

Optimal Transport: Theory and Applications

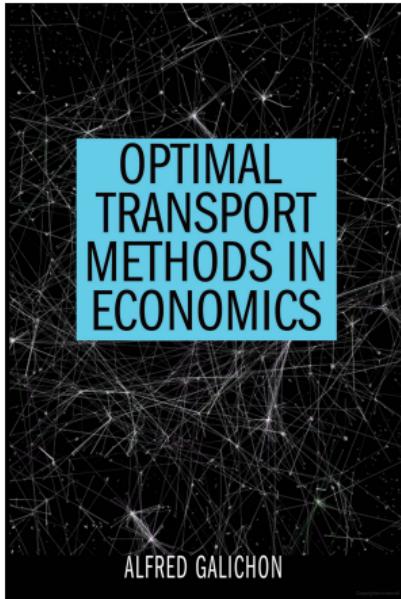
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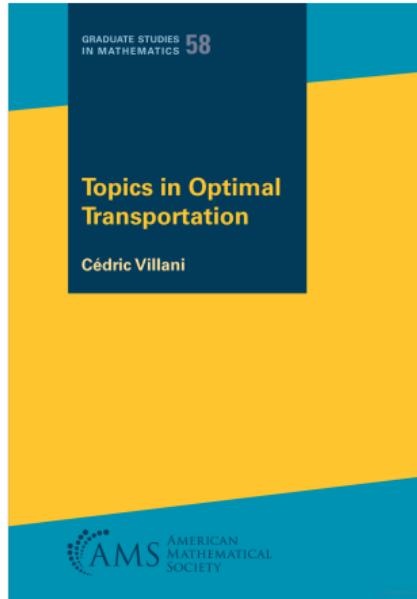


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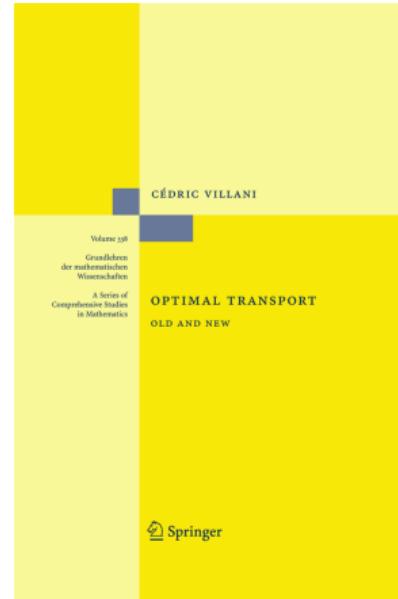
Books



(a)



(b)



(c)

Road map

1 Optimal Transport Preliminaries

- Definition
- Discrete Distributions
- Continuous Distributions

2 Optimization of Optimal Transport

3 Statistics of Optimal Transport

- Curse of Dimensionality
- Projection
- Smoothness

4 Applications

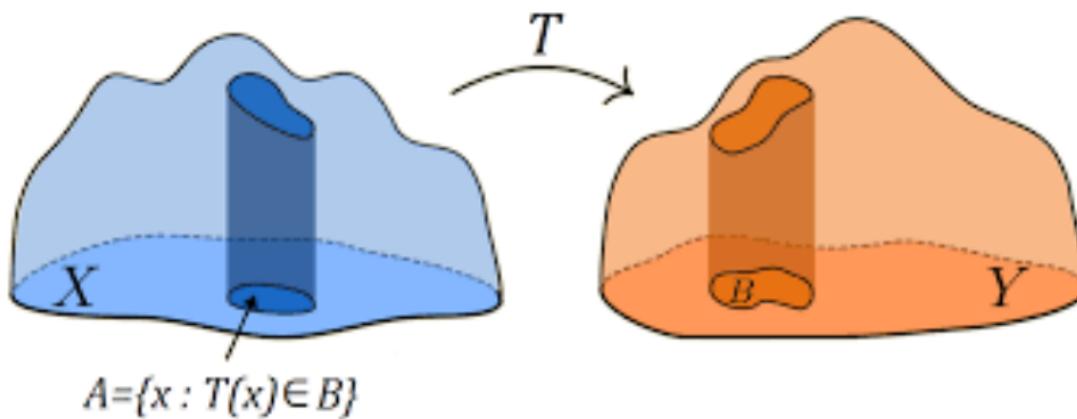
- Wasserstein GANs
- Distributionally Robust Optimization

