

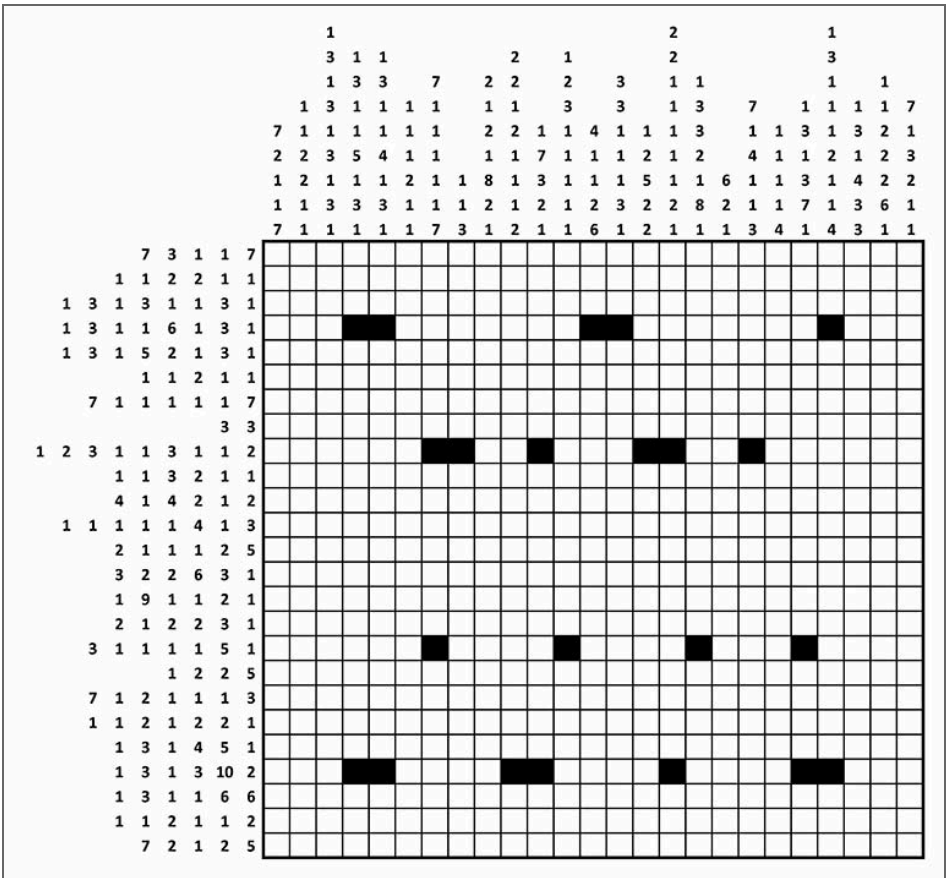
Linear Programming

CS5491: Artificial Intelligence
ZHICHAO LU

Content Credits: **Prof. Wei**'s CS4486 Course
and **Prof. Boddeti**'s AI Course

TODAY

- Solving Linear Program
- Integer Program
- Branch and Bound



GCHQ

LINEAR PROGRAMMING: WHAT TO EAT?

We are trying healthy by finding the optimal amount of food to purchase.

We can choose the amount of **stir-fry** (ounce) and **boba** (fluid ounces).

| Health Goals | Food | Cost | Calories | Sugar | Calcium |
|---|--------------------------|------|----------|-------|---------|
| • $2000 \leq \textit{Calories} \leq 2500$ | stir-fry (per oz) | 1 | 100 | 3 | 20 |
| | Boba (per fl oz) | 0.5 | 50 | 4 | 70 |
| • $\textit{Sugar} \leq 100g$ | | | | | |
| • $\textit{Calcium} \geq 700mg$ | | | | | |

What is the cheapest way to stay "healthy" with this menu?

How much **stir-fry** (ounce) and **boba** (fluid ounces) should we buy?

