 19876543 21212121]
     Graded: ["A"      "C"      "B"]
```

5: Structures

We can access or set a struct's entries using the DOT operator:

```
1 >> Module.code='2223-CS319'
Module =
3   struct with fields:
        code: '2122-CS319'
        name: 'Scientific Computing'
        Students: [20123456 19876543 21212121]
7        Graded: ["A"      "C"      "B"]
```

There is lots more one can do with structs, such as creating arrays of structures, or stuctures with structures... but the main point today is as an introduction to user-defined **composite data types**.

However, it is also worth noting that some MATLAB functions return **stucts**, including piecewise interpolation functions.

One can use the `pchip()` function to compute the piecewise cubic Hermite interpolant to a data set. E.g.,

```
1  x=[0 .1 .5 1]
   y=[1, 0, 0.2, .3]
3  Y = pchip(x,y, X); % Some big vector X
```

The $Y(i)$ is the PCHIP interpolant to (x,y) evaluated at $X(i)$. But we can also compute the interpolant itself.

```
1  >> p = pchip(x,y)
   p =
3  struct with fields:
       form: 'pp'
5  breaks: [0 0.1000 0.5000 1]
       coefs: [3x4 double]
7  pieces: 3
       order: 4
9  dim: 1
```