5 Some New Advances and Open Problems

Monge Map

- Let $P \in \mathcal{P}(S)$ and $Q \in \mathcal{P}(S)$ be two probability distributions defined on a space S ;
 $c : S \times S \rightarrow [0, \infty]$ is a cost function.
- Monge problem:

$$\inf_{T(\cdot)} \mathbb{E}_P[c(X, T(X)) | T_{\sharp} P = Q].$$

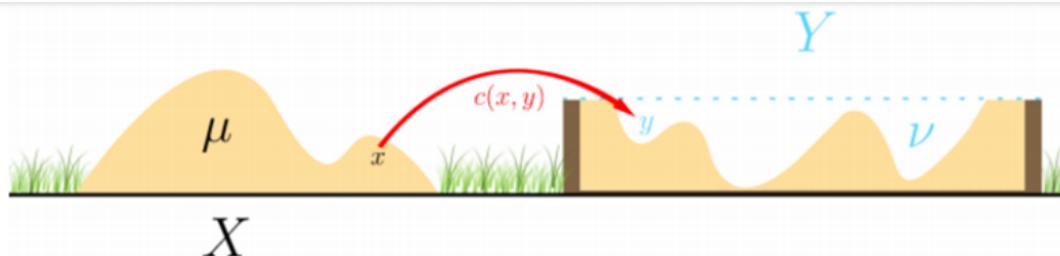


Definition

Definition (Optimal Transport Cost)

Let $P \in \mathcal{P}(S)$ and $Q \in \mathcal{P}(S)$ be two probability distributions defined on a space S ; $c : S \times S \rightarrow [0, \infty]$ is a cost function. Then, the optimal transport cost is defined as

$$D_c(P, Q) := \inf_{\pi} \left\{ \mathbb{E}_{\pi}[c(U, V)] \mid \pi \in \mathcal{P}(S \times S), \right.$$
$$\left. \pi(A \times S) = P(A), \pi(S \times B) = Q(B) \text{ for every subsets } A, B \text{ of } S \right\}$$



Wasserstein Distance, Earth Moving Distance

Let $S = \mathbb{R}^d$ and $c(x, y) = d(x, y)^\rho$ for some metric function $d(\cdot, \cdot)$ and $\rho \geq 1$,

$$W_\rho(P, Q) = D_c(P, Q)^{1/\rho}$$

is a metric on the probability space. We call it the type- ρ Wasserstein distance. In particular, if $\rho = 1$, it is also called the earth moving distance.

Discrete Distributions: Duality

- Let P support $\{x_1, x_2, \dots, x_N\}$ and Q support $\{y_1, y_2, \dots, y_m\}$. Let $P = \{p_x\}$ and $Q = \{q_y\}$ with $\sum_{x=1}^N p_x = \sum_{y=1}^M q_y = 1, p_x \geq 0, q_y \geq 0$.
- Optimal transport cost is the optimal value of the linear programming:

$$D_c(P, Q) = \min_{\pi_{xy} \geq 0} \sum_{x=1, y=1}^{N, M} c(x, y) \pi_{xy} \quad (1)$$

$$\text{s.t. } \sum_{y=1}^N \pi_{xy} = p_x \text{ and } \sum_{x=1}^N \pi_{xy} = q_y. \quad (2)$$

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- Duality:

$$D_c(P, Q) = \max_{u, v} \sum_{x=1}^N p_x u_x + \sum_{y=1}^M q_y v_y \quad (3)$$

$$\text{s.t. } u_x + v_y \leq c(x, y) \quad (4)$$

Discrete Distributions: Optimal Solutions

- Primal and dual:

$$\min_{\pi_{xy} \geq 0} \sum_{x=1, y=1}^{N, M} c(x, y) \pi_{xy} \text{ s.t. } \sum_{y=1}^N \pi_{xy} = p_x \text{ and } \sum_{x=1}^N \pi_{xy} = q_y. \quad (\text{P})$$

$$\max_{u, v} \sum_{x=1}^N p_x u_x + \sum_{y=1}^M q_y v_y \text{ s.t. } u_x + v_y \leq c(x, y) \quad (\text{D})$$

