

Def'n 1: A k-colouring of a graph G is a partition of $V(G)$ into at most k independent sets.

Def'n 2: A k-colouring of a graph G is a map $f: V(G) \rightarrow [k]$ such that $\forall e = uv \in E(G)$, $f(u) \neq f(v)$

Def'n 3: A k-colouring of a graph G is a graph homomorphism to K_k .

Takeaway: There are multiple ways to view what a colouring is.

Weakenings, Generalizations, and Variants of Colouring:

Variants: ("Changing what you color")

- Edge Colouring: A k-edge-colouring of a graph G is

Def 1: A partition of $E(G)$ into at most k matchings

2) A k-colouring of $L(G)$ (the line graph of G)

"Changing what you color"

- Total Colouring: A k-total-colouring of a graph G is a map

$f: V(G) \cup E(G) \rightarrow [k]$ s.t.

- $f(u) \neq f(w) \quad \forall uw \in E(G)$
- $f(u) \neq f(e) \quad \forall u \sim e \in E(G)$
- $f(e) \neq f(e') \quad \forall e \sim e' \in E(G)$

Generalizations: ("Changing what you are allowed to color")

- List Colouring: (idea: Lists of available colors to vertices)

Def: A k-list-colouring k-list-assignment of a graph G is an assignment of lists $(L_v: v \in V(G))$ such that

$$|L(v)| \geq k \quad \forall v \in V(G)$$

An L-colouring of a graph G is a colouring ϕ of G such that $\phi(v) \in L(v) \quad \forall v \in V(G)$.

"~~Remember what you are allowed to colour!~~"

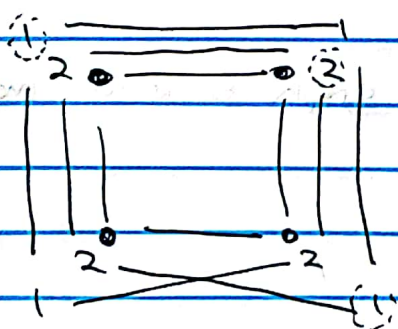
- Correspondence Colouring:

Def: A k-correspondence-assignment is a pair $((L(v) : v \in V(G)), (M(u) : u \in E(G)))$, where $M(u)$ is a matching from $L(u)$ to $L(v)$.

An (L, M) -colouring of G is a colouring ϕ of G such that:

- $\phi(v) \in L(v) \quad \forall v \in V(G)$
- $\phi(u)$ is not matched to $\phi(v)$ in $M(u) \quad \forall uv \in E(G)$

Ex: $G = C_4$



$$L(v) = \{1, 2\} \quad \forall v$$

Q: Is there an (L, M) -colouring of G for this L, M ?

No!

And we run into trouble here:

Remarks:

- We may as well assume $L(v) = \{1, \dots, |L(v)|\} \forall v \in V(G)$
- Correspondence has a "local notion of colour", while list colouring / ordinary colouring have a "global notion of colour"

Weakenings: ("Change what a ^{colour} graph is")

- No Restrictions: Improper colourings, i.e. mappings of graphs
- d -defective colouring: Each colour has maximum degree d
- C -clustered Colouring: Every monochromatic component has size (# of vertices) $\leq C$
- No mono. path of length $> L$
- Every colour is d -degenerate
- Every colour is triangle-free (or more generally, bounded clique #)

or Fractional Colouring: ("Changing how you colour")

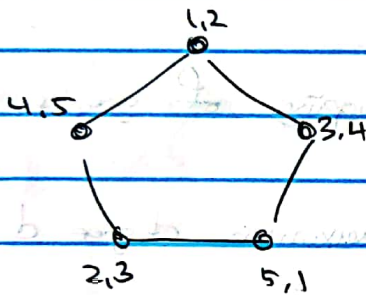
Def: An (a, b) -colouring of a graph G is a map ϕ such that $\phi(v)$ is a subset of $[a]$ of size b and $\forall uv \in E(G), \phi(u) \cap \phi(v) = \emptyset$

Remark: If G has a k -colouring, then G has a (kb, b) -colouring $\forall b$

(i.e. Graph homomorphism to known graph on a, b).

The fractional chromatic #': $\chi_f(G) = \inf \left\{ \frac{a}{b} : G \text{ has an } (a, b)\text{-colouring} \right\}$.

Ex. $\chi_f(C_5) = 5/2$



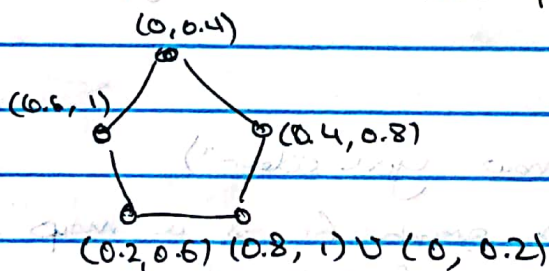
Remark:

- There is also an LP-formulation and its dual. (its dual asks about weighted independent set)
- An assignment of measurable subsets $\Phi(v)$ of $[0, 1] \times V(G)$ such that:

$$u \in \Phi(v) \cap \Phi(w) = \emptyset$$

$$\chi_f(G) = \frac{1}{\sup \left\{ \epsilon : G \text{ has a colouring as above, with } \mu(\Phi(v)) = \epsilon \text{ } \forall v \right\}}$$

Ex:



Let's call the above an f-colouring if $\mu(\phi(v)) = f(v)$ $\forall v$.

Proposition: $\chi_f(G) = k$ iff G has a $(\frac{1}{k})$ -colouring

Proposition: G has an f -colouring iff the vector $(f(v) : v \in V(G))$ is in the independent set polytope (i.e. \exists probability distribution on independent sets such that $\Pr[v \in I] = f(v)$)

Combinatorics:

- Defective Colouring \times
- List Defective \checkmark
- Prop. Defective \checkmark
- List Edge \checkmark
- Prop. Edge \checkmark
- Edge Total \times
- List Correspondence \times (Correspondence is already list).
- Fractional List \checkmark
 - Find b -colouring from a list-assignment. (multicolouring)

$$\chi_{f, \text{list}}(G) = \inf \left\{ \frac{a}{b} : \forall \text{ a-list-assign } G \text{ has a } b\text{-colouring} \right\}$$

Theorem: $\chi_f(G) = \chi_{f, \text{list}}(G)$

- Fractional Correspondence.