OPTIMIZATION FORMULATION

Diet Problem

$$\begin{array}{ll}\min_{\mathbf{x}} & \mathbf{c}^T \mathbf{x} \\ s.t. & \mathbf{A}\mathbf{x} \leq \mathbf{b}\end{array}$$

$$A = \begin{bmatrix} -100 & -50 \\ 100 & 50 \\ 3 & 4 \\ -20 & -70 \end{bmatrix} \quad b = \begin{bmatrix} -2000 \\ 2500 \\ 100 \\ -700 \end{bmatrix} \quad \text{and} \quad c = \begin{bmatrix} 1 \\ 0.5 \end{bmatrix}$$

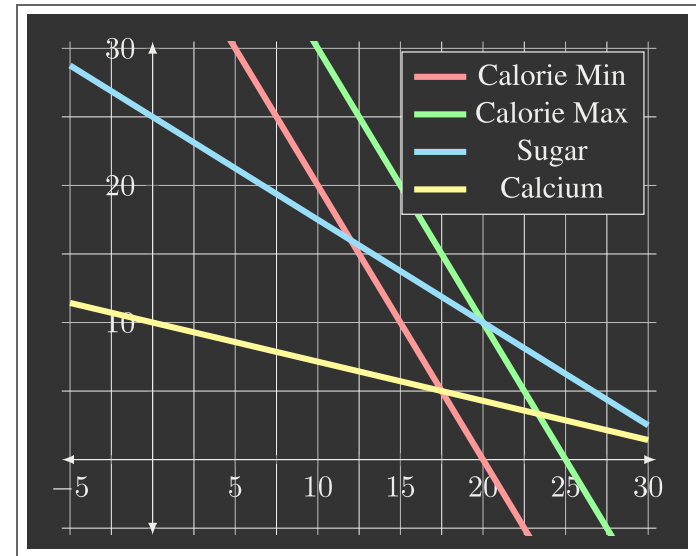
REPRESENTATION AND PROBLEM SOLVING

Problem Description

Optimization Representation

$$\begin{array}{ll}\min_{\mathbf{x}} & \mathbf{c}^T \mathbf{x} \\ s.t. & \mathbf{Ax} \leq \mathbf{b}\end{array}$$

Graphical Representation

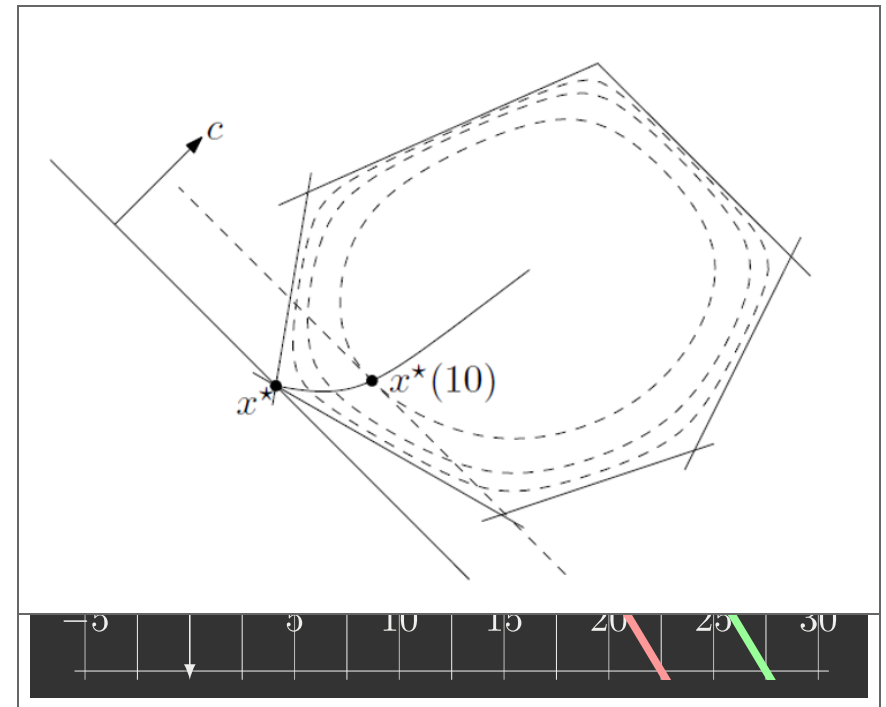


SOLVING A LP

Solutions are at feasible intersections of constraint boundaries.

