

## Solutions to CS511 Homework 01

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**Exercise 1** Go to page 9 in Lecture Slides 06. Your task is to carefully write all the details of the proof by structural induction. These details are not included in the slides.

### Solution

#### Concise Proof by Structural Induction

**Proposition:**  $\forall s, t \in A^*, \text{reverse}(s \cdot t) = \text{reverse}(t) \cdot \text{reverse}(s)$

Let  $P(t) := \forall s \in A^*, \text{reverse}(s \cdot t) = \text{reverse}(t) \cdot \text{reverse}(s)$

**Proof by structural induction on  $t$ :**

1. Base case:  $t = \varepsilon$

$$\forall s \in A^*, \text{reverse}(s \cdot \varepsilon) = \text{reverse}(s) = \varepsilon \cdot \text{reverse}(s) = \text{reverse}(\varepsilon) \cdot \text{reverse}(s)$$

2. Inductive step: Assume  $P(t)$  holds for  $t \in A^*$ . Show  $P(a \cdot t)$  for  $a \in A$ .

$\forall s \in A^*$ :

$$\begin{aligned} \text{reverse}(s \cdot (a \cdot t)) &= \text{reverse}((s \cdot a) \cdot t) && [\text{associativity}] \\ &= \text{reverse}(t) \cdot \text{reverse}(s \cdot a) && [\text{I.H.}] \\ &= \text{reverse}(t) \cdot (\text{reverse}(a) \cdot \text{reverse}(s)) && [\text{def. of reverse}] \\ &= (\text{reverse}(t) \cdot \text{reverse}(a)) \cdot \text{reverse}(s) && [\text{associativity}] \\ &= \text{reverse}(a \cdot t) \cdot \text{reverse}(s) && [\text{def. of reverse}] \end{aligned}$$

By structural induction,  $P(t)$  holds  $\forall t \in A^*$ , proving the proposition.

**Exercise 2 [LCS, page 87]: Exercise 1.4.15. Hint: You may find it helpful to review pages 20 and 21 in Lecture Slides 02.**

## Solution

### Concise Proof by Mathematical Induction

**Theorem:** For  $n \geq 1$ ,

$$((\varphi_1 \wedge (\varphi_2 \wedge (\cdots \wedge \varphi_n) \cdots) \rightarrow \psi) \rightarrow (\varphi_1 \rightarrow (\varphi_2 \rightarrow (\cdots (\varphi_n \rightarrow \psi) \cdots))))$$

Let  $P(n)$  denote the theorem statement.

**Proof:**

1. Base case ( $n = 1$ ):

$$P(1) : ((\varphi_1 \rightarrow \psi) \rightarrow (\varphi_1 \rightarrow \psi)) \text{ [Trivially true]}$$

2. Inductive step: Assume  $P(k)$  holds for some  $k \geq 1$ . To prove  $P(k+1)$ :

LHS of  $P(k+1)$ :

$$\begin{aligned} & (\varphi_1 \wedge (\varphi_2 \wedge (\cdots \wedge \varphi_{k+1}) \cdots) \rightarrow \psi) \\ & \equiv ((\varphi_1 \wedge (\varphi_2 \wedge (\cdots \wedge \varphi_k) \cdots)) \wedge \varphi_{k+1} \rightarrow \psi) \\ & \equiv (\varphi_1 \wedge (\varphi_2 \wedge (\cdots \wedge \varphi_k) \cdots) \rightarrow (\varphi_{k+1} \rightarrow \psi)) \quad \text{[Deduction theorem]} \end{aligned}$$

Applying  $P(k)$  to this:

$$(\varphi_1 \rightarrow (\varphi_2 \rightarrow (\cdots (\varphi_k \rightarrow (\varphi_{k+1} \rightarrow \psi)) \cdots)))$$

This is the RHS of  $P(k+1)$ .

Therefore, by mathematical induction,  $P(n)$  holds for all  $n \geq 1$ .

