

# Bankruptcy Problem from the Talmud

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# Original paper

- Game Theoretic Analysis of a Bankruptcy Problem from the Talmud
- Robert J. Aumann and Michael Maschler
- Journal of Economic Theory **36**. 195-213 (1985)

# Outline

## 1 Introduction

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2 Consistent Solution

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# Bankruptcy Problem

- A man dies, leaving debts  $d_1, d_2, \dots, d_n$  totalling more than his estate  $E$ . How should the estate be divided among the creditors?

# Bankruptcy Problem from the Talmud

		Debt		
		100	200	300
Estate	100	$\frac{100}{3}$	$\frac{100}{3}$	$\frac{100}{3}$
	200	50	75	75
	300	50	100	150



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- The lesser claimant concedes half the garment to the greater one. The remaining half is divided equally.

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- We will say that this division is prescribed by the **CG(contested garment) principle**.

# Consistency

- A bankruptcy problem is defined as a pair  $(E; d)$ , where  $d = (d_1, d_2, \dots, d_n)$ ,  $0 \leq d_1 \leq d_2 \leq \dots \leq d_n$  and  $0 \leq E \leq D := d_1 + \dots + d_n$ .

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- A solution to such a problem is an  $n$ -tuple  $x = (x_1, \dots, x_n)$  of real numbers with

$$x_1 + x_2 + \dots + x_n = E.$$



# Consistency

- A solution is called **CG-consistent or consistent** if for all  $i \neq j$ , the division of  $x_i + x_j$  prescribed by the CG-principle for claims  $d_i, d_j$  is  $(x_i, x_j)$ .

## Remark

The bankruptcy problem from the Talmud are consistent solution.

# Consistent solution of bankruptcy problem

## Theorem 1

*Each bankruptcy problem has the unique consistent solution.*

## Uniqueness of consistent solution

Let  $x, y$  be two consistent solutions of given bankruptcy problem.  
We can find  $i \neq j$  such that

$$y_i > x_i, \quad y_j < x_j, \quad y_i + y_j \geq x_i + x_j.$$

Since CG principle has monotonicity with respect to estate,  
 $y_j \geq x_j$ . It is contradiction.

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# Existence of consistent solution

- This, in turn, continues until 2 has received  $d_2/2$ , at which point she stops receiving payment for the time being, and each additional dollar is divided equally between the remaining  $n - 2$  claimants.

# Existence of consistent solution

Case2)  $E > D/2$ .

- The process is the mirror image of the above. Instead of thinking  $i$ 's award  $x_i$ , one thinks in terms of her loss  $d_i - x_i$ .

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- When  $D - E$  is small, the total loss is shared equally between all creditors.
- This continues until 1 has lost  $d_1/2$ . For the time being she then stops losing and each additional loss is divided equally between the remaining  $n - 1$  claimants.

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## Remarks

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- 1 Above gives the consistent solution.
- 2 Let us denote  $f(E; d)$  be the unique consistent solution of bankruptcy problem  $(E; d)$ . Then the following duality holds;

$$f(E; d) = d - f(D - E; d).$$

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# Definition of bankruptcy game

## Definition

We define the bankruptcy game  $v_{E;d}$  corresponding to the bankruptcy problem  $(E; d)$  by

$$v_{E;d} := (E - d(N \setminus S))_+.$$

## Remark

$v_{E;d}$  is a coalition game with superadditivity.

$v_{E;d}(S)$ : Accepting either nothing, or what is left of the estate  $E$  after each member  $i$  of the coalition  $N \setminus S$  is paid his complete claim  $d_i$ .

# Main theorem

## Theorem 2

*The consistent solution of a bankruptcy problem is the nucleolus of the corresponding game.*

From now on, let  $v$  be a game,  $S$  a coalition and  $x$  a payoff vector.

# The reduced game

## Definition

The reduced game  $v^{S,x}$  is defined on the player space  $S$  as

$$v^{S,x}(T) = \begin{cases} x(T) & \text{if } T = \emptyset \text{ or } T = S \\ \max_{Q \subset N \setminus S} \{v(Q \cup T) - x(Q)\} & \text{if } \emptyset \subsetneq T \subsetneq S \end{cases}$$

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In the reduced game, the players of  $S$  consider how to divide the total amount assigned to them by  $x$  under the assumption that the players  $i$  outside  $S$  get exactly  $x_i$ .



# The reduced game

## Lemma 3

*Let  $x$  be a solution of the bankruptcy problem  $(E; d)$  such that  $0 \leq x_i \leq d_i$  for all  $i$ . Then for any coalition  $S$  the following holds;*

$$v_{E;d}^{S,x} = v_{x(S);d|S}.$$

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proof. Set  $v := v_{E;d}$ ,  $v^S := v_{E;d}^{S,x}$ . First let  $\emptyset \subsetneq T \subsetneq S$  and let the maximum in the definition of  $v^S(T)$  be attained at  $Q$ .

