

Theorem (Bollobás, Thomason '96)

If G is $2k$ -connected, then G is k -linked.

(Strengthening of the crt: Thomason-Wolfe '07: $10k$ -connected $\Rightarrow k$ -linked)

Theorem (Kostochka, Thomason '80's)

If G has no k -minors, then $\overline{ad}(G) \leq O(k \log k)$
 $\overline{ad}(G)$ = average degree

Idea: If G has large avg. degree, then G contains a minor of large avg. degree and a small # of vertices.

Lemma:

If $e(G) \geq k \cdot |V(G)|$, then G contains a minor H s.t. $|V(H)| \leq 2k+1$ and $\delta(H) \geq k+1$.

Proof:

We may assume that $\delta(G) \geq k+1$, else we can delete a min. degree vertex and apply induction.

Similarly, we may assume that $e(G) = k \cdot |V(G)|$, else we can delete edges and apply induction.

Thus, \exists a vertex $v \in V(G)$ with $d(v) \leq 2k$ (and $d(v) \geq 1$).
 Let $H = N[v]$. If $\exists u \in N(v)$ s.t. $|N(u) \cap N(v)| \leq k-1$, then induct on $G - uv$. So $\forall u \in N(v)$, $|N(u) \cap N(v)| \geq k$ and hence $\delta(H) \geq k+1$ and $|V(H)| \leq 2k+1$, as desired.



In fact, something stronger is true!

Lemma (Thomason, '84)

Let $0 < \beta < 1$ be the root of the equation $1 = \beta(1 + \ln(\frac{2}{\beta}))$ (Note: $\beta \approx .37$) and let $k \geq 3$ be an integer. If G is a graph s.t. $e(G) \geq kv(G)$, then G contains a minor H s.t. $V(H) \leq k\beta + 2$ and $\delta(H) \geq v(H) + \lfloor \beta k \rfloor - 1$.

Why is this nice? The guarantee $\delta(H) \geq v(H) + \lfloor \beta k \rfloor - 1$ means we get min. degree at least ~~at~~ $\sim 60\%$ of the original degree. For min. degrees higher than 50% , we get a linear amount of common neighbors!

Proof (of Kosterlka-Thomason)

Idea: We will prove the contrapositive, if we have high average degree, then we can use the previous lemma to find small minors, and then we finish with probabilistic method.

We'll prove that if n is large enough and $ed(G) \geq c \sqrt{\log n}$, then G contains a k -minor. $c := 2k$.

Using the earlier lemma, if minor H of G s.t. $V(H) \leq 2k\beta + 2$ and $\delta(H) \geq k\beta$.

(Note: We will try to bag vertices up to use for k -libed theorems, without worrying about the connectivity of the bags)

Proof (con't)

Example: We want graph like:



And we'll worry about individual connectivity in a bit

Let $p = \frac{E(H)}{V(H)}$, the density.
Let X_1, \dots, X_t be vertex-disjoint subsets of H of size $\frac{c}{200} \log t$ chosen u.a.r.

$$\Pr[e(X_i, X_j) = 0] = (1-p)^{|X_i||X_j|} \\ = (1-p)^{\frac{c^2}{40000} \log^2 t}$$

and we claim this is $< t^{-2}$, and so:

$$\mathbb{E}[\sum_{i \neq j \in [t]} \mathbb{1}_{e(X_i, X_j) = 0}] < t^2 \cdot \frac{1}{t^2} < 1.$$

So, there exist a choice of X s s.t. $\sum_{i \neq j \in [t]} e(X_i, X_j) > 0$
 $\forall i \neq j \in [t]$.

Now, we're almost done, but each X_i may not be connected! (So, we may not be able to actually extract them)

→ i.e. (Thomason's lemma).

Finish #1: Use better lemma to guarantee that $|N(x_i) \cap N(x_j)| \geq .05K$ $\forall x_i \in V(H)$, then greedily pick common neighbors for each x_i outside of $\bigcup_{i=1}^t X_i$ to connect each x_i , avoiding previously picked common neighbors.

