

# GUARDED RECURSIVE TYPE THEORY VIA SIZED TYPES

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## 1. INTRODUCTION

[?]

## 2. PRELIMINARIES

**postulate** funext :  $\forall \{\ell \ell'\} \rightarrow \text{Extensionality } \ell \ell'$

**uip** :  $\forall \{\ell\} \{A : \text{Set } \ell\} \rightarrow \{a \ a' : A\}$   
 $\rightarrow \{p \ p' : a \equiv a'\} \rightarrow p \equiv p'$   
**uip**  $\{p = \text{refl}\} \{\text{refl}\} = \text{refl}$

**data** tag : **Set** **where**  
**set** : tag  
**tot** : tag

$\Sigma \equiv : \{A : \text{Set}\} \{P : A \rightarrow \text{Set}\}$   
 $\rightarrow \{a \ a' : A\} \{p : P \ a\} \{p' : P \ a'\}$   
 $\rightarrow (e : a \equiv a') \rightarrow \text{subst } P \ e \ p \equiv p' \rightarrow (a \ , \ p) \equiv (a' \ , \ p')$   
 $\Sigma \equiv \text{refl refl} = \text{refl}$

**isProp** : **Set**  $\rightarrow$  **Set**  
**isProp**  $P = \{x \ y : P\} \rightarrow x \equiv y$

$\Sigma \equiv \text{-uip} : \{A : \text{Set}\} \{P : A \rightarrow \text{Set}\}$   
 $\rightarrow (\{a : A\} \rightarrow \text{isProp } (P \ a))$   
 $\rightarrow \{a \ a' : A\} \{p : P \ a\} \{p' : P \ a'\}$   
 $\rightarrow (e : a \equiv a') \rightarrow (a \ , \ p) \equiv (a' \ , \ p')$   
 $\Sigma \equiv \text{-uip } q \ \text{refl} = \text{cong } (, \ ) \ q$

Dependent functions preserve equality

**cong-dep** :  $\{A : \text{Set}\} \{P : A \rightarrow \text{Set}\}$   
 $\rightarrow (f : (a : A) \rightarrow P \ a)$   
 $\rightarrow \{x \ y : A\}$   
 $\rightarrow (e : x \equiv y) \rightarrow \text{subst } P \ e \ (f \ x) \equiv f \ y$   
**cong-dep**  $f \ \text{refl} = \text{refl}$

Functions with two arguments preserve equality

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cong2-dep : {A B : Set}{P : A → Set}
  → (f : (a : A) (p : P a) → B)
  → {a a' : A} {p : P a} {p' : P a'}
  → (e : a ≡ a') → subst P e p ≡ p' → f a p ≡ f a' p'
cong2-dep f refl refl = refl

```

Transport of a composition

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subst-trans : {A : Set}{P : A → Set}
  → {x y z : A}(e : x ≡ y)(e' : y ≡ z)
  → {p : P x}
  → subst P e' (subst P e p) ≡ subst P (trans e e') p
subst-trans refl refl = refl

```

```

record PSh : Set1 where
  field
    Obj : Size → Set
    Mor : (i : Size) (j : Size< (↑ i))
      → Obj i → Obj j
    MorId : {i : Size} {x : Obj i}
      → Mor i i x ≡ x
    MorComp : {i : Size} {j : Size< (↑ i)} {k : Size< (↑ j)}
      → {x : Obj i}
      → Mor i k x ≡ Mor j k (Mor i j x)

```

```

ConstObj : Size → Set
ConstObj = A

```

```

ConstMor : (i : Size) (j : Size< (↑ i))
  → ConstObj i → ConstObj j
ConstMor x = x

```

```

ConstMorId : {i : Size} {x : A}
  → ConstMor i i x ≡ x
ConstMorId = refl

```

```

ConstMorComp : {i : Size} {j : Size< (↑ i)} {k : Size< (↑ j)}
  → {x : ConstObj i}
  → ConstMor i k x ≡ ConstMor j k (ConstMor i j x)
ConstMorComp = refl

```

```

Const : PSh
Const = record
  { Obj = ConstObj
  ; Mor = ConstMor
  ; MorId = ConstMorId

```

```

; MorComp = ConstMorComp
}

```

```

Terminal : PSh
Terminal = Const  $\top$ 

```

```

ProdObj : Size  $\rightarrow$  Set
ProdObj  $i = \text{P.Obj } i \times \text{Q.Obj } i$ 

```

```

ProdMor : ( $i : \text{Size}$ ) ( $j : \text{Size} < (\uparrow i)$ )
 $\rightarrow \text{ProdObj } i \rightarrow \text{ProdObj } j$ 
ProdMor  $i j = \text{map } (\text{P.Mor } i j) (\text{Q.Mor } i j)$ 

```

```

ProdMorId : { $i : \text{Size}$ } { $x : \text{ProdObj } i$ }
 $\rightarrow \text{ProdMor } i i x \equiv x$ 
ProdMorId { $i$ } { $x$ } =
begin
  ( $\text{P.Mor } i i (\text{proj}_1 x), \text{Q.Mor } i i (\text{proj}_2 x)$ )
 $\equiv \langle \text{cong } (\lambda z \rightarrow (z, \text{Q.Mor } i i (\text{proj}_2 x))) \text{ P.MorId } \rangle$ 
  ( $\text{proj}_1 x, \text{Q.Mor } i i (\text{proj}_2 x)$ )
 $\equiv \langle \text{cong } (\lambda z \rightarrow (\text{proj}_1 x, z)) \text{ Q.MorId } \rangle$ 
   $x$ 
■