Algorithms

- Check objective at all feasible intersections.
- Simplex
- Interior Point



SOLVING AN LP

But, how do we find the intersection between boundaries?

$$\begin{array}{ll} \min_{\mathbf{x}} & \mathbf{c}^T \mathbf{x} \\ \text{s.t.} & \mathbf{Ax} \leq \mathbf{b} \end{array} \quad A = \begin{bmatrix} -100 & -50 \\ 100 & 50 \\ 3 & 4 \\ -20 & -70 \end{bmatrix} \quad b = \begin{bmatrix} -2000 \\ 2500 \\ 100 \\ -700 \end{bmatrix} \quad \text{and } c = \begin{bmatrix} 1 \\ 0.5 \end{bmatrix}$$

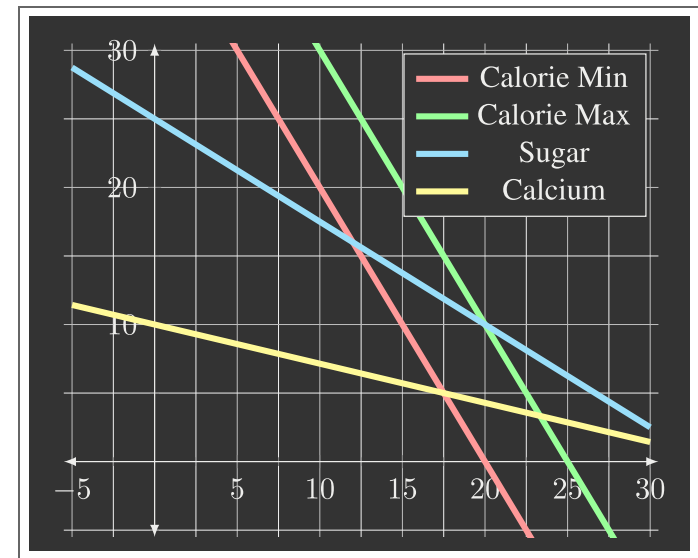
WHAT ABOUT HIGHER DIMENSIONS?

Problem Description

Optimization Representation

$$\begin{array}{ll}\min_{\mathbf{x}} & \mathbf{c}^T \mathbf{x} \\ s.t. & \mathbf{Ax} \leq \mathbf{b}\end{array}$$

Graphical Representation



SHAPES IN HIGHER DIMENSIONS

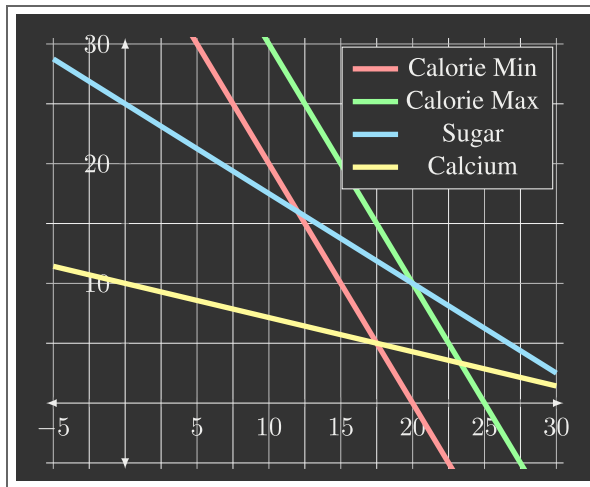
How do these linear shapes extend to 3-D, N-D?

