

ISIT312 Big Data Management

HBase Data Model

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HBase Data Model

Outline

Background

Logical view of data

Design fundamentals

Physical implementation

Background

Hbase is open source distributed database based on a data model of Google's BigTable

HBase provides a BigTable view of data stored in HDFS

HBase is also called as Hadoop DataBase

HBase still provides a tabular view of data however it is also very different from the traditional relational data model

HBase data model is a sparse, distributed, persistent multidimensional sorted map

It is indexed by a row key, column key, and timestamp

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Logical view of data

HBase organizes data into **tables**

HBase **table** consists of **rows**

Each **row** is uniquely identified by a **row key**

Data within a **row** is grouped by a **column family**

Column families have an important impact on the **physical implementation** of HBase **table**

Every **row** has the same **column families** although some of them can be **empty**

Data within a **column family** is addressed via its **column qualifier**, or simply, **column name**

Hence, a combination of **row key**, **column family**, and **column qualifier** uniquely identifies a **cell**

Values in cells do not have a data type and are always treated as

[TOP](#) **sequences of bytes**

Logical view of data

Each cell has multiple versions,
typically represented by the timestamp
of when they were inserted into the table

Timestamp1 Timestamp2

The table is lexicographically sorted on the row keys

Row Key	Column Family - Personal			Column Family - Office	
	Name	Residence	Phone	Phone	Address
00001	John	415-111-1234	415-212-5544	1021 Market St	
00002	Paul	408-432-9922	415-212-5544	1021 Market St	
00003	Ron	415-993-2124	415-212-5544	1021 Market St	
00004	Rob	818-243-9988	408-998-4322	4455 Bird Ave	
00005	Carly	206-221-9123	408-998-4325	4455 Bird Ave	
00006	Scott	818-231-2566	650-443-2211	543 Dale Ave	

Cells

Values within a **cell** have **multiple versions**

Versions are identified by their **version number**, which by default is a **timestamp** when the cell was written

Logical view of data

If a timestamp is not determined at write time, then the current timestamp is used

If a timestamp is not determined during a read, the latest one is returned

The maximum allowed number of cell value versions is determined for each column family

The default number of cell versions is three

Logical view of data

A view of HBase table as a nested structure

```
{"Row-0001":  
    {"Home":  
        {"Name":  
            {"timestamp-1": "James"}  
        "Phones":  
            {"timestamp-1": "2 42 214339"  
             "timestamp-2": "2 42 213456"  
             "timestamp-3": "+61 2 4567890"}  
        }  
    "Office":  
        {"Phone":  
            {"timestamp-4": "+64 345678"}  
        "Address":  
            {"timestamp-5": "10 Ellenborough Pl"}  
        }  
    }  
}
```

HBase Table

Logical view of data

A view of HBase table as a nested structure

```
{"Row-0002":  
    {"Home":  
        {"Name":  
            {"timestamp-6": "Harry"}  
        "Phones":  
            {"timestamp-7": "2 42 214234"}  
        }  
    "Office":  
        {"Phone":  
            {"timepstamp-8": "+64 345678"}  
        "Address":  
            {"timestamp-9": "10 Bong Bong Rd"  
             "timestamp-10": "23 Victoria Rd"}  
        }  
    }  
}
```

HBase Table

Logical view of data

A **key** can be **row key** or a combination of a **row key**, **column family**, **qualifier**, and **timestamp** depending on what supposed to be retrieved

If all the **cells** in a row are of interest then a **key** is a **row key**

If only specific **cells** are of interest, the appropriate **column families** and **qualifiers** are a part of a **key**

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Design Fundamentals

When designing **Hbase table** we have to consider the following questions:

- What should be a **row key** and what should it contain?
- How many **column families** should a **table** have?
- What **columns (qualifiers)** should be included in each **column family**?
- What information should go into the **cells**?
- How many **versions** should be stored for each **cell**?

In fact **HBase table** is a four level **hierarchical structure** where a **table** consists of **rows**, **rows** consists of **column families**, **column families** consist of **columns** and **columns** consists of **versions**

If cells contain the keys then **HBase table** becomes a **network/graph structure**

Design Fundamentals

Important facts to remember:

- Indexing is performed only for a row key
- Hbase tables are sorted based on a row key
- Everything in Hbase tables is stored as untyped sequence of bytes
- Atomicity is guaranteed only at a row level and there are no multi-row transactions
- Column families must be defined at Hbase table creation time
- Column qualifiers are dynamic and can be defined at write time
- Column qualifiers are stored as sequences of bytes such that they can represent data

Design Fundamentals

Implementation of Entity type

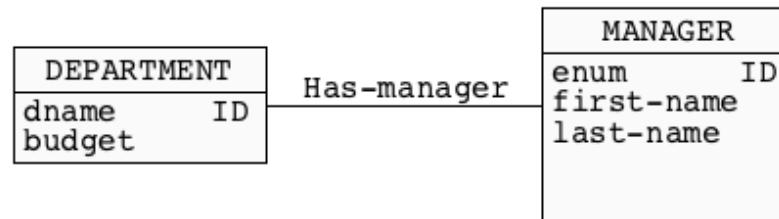
CUSTOMER	
cnumber	ID
first-name	
last-name	
phone	
email	

```
{"007":  
  { "CUSTOMER":  
      {"first-name": {"timestamp-1": "James"},  
       "last-name": {"timestamp-2": "Bond"},  
       "phone": {"timestamp-1": "007-007"},  
       "email": {"timestamp-1": "jb@mi6.com"}  
    }  
}
```