```

```

ProdMorComp : { $i : \text{Size}$ } { $j : \text{Size} < (\uparrow i)$ } { $k : \text{Size} < (\uparrow j)$ }
 $\rightarrow \{x : \text{ProdObj } i\}$ 
 $\rightarrow \text{ProdMor } i k x \equiv \text{ProdMor } j k (\text{ProdMor } i j x)$ 
ProdMorComp { $i$ } { $j$ } { $k$ } { $x$ } =
begin
  ( $\text{P.Mor } i k (\text{proj}_1 x), \text{Q.Mor } i k (\text{proj}_2 x)$ )
 $\equiv \langle \text{cong } (\lambda z \rightarrow (z, \text{Q.Mor } i k (\text{proj}_2 x))) \text{ P.MorComp } \rangle$ 
  ( $\text{P.Mor } j k (\text{P.Mor } i j (\text{proj}_1 x)), \text{Q.Mor } i k (\text{proj}_2 x)$ )
 $\equiv \langle \text{cong } (\lambda z \rightarrow (\text{P.Mor } j k (\text{P.Mor } i j (\text{proj}_1 x)), z)) \text{ Q.MorComp } \rangle$ 
  ( $\text{P.Mor } j k (\text{P.Mor } i j (\text{proj}_1 x)), \text{Q.Mor } j k (\text{Q.Mor } i j (\text{proj}_2 x))$ )
■

```

```

Prod : PSh
Prod = record
{ Obj = ProdObj
; Mor = ProdMor
; MorId = ProdMorId
; MorComp = ProdMorComp
}

```

```

SumObj : Size  $\rightarrow$  Set
SumObj  $i = \text{P.Obj } i \uplus \text{Q.Obj } i$ 

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SumMor : (i : Size) (j : Size< (↑ i))
  → SumObj i → SumObj j
SumMor i j = map (P.Mor i j) (Q.Mor i j)

SumMorId : {i : Size} {x : SumObj i}
  → SumMor i i x ≡ x
SumMorId {i} {inj1 p} =
  begin
    inj1 (P.Mor i i p)
  ≡⟨ cong inj1 P.MorId ⟩
    inj1 p
  ■

SumMorId {i} {inj2 q} =
  begin
    inj2 (Q.Mor i i q)
  ≡⟨ cong inj2 Q.MorId ⟩
    inj2 q
  ■

SumMorComp : {i : Size} {j : Size< (↑ i)} {k : Size< (↑ j)}
  → {x : SumObj i}
  → SumMor i k x ≡ SumMor j k (SumMor i j x)
SumMorComp {i} {j} {k} {inj1 p} =
  begin
    inj1 (P.Mor i k p)
  ≡⟨ cong inj1 P.MorComp ⟩
    inj1 (P.Mor j k (P.Mor i j p))
  ■

SumMorComp {i} {j} {k} {inj2 q} =
  begin
    inj2 (Q.Mor i k q)
  ≡⟨ cong inj2 Q.MorComp ⟩
    inj2 (Q.Mor j k (Q.Mor i j q))
  ■

Sum : PSh
Sum = record
  { Obj = SumObj
  ; Mor = SumMor
  ; MorId = SumMorId
  ; MorComp = λ {}{}{}{x} → SumMorComp {x = x}
  }

ExpObj : Size → Set
ExpObj i =
  Σ ((j : Size< (↑ i)) → P.Obj j → Q.Obj j)

```

$$\begin{aligned}
& (\lambda f \rightarrow (j : \text{Size} < (\uparrow i)) (k : \text{Size} < (\uparrow j)) \\
& \quad (x : \text{P.Obj } j) \\
& \quad \rightarrow \text{Q.Mor } j \ k (f \ j \ x) \\
& \quad \equiv \\
& \quad f \ k (\text{P.Mor } j \ k \ x))
\end{aligned}$$

$$\begin{aligned}
\text{ExpMor} & : (i : \text{Size}) (j : \text{Size} < (\uparrow i)) \\
& \rightarrow \text{ExpObj } i \rightarrow \text{ExpObj } j \\
\text{ExpMor } i \ j (f, p) & = (\lambda \rightarrow f), (\lambda \rightarrow p)
\end{aligned}$$

$$\begin{aligned}
\text{ExpMorId} & : \{i : \text{Size}\} \{x : \text{ExpObj } i\} \\
& \rightarrow \text{ExpMor } i \ i \ x \equiv x \\
\text{ExpMorId} & = \text{refl}
\end{aligned}$$

$$\begin{aligned}
\text{ExpMorComp} & : \{i : \text{Size}\} \{j : \text{Size} < (\uparrow i)\} \{k : \text{Size} < (\uparrow j)\} \\
& \rightarrow \{x : \text{ExpObj } i\} \\
& \rightarrow \text{ExpMor } i \ k \ x \equiv \text{ExpMor } j \ k (\text{ExpMor } i \ j \ x) \\
\text{ExpMorComp} & = \text{refl}
\end{aligned}$$

$$\begin{aligned}
\text{Exp} & : \text{PSh} \\
\text{Exp} & = \text{record} \\
& \{ \text{Obj} = \text{ExpObj} \\
& \ ; \ \text{Mor} = \text{ExpMor} \\
& \ ; \ \text{MorId} = \text{ExpMorId} \\
& \ ; \ \text{MorComp} = \text{ExpMorComp} \\
& \}
\end{aligned}$$

$$\begin{aligned}
\text{Ctx} & : \text{tag} \rightarrow \text{Set}_1 \\
\text{Ctx set} & = \text{Set} \\
\text{Ctx tot} & = \text{PSh}
\end{aligned}$$