| | 2-D | 3-D | N-D |
|------------------------------------|-----------|------------|------------|
| $a_1x_1 + a_2x_2 = b_1$ | line | plane | hyperplane |
| $a_1x_1 + a_2x_2 \leq b_1$ | halfplane | halfspace | halfspace |
| $a_{1,1}x_1 + a_{1,2}x_2 \leq b_1$ | polygon | polyhedron | polytope |
| $a_{2,1}x_1 + a_{2,2}x_2 \leq b_2$ | | | |
| $a_{3,1}x_1 + a_{3,2}x_2 \leq b_3$ | | | |
| $a_{4,1}x_1 + a_{4,2}x_2 \leq b_4$ | | | |

INTERSECTIONS IN HIGHER DIMENSIONS

How do these linear shapes extend to 3-D, N-D?

$$\begin{array}{ll}\min_{\mathbf{x}} & \mathbf{c}^T \mathbf{x} \\ \text{s.t.} & \mathbf{Ax} \leq \mathbf{b}\end{array}$$



$$\mathbf{A} = \begin{bmatrix} -100 & -50 \\ 100 & 50 \\ 3 & 4 \\ -20 & -70 \end{bmatrix}$$

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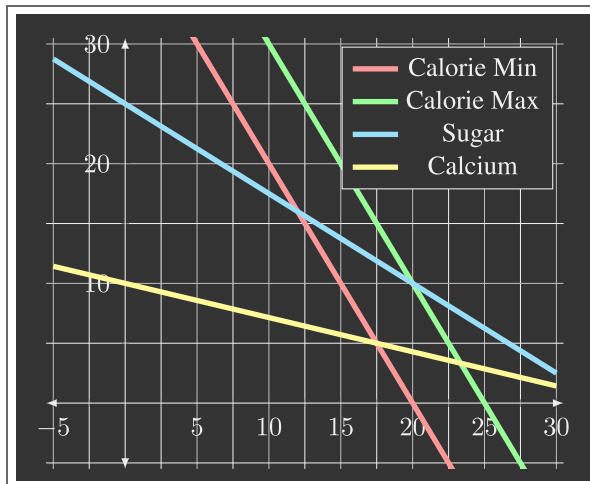
INTERSECTIONS IN HIGHER DIMENSIONS

Still looking at subsets of **A** matrix

$$\begin{array}{ll}\min & \mathbf{c}^T \mathbf{x} \\ \mathbf{x} & \\ s.t. & \mathbf{Ax} \leq \mathbf{b}\end{array}$$

$$\mathbf{A} = \begin{bmatrix} -100 & -50 \\ 100 & 50 \\ 3 & 4 \\ -20 & -70 \end{bmatrix}$$

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LINEAR PROGRAMMING

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LINEAR \rightarrow INTEGER

We are trying healthy by finding the optimal amount of food to purchase.

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LINEAR VS INTEGER PROGRAMMING

Linear objective with linear constraints, but now with additional constraint that all values in \mathbf{x} must be integers

$$\begin{array}{ll}\min_{\mathbf{x}} & \mathbf{c}^T \mathbf{x} \\ s.t. & \mathbf{Ax} \leq \mathbf{b}\end{array}$$

$$\begin{array}{ll}\min_{\mathbf{x}} & \mathbf{c}^T \mathbf{x} \\ s.t. & \mathbf{Ax} \leq \mathbf{b} \\ & \mathbf{x} \in \mathbb{Z}^N\end{array}$$

We could also do:

- Even more constrained: Binary Integer Programming
- A hybrid: Mixed Integer Linear Programming

INTEGER PROGRAMMING: GRAPHICAL REPRESENTATION

Just add a grid of integer points onto our LP representation

$$\begin{array}{ll}\min_{\mathbf{x}} & \mathbf{c}^T \mathbf{x} \\ s.t. & \mathbf{Ax} \leq \mathbf{b} \\ & \mathbf{x} \in \mathbb{Z}^N\end{array}$$

RELAXATION

Relax IP to LP by dropping integer constraints

$$\begin{array}{ll}\min_{\mathbf{x}} & \mathbf{c}^T \mathbf{x} \\ s.t. & \mathbf{Ax} \leq \mathbf{b} \\ & \mathbf{x} \in \mathbb{Z}^N\end{array}$$

Remember Heuristics?