HBase Table

Design Fundamentals

Implementation of one-to-one relationship



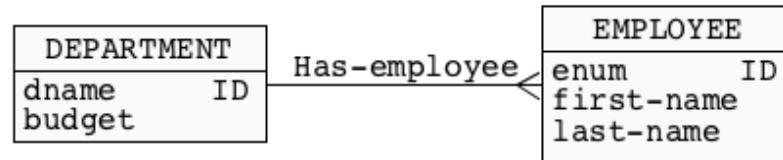
```
{"Sales":  
  {"DEPARTMENT":  
    {"dname": {"timestamp-1": "Sales"},  
     "budget": {"timestamp-1": "1000"}  
    }  
  },  
  {"MANAGER":  
    {"enumber": {"timestamp-2": "007"},  
     "first-name": {"timestamp-3": "James"},  
     "last-name": {"timestamp-4": "Bond"}  
    }  
  }  
}
```

HBase Table

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Design Fundamentals

Implementation of one-to-many relationship

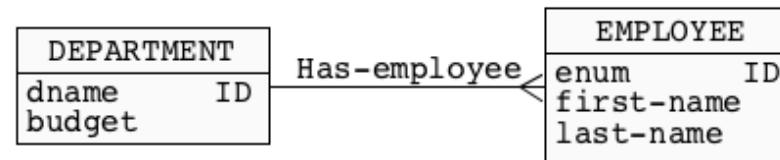


```
{"007":  
  { "EMPLOYEE":  
      {"enum": {"timestamp-1": "007"},  
       "first-name": {"timestamp-2": "James"},  
       "last-name": {"timestamp-3": "Bond"},  
       "department": {"timestamp-4": "Sales"}  
    }  
}
```

HBase Table

Design Fundamentals

Another implementation of one-to-many relationship

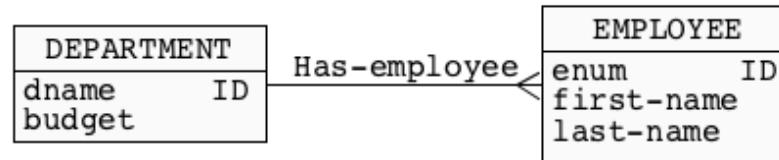


```
{"Sales":  
  { "DEPARTMENT":  
      { "dname": {"timestamp-1": "Sales"},  
       "budget": {"timestamp-1": "1234567"}  
     }  
  { "EMPLOYEE":  
      {"007": {"timestamp-2": "James Bond"},  
       "008": {"timestamp-3": "Harry Potter"},  
       "009": {"timestamp-4": "Robin Hood"}  
      ...  
    }  
}
```

HBase Table

Design Fundamentals

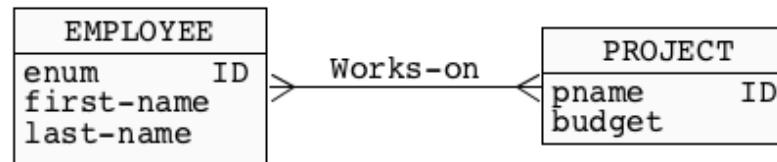
Yet another implementation of one-to-many relationship



```
{"Sales":  
  {"DEPARTMENT":  
    {"dname": {"timestamp-1": "Sales"},  
     "budget": {"timestamp-1": "1000"}  
    },  
   {"HAS-EMPLOYEES":  
     {"employees": {"timestamp-2": "007 James Bond"},  
      {"timestamp-3": "008 Harry Potter"},  
      {"timestamp-4": "009 Robin Hood"}  
     ...  
    }  
  }  
}
```

Design Fundamentals

Implementation of many-to-many relationship

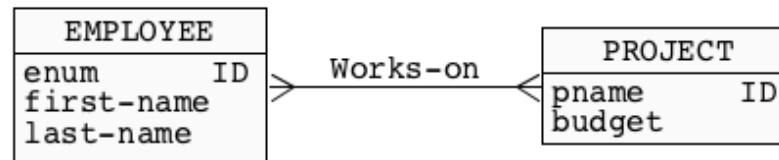


```
{"participation-1":  
    {"EMPLOYEE":  
        {"enum": {"timestamp-1": "007"},  
         "first-name": {"timestamp-2": "James"},  
         "last-name": {"timestamp-3": "Bond"},  
         "pnumber": {"timestamp-4": "project-1"},  
         {"timestamp-5": "project-2"}  
        ...  
    }  
}
```

