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Solving linear systems of equations

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Annotated §3.6 Condition Numbers

Annotated

Slides

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MA385 – Numerical Analysis 1

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Numerical solutions to some linear systems are adversely affected by round-off errors.

This phenomenon is related the matrices in the linear systems. Those matrices for which the issue is particularly prevalent are referred to as being *ill-conditioned*.

For any matrix, we can assign a numerical score that gives an indication of whether it is ill-conditioned. That score is called the *condition number*, and is the subject of these section.

The condition number is defined in terms of matrix norms.

Suppose we have a vector norm, $\|\cdot\|$ and associated subordinate matrix norm. It is not hard to see that

$$\|Au\| \leq \|A\| \|u\| \quad \text{for any } u \in \mathbb{R}^n, A \in \mathbb{R}^{n \times n}.$$

Here is why:

Recall that $\|A\| := \max_{v \in \mathbb{R}^n / \{0\}} \frac{\|Av\|}{\|v\|}.$

So, for an arbitrary vector u ,

$$\|A\| \geq \frac{\|Au\|}{\|u\|}.$$

So $\|Au\| \leq \|A\| \cdot \|u\|.$

There is an analogous statement for the product of two matrices:

Definition 3.26 (Consistent matrix norm)

A matrix norm $\| \cdot \|$ is **consistent** (or “*sub-multiplicative*”) if

$$\|AB\| \leq \|A\|\|B\|, \quad \text{for all } A, B \in \mathbb{R}^{n \times n}.$$

Theorem 3.27

Any subordinate matrix norm is consistent.

The proof is left to Exercise 3.17. That exercises also demonstrates that there are matrix norms which are *not* consistent.

[Please read this slide in your own time!]

Modern computers don't store numbers in decimal (base 10), but in binary (base 2) “floating point numbers” of the form :

$$x = \pm a \times 2^{b-M}.$$

Most use *double precision*, where 8 bytes (64 bits or *binary digits*) are used to store

- the sign (1 bit),
- a , called the “significand” or “mantissa” (52 bits)
- and the exponent, $b - 1023$ (11 bits)

Note that a has roughly 16 decimal digits.

(Some older computer systems sometimes use *single precision* where a has 23 bits — giving 8 decimal digits — and b has 7; so too do many new GPU-based systems).

[OK, you can start reading again!] When we try to store a real number x on a computer, we actually store the nearest floating-point number. That is, we end up storing $x + \delta x$, where δx is the “round-off” error.

But the quantity we are mainly interested in is the **relative error**: $|\delta x|/|x|$.

Since this is not a course on computer architecture, we'll simplify a little and just take it that single and double precision systems lead to a relative error of 10^{-8} and 10^{-16} respectively.

(Sew p68–70 of Süli and Mayers for a thorough development of the concept of a condition number).

Suppose we use, say, LU -factorization and back-substitution on a computer to solve

$$Ax = b.$$

Because of the “round-off error” we actually solve

$$A(x + \delta x) = (b + \delta b).$$

Our problem now is, for a given A , if we know the (relative) error in b , can we find an upper-bound on the relative error in x ?

Definition 3.28

The *condition number* of a matrix, with respect to a particular matrix norm $\|\cdot\|_*$ is

$$\kappa_*(A) = \|A\|_* \|A^{-1}\|_*.$$

If $\kappa_*(A) \gg 1$ then we say A is *ill-conditioned*.

$$\left\{ \begin{array}{l} \kappa_*(I) = 1. \\ \kappa_*(A) \geq 1 \text{ for} \\ \text{all other } A. \end{array} \right.$$

Example: Find the condition number κ_∞ of

$$A = \begin{pmatrix} 10 & 12 \\ 0.08 & 0.1 \end{pmatrix}.$$

$$A^{-1} = \begin{pmatrix} 2.5 & 300 \\ -2 & 250 \end{pmatrix}.$$

$$\text{So } \|A\|_\infty = 22. \quad \|A^{-1}\|_\infty = 302.5$$

$$\text{So } \kappa_\infty(A) = 6,655.$$

Theorem 3.29

Suppose that $A \in \mathbb{R}^{n \times n}$ is nonsingular and that $\mathbf{b}, \mathbf{x} \in \mathbb{R}^n$ are non-zero vectors. If $A\mathbf{x} = \mathbf{b}$ and $A(\mathbf{x} + \delta\mathbf{x}) = (\mathbf{b} + \delta\mathbf{b})$ then

$$\frac{\|\delta\mathbf{x}\|}{\|\mathbf{x}\|} \leq \kappa(A) \frac{\|\delta\mathbf{b}\|}{\|\mathbf{b}\|}.$$

Proof: Since $A\mathbf{x} = \mathbf{b}$ and $A(\mathbf{x} + \delta\mathbf{x}) = \mathbf{b} + \delta\mathbf{b}$,

$$\text{so } A\delta\mathbf{x} = \delta\mathbf{b}.$$

Then $\mathbf{b} = A\mathbf{x}$, so $\|\mathbf{b}\| = \|A\mathbf{x}\| \leq \|A\| \cdot \|\mathbf{x}\|$.