SOLUTION OF LINEAR VS INTEGER PROGRAM

Let y_{IP}^* be the optimal objective of an integer program P .

$$\begin{aligned} y_{IP}^* = \min_{\mathbf{x}} \quad & \mathbf{c}^T \mathbf{x} \\ \text{s.t.} \quad & \mathbf{Ax} \leq \mathbf{b} \\ & \mathbf{x} \in \mathbb{Z}^N \end{aligned}$$

Let \mathbf{x}_{IP}^* be the optimal point of an integer program P .

Let y_{LP}^* be the optimal objective of the LP-relaxed version of P .

$$\begin{aligned} y_{LP}^* = \min_{\mathbf{x}} \quad & \mathbf{c}^T \mathbf{x} \\ \text{s.t.} \quad & \mathbf{Ax} \leq \mathbf{b} \end{aligned}$$

Let \mathbf{x}_{LP}^* be the optimal point of the LP-relaxed version of P .

Assume that P is a minimization problem.

Which of the following are true?

- $\mathbf{x}_{IP}^* = \mathbf{x}_{LP}^*$

- $y_{IP}^* \leq y_{LP}^*$

- $y_{IP}^* \geq y_{LP}^*$

SOLVING INTEGER PROGRAMS: A NAIVE SOLUTION

True/False: It is sufficient to consider the integer points around the corresponding LP solution?

