

# (1) Introduction

## 3 Motivating Examples:

### 1) Scheduling in CPU.

We want to be able to schedule the various processes which are running concurrently on a computer i.e. Allocate chunks of CPU time towards each process.

We can have multiple objectives:

- Minimize average waiting time, or
- Maximize throughput.

Our tasks i.e. The processes have multiple properties:

- Processing time - Time needed to finish the proc
- Priorities - Weights indicating importance

Further,

- Tasks need not be known beforehand; in fact, neither does the processing time

(These are known as real-time problems, we will mainly focus on problems with known quantities)

- Tasks need not run in contiguous chunks. They can be interrupted, and resumed at a later time. (This is known as preemption)

→ Ex. 2, 3

## 2) Production Planning.

The Cartz-er-U factory produces 3 types of carts:

- Shopping Carts
- Airport Trolleys
- Cargo Carts

Production of each type of cart goes through the same stages

- Producing metal skeleton
- Coating Metal
- Attaching parts to the skeleton
- Packaging of carts

Each stage consists of the same machines, not necessarily identical (i.e. Could differ in speed or other capabilities)

We could have multiple objectives:

- Minimize # of late demands

(The "demand" for each cart is the # required and due date)

- Minimize the "total penalty" of late demands

(i.e. Some demands may not be met, but we want to minimize lateness)

- Minimize maximum lateness of demands.

## 3) Gate Assignments at an airport

An airport has  $m$  gates where planes land and take off. We assume all gates are identical, and there is a hanger where planes can ~~wait~~ wait.

Suppose plane  $j$  lands at time  $r_j$ , departs at time  $d_j$  and takes time  $\tau_{j,i,i'}$  to go from gate  $i$  to  $i'$ . How should planes be assigned to gates to minimize total airline delay?



## 1.1 Def'n (Job, Machine)

Every scheduling problem consists of:

- Jobs - Entities that require resources. We will commonly use  $j, k, l$  to denote jobs.
- Machines (M/C) - Resources that jobs require. We will use  $i, i', i'', \text{etc.}$  to denote M/Cs.

### - Jobs:

A job  $j$  typically is associated with:

#### - Processing time ( $P_{ij}$ )

This is the time taken by machine  $i$  to process job  $j$ .

Note:

- 1) Always  $\geq 0$
- 2) Can be infinity (i.e. This machine cannot process job  $j$ )
- 3) If machines are identical, then  $j$ 's processing time is independent of  $i$ , and we will use  $P_j$  to denote  $j$ 's processing time.

#### - Release Date ( $r_j$ )

Time when  $j$  enters the system.

Note:  $j$  can only be scheduled on a machine at or after its release date

#### - Due Date ( $d_j$ )

Time by which  $j$  must be completed

Note: Could be a soft deadline that can incur penalty in our objective fn., or hard deadline

#### - Weight ( $w_j$ )

Indicates the importance / priority of job  $j$ .

$w_j \geq 0$  unless otherwise stated.

← Machines  
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## - Machines:

We can have multiple settings/environments:

- Single M/C
- Multiple M/C
  - Identical machines
  - Non-identical machines
- Shop Scheduling: Multiple machines and a job consists of various operations that need to be performed on various machines (Could be a sequence - See Ex. 2)

## - Objective

1) Sum of completion times:

Let  $C_j$  denote the completion time of job  $j$ . We write:

$$\min \sum_j C_j$$

2) Weighted Completion time:

$$\min \sum_j w_j C_j$$

3) Total Waiting time:

$$\min \sum (C_j - r_j)$$

4) Maximum lateness:

Let  $L_j := C_j - d_j$  be the lateness of  $j$ .

$$\min \max_j L_j$$

5) No. of late jobs:

$$\text{Define } u_j := \begin{cases} 1 & C_j > d_j \\ 0 & \text{otherwise} \end{cases}$$

$$\min \sum u_j$$