- The optimal solution satisfies

$$\pi_{xy}^* > 0 \Rightarrow u_x^* + v_y^* = c(x, y)$$

$$u_x^* = \min_{y \in \{1, 2, \dots, M\}} c(x, y) - v_y^* \text{ and } v_y^* = \min_{x \in \{1, 2, \dots, N\}} c(x, y) - u_x^*.$$

An Economic Interpretation

- Primal and dual:

$$\min_{\pi_{xy} \geq 0} \sum_{x=1, y=1}^{N, M} c(x, y) \pi_{xy} \text{ s.t. } \sum_{y=1}^N \pi_{xy} = p_x \text{ and } \sum_{x=1}^N \pi_{xy} = q_y. \quad (\text{P})$$

$$\max_{u, v} \sum_{x=1}^N p_x u_x + \sum_{y=1}^M q_y v_y \text{ s.t. } u_x + v_y \leq c(x, y) \quad (\text{D})$$

- Transfer coal from mines in $\{x_1, x_2, \dots, x_N\}$ to factories in $\{y_1, y_2, \dots, y_m\}$:
 - Transportation cost is $c(x, y)$;
 - u_x, v_y are shadow prices: u_x is the price of loading one ton of coal at place x ; and v_y is the price of unloading it at destination y .

Pure Assignments and Monge Map

- Consider $N = M$ and $p_x = q_y = 1/N$.
- Then, the optimal solution is a permutation σ : an invertible map from $\{1, 2, \dots, N\}$ onto itself.

$$\pi_{xy}^* = \mathbb{I}\{y = \sigma(x)\}.$$

- The optimal transport problem is equivalent to the Monge problem.

Continuous Distributions: Duality

- Recall

$$D_c(P, Q) = \inf_{\pi} \left\{ \mathbb{E}_{\pi}[c(U, V)] \mid \pi \in \mathcal{P}(S \times S), \right.$$

$\pi(A \times S) = P(A), \pi(S \times B) = Q(B)$ for every subsets A, B of $S\}$

- Duality:

$$\begin{aligned} D_c(P, Q) &= \sup_{\varphi, \psi} \int \varphi dP + \int \psi dQ \\ &\quad s.t. \varphi(x) + \psi(y) \leq c(x, y). \end{aligned}$$

2-Wasserstein Distance Between Gaussian Distributions

- Cost function $c(x, y) = \|x - y\|_2^2$, $P = \mathcal{N}(\mu_1, \Sigma_1)$ and $Q = \mathcal{N}(\mu_2, \Sigma_2)$.
- Then, the 2-Wasserstein distance between P and Q is

$$W_2^2(P, Q) = \|\mu_1 - \mu_2\|_2^2 + \text{tr}(\Sigma_1) + \text{tr}(\Sigma_2) - 2\text{tr}\left[(\Sigma_1^{1/2}\Sigma_2\Sigma_1^{1/2})^{1/2}\right].$$

1-Wasserstein Distance: Duality

- $c(x, y) = d(x, y)$;
- Duality:

$$W_1(P, Q) = D_c(P, Q) = \sup_{\varphi} \int \varphi dP - \int \varphi dQ$$

s.t. $\varphi(x)$ is 1-Lipschitz with respect to $d(\cdot, \cdot)$

Total Variation Distance

- $c(x, y) = \mathbb{I}(x \neq y);$
- $D_c(P, Q) = TV(P, Q).$

One-Dimensional Case

- $d = 1$, we have

$$W_\rho(P, Q) = \left(\int_0^1 |F_P^{-1}(s) - F_Q^{-1}(s)|^\rho ds \right)^{1/\rho},$$

where F_P and F_Q are CDFs of measures P and Q .