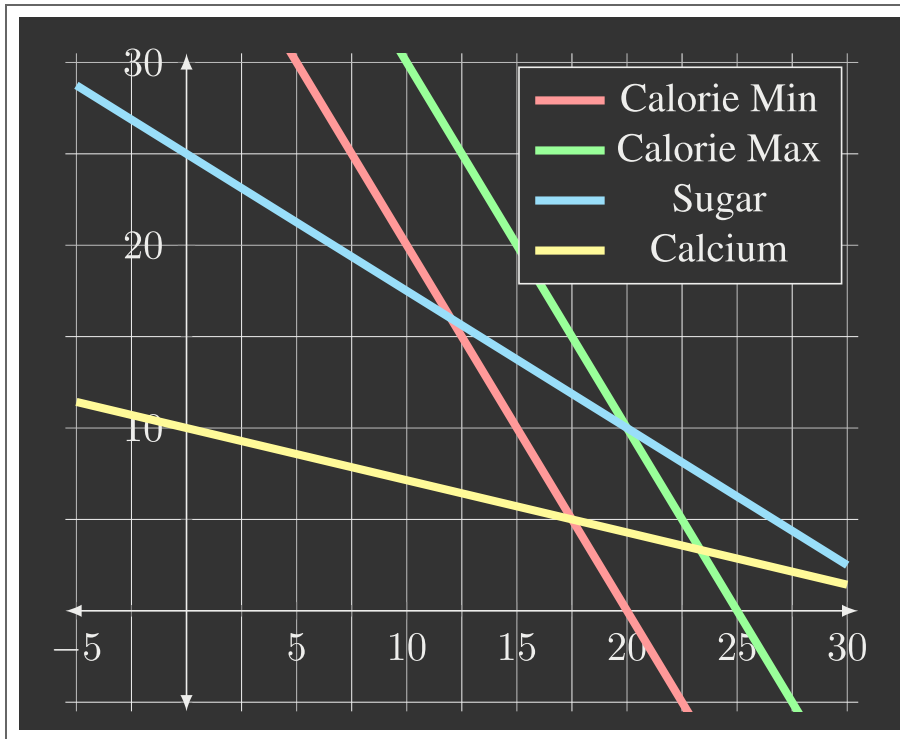
SOLVING AN IP

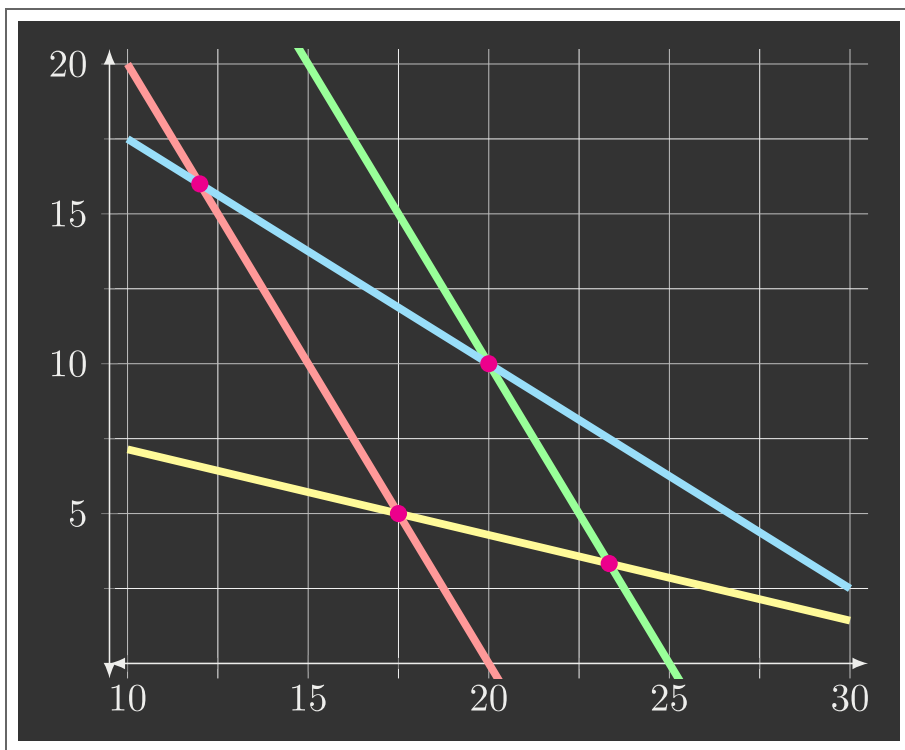
Branch and Bound algorithm

- Start with LP-relaxed version of IP
- If solution \mathbf{x}_{LP}^* has non-integer value at \mathbf{x}_i
- Consider two branches with two different slightly more constrained LP problems:
 - Left branch: Add constraint $\mathbf{x}_i \leq \text{floor}(\mathbf{x}_{LP}^*)$
 - Right branch: Add constraint $\mathbf{x}_i \geq \text{ceil}(\mathbf{x}_{LP}^*)$
- Recursion. Stop going deeper:
 - When the LP returns a worse objective than the best feasible IP objective you have seen before.
 - When you hit an integer result from the LP