**PROBLEM 1** Show that any of the three rules (LEM),(PBC),( $\neg\neg$ E) are interderivable.

## Solution

### Interderivability of LEM, PBC, and $\neg\neg$ E

We will show that the three rules Law of Excluded Middle (LEM), Proof by Contradiction (PBC), and Double Negation Elimination ( $\neg\neg$ E) are interderivable.

**(a) (PBC) is derivable from ( $\neg\neg$ E)**

1. $\neg\varphi \rightarrow \perp$	given
2. $\neg\varphi$	assumption
3. $\perp$	$\rightarrow$ E 1, 2
4. $\neg\neg\varphi$	$\neg$ I 2-3
5. $\varphi$	$\neg\neg$ E 4

**(b) (LEM) is derivable from (PBC)**

1. $\neg(\varphi \vee \neg\varphi)$	assumption
2. $\varphi$	assumption
3. $\varphi \vee \neg\varphi$	$\vee\text{I } 2$
4. $\perp$	$\neg\text{E } 1, 3$
5. $\neg\varphi$	$\neg\text{I } 2-4$
6. $\varphi \vee \neg\varphi$	$\vee\text{I } 5$
7. $\perp$	$\neg\text{E } 1, 6$
8. $\varphi \vee \neg\varphi$	PBC 1-7

**(c)  $(\neg\neg\text{E})$  is derivable from (LEM)**

1.	$\neg\neg\varphi$	premise
2.	$\varphi \vee \neg\varphi$	LEM
3.	$\varphi$	assumption
4.	$\varphi$	reiteration 3
5.	$\neg\varphi$	assumption
6.	$\perp$	$\neg\text{E}$ 1, 5
7.	$\varphi$	$\perp\text{E}$ 6
8.	$\varphi$	$\vee\text{E}$ 2, 3-4, 5-7

Therefore, we have shown that  $(\neg\neg\text{E}) \Rightarrow (\text{PBC}) \Rightarrow (\text{LEM}) \Rightarrow (\neg\neg\text{E})$ , proving that these three rules are interderivable.

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**Solutions in one file at:** [https://github.com/nich-ikech/CS511-hw-macbeth/blob/main/cs511HwSolutions/hw02/hw02\\_nicholas\\_ikechukwu.lean](https://github.com/nich-ikech/CS511-hw-macbeth/blob/main/cs511HwSolutions/hw02/hw02_nicholas_ikechukwu.lean)

**Exercise 3** For each of the three examples in the following three sections of Macbeth's book, your task is to remove 'sorry' and insert appropriate Lean 4 tactics

### **Solution**

[https://github.com/nich-ikech/CS511-hw-macbeth/blob/main/cs511HwSolutions/hw02/hw02\\_nicholas\\_ikechukwu.lean](https://github.com/nich-ikech/CS511-hw-macbeth/blob/main/cs511HwSolutions/hw02/hw02_nicholas_ikechukwu.lean)

**Exercise 4** For each of the three examples in the following three sections of Macbeth's book, your task is to remove 'sorry' and insert appropriate Lean 4 tactics.

### **Solution**

[https://github.com/nich-ikech/CS511-hw-macbeth/blob/main/cs511HwSolutions/hw02/hw02\\_nicholas\\_ikechukwu.lean](https://github.com/nich-ikech/CS511-hw-macbeth/blob/main/cs511HwSolutions/hw02/hw02_nicholas_ikechukwu.lean)



**PROBLEM 2** For each of the three examples in the following three sections of Macbeth's book, your task is to remove 'sorry' and insert appropriate Lean 4 tactics

### **Solution**

[https://github.com/nich-ikech/CS511-hw-macbeth/blob/main/cs511HwSolutions/hw02/hw02\\_nicholas\\_ikechukwu.lean](https://github.com/nich-ikech/CS511-hw-macbeth/blob/main/cs511HwSolutions/hw02/hw02_nicholas_ikechukwu.lean)