- if $\rho = 1$ and $d = 1$, we have

$$W_1(P, Q) = \int_{\mathbb{R}} |F_P(s) - F_Q(s)| ds,$$

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Optimization of Optimal Transport

- Discrete case, linear programming: for simplicity, assume $N = M$

$$\min_{\pi_{xy} \geq 0} \sum_{x=1, y=1}^N c(x, y) \pi_{xy} \text{ s.t. } \sum_{y=1}^N \pi_{xy} = p_x \text{ and } \sum_{x=1}^N \pi_{xy} = q_y. \quad (\text{P})$$

- Linear programming sample complexity $O(N^{3.5} \log(1/\epsilon))$.
- Sinkhorn method [Cuturi, 2013] with sample complexity $\tilde{O}(N^2/\epsilon^2)$ [Dvurechensky et al., 2018].

Sinkhorn Method

- We optimize the following program:

$$\min_{\pi_{xy} \geq 0} \sum_{x=1, y=1}^N c(x, y) \pi_{xy} + \frac{1}{\lambda} \sum_{i, j=1}^N \pi_{ij} \log(\pi_{ij}) \text{ s.t. } \sum_{y=1}^N \pi_{xy} = p_x \text{ and } \sum_{x=1}^N \pi_{xy} = q_y.$$

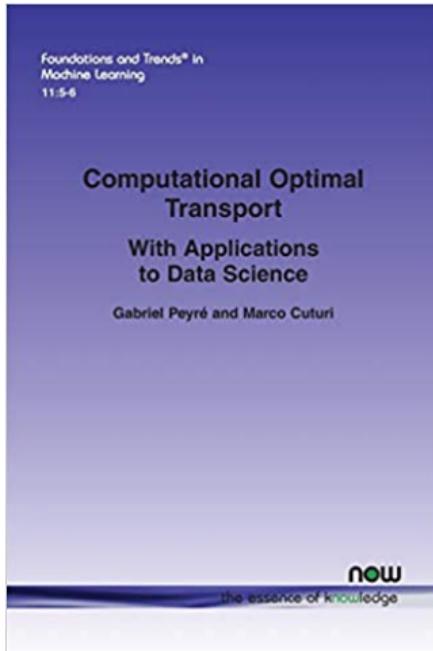
- The solution admits the form:

$$\pi_{ij}^\lambda = u_i \exp(-\lambda c(i, j)) v_j.$$

- By Sinkhorn and Knopp's algorithm [Sinkhorn and Knopp, 1967], we can use a fixed-point iteration to arrive

$$\pi^\lambda \mathbf{1} = P, (\pi^\lambda)^\top \mathbf{1} = Q.$$

Computational Optimal Transport



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Curse of Dimensionality

- Let P^* be a measure on \mathbb{R}^d and let P_n be the associated empirical measure, i.e., for i.i.d. sample X_1, X_2, \dots, X_n ,

$$P_n = \frac{1}{n} \sum_{i=1}^n \delta_{X_i}.$$

- Consistency: $W_\rho(P_n, P^*) \rightarrow 0$.
- Curse of Dimensionality: $\mathbb{E}[W_\rho(P_n, P^*)] = O(n^{-1/d})$ [Fournier and Guillin, 2015].
- If P^* supports on an m -dimensional manifold of \mathbb{R}^d , we have $\mathbb{E}[W_\rho(P_n, P^*)] = O(n^{-1/m})$ [Weed and Bach, 2019].