HBase Table

Design Fundamentals

Another implementation of many-to-many relationship

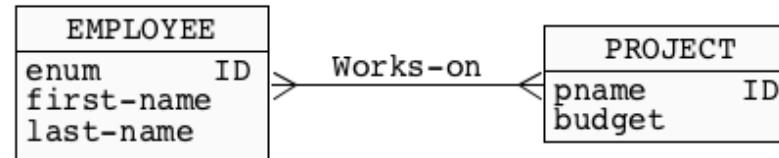


```
{"participation-1":  
    {"PROJECT":  
        {"pnumber": {"timestamp-1": "project-1"},  
         "budget": {"timestamp-2": "12345.25"},  
         "employee": {"timestamp-3": "007"},  
                   {"timestamp-4": "008"},  
                   {"timestamp-5": "009"},  
         "..."},  
        }},  
    }  
}
```

HBase Table

Design Fundamentals

Another implementation of many-to-many relationship



```
{"participation-1":  
    {"PARTICIPATION":  
        {"pnumber": {"timestamp-1": "project-1"},  
         "employee": {"timestamp-2": "employee-007"}  
    }  
}  
  
{"participation-2":  
    {"PARTICIPATION":  
        {"pnumber": {"timestamp-1": "project-1"},  
         "employee": {"timestamp-2": "employee-008"}  
    }  
}}
```

HBase Table

HBase Table

Design Fundamentals

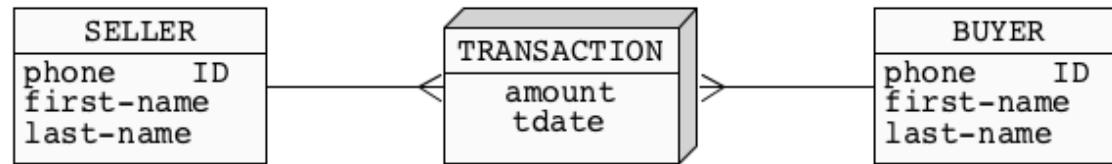
Note, that it is possible to group in one **Hbase table** rows of different types

```
{ "employee-007":  
    { "EMPLOYEE":  
        { "enumber": { "timestamp-1": "007" },  
          "first-name": { "timestamp-2": "James" },  
          "last-name": { "timestamp-3": "Bond" }  
        }  
    }  
,  
{ "project-1":  
    { "PROJECT":  
        { "pnumber": { "timestamp-4": "1" },  
          "budget": { "timestamp-5": "12345.25" }  
        }  
    }  
,  
{ "participation-2":  
    { "PARTICIPATION":  
        { "pnumber": { "timestamp-1": "project-1" },  
          "employee": { "timestamp-2": "employee-007" }  
        }  
    }  
}
```

HBase Table

Design Fundamentals

Implementation of fact with dimensions

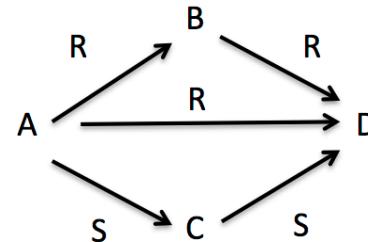


```
{"1234567":  
    {"MEASURE":  
        {"amount": {"timestamp-1": "1000000"}  
     },  
    "BUYER":  
        {"phone": {"timestamp-1": "242214339"},  
         "first-name": {"timestamp-1": "James"},  
         "last-name": {"timestamp-1": "Bond"}  
     },  
    "SELLER":  
        {"phone": {"timestamp-1": "242215612"},  
         "first-name": {"timestamp-1": "Harry"},  
         "last-name": {"timestamp-1": "potter"}  
     }  
}
```

HBase Table

Design Fundamentals

Implementation of graph structure



```
{"A":  
  {"R":  
    {"1": {"timestamp-1": "B"},  
     "2": {"timestamp-1": "D"}  
    },  
  "S":  
    {"1": {"timestamp-1": "C"}  
  }  
}  
}
```

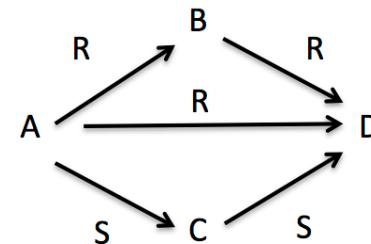
HBase Table

```
{"B":  
  {"R":  
    {"1": {"timestamp-1": "D"}  
  }  
}  
}
```

HBase Table

Design Fundamentals

Implementation of graph structure



```
{"C":  
  {"S":  
    {"1": {"timestamp-1": "D"}  
    }  
  }  
}  
  
{ "D":  
}
```

HBase Table

HBase Table

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HBase is a database built on top of HDFS

HBase tables can scale up to billions of rows and millions of columns

Because Hbase tables can grow up to terabytes or even petabytes, Hbase tables are split into smaller chunks of data that are distributed across multiple servers

Chunks of data are called as regions and servers that host regions are called as region servers

Region servers are usually collocated with data nodes of HDFS

The splits of Hbase tables are usually horizontal, however, it is also possible to benefit from vertical splits separating column families

Region assignments happen when Hbase table grows in size or when a region server is malfunctioning or when a new region server is added

References

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