Proof (cont)

Finish #2: Using Mader, I subgraph H' of H with connectivity $\geq k/2$ and hence $\min \deg \geq k/2$ and so $p \geq 1/k$ for each H' and $V(H') \leq V(H) \leq 2k\epsilon$.

So, apply random partition (of X 's) to H' instead of H .

Now, by Bollobás, Thomason, H' is $k/2\epsilon$ -linked. Take a

matching from $\{X_i\}$ to a set of vertices $\{v_i\}$ in $G - U_i$.

(Such exists since $\delta(H') \geq k/2$ and $|U_i| \leq k\epsilon$). Now, use

the matched vertices ~~and~~ vertices of X as terminals to find vertex disjoint paths meeting each X_i .

□

Proof (Thomason's Lemma)

By the other lemma, there exists a minor H' of G s.t.

$V(H') \leq 2k\epsilon$ and $\delta(H') \geq k/2$. (We can drop the ϵ 's)

We define a \mathbb{R} -valued function $f(G)$ on graphs as

$$f(G) = \frac{BkV(G)}{2} \left(\log \left(\frac{V(G)}{Bk} \right) + 1 \right)$$

and let:

$$D := \{ G : V(G) \geq Bk \text{ and } e(G) > f(G) \}$$

Now

Claim: $H' \in D$

Proof: $V(H') \geq k > Bk$ and $e(H') \geq \frac{kV(H')}{2} \cdot 1$

$$\geq \frac{kV(H')}{2} \left(B \left(1 + \log \left(\frac{2}{B} \right) \right) \right)$$

□

Proof: (cont)

And since $v(H'') \leq v(H') \leq 2k$:

$$\log\left(\frac{v(H'')}{2k}\right) \leq \log\left(\frac{2}{\beta}\right) \text{ and hence}$$

$$v(H) \leq \beta k (\log\left(\frac{2}{\beta}\right) + 1) = k \quad (\text{Since } \beta(\log\left(\frac{2}{\beta}\right) + 1) = 1)$$

It remains to show that $\delta(H) \geq v(H) + \lfloor \beta k \rfloor - 1$.

Note:

$$\delta(H) \geq \min_{v \in H''(u)} |N_{H''}(u) \cap N_{H''}(v)|.$$

(Now, we have to show that the # of common neighbors is actually high)

By previous obs., we have that

$$e(H'') - e(H''/uv) \geq f(H'') - f(H''/uv)$$

So,

$$|N_{H''}(u) \cap N_{H''}(v)| \geq f(H'') - f(H''/uv) - 1$$

$\xrightarrow{\quad} \text{For } uv \text{ itself.}$

Then

$$\delta(H) - v(H)$$

$$\geq 2(f(H'') - f(H''/uv) - 1) - v(H)$$

$$\geq 2\left(\frac{\beta k v(H'')}{2} \left(\log\left(\frac{v(H'')}{\beta k}\right) + 1\right) - \frac{\beta k \frac{v(H''/uv)}{2}}{2} \left(\log\left(\frac{v(H''/uv)}{\beta k}\right) + 1\right) - 1\right)$$

$$= \beta k \left(v(H'') \left(\log\left(\frac{v(H'')}{\beta k}\right) + 1\right) - (v(H'') - 1) \left(\log\left(\frac{v(H'')}{\beta k}\right) + 1\right) - 2 - v(H) \right)$$

Proof (c.w.1)

$$= \beta^k \left(v(k'') \left(\log\left(\frac{v(k'')}{\beta^k}\right) - \log\left(\frac{v(k'')-1}{\beta^k}\right) + 1 + \log\left(\frac{v(k'')-1}{\beta^k}\right) \right) \right.$$

$$\left. - 2 - v(k) \right)$$

Recall that

$$v(k) \leq \beta^k \left(\log\left(\frac{v(k'')}{\beta^k}\right) + 1 \right)$$