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- CLT [Del Barrio and Loubes, 2019]:

$$\sqrt{n}(W_2^2(P_n, P^*) - \mathbb{E}[W_2^2(P_n, P^*)]) \Rightarrow \mathcal{N}(0, \sigma^2).$$

Beating Curse of Dimensionality: Projection

- Sliced Wasserstein distance [Bonneel et al., 2015]:

$$SW_{\rho}^{\rho}(P, Q) = \int_{\mathbb{S}^{d-1}} W_{\rho}^{\rho}(\theta_{\sharp} P, \theta_{\sharp} Q) d\theta,$$

where $\theta_{\sharp} P$ is the push-forward measure:

$$\theta_{\sharp} P(A) = P(\{x : \theta^{\top} x \in A\}), \text{ for any Borel set } A \in \mathbb{R}.$$

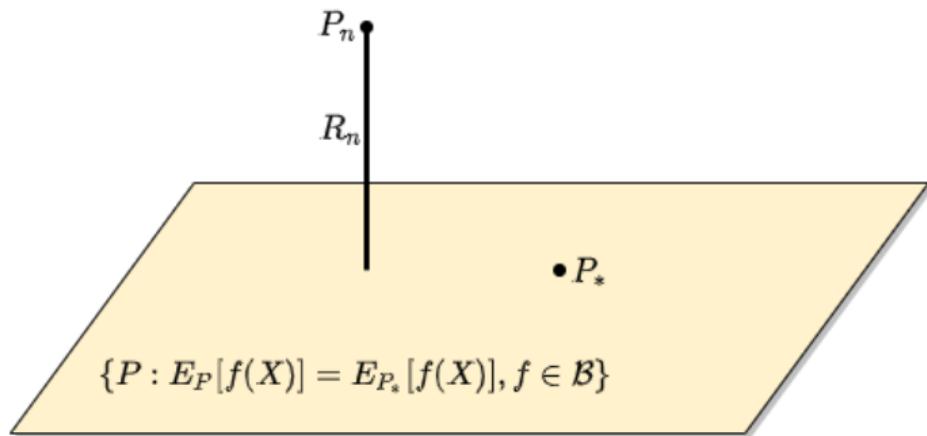
- Convergence rate: $SW_{\rho}(P_n, P^*) = O(\sqrt{n})$ [Nadjahi et al., 2019].
- Another variance: max-sliced Wasserstein distance [Deshpande et al., 2019]

$$MSW_{\rho}^{\rho}(P, Q) = \max_{\theta \in \mathbb{S}^{d-1}} W_{\rho}^{\rho}(\theta_{\sharp} P, \theta_{\sharp} Q),$$

Beating Curse of Dimensionality: Projection

- Robust Wasserstein profile function [Si et al., 2020]: consider a function class \mathcal{B}

$$R_n(P_*, P_n) := \inf_P \{D_c(P, P_n) : \mathbb{E}_P[f(X)] = \mathbb{E}_{P_*}[f(X)] \text{ for all } f \in \mathcal{B}\}.$$



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- Duality:

$$R_n(P_*, P_n) = \sup_{f \in \mathcal{LB}} \{\mathbb{E}_{P_*}[f(X)] - \mathbb{E}_{P_n}[f^c(X)]\},$$

where $f^c(x) = \sup_z \{f(z) - c(z, x)\}$ and \mathcal{LB} is a linear space spanned by the function class \mathcal{B} :

$$\mathcal{LB} = \left\{ f(\cdot) = \sum_{i=1}^m \lambda_i f_i(\cdot) : \{f_i(\cdot)\}_{i=1}^m \subset \mathcal{B}, \lambda \in \mathbb{R}^m, \text{ and } m \in \mathbb{Z}_+ \right\}.$$

- $R_n = O(n^{-1/2})$ under some assumptions of \mathcal{B} .

Beating Curse of Dimensionality: Smoothness

- If P^* is sufficient smooth: the density of P^* is in the Besov space $B_{p,q}^s$, then we can construct a wavelet estimator based on data such that $\mathbb{E} W_\rho(\hat{P}_n^w, P^*) = O\left(n^{-\frac{1+s}{d+2s}}\right)$ [Weed and Berthet, 2019].

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- σ -smooth Wasserstein distance:

$$W_\rho^{(\sigma)}(P, Q) = W_\rho(P * \mathcal{N}_\sigma, Q * \mathcal{N}_\sigma),$$

where $P * \mathcal{N}_\sigma(A) = \int_{-\infty}^{\infty} P(A - t)\phi_\sigma(t)dt$ and $\phi_\sigma(t)$ is the PDF of the Gaussian distribution \mathcal{N}_σ .