# The reduced game

$$\begin{aligned}
 v^S(T) &= v(T \cup Q) - x(Q) \\
 &= [E - d(N \setminus (Q \cup T))]_+ - x(Q)_+ \\
 &\leq [x(N) - d(N \setminus (Q \cup T)) - x(Q)]_+ \\
 &= [x(S) - d(S \setminus T) - (d - x)(N \setminus (S \cup Q))]_+ \\
 &\leq (x(S) - d(S \setminus T))_+
 \end{aligned}$$

Setting  $Q = N \setminus S$

$$\begin{aligned}
 v^S(T) &\geq v(T \cup (N \setminus S)) - x(N \setminus S) \\
 &\geq [E - d(N \setminus (T \cup (N \setminus S)))]_+ - (x(N) - x(S)) \\
 &\geq (E - d(S \setminus T)) - (E - x(S)) = x(S) - d(S \setminus T).
 \end{aligned}$$

# The reduced game

and setting  $Q = \emptyset$

$$v^S(T) \geq v(T \cup \emptyset) - x(\emptyset) = v(T) = (E - d(N \setminus T))_+ \geq 0.$$

Therefore  $v^S(T) = (x(S) - d(S \setminus T))_+ = v_{x(S); d|_S}(T)$ .

# Kernel and pre-kernel

Let  $v$  be a game. For each payoff vector  $x$  and players  $i, j$ , define

$$S_{ij}(x) = \max\{v(S) - x(S) \mid S \text{ contains } i \text{ but not } j\}.$$

## Definition

The pre-kernel of  $v$  is the set of all payoff vectors  $x$  with  $x(N) = v(N)$  and  $S_{ij}(x) = S_{ji}(x)$  for all  $i \neq j$ .

## Definition

The kernel of  $v$  is the set of all payoff vectors  $x$  with  $x(N) = v(N)$ ,  $x_i \geq v(i)$  for all  $i$  and for all  $i \neq j$

$$S_{ij}(x) > S_{ji}(x) \text{ implies } x_i = v(j).$$

# Standard solution of 2 person game

- For a given 2 person game  $v$ , the payoff vector

$$x_i = \frac{v(12) - v(1) - v(2)}{2} + v(i)$$

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### Remark

The nucleolus and kernel, the pre-kernel and the Shapley value of 2 person game all coincide with its standard solution.



# Pre-kernel and standard solution

## Lemma 4

*Let  $x$  be in the pre-kernel of a game  $v$  and let  $S$  be a coalition with exactly two person. Then  $x|_S$  is the standard solution of  $v^{S,x}$ .*

proof) Let  $S = \{i, j\}$  Then

$$\begin{aligned} S_{ij}(x) &= \max_{Q \subset N \setminus S} (v(Q \cup i) - x(Q \cup i)) = \max_{Q \subset N \setminus S} (v(Q \cup i) - x(Q)) - x_i \\ &= v^{S,x}(i) - x_i. \end{aligned}$$

Similarly, we know that  $S_{ji}(x) = v^{S,x}(j) - x_j$ . Since  $x$  is in the pre-kernel  $S_{ij}(x) = S_{ji}(x)$ . This implies

$$x_i - x_j = v^{S,x}(i) - v^{S,x}(j)$$

also by definition of reduced game

$$x_i + x_j = x(i, j) = v^{S,x}(i, j).$$

# Kernel and the consistent solution of bankruptcy problem

## Lemma 5

*The contested garment solution of 2 person bankruptcy problem is the standard solution of the corresponding game.*

## Proof.

Directly from Lemma 4 and definition of bankruptcy game and the standard solution. □

# Kernel and the consistent solution of bankruptcy problem

## Proposition 6

*The kernel of a bankruptcy game  $v_{E;d}$  consists of a single point, namely the consistent solution of the problem  $(E; d)$ .*

Proof) Set  $v := v_{E;d}$ .  $v$  is superadditive, hence 0-monotonic, i.e.  $S \subset T$  implies  $v(S) + \sum_{i \in T \setminus S} v(i) \leq v(T)$ . In 0-monotonic game, the kernel coincides with the pre-kernel. Hence  $x$  is in the pre-kernel of  $v$ . Let  $S$  be an arbitrary 2 person coalition. By lemma 4,  $x|_S$  is the standard solution of  $v^{S,x}$  and hence, of  $v_{x(S);d|_S}$ . Therefore  $x|_S$  is the CG-solution of  $(x(S); d|_S)$ , but that means  $x$  is the consistent solution of  $(E; d)$ .

# Proof of main theorem

## Theorem

*The consistent solution of a bankruptcy problem is the nucleolus of the corresponding game.*

## Proof.

Since the nucleolus is always in the kernel, by proposition 6, it coincides with the consistent solution of the bankruptcy problem. □