$$\begin{aligned}
\text{Ty} & : \text{tag} \rightarrow \text{Set}_1 \\
\text{Ty set} & = \text{Set} \\
\text{Ty tot} & = \text{PSh}
\end{aligned}$$

$$\begin{aligned}
\text{Tm} & : (b : \text{tag}) ( : \text{Ctx } b) (A : \text{Ty } b) \rightarrow \text{Set} \\
\text{Tm set } A & = \rightarrow A \\
\text{Tm tot } A & = \\
& \Sigma ((i : \text{Size}) \rightarrow \text{PSh.Obj } i \rightarrow \text{PSh.Obj } A \ i) \\
& (\lambda f \rightarrow (i : \text{Size}) (j : \text{Size} < (\uparrow i)) (x : \text{PSh.Obj } i) \\
& \quad \rightarrow \text{PSh.Mor } A \ i \ j (f \ i \ x) \equiv f \ j (\text{PSh.Mor } i \ j \ x))
\end{aligned}$$

- : (b : tag) → Ctx b
- set =  $\top$
- tot = Terminal

$\text{,,} : \{b : \text{tag}\} \rightarrow \text{Ctx } b \rightarrow \text{Ty } b \rightarrow \text{Ctx } b$   
 $\text{,,} \{ \text{set} \} A = \times A$   
 $\text{,,} \{ \text{tot} \} A = \text{Prod } A$

$\text{var} : \{b : \text{tag}\} ( : \text{Ctx } b) (A : \text{Ty } b) \rightarrow \text{Tm } b ( \text{,, } A) A$   
 $\text{var} \{ \text{set} \} A = \text{proj}_2$   
 $\text{proj}_1 (\text{var} \{ \text{tot} \} A) i (y, x) = x$   
 $\text{proj}_2 (\text{var} \{ \text{tot} \} A) i j (y, x) = \text{refl}$

$\text{weaken} : \{b : \text{tag}\} ( : \text{Ctx } b) (A B : \text{Ty } b)$   
 $\rightarrow \text{Tm } b B \rightarrow \text{Tm } b ( \text{,, } A) B$   
 $\text{weaken} \{ \text{set} \} A B t (x, ) = t x$   
 $\text{proj}_1 (\text{weaken} \{ \text{tot} \} A B (t, p)) i (x_1, x_2) = t i x_1$   
 $\text{proj}_2 (\text{weaken} \{ \text{tot} \} A B (t, p)) i j (x_1, x_2) = p i j x_1$

$\text{subst-Tm} : \{b : \text{tag}\} \{ : \text{Ctx } b \} \{A B : \text{Ty } b\}$   
 $\rightarrow (t : \text{Tm } b ( \text{,, } A) B) (x : \text{Tm } b A)$   
 $\rightarrow \text{Tm } b B$   
 $\text{subst-Tm} \{ \text{set} \} t x y = t (y, (x y))$   
 $\text{proj}_1 (\text{subst-Tm} \{ \text{tot} \} (t, p) (x, q)) i y = t i (y, x i y)$   
 $\text{proj}_2 (\text{subst-Tm} \{ \text{tot} \} \{ \} \{A\} \{B\} (t, p) (x, q)) i j y =$   
 $\text{begin}$   
 $\text{PSh.Mor } B i j (t i (y, x i y))$   
 $\equiv \langle p i j (y, x i y) \rangle$   
 $t j (\text{PSh.Mor } ( \text{,, } A) i j (y, x i y))$   
 $\equiv \langle \text{cong } (\lambda z \rightarrow t j ( , z)) (q i j y) \rangle$   
 $t j (\text{PSh.Mor } i j y, x j (\text{PSh.Mor } i j y))$

■

$\text{Unit} : (b : \text{tag}) \rightarrow \text{Ty } b$   
 $\text{Unit set} = \top$   
 $\text{Unit tot} = \text{Terminal}$

$\otimes : \{b : \text{tag}\} (A B : \text{Ty } b) \rightarrow \text{Ty } b$   
 $\otimes \{ \text{set} \} A B = A \times B$   
 $\otimes \{ \text{tot} \} A B = \text{Prod } A B$

$\text{pair} : \{b : \text{tag}\} ( : \text{Ctx } b) (A B : \text{Ty } b) (x : \text{Tm } b A) (y : \text{Tm } b B)$   
 $\rightarrow \text{Tm } b (A \otimes B)$   
 $\text{pair} \{ \text{set} \} A B x y t = x t, y t$   
 $\text{proj}_1 (\text{pair} \{ \text{tot} \} A B (x, p) (y, q)) i t = (x i t), (y i t)$   
 $\text{proj}_2 (\text{pair} \{ \text{tot} \} A B (x, p) (y, q)) i j t =$   
 $\text{begin}$   
 $(\text{PSh.Mor } A i j (x i t), \text{PSh.Mor } B i j (y i t))$   
 $\equiv \langle \text{cong } (\lambda z \rightarrow (z, )) (p i j t) \rangle$   
 $(x j (\text{PSh.Mor } i j t), \text{PSh.Mor } B i j (y i t))$   
 $\equiv \langle \text{cong } (\lambda z \rightarrow ( , z)) (q i j t) \rangle$

$$(x\ j\ (\text{PSh.Mor } i\ j\ t) ,\ y\ j\ (\text{PSh.Mor } i\ j\ t))$$

■

$\text{pr}_1 : \{b : \text{tag}\} ( : \text{Ctx } b) (A\ B : \text{Ty } b) \rightarrow \text{Tm } b\ (A \otimes B) \rightarrow \text{Tm } b\ A$   
 $\text{pr}_1 \{\text{set}\} A\ B\ x\ t = \text{proj}_1 (x\ t)$   
 $\text{proj}_1 (\text{pr}_1 \{\text{tot}\} A\ B\ (x ,\ p))\ i\ t = \text{proj}_1 (x\ i\ t)$   
 $\text{proj}_2 (\text{pr}_1 \{\text{tot}\} A\ B\ (x ,\ p))\ i\ j\ t = \text{cong } \text{proj}_1 (p\ i\ j\ t)$