Similarly $\delta\mathbf{x} = A^{-1}\delta\mathbf{b}$, so $\|\delta\mathbf{x}\| \leq \|A^{-1}\| \cdot \|\delta\mathbf{b}\|$.

$$\text{So } \|\mathbf{b}\| \cdot \|\delta\mathbf{x}\| \leq \underbrace{\|A\| \cdot \|A^{-1}\|}_{\kappa(A)} \|\mathbf{x}\| \|\delta\mathbf{b}\|$$

This gives $\|\delta\mathbf{x}\|/\|\mathbf{x}\| \leq \kappa(A) \|\delta\mathbf{b}\|/\|\mathbf{b}\|$.

Example 3.30

Suppose we are using a computer to solve $Ax = b$ where

$$A = \begin{pmatrix} 10 & 12 \\ 0.08 & 0.1 \end{pmatrix} \quad \text{and} \quad b = \begin{pmatrix} 1 \\ 1 \end{pmatrix}.$$

But, due to round-off error, right-hand side has a relative error (in the ∞ -norm) of 10^{-6} . Give a bound for the relative error in x in the ∞ -norm.

We already saw that $\kappa_\infty(A) = 6655$.

So, since $\frac{\| \delta b \|}{\| b \|} = 10^{-6}$,

$$\frac{\| \delta x \|}{\| x \|} \leq \kappa_\infty(A) \times 10^{-6} = 0.006655.$$

For every matrix norm we get a different condition number.

Example 3.31

Let A be the $n \times n$ matrix

$$A = \begin{pmatrix} 1 & 0 & 0 & \dots & 0 \\ 1 & 1 & 0 & \dots & 0 \\ 1 & 0 & 1 & \dots & 0 \\ \vdots & & & & \\ 1 & 0 & 0 & \dots & 1 \end{pmatrix}.$$

What are $\kappa_1(A)$, and $\kappa_\infty(A)$?

First we compute $\|A\|_1$ and $\|A\|_\infty$.

$$\|A\|_1 = 2, \quad \|A\|_\infty = n.$$

For this very special example, it is easy to write down the inverse of A :

$$A^{-1} = \begin{pmatrix} 1 & 0 & 0 & \dots & 0 \\ -1 & 1 & 0 & \dots & 0 \\ -1 & 0 & 1 & \dots & 0 \\ \vdots & & & & \\ -1 & 0 & 0 & \dots & 1 \end{pmatrix}.$$

$$\|A^{-1}\|_1 = 2 \qquad \|A^{-1}\|_\infty = n.$$

$$\text{So } \kappa_1(A) = 4 \qquad \kappa_\infty(A) = n^2.$$

To compute $\kappa_1(A)$ and $\kappa_\infty(A)$, we need to know A^{-1} , which is usually not practical. However, for κ_2 , we are able to *estimate* the condition number of A without knowing A^{-1} .

Recall that $\|A\|_2 = \sqrt{\lambda_n}$ where λ_n is the largest eigenvalue of $B = A^T A$.

We can also show that $\|A^{-1}\|_2 = \frac{1}{\sqrt{\lambda_1}}$ where λ_1 is the smallest eigenvalue of B (see Section 3.6.5 of notes). So

$$\kappa_2(A) = \left(\lambda_n / \lambda_1 \right)^{1/2}.$$

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Motivated by this, we'll finish MA385, by studying an easy way of estimating the eigenvalues of a matrix.

Exercise 3.17

- (i) Prove that, if $\|\cdot\|$ is a subordinate matrix norm, then it is *consistent*, i.e., for any pair of $n \times n$ matrices, A and B , we have $\|AB\| \leq \|A\|\|B\|$.
- (ii) One might think it intuitive to define the “max” norm of a matrix as follows:

$$\|A\|_{\infty}^{\sim} = \max_{i,j} |a_{ij}|.$$

Show that this is indeed a norm on $\mathbb{R}^{n \times n}$. Show that, however, it is not consistent.

Exercise 3.18

Let A be the matrix

$$A = \begin{pmatrix} 0.1 & 0 & 0 \\ 10 & 0.1 & 10 \\ 0 & 0 & 0.1 \end{pmatrix}$$

Compute $\kappa_{\infty}(A)$. Suppose we wish to solve the system of equations $Ax = b$ on *single precision* computer system (i.e., the relative error in any stored number is approximately 10^{-8}). Give an upper bound on the relative error in the computed solution x .

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