

Optimization

CS5491: Artificial Intelligence
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Content Credits: Prof. Wei's CS4486 Course
and Prof. Boddeti's AI Course

TODAY

Optimization Setup

OPTIMIZATION PROBLEM: DEFINITION

Optimization Problem: Determine value of optimization variable within feasible region/set to optimize optimization objective

$$\begin{aligned} & \min_{\mathbf{x}} f(\mathbf{x}) \\ & \text{s.t. } \mathbf{x} \in \mathcal{F} \end{aligned} \tag{1}$$

- › Optimization variable $\mathbf{x} \in \mathbb{R}^n$
- › Feasible region/set $\mathcal{F} \subseteq \mathbb{R}^n$
- › Optimization objective: $f : \mathcal{F} \rightarrow \mathbb{R}$

Optimal solution: $\mathbf{x}^* = \operatorname{argmin}_{\mathbf{x} \in \mathcal{F}} f(\mathbf{x})$

Optimal objective value $f^* = \min_{\mathbf{x} \in \mathcal{F}} f(\mathbf{x}) = f(\mathbf{x}^*)$

OPTIMIZATION PROBLEM: DEFINITION

$$\begin{aligned} & \min_{\mathbf{x}} f(\mathbf{x}) \\ & \text{s.t. } \mathbf{x} \in \mathcal{F} \end{aligned} \tag{2}$$

Optimization variable $\mathbf{x} \in \mathbb{R}^n$

- › Discrete variables: Combinatorial optimization
- › Continuous variables: Continuous optimization
- › Mixed: Some variables are discrete, and some are continuous

OPTIMIZATION PROBLEM: DEFINITION

$$\begin{aligned} & \min_{\mathbf{x}} f(\mathbf{x}) \\ & \text{s.t. } \mathbf{x} \in \mathcal{F} \end{aligned} \tag{3}$$

Feasible region/set $\mathcal{F} \subseteq \mathbb{R}^n$

- › Unconstrained optimization: $\mathcal{F} = \mathbb{R}^n$
- › Constrained optimization: $\mathcal{F} \subset \mathbb{R}^n$
- › Finding a feasible point $\mathbf{x} \in \mathcal{F}$ can already be difficult

OPTIMIZATION PROBLEM: DEFINITION

$$\begin{aligned} & \min_{\mathbf{x}} f(\mathbf{x}) \\ & \text{s.t. } \mathbf{x} \in \mathcal{F} \end{aligned} \tag{4}$$

Optimization objective $f : \mathcal{F} \rightarrow \mathbb{R}$

- › $f(\mathbf{x}) = 1$: Feasibility problem
- › Simple functions
- › Linear function $f(\mathbf{x}) = \mathbf{a}^T \mathbf{x}$
- › Convex function (next lecture)
- › Complicated functions

- Can be implicitly represented through an algorithm which takes $\mathbf{x} \in \mathcal{F}$ as input, and outputs a value

OPTIMIZATION PROBLEM: DEFINITION

$$\begin{aligned} & \min_{\mathbf{x}} f(\mathbf{x}) \\ & \text{s.t. } \mathbf{x} \in \mathcal{F} \end{aligned} \tag{5}$$

Minimization can be converted to maximization (and vice versa)

$$\begin{aligned} & \max_{\mathbf{x}} g(\mathbf{x}) = -f(\mathbf{x}) \\ & \text{s.t. } \mathbf{x} \in \mathcal{F} \end{aligned} \tag{6}$$

Same optimal solution and optimal objective value $g^* = -f^*$

OPTIMIZATION PROBLEM: EXAMPLE

Example: Traveling Salesman Problem (TSP)

- › Problem: n cities, distance from city i to city j is $d(i, j)$, find a tour (a closed path that visits every city exactly once) with minimal total distance.
- › Variable \mathbf{x} : ordered list of cities being visited
- › x_i is the index of the i -th city being visited
- › Feasible set $\mathcal{F} = \{\mathbf{x} : \text{"each city visited exactly once"}\}$

$$\mathcal{F} = \{\mathbf{x} : \mathbf{x} \in \{1, \dots, n\}^n; \sum_k \mathbb{I}(x_k = i) = 1, \forall i \in \{1, \dots, n\}\} \quad (7)$$

