

* All p 's are probs.

Recall:

$k_0 = c \log n$, $k_1 = n^{2/3}$. Prove a.s.

(i) No components of order (k_0, k_1)

(ii) Only one component of order $\geq k_1$.

(iii) # vertices in small component $\sim n p^n$, where $p e^c = e^{p c}$ ($0 < p < 1$)

We will prove (iii) this class!

(Let Y denote the # of vertices lying in small components (i.e. of order $< k_0$). We will prove that the probability for a random uniformly random vertex to be in a small component is $p + o(1)$)

Lemma: Fix $c > 1$ and let $Z_n \sim \text{Bin}(n, \frac{c}{n})$ where $\frac{c}{n} \in \mathcal{P}$. Then $P_{Z_n} \rightarrow P$ as $n \rightarrow \infty$ where $p e^c = e^{p c}$, ($0 < p < 1$)

Proof:

(Let $G(x) = E(x^{Z_n}) = (1 - p + p x)^n = e^{-c + c x + o(\frac{c^2}{n}) + o(1)}$, which converges to $e^{-c + c x}$ pointwise as $n \rightarrow \infty$.)

By the proof of the theorem for the branching process, part (iii):

P_n is the unique solution of $G(x) = x$, where $0 < x < 1$.

As $G(x) \rightarrow \underbrace{e^{-c + c x}}_{F(x)}$ pointwise and P is the unique solution of

$x = F(x)$, $0 < x < 1$, it follows immediately $P_n \rightarrow P$ as $n \rightarrow \infty$.

□

Lemma $\mathbb{E}Y = (1+o(1))Pn$.

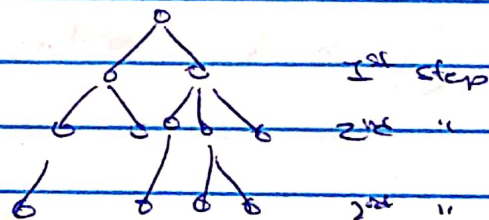
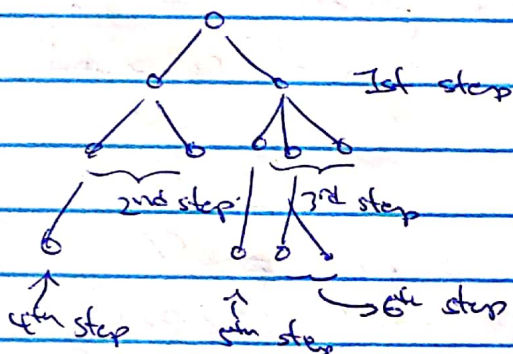
Proof:

Consider the graph exploration process starting from u , $\hat{\mathcal{U}}_u$.

Let \mathcal{E}_t denote the event that $\hat{\mathcal{U}}_u$ lasts $\leq t$ steps.

Consider the slowed down version of the branching process where $Z_n \sim \text{Bin}(n, (1+\xi)\frac{c}{n})$, where $\xi = 0 < \xi = o(1)$.

(Slowed-down Branching process.)



i.e. Parent's children are added at a time

$\text{Bin}(n, (1+\xi)\frac{c}{n})$ stochastically dominates $\text{Bin}(n - (\frac{n}{2} + t + 1), \frac{c}{n})$,
so we can couple $(\hat{\mathcal{U}}_u, \mathcal{U})$ such that $\hat{\mathcal{U}}_u$ terminates no
later than \mathcal{U} . (Exercise!)

Let P_t denote the probability that \mathcal{U} terminates during
the first t steps. We know $P_t \rightarrow P$.

That is $\forall \epsilon > 0, \exists n_0 > n \forall n > n_0: P_{k_0} \geq P - \epsilon$

And also, $P_{2n} \rightarrow P$ as $n \rightarrow \infty \Rightarrow P_{2n} \geq P - \epsilon$

$\Rightarrow P_{k_0} \geq P - 2\epsilon$

\hookrightarrow

By the coupling (\tilde{Z}_v, φ) , we have $P(Z_v) \geq P(\varphi \text{ terminates before } k_0 \text{ steps}) = P_{\text{iso}} \geq P - 2\epsilon \Rightarrow P(Z_v) \geq P - o(1)$

For the upper bound, consider the stand-down ^{branching} process with $Z_v \sim \text{Bin}(C_n - k_1, (1-\epsilon)\frac{c}{n})$. Z_v is stochastically dominated by $\text{Bin}(\frac{c}{n} - (Z_t + t), \frac{c}{n})$ for $t < k_0$

Again, we can couple these two processes such that during the first k_0 steps the graph exploration process will terminate no earlier than the branching process

$$\Rightarrow P(Z_v) \leq P(\text{Branching process}^{\text{last}} \leq k_0 \text{ steps}) = P + o(1), \text{ using similar argument as above}$$

$$\Rightarrow P(Z_v) = P + o(1)$$

$$\Rightarrow \mathbb{E}Y = (1 + o(1))Pn$$

□

$\mathbb{E}Y(Y-1)$. Pick a pair of vertices u and v .

$$P(C_u \text{ is in a small component}) = P + o(1)$$

Condition on that. The graph exploration process only "exposes" $\leq k_0$ vertices.