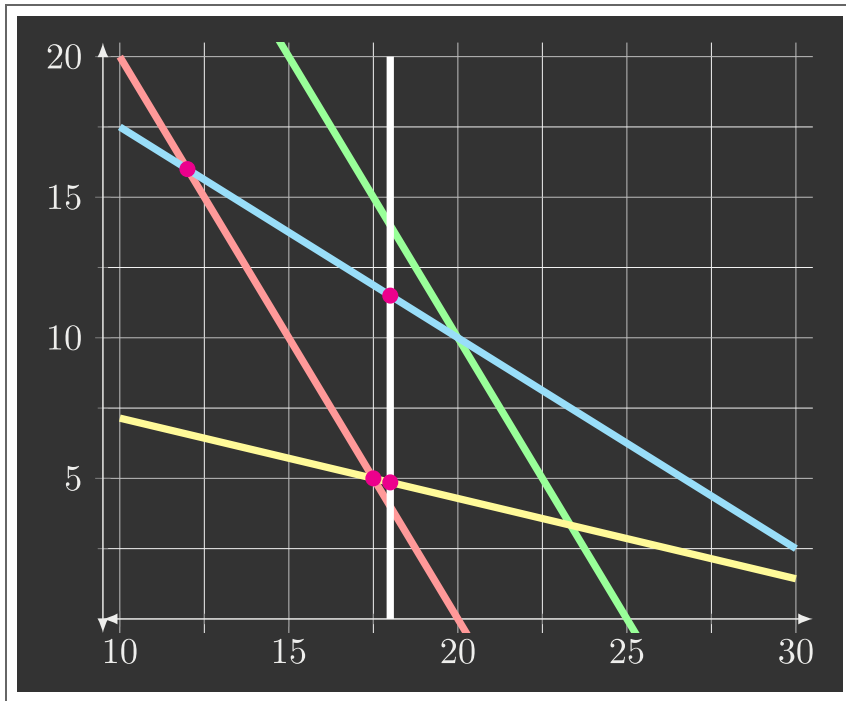
- When LP is infeasible

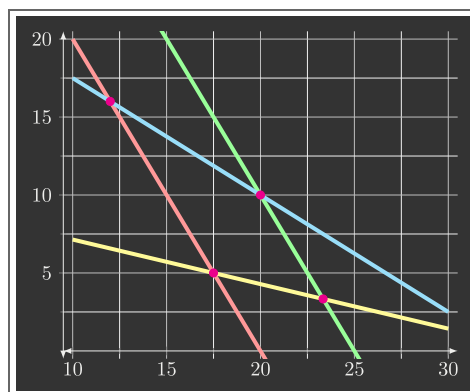
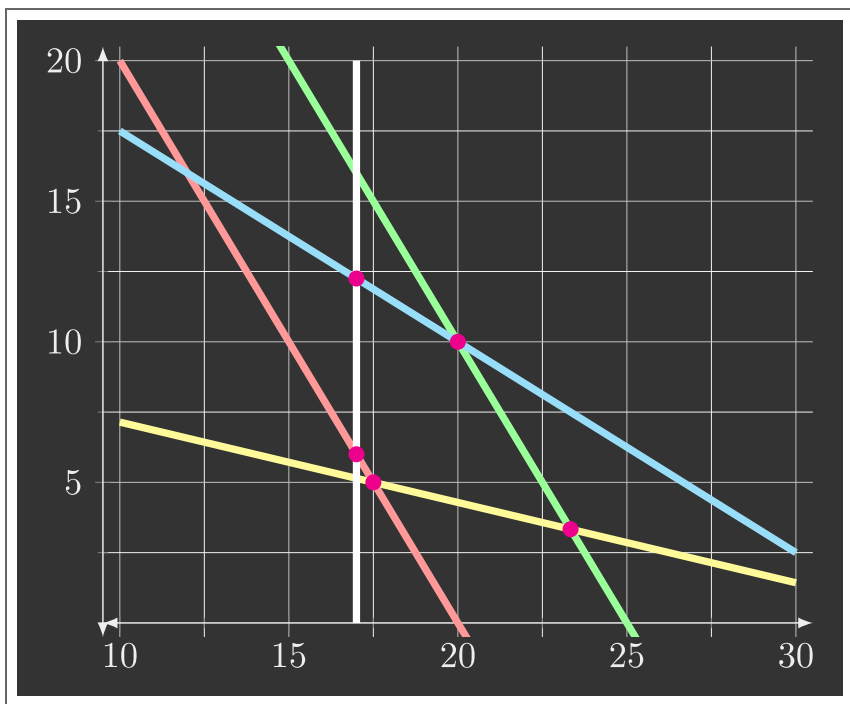
BRANCH AND BOUND EXAMPLE