$\text{pr}_2 : \{b : \text{tag}\} ( : \text{Ctx } b) (A\ B : \text{Ty } b) \rightarrow \text{Tm } b\ (A \otimes B) \rightarrow \text{Tm } b\ B$   
 $\text{pr}_2 \{\text{set}\} A\ B\ x\ t = \text{proj}_2 (x\ t)$   
 $\text{proj}_1 (\text{pr}_2 \{\text{tot}\} A\ B\ (x ,\ p))\ i\ t = \text{proj}_2 (x\ i\ t)$   
 $\text{proj}_2 (\text{pr}_2 \{\text{tot}\} A\ B\ (x ,\ p))\ i\ j\ t = \text{cong } \text{proj}_2 (p\ i\ j\ t)$

$\text{pr}_1\text{-pair} : \{b : \text{tag}\} ( : \text{Ctx } b) (A\ B : \text{Ty } b) (x : \text{Tm } b\ A) (y : \text{Tm } b\ B)$   
 $\rightarrow \text{def-eq } A$   
 $(\text{pr}_1\ A\ B\ (\text{pair } A\ B\ x\ y))$   
 $x$   
 $\text{pr}_1\text{-pair} \{\text{set}\} A\ B\ x\ y\ t = \text{refl}$   
 $\text{pr}_1\text{-pair} \{\text{tot}\} A\ B\ x\ y\ i\ t = \text{refl}$

$\text{pr}_2\text{-pair} : \{b : \text{tag}\} ( : \text{Ctx } b) (A\ B : \text{Ty } b) (x : \text{Tm } b\ A) (y : \text{Tm } b\ B)$   
 $\rightarrow \text{def-eq } B$   
 $(\text{pr}_2\ A\ B\ (\text{pair } A\ B\ x\ y))$   
 $y$   
 $\text{pr}_2\text{-pair} \{\text{set}\} A\ B\ x\ y\ t = \text{refl}$   
 $\text{pr}_2\text{-pair} \{\text{tot}\} A\ B\ x\ y\ i\ t = \text{refl}$

$\text{prod-eta} : \{b : \text{tag}\} ( : \text{Ctx } b) (A\ B : \text{Ty } b) (x : \text{Tm } b\ (A \otimes B))$   
 $\rightarrow \text{def-eq } (A \otimes B)$   
 $(\text{pair } A\ B\ (\text{pr}_1\ A\ B\ x) (\text{pr}_2\ A\ B\ x))$   
 $x$   
 $\text{prod-eta} \{\text{set}\} A\ B\ x\ t = \text{refl}$   
 $\text{prod-eta} \{\text{tot}\} A\ B\ x\ i\ t = \text{refl}$

$\Rightarrow : \{b : \text{tag}\} (A\ B : \text{Ty } b) \rightarrow \text{Ty } b$   
 $\Rightarrow \{\text{set}\} A\ B = A \rightarrow B$   
 $\Rightarrow \{\text{tot}\} A\ B = \text{Exp } A\ B$

$\text{lambda} : \{b : \text{tag}\} ( : \text{Ctx } b) (A\ B : \text{Ty } b) (t : \text{Tm } b\ ( ,\ A)\ B) \rightarrow \text{Tm } b\ (A \Rightarrow B)$   
 $\text{lambda} \{\text{set}\} A\ B\ t\ x\ y = t\ (x ,\ y)$   
 $\text{proj}_1 (\text{proj}_1 (\text{lambda} \{\text{tot}\} A\ B\ (t ,\ p))\ i\ x)\ j\ z = t\ j\ (\text{Mor } i\ j\ x ,\ z)$   
 $\text{where open PSh renaming (Mor to Mor; MorComp to MorComp)}$   
 $\text{proj}_2 (\text{proj}_1 (\text{lambda} \{\text{tot}\} A\ B\ (t ,\ p))\ i\ x)\ j\ k\ y =$   
 $\text{begin}$   
 $\text{PSh.Mor } B\ j\ k\ (t\ j\ (\text{Mor } i\ j\ x ,\ y))$   
 $\equiv \langle p\ j\ k\ (\text{Mor } i\ j\ x ,\ y) \rangle$   
 $t\ k\ (\text{PSh.Mor } ( ,\ A)\ j\ k\ (\text{Mor } i\ j\ x ,\ y))$

```

≡ ( cong (λ z → t k (z , )) (sym MorComp) )
  t k (Mor i k x , PSh.Mor A j k y)
■

where open PSh renaming (Mor to Mor; MorComp to MorComp)
proj2 (lambda {tot} A B (t , p)) i j x =
  Σ≡-uip (funext (λ → funext (λ → funext (λ → uip))))
    (funext (λ k → (funext (λ z → cong (λ z → t k (z , )) MorComp))))
where open PSh renaming (Mor to Mor; MorComp to MorComp)

app : {b : tag} { : Ctx b} {A B : Ty b} (f : Tm b (A ⇒ B)) (t : Tm b A) → Tm b B
app {set} f t x = f x (t x)
proj1 (app {tot} {} (f , p) (t , q)) i x = let (f' , ) = f i x in f' (t i x)
proj2 (app {tot} {} {A} {B} (f , p) (t , q)) i j x =
  let (f' , p') = f i x in
begin
  PSh.Mor B i j (proj1 (f i x) (t i x))
≡⟨ p' i j (t i x) ⟩
  proj1 (f i x) j (PSh.Mor A i j (t i x))
≡⟨ cong2 (λ z g → proj1 g z) (q i j x) (p i j x) ⟩
  proj1 (f j (PSh.Mor i j x)) (t j (PSh.Mor i j x))
■

beta : {b : tag} { : Ctx b} {A B : Ty b} (t : Tm b ( , A) B) (x : Tm b A)
  → def-eq B (app {b} {} {A} {B} (lambda A B t) x) (subst-Tm {} {} {A} {B} t x)
beta {set} t x = refl
beta {tot} {} (t , p) (x , q) z = cong (λ z → t (z , )) World
  where open PSh renaming (MorId to MorId)

eta : {b : tag} { : Ctx b} {A B : Ty b} (t : Tm b (A ⇒ B))
  → def-eq (A ⇒ B)
    (lambda A B (app {} {} , A) {A} {B} (weaken A (A ⇒ B) t) (var A)))
    t
eta {set} t x = refl
eta {tot} (t , p) x =
  Σ≡-uip (funext (λ → funext (λ → funext (λ → uip))))
    (funext (λ ' → funext (λ z → sym (cong (λ h → proj1 h z) (p ' x)))))

id-tm : {b : tag} ( : Ctx b) (A : Ty b) → Tm b (A ⇒ A)
id-tm A = lambda A A (var A)

comp-tm : {b : tag} ( : Ctx b) (A B C : Ty b)
  → Tm b ((B ⇒ C) ⇒ ((A ⇒ B) ⇒ (A ⇒ C)))
comp-tm A B C = lambda (B ⇒ C) ((A ⇒ B) ⇒ (A ⇒ C))
  (lambda ( , (B ⇒ C)) (A ⇒ B) (A ⇒ C)
    (lambda (( , (B ⇒ C)) , (A ⇒ B)) A C (app
      (weaken (( , (B ⇒ C)) , (A ⇒ B)) A (

```



$$(\text{app } (\text{weaken } ((\text{,, } (B \Rightarrow C)) \text{,, } (A \Rightarrow B))) A (A \Rightarrow B) (\text{var } (\text{,, } (B \Rightarrow C)) \text{,, } (A \Rightarrow B))) A)))))$$
 $\square : \text{Ty tot} \rightarrow \text{Ty set}$ 
 $\square A =$ 

$$\begin{aligned} & \Sigma ((i : \text{Size}) \rightarrow \text{PSh.Obj } A \ i) \\ & (\lambda x \rightarrow (i : \text{Size}) (j : \text{Size} < (\uparrow i)) \\ & \rightarrow \text{PSh.Mor } A \ i \ j (x \ i) \equiv x \ j) \end{aligned}$$
 $\text{data SizeLt } (i : \text{Size}) : \text{Set where}$ 

$$[] : (j : \text{Size} < i) \rightarrow \text{SizeLt } i$$
 $\text{size} : \forall \{i\} \rightarrow \text{SizeLt } i \rightarrow \text{Size}$ 
 $\text{size } [j] = j$ 
 $\text{elimLt} : \forall \{\ell\} \{A : \text{Size} \rightarrow \text{Set } \ell\} \{i : \text{Size}\} (j : \text{SizeLt } i)$ 

$$\rightarrow ((j : \text{Size} < i) \rightarrow A \ j) \rightarrow A \ (\text{size } j)$$
 $\text{elimLt } [j] f = f \ j$ 
 $\text{Later} : (\text{Size} \rightarrow \text{Set}) \rightarrow \text{Size} \rightarrow \text{Set}$ 
 $\text{Later } A \ i = (j : \text{SizeLt } i) \rightarrow A \ (\text{size } j)$ 
 $\text{module } (A : \text{Size} \rightarrow \text{Set}) (m : (i : \text{Size}) (j : \text{Size} < (\uparrow i)) \rightarrow A \ i \rightarrow A \ j) \text{ where}$ 
 $\text{LaterLim} : (i : \text{Size}) (x : \text{Later } A \ i) \rightarrow \text{Set}$ 
 $\text{LaterLim } i \ x = (j : \text{SizeLt } i)$ 

$$\rightarrow \text{elimLt } j (\lambda \{j' \rightarrow (k : \text{SizeLt } (\uparrow j'))$$

$$\rightarrow \text{elimLt } k (\lambda k' \rightarrow m \ j' \ k' (x \ [j']) \equiv x \ [k']) \}$$
 $\text{LaterLimMor} : (i : \text{Size}) (j : \text{Size} < (\uparrow i)) (x : \text{Later } A \ i)$ 

$$\rightarrow \text{LaterLim } i \ x \rightarrow \text{LaterLim } j \ x$$
 $\text{LaterLimMor } i \ j \ x \ p \ [k] \ [l] = p \ [k] \ [l]$ 
 $\text{module } (A : \text{Ty tot}) \text{ where}$ 
 $-- \text{ 3. Object part}$ 
 $\triangleright \text{Obj} : (i : \text{Size}) \rightarrow \text{Set}$ 
 $\triangleright \text{Obj } i = \Sigma (\text{Later } (\text{PSh.Obj } A) \ i) (\text{LaterLim } (\text{PSh.Obj } A) (\text{PSh.Mor } A) \ i)$ 
 $-- \text{ 4. Morphism part}$ 
 $\triangleright \text{Mor} : (i : \text{Size}) (j : \text{Size} < (\uparrow i))$ 

$$\rightarrow \triangleright \text{Obj } i \rightarrow \triangleright \text{Obj } j$$
 $\triangleright \text{Mor } i \ j \ (x, p) = x, \text{LaterLimMor } (\text{PSh.Obj } A) (\text{PSh.Mor } A) \ i \ j \ x \ p$ 
 $\text{where}$ 

$$p' : \text{LaterLim } (\text{PSh.Obj } A) (\text{PSh.Mor } A) \ j \ x$$

$$p' \ [j] \ [k] = p \ [j] \ [k]$$

```

-- 5. Preservation of identity
▷MorId : {i : Size} {x : ▷Obj i}
        → ▷Mor i i x ≡ x
▷MorId = Σ≡-uip (funext (λ { [ j ] → funext (λ { [ k ] → uip } ) } ) refl

-- 6. Preservation of composition
▷MorComp : {i : Size} {j : Size< (↑ i)} {k : Size< (↑ j)} {x : ▷Obj i}
        → ▷Mor i k x ≡ ▷Mor j k (▷Mor i j x)
▷MorComp = Σ≡-uip (funext (λ { [ j ] → funext (λ { [ k ] → uip } ) } ) refl

▷ : Ty tot
▷ = record
  { Obj = ▷Obj
  ; Mor = ▷Mor
  ; MorId = ▷MorId
  ; MorComp = ▷MorComp
  }

pure : ( : Ctx tot) (A : Ty tot) (t : Tm tot A) → Tm tot (▷ A)
proj₁ (proj₁ (pure A (t , ) i x) [ j ] = t j (PSh.Mor i j x)
proj₂ (proj₁ (pure A (t , p)) i x) [ j ] [ k ] =
begin
  PSh.Mor A j k (t j (PSh.Mor i j x))
≡⟨ p j k (PSh.Mor i j x) ⟩
  t k (PSh.Mor j k (PSh.Mor i j x))
≡⟨ cong (t k) (sym (PSh.MorComp)) ⟩
  t k (PSh.Mor i k x)
■
proj₂ (pure A (t , p)) i j x =
  Σ≡-uip (funext (λ { [ ] → funext (λ { [ ] → uip } ) } )
    (funext (λ { [ k ] → cong (t k) (PSh.MorComp ) } ) )

fmap : ( : Ctx tot) (A B : Ty tot)
      → (f : Tm tot (▷ (A ⇒ B))) (t : Tm tot (▷ A))
      → Tm tot (▷ B)
proj₁ (proj₁ (fmap A B (f , ) (t , ) i x) [ j ] = proj₁ (proj₁ (f i x) [ j ]) j (proj₁ (t i x) [ j ])
proj₂ (proj₁ (fmap A B (f , p) (t , q)) i x) [ j ] [ k ] =
begin
  PSh.Mor B j k (proj₁ (proj₁ (f i x) [ j ]) j (proj₁ (t i x) [ j ]))
≡⟨ proj₂ (proj₁ (f i x) [ j ]) j k (proj₁ (t i x) [ j ]) ⟩
  proj₁ (proj₁ (f i x) [ j ]) k (PSh.Mor A j k (proj₁ (t i x) [ j ]))
≡⟨ cong (proj₁ (proj₁ (f i x) [ j ]) k) (proj₂ (t i x) [ j ] [ k ]) ⟩
  proj₁ (proj₁ (f i x) [ j ]) k (proj₁ (t i x) [ k ])
≡⟨ cong (λ z → proj₁ z k (proj₁ (t i x) [ k ])) (sym (proj₂ (f i x) [ j ] [ j ])) ⟩

```

$\text{proj}_1 (\text{PSh.Mor } (A \Rightarrow B) j j (\text{proj}_1 (f i x) [j])) k (\text{proj}_1 (t i x) [k])$   
 $\equiv \langle \text{cong } (\lambda z \rightarrow \text{proj}_1 z k (\text{proj}_1 (t i x) [k])) (\text{proj}_2 (f i x) [j] [k]) \rangle$   
 $\text{proj}_1 (\text{proj}_1 (f i x) [k]) k (\text{proj}_1 (t i x) [k])$

■

$\text{proj}_2 (\text{fmap } A B (f, p) (e, q)) i j x =$   
 $\Sigma \equiv \text{-uip } (\text{funext } (\lambda \{ [ ] \} \rightarrow \text{funext } (\lambda \{ [ ] \} \rightarrow \text{uip } \{ \})))$   
 $(\text{funext } (\lambda \{ [k] \} \rightarrow \text{cong}_2 (\lambda a b \rightarrow \text{proj}_1 (\text{proj}_1 a [k]) k (\text{proj}_1 b [k])) (p i j x) (q i j x) \{ \})))$

$\text{pure-fmap-pure} : ( : \text{Ctx tot}) (A B : \text{Ty tot})$   
 $\rightarrow (f : \text{Tm tot } (A \Rightarrow B)) (t : \text{Tm tot } A)$   
 $\rightarrow \text{def-eq } (\triangleright B) (\text{fmap } A B (\text{pure } (A \Rightarrow B) f) (\text{pure } A t)) (\text{pure } B (\text{app } \{ \text{tot} \} \{ \} \{ A \} \{ B \} f t))$   
 $\text{pure-fmap-pure } A B (f, p) (t, q) i x =$   
 $\Sigma \equiv \text{-uip } (\text{funext } (\lambda \{ [ ] \} \rightarrow \text{funext } (\lambda \{ [ ] \} \rightarrow \text{uip } \{ \})))$   
 $(\text{funext } (\lambda \{ [j] \} \rightarrow \text{refl } \{ \})))$