- $\mathbb{E}[W_\rho^{(\sigma)}(P_n, P)] = O(n^{-1/2})$.

More Properties of Smooth Wasserstein Distance

- $W_\rho^{(\sigma)}$ is continuous and monotonically non-increasing in $\sigma \in [0, +\infty)$;
- $\lim_{\sigma \rightarrow 0} W_\rho^{(\sigma)}(P, Q) = W_\rho(P, Q)$;
- $\lim_{\sigma \rightarrow +\infty} W_\rho^{(\sigma)}(P, Q) = |\mathbb{E}[X] - \mathbb{E}[Y]|$, for $X \sim P$ and $Y \sim Q$ sub-Gaussian.
- The constants in $\mathbb{E}[W_\rho^{(\sigma)}(P_n, P)]$ exhibit an exponential dependence on dimension d .

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Wasserstein Generative Adversarial Networks (GANs) [Arjovsky et al., 2017]

Stanford
University

- Goal: learn a generative model $g_\theta(\cdot)$ from data X_1, X_2, \dots, X_n sampled from a real data distribution P_r . We let P_θ be the distribution induced by the generative model $g_\theta(\cdot)$.

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- Minimize Wasserstein distance:

$$\min_{\theta} W_1(P_r, P_\theta) = \min_{\theta} \sup_{\|f\|_L \leq 1} \mathbb{E}_{P_r}[f(x)] - \mathbb{E}_{P_\theta}[f(x)]$$

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- Parametric $f(\cdot)$ to be a neural network:

$$\min_{\theta} W_1(P_r, P_\theta) = \min_{\theta} \max_w \left\{ \mathbb{E}_{P_r}[f_w(x)] - \mathbb{E}_{z \sim p(z)}[f_w(g_\theta(z))] \right\}$$

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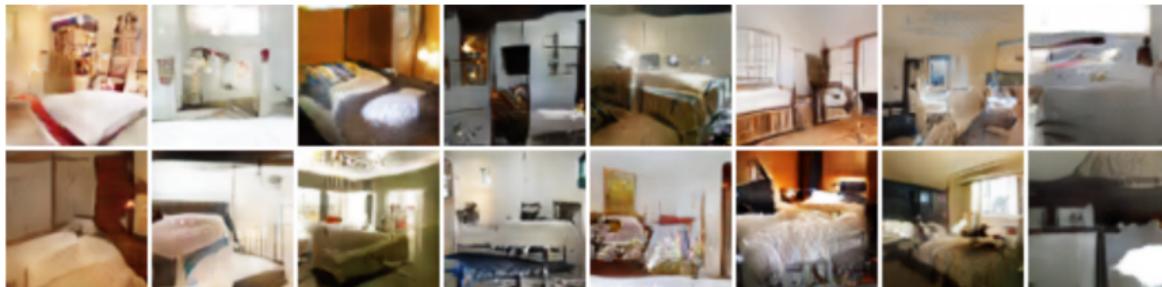
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- Adversarial.

Wasserstein GANs Results



(d)



(e)

Distributionally Robust Optimization Formulation

Distributionally Robust Optimization (DRO):

$$\inf_{\beta \in \mathbb{R}^d} \underbrace{\sup_{P \in \mathcal{U}} \mathbb{E}_P[\ell(X; \beta)]}_{\text{worst case expectation}},$$

\mathcal{U} : distributional uncertainty set.

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\mathcal{U} : distributional uncertainty set.

Construction of distributional uncertainty set \mathcal{U} :

$$\mathcal{U} = \mathcal{U}_\delta(P_n) = \{P \in \mathcal{P}(S) : D_c(P, P_n) \leq \delta\}$$

Why DRO?

- Statistical errors and overfitting;
- Distributional shifts.