- › Objective function $f(\mathbf{x}) = \text{"total distance when following } \mathbf{x}\text{"}$

$$f(\mathbf{x}) = d(x_n, x_1) + \sum_{k=1}^{n-1} d(x_k, x_{k+1}) \quad (8)$$

OPTIMIZATION PROBLEM: EXAMPLE

$$\begin{aligned} & \min_{\mathbf{x}} f(\mathbf{x}) \\ & \text{s.t. } \mathbf{x} \in \mathcal{F} \end{aligned}$$

Example: 8-Queens Problem (Solution 1)

- › Variable \mathbf{x} : location of the queen in each column
- › x_i is the row index of the queen in i -th column
- › Feasible set $\mathcal{F} = \{\mathbf{x} : \text{"no queens in the row, col, diag"}\}$

$$\mathcal{F} = \{\mathbf{x} : \mathbf{x} \in \{1, \dots, 8\}^8; x_i \neq x_j, |x_i - x_j| \neq |i - j|, \forall i, j \in \{1, \dots, 8\}\}$$

- › Objective function $f(\mathbf{x}) = 1$ (dummy)

OPTIMIZATION PROBLEM: EXAMPLE

$$\begin{aligned} & \min_{\mathbf{x}} f(\mathbf{x}) \\ & \text{s.t. } \mathbf{x} \in \mathcal{F} \end{aligned}$$

Example: 8-Queens Problem (Solution 2)

- Variable x_i, y_i : index of row and column of the i -th queen
- Feasible set $F = \{x, y : \text{"no queens in the row, col, diag"}\}$

$$\begin{aligned} F = \{x, y : x, y \in \{1, \dots, 8\}^8; \sum_i \mathbb{I}(x_i = k) = 1, \forall k \in \{1, \dots, 8\}; \\ \sum_i \mathbb{I}(y_i = k) = 1, \forall k \in \{1, \dots, 8\}; |x_i - x_j| \neq |y_i - y_j|, \forall i, j \in \{1, \dots, 8\} \} \end{aligned}$$

- Objective function $f(x) = 1$ (dummy)

OPTIMIZATION PROBLEM: EXAMPLE

x_i	1.0	2.0	3.0
y_i	2.1	3.98	7.0

Example: Linear Regression

- › Problem: Find a such that $y_i \approx ax_i, \forall i = \{1, 2, 3\}$
- › Variable a
- › Feasible region \mathbb{R}
- › Objective function $f(a)$?

$$\min_a \sum_{i=1}^3 |y_i - ax_i| \quad (9)$$

s.t. $a \in \mathbb{R}$

$$\min_a \sum_{i=1}^3 (y_i - ax_i)^2 \quad (10)$$

s.t. $a \in \mathbb{R}$

OPTIMIZATION PROBLEM: HOW TO SOLVE?

No general way to solve

Many algorithms developed for special classes of optimization problems (i.e., when $f(x)$ and \mathcal{F} satisfy certain constraints)

- › Convex optimization problem (CO)
- › Linear Program (LP)
- › (Mixed) Integer Linear Program (MILP)
- › Quadratic program (QP), (Mixed) Integer Quadratic program (MIQP), Semidefinite program (SDP), Second-order cone program (SOCP), . . .

Existing solvers and code packages for these problems

- › Cplex (LP, MILP, QP), Gurobi (LP, MILP, MIQP), GLPK (LP, MILP), Cvxopt (CO), DSDP5 (SDP), MOSEK (QP, SOCP), Yalmip (SDP), . . .

OPTIMIZATION PROBLEM: WHY USEFUL?