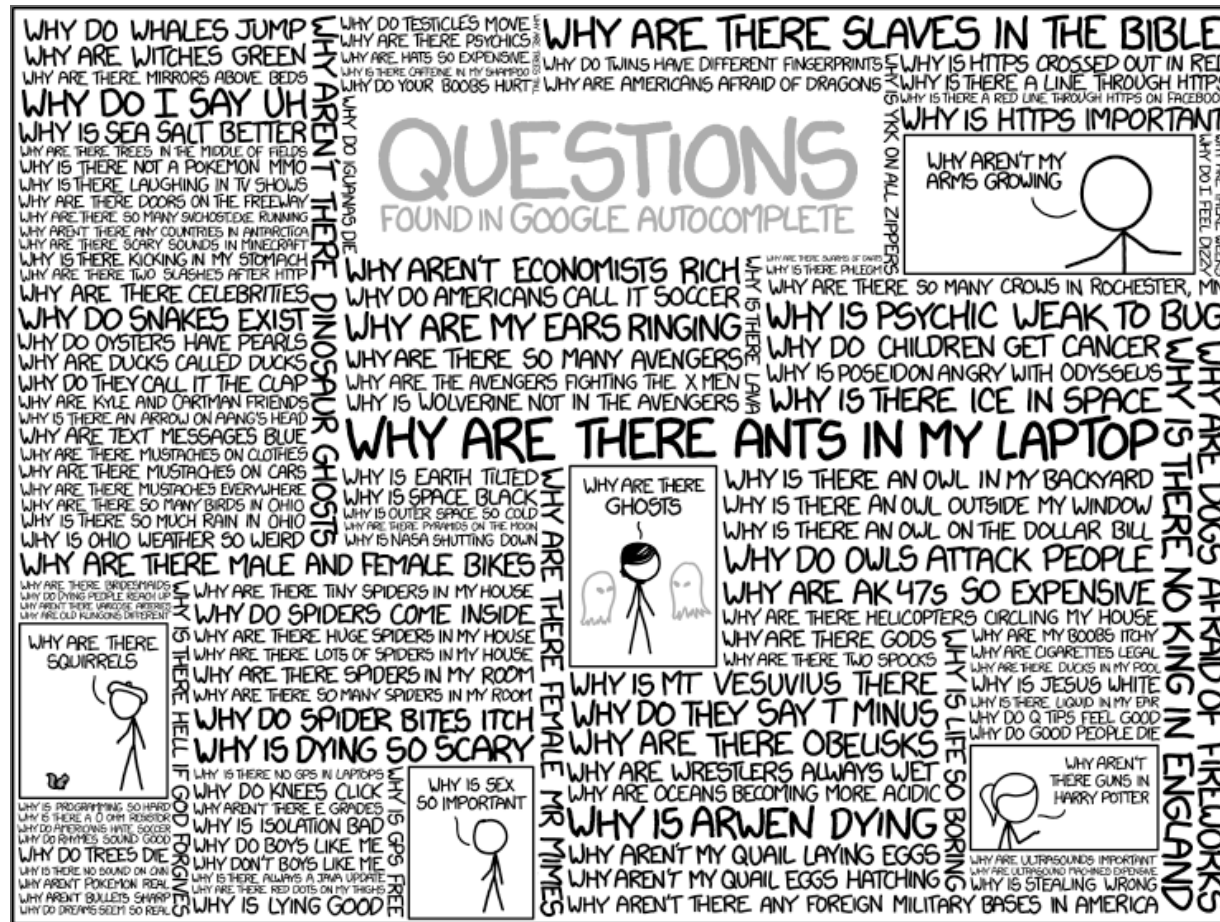


BRANCH AND BOUND EXAMPLE





Q & A



XKCD

Speaker notes