(Because: Easy to show that a.s. v is not exposed)

→ i.e. In the

same component that u is in.

Run the graph exploration process starting from u , condition on C_u , the component containing u . This is equivalent to running the graph exploration process in $G(n, p)$ where $n' = n - O(k)$.

Repeat the proof as for $\mathbb{E}Y$. We have \dots

②

We have $R(V)$ is in a small component / a is in a small component
 $= P + o(1)$.

$$\sum_{i=1}^n (d_i - 1) = n(n-1)(P^2 + o(1)) \sim (n-1)^2$$

By Chebyshev's Inequality: $\forall \epsilon > 0$

$$Pr(|R(V) - \sum_{i=1}^n (d_i - 1)| > \epsilon \sum_{i=1}^n (d_i - 1)) \leq \frac{\sum_{i=1}^n (d_i - 1)^2}{\epsilon^2 (\sum_{i=1}^n (d_i - 1))^2} = \frac{o(1) \cdot (n-1)^2}{\epsilon^2 (n-1)^2} = o(1)$$

So a.a.s $|R(V) - (P + o(1))n| = o(n)$. □

Hamiltonian Cycles in $G(n, p)$:

↓
 p 's are new
 p 's

Theorem:

$$\text{Let } p = \frac{\log n + \log \log n + o(1)}{n}.$$

(a) If $\frac{x(n)}{n} \rightarrow \infty$, then a.a.s $G(n, p)$ is Hamiltonian

(b) If $\frac{x(n)}{n} \rightarrow -\infty$, then $G(n, p)$ is not Hamiltonian

(Note: When degree 1 vertices disappear, then our graph becomes Hamiltonian (w.h.p)).

Proof:

Let D denote the property that min. degree is ≥ 2

$HAM \implies D$ Graph is Hamiltonian

$$G \in HAM \implies G \in D$$

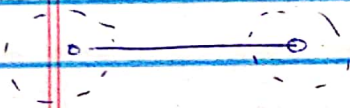
Lemma (Exercise)

(a) If $p_n - (\log n + \log \log n) \rightarrow -\infty$ then a.a.s $G(n, p) \in D_2$

(b) If $p_n - (\log n + \log \log n) \rightarrow \infty$ then a.a.s $G(n, p) \in D_2$

Lemma \Rightarrow Part (b) of the theorem

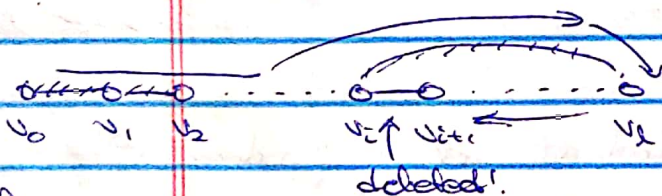
Idea for proof:



We try to build a layer and layer path until we build a Hamiltonian cycle.

What is the probability of us creating a layer path? \rightarrow But we can only build from the 2 ends

Rosé Rotation:



$v_0 v_1 v_2 \dots v_{i-1} v_i v_{i+1} \dots v_{l-1} v_l$ is a Rose Rotation

Def'n:

Let $P = v_0 v_1 \dots v_l$ be a longest path in G . So assume $v_i v_{i+1}$ is an edge. Then, the operation of deleting edge $v_i v_{i+1}$ from P and adding edge $v_i v_l$ to P is called a Rose Rotation.

The Rose Rotation converts P to another longest path



Given a longest path $P = v_0 v_1 \dots v_k$. Let \mathcal{P} be the set of longest paths obtained by fixing end v_0 and repeatedly performing Pivot Rotations.

Let $\text{End}(v_0)$ be the set of ends of paths in \mathcal{P} other than v_0 .

Given a subset S of vertices, let

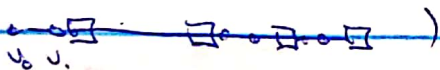
$$N(S) = \{u \in V : \exists v \in S \text{ s.t. } v \sim u\}$$

Lemma: $|N(\text{End}(v_0))| \leq 2|\text{End}(v_0)|$

Proof:

It is sufficient to show that if $v_i \in N(\text{End}(v_0))$, then one of v_{i-1}, v_{i+1} must be in $\text{End}(v_0)$.

(Extremal Case)



Consider a sequence of Pivot rotations which produces P' with $v_i v_x$, where $x \neq v_0$ is an end of P' . Let y and z be the left and right neighbors of v_i on P' . If $\{y, z\} = \{v_{i-1}, v_{i+1}\}$, the one of them can be added to $\text{End}(v_0)$ by performing a Pivot Rotation on P' . Otherwise, one of the edges $v_i v_x, v_i v_y$ has been deleted in a previous Pivot Rotation \Rightarrow one of v_{i-1}, v_{i+1} has been added to $\text{End}(v_0)$.

□