$\text{pure-id-fmap} : ( : \text{Ctx tot}) (A B : \text{Ty tot}) (t : \text{Tm tot } (\triangleright A))$   
 $\rightarrow \text{def-eq } (\triangleright A) (\text{fmap } A A (\text{pure } (A \Rightarrow A) (\text{id-tm } A)) t) t$   
 $\text{pure-id-fmap } A B (t, p) i y =$   
 $\Sigma \equiv \text{-uip } (\text{funext } (\lambda \{ [ ] \} \rightarrow \text{funext } (\lambda \{ [ ] \} \rightarrow \text{uip } \{ \})))$   
 $(\text{funext } (\lambda \{ [j] \} \rightarrow \text{refl } \{ \})))$

$\text{pure-comp-fmap} : ( : \text{Ctx tot}) (A B C : \text{Ty tot})$   
 $\rightarrow (g : \text{Tm tot } (\triangleright (B \Rightarrow C))) (f : \text{Tm tot } (\triangleright (A \Rightarrow B))) (t : \text{Tm tot } (\triangleright A))$   
 $\rightarrow \text{def-eq}$   
 $(\triangleright C)$   
 $(\text{fmap } A C (\text{fmap } (A \Rightarrow B) (A \Rightarrow C) (\text{fmap } (B \Rightarrow C) ((A \Rightarrow B) \Rightarrow (A \Rightarrow C)) (\text{pure } ((B \Rightarrow C) \Rightarrow (A \Rightarrow B) \Rightarrow (A \Rightarrow C))$   
 $(\text{fmap } B C g (\text{fmap } A B f t))))$

$\text{pure-comp-fmap } A B C g f t i y =$   
 $\Sigma \equiv \text{-uip } (\text{funext } (\lambda \{ [ ] \} \rightarrow \text{funext } (\lambda \{ [ ] \} \rightarrow \text{uip } \{ \})))$   
 $(\text{funext } (\lambda \{ [j] \} \rightarrow \text{refl } \{ \})))$

$\text{fmap-pure-fun} : ( : \text{Ctx tot}) (A B : \text{Ty tot})$   
 $\rightarrow (f : \text{Tm tot } (\triangleright (A \Rightarrow B))) (t : \text{Tm tot } A)$   
 $\rightarrow \text{def-eq}$

$(\triangleright B)$   
 $(\text{fmap } A B f (\text{pure } A t))$   
 $(\text{fmap } (A \Rightarrow B) B (\text{pure } ((A \Rightarrow B) \Rightarrow B) (\text{lambda } (A \Rightarrow B) B (\text{app } \{ \} \{ \text{,, } (A \Rightarrow B) \} \{ A \} \{ B \} (\text{var } (A \Rightarrow B) \Rightarrow (A \Rightarrow B) \Rightarrow (A \Rightarrow B))$   
 $\text{fmap-pure-fun } A B (f, p) (t, q) i y =$   
 $\Sigma \equiv \text{-uip } (\text{funext } (\lambda \{ [ ] \} \rightarrow \text{funext } (\lambda \{ [ ] \} \rightarrow \text{uip } \{ \})))$   
 $(\text{funext } (\lambda \{ [j] \} \rightarrow \text{cong } (\lambda a \rightarrow \text{proj}_1 (\text{proj}_1 (f i y) [j]) j (t j a)) (\text{sym } (\text{PSh.MorId } \{ \}))))$

$\text{WCObj} : \text{Size} \rightarrow \text{Set}$   
 $\text{WCObj } = A$

$\text{WCMor} : (i : \text{Size}) (j : \text{Size} < (\uparrow i))$   
 $\rightarrow \text{WCObj } i \rightarrow \text{WCObj } j$

$\text{WCMor } x = x$

$\text{WCMorId} : \{i : \text{Size}\} \{x : \text{WCObj } i\}$

$\rightarrow \text{WCMor } i \ i \ x \equiv x$

$\text{WCMorId} = \text{refl}$

$\text{WCMorComp} : \{i : \text{Size}\} \{j : \text{Size} < (\uparrow i)\} \{k : \text{Size} < (\uparrow j)\}$

$\{x : \text{WCObj } i\}$

$\rightarrow \text{WCMor } i \ k \ x \equiv \text{WCMor } j \ k \ (\text{WCMor } i \ j \ x)$

$\text{WCMorComp} = \text{refl}$

$\text{WC} : \text{Ty tot}$

$\text{WC} = \text{record}$

$\{ \text{Obj} = \text{WCObj}$

$; \text{Mor} = \text{WCMor}$

$; \text{MorId} = \text{WCMorId}$

$; \text{MorComp} = \text{WCMorComp}$

$\}$

$\text{WC-fun} : ( : \text{Ctx set}) (A : \text{Ty set}) \rightarrow \text{Tm set } A \rightarrow \text{Tm tot } (\text{WC}) (\text{WC } A)$

$\text{proj}_1 (\text{WC-fun } A \ t) = t$

$\text{proj}_2 (\text{WC-fun } A \ t) = \text{refl}$

$\text{WC-unfun} : ( : \text{Ctx set}) (A : \text{Ty set}) \rightarrow \text{Tm tot } (\text{WC}) (\text{WC } A) \rightarrow \text{Tm set } A$

$\text{WC-unfun } A \ (t, p) = t \infty$

$\text{dfix}_1 : (A : \text{Ty tot}) (i : \text{Size}) \rightarrow \text{ExpObj } (\triangleright A) \ A \ i \rightarrow \triangleright \text{Obj } A \ i$

