

Class activity, Week 9

Primitive Roots and Quadratic Residues

In this activity, we first investigate some of the properties of the *primitive roots* you discovered in the pre-class. We will also see why having a primitive root modulo n allows us to put a specific structure on the elements modulo n , and when this happens. Then at the end, we will use the primitive roots to obtain some information about quadratic residues modulo p , i.e. those elements that are squares modulo p .

Testing for Primitive Roots

Using the definition, to determine if g is a primitive root modulo n , we need to check that none of $g^2, g^3, \dots, g^{\phi(n)-1}$ is congruent to 1 modulo n . But that is a lot of work! Instead, using the next proposition, we can significantly reduce the number of powers of g that we need to compute to find whether g is a primitive root.

Proposition 1: An integer g is a primitive root modulo n if and only if for every p dividing $\phi(n)$

$$g^{\frac{\phi(n)}{p}} \not\equiv 1 \pmod{n}.$$

proof: Suppose g is a primitive root. Then, by definition, $g^k \not\equiv 1 \pmod{n}$ for all $1 \leq k < \phi(n) - 1$. Since $\frac{\phi(n)}{p} < \phi(n)$, $g^{\frac{\phi(n)}{p}} \not\equiv 1 \pmod{n}$.

Suppose now that g is an element for which

$$g^{\frac{\phi(n)}{p}} \not\equiv 1 \pmod{n}$$

for every p dividing $\phi(n)$, and assume that the order of g is $r < \phi(n)$. We will reach a contradiction.

From Corollary 2 of the week 8 class activity, we know that $r | \phi(n)$. Consider $\frac{\phi(n)}{r}$. There is a prime factor p of $\frac{\phi(n)}{r}$ and an integer m such that $pm = \frac{\phi(n)}{r}$. Then $rm = \phi(n)$. But this is contradiction because

$$g^{\frac{\phi(n)}{p}} \equiv (g^r)^m \equiv 1 \pmod{n}.$$

Therefore, $r = \phi(n)$ and g is a primitive root. □

1. Check that 2 is a primitive root modulo 101 in an efficient way. (Note: If you need to calculate large powers modulo 101, say 2^{60} , you can do this in two steps: 2^{15} and reduce modulo 101, and then take the fourth power of the reduced expression. You will need different powers than 60; I used 60 as an example.)

✓

$$\begin{aligned} \phi(101) &= 100 = 2^2 \cdot 5^2 \\ 2^{100/2} &= 2^{50} = (2^{10})^5 = (14)^5 \equiv 100 \pmod{101} \neq 1 \\ 2^{100/5} &= 2^{20} = (2^{10})^2 = (14)^2 \equiv 95 \pmod{101} \neq 1 \end{aligned}$$

2. Find a primitive root modulo 18. (Hint: A primitive root g modulo n has to be relatively prime to n for otherwise order of g is not defined. So that cuts down the number of g 's to try.)

$$\phi(18) = 6$$

(5) ✓

$$5^2 = 7$$

$$5^3 = -1$$

$$\pmod{18}$$

$$\pmod{18}$$

Structure of \mathbf{Z}_n^\times

Given a modulus n , we let \mathbf{Z}_n^\times denote the set of all residue classes with integers relatively prime to n . These integers are those for which we can find the order and there are $\phi(n)$ such integers. In this section, we will show that the set \mathbf{Z}_n^\times has a specific structure when there is a primitive root, as you guessed in problem 2 in the pre-class.

3. Suppose a has order k modulo n .

a. If $a^i \equiv a^j \pmod{n}$, what is the relation between i and j ?

$$i \equiv j \pmod{k}$$

b. Justify why $a^0, a^1, a^2, \dots, a^{k-1}$ are all different modulo n .

dis. res. classes $\neq K$

Using the above problem with the special case of a being a primitive root, we obtain the first half of the following theorem:

Theorem 1: If g is a primitive root modulo n , then any integer relatively prime to n is congruent to g^i modulo n for some i . Conversely, if g is an integer for which any integer relatively prime to n is congruent to g^i for some i , then g is a primitive root.

proof: Suppose g is a primitive root. Then g has order $\phi(n)$. Using the above problem, this says that $g^0, g^1, \dots, g^{\phi(n)-1}$ are all different modulo n . Also note that g^i are all relatively prime to n if g is. But there are only $\phi(n)$ residue classes which are relatively prime to n , so the powers of g must cover all these residue classes.

Suppose now that g is an integer such that every integer relatively prime to n is congruent to g^i for some i . Suppose the order of g is k . Then, from the above problem, there are \underline{k} different integers modulo n which are of the form g^i . But there are $\underline{\phi(n)}$ integers modulo n which are relatively prime to n . Therefore, $\underline{k = \phi(n)}$ and g is a primitive root. □

This theorem motivates another definition for a primitive root: An element g of \mathbf{Z}_n^\times is a *primitive root* if every element of \mathbf{Z}_n^\times is congruent to g^k for some power k . For this reason, a primitive root is also called a *generator* of \mathbf{Z}_n^\times . For those of you familiar with groups, this basically translates into saying that “if there is a primitive root, \mathbf{Z}_n^\times is a cyclic group with any primitive root being a generator.”

4. Check that 3 and 7 are primitive roots modulo 10, and express each in terms of the other.

$$\phi(10) = 4$$

$$3^2 = 9 \not\equiv 1$$

$$7^2 = 49 \not\equiv 1$$

$$7 \equiv 3^3 \pmod{10}$$

$$3 \equiv 7^3 \pmod{10}$$

Number of Primitive Roots

The next question to consider is that in cases when a primitive root exists, how many primitive roots are there total? If g is a primitive root, then every element is of the form g^k . So the question then becomes, “for which powers k is g^k still a primitive root?”

5. Given: 2 is a primitive root modulo 11.

a. Describe what this means.

order of 2 is 10

b. Find the order of $2^4, 2^7, 2^9$ modulo 11, using the information on the order of 2.

5, 10, 10

c. Generalizing your observations from part b, determine for which k the order of 2^k is 10, i.e. how many primitive roots are there modulo 11?

$$\phi(10) = 4$$

$$k = 1, 3, 7, 9$$

4 Prim roots

We consider now another example with a larger modulus $n = 101$. We know from earlier that 2 is a primitive root modulo 101. We know that $2^4, 2^{25}, 2^{75}$ are not primitive roots because

$$(2^4)^{25} \equiv 1 \pmod{101}, (2^{25})^4 \equiv 1 \pmod{101}, (2^{75})^4 \equiv 1 \pmod{101}.$$

But $2^3, 2^7, 2^9, 2^{11}, \dots$ are primitive roots. More generally, 2^k is a primitive root as long as $\gcd(k, 100) = 1$ (where $100 = \phi(101)$ is the order of 2). There are $\phi(100)$ such k 's.

Theorem 2: Suppose there is a primitive root g modulo n . Then g^k is also a primitive root if and only if $\gcd(k, \phi(n)) = 1$. Therefore, there are $\phi(\phi(n))$ primitive roots modulo n .

The following theorem (whose proof is omitted) describes for which n we can expect to find a primitive root in general.

Theorem 3: A primitive root modulo n exists if and only if n is of the form $2, 4, p^r, 2p^r$ where p is an odd prime and $r \geq 1$.

Quadratic residues

Recall how we defined quadratic residues modulo p a prime: If $x^2 \equiv a \pmod{p}$ has a solution, then a is a quadratic residue. If $x^2 \equiv a \pmod{p}$ does not have a solution, then a is a *quadratic non-residue*. We let R denote the set of quadratic residues, and N denote the set of quadratic nonresidues.

Theorem 3: If p is an odd prime, $|R| = |N| = (p-1)/2$, where $|\cdot|$ denotes set cardinality.

proof: In order to determine how many quadratic residues modulo p there are, we consider how many different squares we have modulo p . In fact, it is enough to count the number of x^2 's modulo p where $x = 1, 2, \dots, (p-1)/2$ because $i^2 = (-i)^2$. Also note that, $x^2 \not\equiv y^2 \pmod{p}$ for $x \not\equiv y \pmod{p}$ and $1 \leq x, y \leq (p-1)/2$. If they were we would have $p | x^2 - y^2 = (x-y)(x+y)$ and neither $x-y$ nor $x+y$ are divisible by p since $1 \leq x-y, x+y \leq p-1$. So, each $x = 1, 2, \dots, (p-1)/2$ produces a unique square modulo p . Therefore, there are $(p-1)/2$ quadratic residues. There are a total of $p-1$ elements relatively prime to p , therefore $|N| = p-1 - (p-1)/2 = (p-1)/2$. \square

6. There is another way to see why Theorem 3 is true using the primitive roots. Recall that for every prime p there is a primitive root modulo p . Let g be a primitive root modulo p .

a. Recall that every element modulo p can be written as g^i for some power i . What is the condition on i for the g^i to be a square? What is the condition for g^i to be a non-square?

$$2|i \rightarrow \square$$

$$2 \nmid i \rightarrow \boxtimes$$

b. Using part a, justify why $|R| = |N| = (p-1)/2$.

$$p-1 \text{ possible } i\text{'s}$$

$$\text{half are even}$$

So, using this last problem, we see that if we have a primitive root at hand, then it is easy to find the squares modulo p . However, finding a primitive root is a lot of work too. If we have to determine whether one element a modulo a large prime p is a quadratic residue, then it is not worth the trouble to find a primitive root. Instead, can we use the idea of the primitive root without finding one explicitly?

7. Suppose g is a primitive root modulo p .

a. Let a be a quadratic residue modulo p . Using g , justify why $a^{(p-1)/2} \equiv 1 \pmod{p}$.

$$a \equiv g^{2i}, \quad a^{(p-1)/2} = g^{2i \cdot (p-1)/2} = g^{i(p-1)} = g^{(p-1)} \equiv 1 \pmod{p}$$

b. Let a be a quadratic nonresidue modulo p . Using g , justify why $a^{(p-1)/2} \not\equiv 1 \pmod{p}$.

$$a = g^{2i+1}, \quad a^{(p-1)/2} = g^{(2i+1)(p-1)/2} \not\equiv 1$$

$$bc \frac{2i+1}{2} \notin \mathbb{Z}$$

c. Noting that $a^{p-1} \equiv 1 \pmod{p}$, justify why $a^{(p-1)/2} \equiv \pm 1 \pmod{p}$. Therefore, using part b, show that $a^{(p-1)/2} \equiv -1 \pmod{p}$ if a is a quadratic nonresidue.

only ± 1 square to 1 mod p

$$a^{(p-1)/2} = \pm 1 \text{ but not } +1 \text{ so } -1$$

Theorem 4: (Euler's criterion) Suppose that p is an odd prime and $p \nmid a$. Then the congruence $x^2 \equiv a \pmod{p}$ has two solutions if $a^{(p-1)/2} \not\equiv -1 \pmod{p}$ and no solution if $a^{(p-1)/2} \equiv -1 \pmod{p}$.

8. Using Euler's criterion, determine whether the following equations have solutions:

a. $x^2 \equiv 5 \pmod{101}$ $5^{50} \equiv 1$ yes, 2

b. $x^2 \equiv 7 \pmod{101}$ $7^{50} \equiv -1$ NO

Note: It should be noted that the Euler's criterion does not tell us what the solution x is. It only tells us that there is a solution. However, the test is actually pretty efficient. For large numbers, there are various fast exponentiation methods which make the testing in Euler's criterion fast.