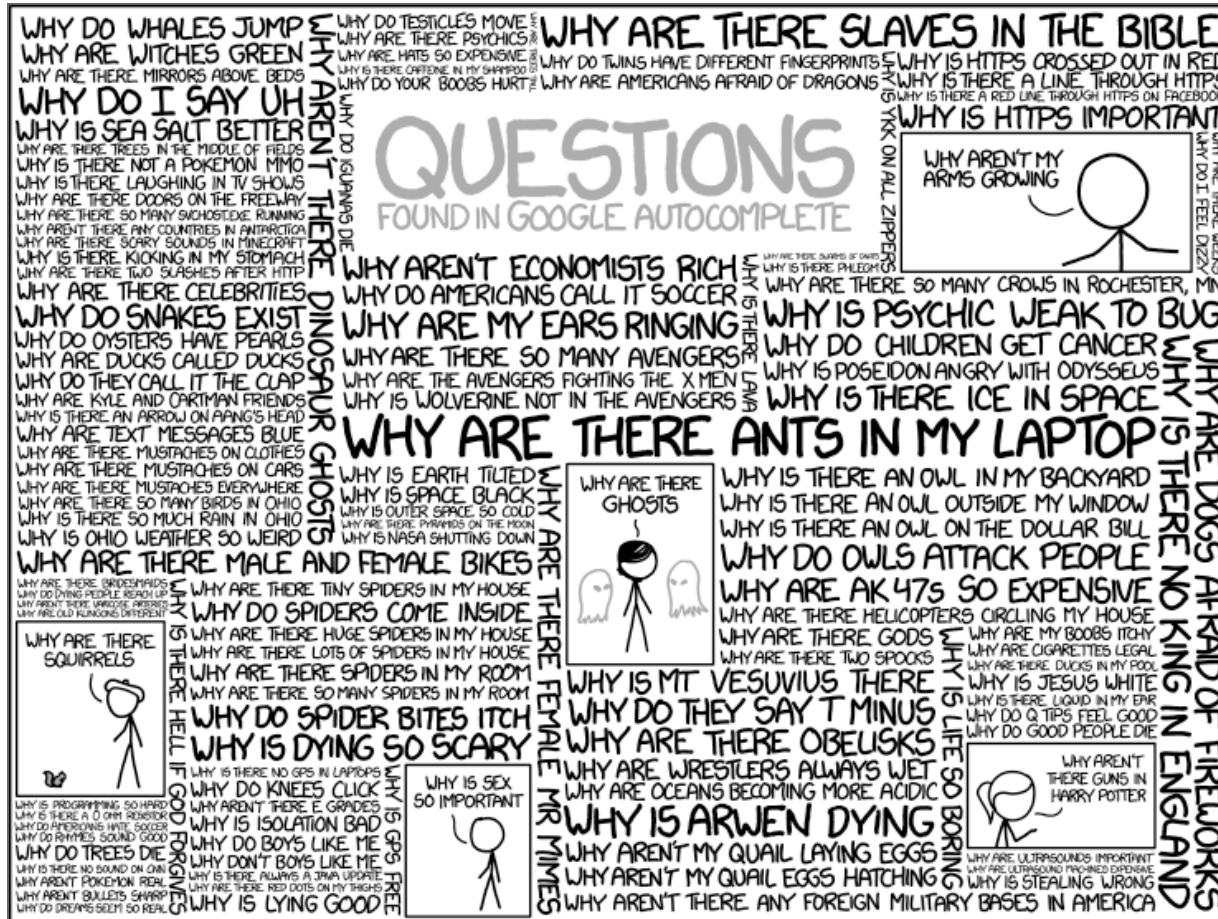
Why formulate problems as optimization problems?

- › For many class of optimization problems, algorithms or algorithmic frameworks have been developed
- › Decouple “representation” and “problem solving”

Lazy mode

- › Formulate a problem as an optimization problem
- › Identify which class the formulation belongs to
- › Call the corresponding solver
- › Done !!

Q & A



XKCD

Speaker notes