$\text{proj}_1 (\text{dfix}_1 \ A \ i \ (f, p)) \ [j] = f \ j \ (\text{dfix}_1 \ A \ j \ (f, p))$

$\text{proj}_2 (\text{dfix}_1 \ A \ i \ (f, p)) \ [j] \ [k] =$

$\text{begin}$

$\text{PSh.Mor } A \ j \ k \ (f \ j \ (\text{dfix}_1 \ A \ j \ (f, p)))$

$\equiv \langle p \ j \ k \ (\text{dfix}_1 \ A \ j \ (f, p)) \rangle$

$f \ k \ (\triangleright \text{Mor } A \ j \ k \ (\text{dfix}_1 \ A \ j \ (f, p)))$

$\equiv \langle \text{cong } (f \ k) \ (\Sigma \equiv \text{uip } (\text{funext } (\lambda \{ [j] \rightarrow \text{funext } (\lambda \{ [k] \rightarrow \text{uip} \} \})) (\text{funext } (\lambda \{ [ ] \rightarrow \text{refl} \})))) \rangle$

$f \ k \ (\text{dfix}_1 \ A \ k \ (f, p))$

■

$\text{dfix} : ( : \text{Ctx tot}) (A : \text{Ty tot}) (f : \text{Tm tot } (\triangleright A \Rightarrow A)) \rightarrow \text{Tm tot } (\triangleright A)$

$\text{proj}_1 (\text{dfix } A \ (f, )) \ i \ y = \text{dfix}_1 \ A \ i \ (f \ i \ y)$

$\text{proj}_2 (\text{dfix } A \ (f, p)) \ i \ j \ y =$

$\Sigma \equiv \text{uip } (\text{funext } (\lambda \{ [j] \rightarrow \text{funext } (\lambda \{ [k] \rightarrow \text{uip} \} \}))$

$(\text{funext } (\lambda \{ [k] \rightarrow \text{cong } (\lambda a \rightarrow \text{proj}_1 \ a \ k \ (\text{dfix}_1 \ A \ k \ (\text{proj}_1 \ a, \text{proj}_2 \ a))) \ (p \ i \ j \ y) \})))$

$\text{fix} : ( : \text{Ctx tot}) (A : \text{Ty tot}) (f : \text{Tm tot } (\triangleright A \Rightarrow A)) \rightarrow \text{Tm tot } A$

$\text{fix } A \ f = \text{app } \{\text{tot}\} \ \{\} \ \{\triangleright A\} \ \{A\} \ f \ (\text{dfix } A \ f)$

$\text{dfix-eq} : ( : \text{Ctx tot}) (A : \text{Ty tot}) (f : \text{Tm tot } (\triangleright A \Rightarrow A))$

$\rightarrow \text{def-eq } \{\text{tot}\} \ (\triangleright A) \ (\text{dfix } A \ f) \ (\text{pure } A \ (\text{fix } A \ f))$

$\text{dfix-eq } A \ (f, p) \ i \ y =$

$\Sigma \equiv \text{uip } (\text{funext } (\lambda \{ [j] \rightarrow \text{funext } (\lambda \{ [k] \rightarrow \text{uip } \} \})))$   
 $(\text{funext } (\lambda \{ [j] \rightarrow \text{cong } (\lambda a \rightarrow \text{proj}_1 a j (\text{dfix}_1 A j (\text{proj}_1 a, \text{proj}_2 a))) (p \ i \ j \ y))))$

$\text{fix-eq} : ( : \text{Ctx tot}) (A : \text{Ty tot}) (f : \text{Tm tot } (\triangleright A \Rightarrow A))$   
 $\rightarrow \text{def-eq } \{\text{tot}\} A (\text{fix } A f) (\text{app } \{\text{tot}\} \{\} \{\triangleright A\} \{A\} f (\text{pure } A (\text{fix } A f)))$   
 $\text{fix-eq } A f \ i \ x = \text{cong } (\text{proj}_1 (\text{proj}_1 f \ i \ x) \ i) (\text{dfix-eq } A f \ i \ x)$

$\text{force-tm} : ( : \text{Ctx set}) (A : \text{Ty tot}) (t : \text{Tm set } (\Box (\triangleright A))) \rightarrow \text{Tm set } (\Box A)$   
 $\text{proj}_1 (\text{force-tm } A \ t \ x) \ j = \text{proj}_1 (\text{proj}_1 (t \ x) \ \infty) [j]$   
 $\text{proj}_2 (\text{force-tm } A \ t \ x) \ i \ j =$   
 $\text{begin}$   
 $\text{PSh.Mor } A \ i \ j (\text{proj}_1 (\text{proj}_1 (t \ x) \ \infty) [i])$   
 $\equiv \langle \text{proj}_2 (\text{proj}_1 (t \ x) \ \infty) [i] [j] \rangle$   
 $\text{proj}_1 (\text{proj}_1 (t \ x) \ \infty) [j]$   
■

### 2.1. Sized Types.

### 2.2. Guarded Recursive Type Theory.

## 3. THE MODEL

### 3.1. Clock Contexts.

### 3.2. Presheaves.

## 4. THE INTERPRETATION

### 4.1. Types, Contexts, Terms.

### 4.2. Operations on Contexts.

### 4.3. Substitution.

### 4.4. Simple Types.

### 4.5. Later.

### 4.6. Clock Quantification.

### 4.7. Fix.

### 4.8. Inductive Types.

## 5. CONCLUSION

## APPENDIX A. OMITTED PROOFS

## REFERENCES

- [1] Robert Atkey and Conor McBride. Productive coprogramming with guarded recursion. In *ACM SIGPLAN Notices*, volume 48, pages 197–208. ACM, 2013.