Strong Duality for DRO

Theorem (Blanchet and Murthy, 2019; Gao and Kleywegt, 2016; Esfahani and Kuhn, 2018)

Suppose $c(\cdot)$ is a nonnegative lower semicontinuous function satisfying $c(x, y) = 0$ if and only if $x = y$ and $\ell(\cdot)$ is upper semicontinuous. Then,

$$\sup_{P: D_c(P, P_n) \leq \delta} \mathbb{E}_P [\ell(X; \beta)] = \inf_{\lambda \geq 0} f(\beta, \lambda),$$

where

$$f(\beta, \lambda) = \lambda \delta + \mathbb{E}_{P_n} [\ell_{rob}(X; \beta, \lambda)], \text{ and}$$

$$\ell_{rob}(X; \beta, \lambda) := \sup_{u \in \mathbb{R}^d} \{\ell(u; \beta) - \lambda c(u, X)\}.$$

Some DRO Estimators

- **Square-root LASSO** [Belloni, Chernozhukov and Wang 2011]:

$$\ell((x, y); \beta) = \|y - \beta^T x\|_2^2$$

$$P_n = \frac{1}{n} \sum_{i=1}^n \delta_{(X_i, Y_i)}(dx, dy)$$

$$c((x, y), (x', y')) = \|x - x'\|_q^2 + \infty \cdot \mathbf{1}\{y \neq y'\}$$

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DRO is equivalent to the square-root LASSO [Blanchet, Kang and Murthy, 2016; Gao, Chen and Kleywegt, 2017],
 $(1/p+1/q = 1)$

$$\sup_{P: D_c(P, P_n) \leq \delta} \mathbb{E}_P [\ell((X, Y); \beta)] = \left(\sqrt{\mathbb{E}_{P_n} [\ell((X, Y); \beta)]} + \sqrt{\delta} \|\beta\|_p \right)^2.$$

Some DRO Estimators

- Regularized logistic regression:

$$\ell((x, y); \beta) = \log(1 + \exp(-y\beta^T x))$$

$$P_n = \frac{1}{n} \sum_{i=1}^n \delta_{(X_i, Y_i)}(dx, dy)$$

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DRO is equivalent to the regularized logistic regression [Blanchet, Kang and Murthy, 2016; Gao, Chen and Kleywegt, 2017; Esfahani and Kuhn, 2015],

$$\sup_{P: D_c(P, P_n) \leq \delta} \mathbb{E}_P [\ell((X, Y); \beta)] = \mathbb{E}_{P_n} [\ell((X, Y); \beta)] + \delta \|\beta\|_p.$$

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Martingale Optimal Transport

$$MD_c(P, Q) := \inf_{\pi} \left\{ \mathbb{E}_{\pi}[c(X, Y)] \mid \pi \in \mathcal{P}(S \times S), \mathbb{E}_{\pi}[Y|X] = X, \right.$$
$$\left. \pi(A \times S) = P(A), \pi(S \times B) = Q(B) \text{ for every subsets } A, B \text{ of } S \right\}$$

Adapted Optimal Transport

Consider two-period case P is the joint distribution of (X_1, X_2) and Q is the joint distribution of (Y_1, Y_2) ,

$$\begin{aligned} AD_c(P, Q) := \inf_{\pi^1} & \left\{ \mathbb{E}_{\pi^1}[c(X_1, Y_1) + D_c(P_{X_1}, Q_{Y_1})] \mid \right. \\ & \left. \pi^1(A \times S) = P^1(A), \pi^1(S \times B) = Q^1(B) \text{ for every subsets } A, B \right\}, \end{aligned}$$

where P^1, Q^1 are the distributions of X_1 and X_2 and P_{X_1}, Q_{Y_1} are the distributions of Y_1 and Y_2 conditional on X_1 and X_2 .

Minimum Wasserstein Distance

- Recall the CLT:

$$\sqrt{n}(W_\rho^\rho(P_n, P^*) - \mathbb{E}[W_\rho^\rho(P_n, P^*)]) \Rightarrow \mathcal{N}(0, \sigma^2).$$

- What can we say about $\hat{\theta}_n$:

$$\hat{\theta}_n = \arg \min_{\theta} D_c(P_n, P_{\theta})$$

Curse of Dimensionality?

Q&A?

Stanford
University

QUESTIONS

Q & A

ANSWERS

References I

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