

# CS487 - Symbolic Computation

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# Lecture 1: Introduction

## 1.1 Course Preview

**Example 1.1** (Simplify Rational Expressions). Suppose we have the two following expressions:

$$f := \frac{x+1}{x-1} - \frac{x^3 - 2x + x^2 + 2}{x^3 + 2x - x^2 - 2} + \frac{x^2 + 3}{x-1} \quad (1.1)$$

$$g := \frac{(x-1)^2 - x^2 - x + 2x}{(x+y+2)^{100}} \quad (1.2)$$

Question: How do we simplify these expressions to a single  $\frac{poly}{poly}$  or return that it is 0?

One idea: Define a “normal” function:

1. If expression is 0, the normal function will be 0
2. If not, the normal function will be the simplest form

(More) Questions: What else do we need to consider?

- How do we represent polynomials (i.e. What data structure do we use?)
- How do we perform polynomial operations computationally?
- Do we need to consider the size of the integers in our computations?

**Example 1.2** (Solving Recurrences). Suppose we have the recurrence:

$$T(n) = \begin{cases} 2T(\frac{n}{2}) + \frac{n}{2} & n > 1 \\ 1 & n = 1 \end{cases} \quad (1.3)$$

We can solve this by hand (using Master theorem or other techniques) to obtain the answer:

$$T(n) = n(1 + \log_2(n)) \quad (1.4)$$

Question: How do we do this computationally?

**Example 1.3.** Consider the following identities:

$$\sum_{k=0}^n k = \frac{n(n-1)}{2} \quad (1.5)$$

$$\sum_{k=0}^n k^4 = \frac{n(n-1)(2n-1)(3n^3-3n-1)}{30} \quad (1.6)$$

Question: Can we return a closed form (without involving the index  $k$ ) for any general expression or report that one doesn't exist?

## 1.2 Representation of Integers

Current computers are based on architecture with 64 bits (We will call this number of bits the word size)

**Example 1.4.** The unsigned long in C represents integers in exactly the range  $[0, 2^{64} - 1]$

Question: How do we represent larger numbers?

Idea: Use an array of word size numbers.

Any integer  $a$  can be expressed as the following summation:

$$a = (-1)^s \sum_{i=0}^n a_i 2^{64i} \quad (1.7)$$

where  $s \in \{0, 1\}$  represents the sign of  $a$  and  $0 \leq a_i \leq 2^{64} - 1$  are the individual elements in the array.

If we assume  $0 \leq n + 1 \leq 2^{63}$ , then we can encode  $a$  as an array:

$$[s \cdot 2^{63} + n + 1, a_0, a_1, \dots, a_n] \quad (1.8)$$

This is sufficient for all practical purposes.

**Note.** The length of  $a$  is given by:  $\lfloor \log_{2^{64}} |a| \rfloor + 1 \in \mathcal{O}(\log |a|)$  words

Addition of Integers:

Suppose our input is  $a : a_0 + a_1\beta + a_2\beta + \dots a_n\beta_n$  and  $b : b_0 + b_1\beta + b_2\beta + \dots b_m\beta_m$  (where  $m \leq n$ ).

Question: How large can each individual nu

## Lecture 2: