

1. (15 points) Use the following table to find the maximum content of the sack using **Dynamic Programming**. The max weight the sack can hold is 15 lb. You May use the program to create the table.

item	1	2	3	4	5
Value(\$)	7	9	5	12	14
Weight(lb)	3	4	2	6	7

```
#include <iostream>
#include <iomanip>
#include <algorithm>

using namespace std;

// returns the maximum value that can be put in a knapsack of capacity W
int knapSack(int W, int wt[], int val[], int n) {
    int i, w;
    int K[n+1][W+1]; // K[n+1][W+1]

    // build table K[][] in bottom up manner
    for (i = 0; i <= n; i++) {
        for (w = 0; w <= W; w++) {
            if (i == 0 || w == 0) K[i][w] = 0;
            else if (wt[i-1] <= w) K[i][w] = max(val[i-1] + K[i-1][w - wt[i-1]],
K[i-1][w]);
            else K[i][w] = K[i - 1][w];
        }
    }

    for (int i = 0; i <= n; ++i) {
        for (int w = 0; w <= W; w++) cout << setw(5) << K[i][w];
        cout << endl;
    }
    return K[n][W];
}

int main() {
    // number of items and capacity of the knapsack
    int n = 5, W = 15;
    int val[5] = { 7, 9, 5, 12, 14 }, wt[5] = { 3, 4, 2, 6, 7 };

    int maxValue = knapSack(W, wt, val, n);
```

```

cout << "\nMaximum value in the Knapsack = " << maxValue << endl;

return 0;
}

```

**WRITTEN WORK:**

1.

Item	1	2	3	4	5
Value (v)	7	9	5	12	14
Weight (w)	3	4	2	6	7

15 lb

Weight

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
bag = {3}	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
bag = {2,3}	0	0	0	7	7	7	7	7	7	7	7	7	7	7	7	7
				{3}	...									...	{3}	
bag = {1,3}	0	0	0	7	9	9	9	16	16	16	16	16	16	16	16	16
				{3}	{2,3}	{2,3}	{2,3}	{1,2,3}	...						{1,2,3}	
bag {1,4,3}	0	0	5	7	9	12	14	16	16	21	21	21	21	21	21	21
			{3}	{1,3}	{2,3}	{1,2,3}	{1,2,3}	{1,2,3}	...	{1,2,3}	...				{1,2,3}	
bag {1,2,3,4}	0	0	5	7	9	12	14	16	17	21	21	24	26	28	28	33
																{1,2,3,4}
bag {1,...,5}	0	0	5	7	9	12	14	16	17	21	21	24	26	28	30	33
																{1,2,3,4,5}

Max Value = 33

2. (15 points) Given the following chain of matrices multiplications

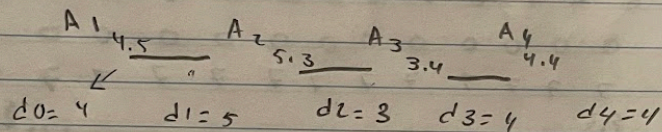
Matrices:     A1 \* A2 \* A3 \* A4  
Dimensions:    4.5     5.3     3.4     4.4

Use Dynamic Programming to determine in what order we should multiply those matrices to obtain minimum number of multiplications. Show all work for maximum credit.



2. Matrices  $A_1 \times A_2 \times A_3 \times A_4$   
 Dimensions: 4,5 5,3 3,4 4,4

Minimum number of multiplications



D	A1	A2	A3	A4	M	A1	A2	A3	A4
A1	0	$d_1=5$	$d_1=5$	$d_3=4$	A1	0	60	108	156
A2		0	$d_2=3$	$d_3=4$	A2		0	60	108
A3			0	$d_3=4$	A3			0	48
A4				0	A4				0

$$M(1,2) = A_1 \cdot A_2, = 4 \times 5 \times 3 = 60 = M(1,1)$$

$$M(2,3) = A_2 \cdot A_3, = 5 \times 3 \times 4 = 60$$

$$M(3,4) = A_3 \cdot A_4, = 3 \times 4 \times 4 = 48$$

$$M(i,j) = \min \sum M(i,k) + M(k+1,j) + d_{i-1} \times d_k \times d_j, i \leq k < j$$

$$M(1,3) = \min_{i=1, j=3} \sum_{k=1} M(1,k) + M(k+1,3) + d_0 \times d_1 \times d_3, M(1,2) + M(3,3) + d_0 \times d_1 \times d_3$$

$$= \min \sum 0 + 60 + 4 \times 5 \times 4, \quad 60 + 0 + 4 \times 3 \times 4 \} \\ = \min \{ 140, 108 \} = 108$$

$M[1,4] = 156$  min num of mul.

$$M(2,4) = \min \sum M(2,k) + M(k+1,4) + d_1 \times d_2 \times d_4, M(2,3) + M(4,4) + d_1 \times d_3 \times d_4$$

$$= \min \sum 0 + 48 + 5 \times 3 \times 4, \quad 60 + 0 + 5 \times 4 \times 4 \}$$

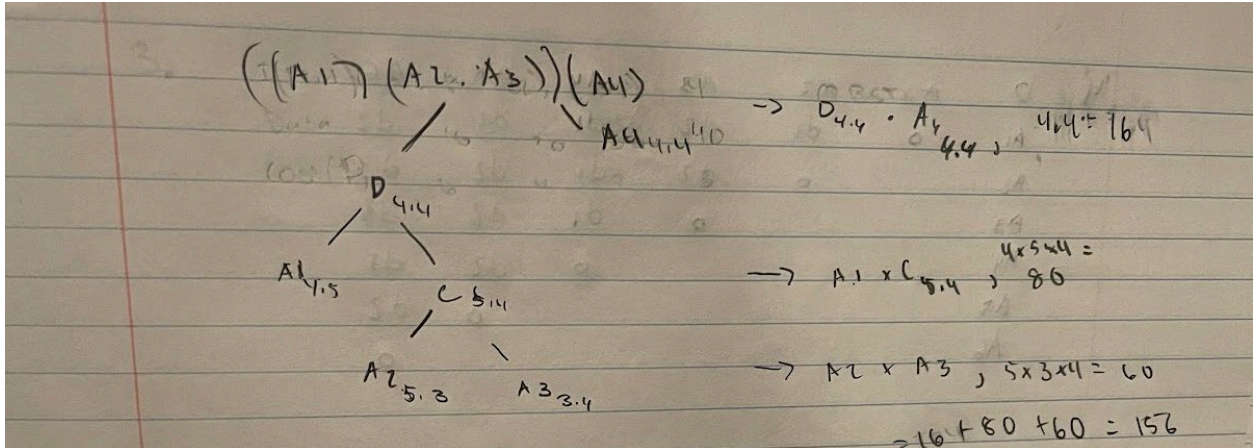
$$= \min \sum 108, \quad 140 \}$$

$$= 108$$

$$M(1,4) = \min \sum M(1,k) + M(k+1,4) + d_0 \times d_1 \times d_4, M(1,2) + M(3,4) + d_0 \times d_2 \times d_4$$

$$= \min \sum 0 + 108 + 4 \times 5 \times 4, \quad 60 + 48 + 4 \times 3 \times 4 \}$$

$$= \min \sum 188, \quad 156 \} = 156$$



3. (15 points) Use the following table to construct an OBST

Items	1	2	3	4
Data	10	20	30	40
Cost(\$)	6	4	2	3

Code to find optimal BST:

```
#include <bits/stdc++.h>
using namespace std;

int sum(int freq[], int i, int j);

int optCost(int freq[], int i, int j)
{
    if (j < i)
        return 0;
    if (j == i)
        return freq[i];

    int fsum = sum(freq, i, j);

    int min = INT_MAX;

    for (int r = i; r <= j; ++r)
    {
        int cost = optCost(freq, i, r - 1) +
                    optCost(freq, r + 1, j);
        if (cost < min)
            min = cost;
    }
    // Return minimum value
    return min + fsum;
}
```

```

}

int optimalSearchTree(int keys[],
                     int freq[], int n)
{
    return optCost(freq, 0, n - 1);
}

int sum(int freq[], int i, int j)
{
    int s = 0;
    for (int k = i; k <= j; k++)
        s += freq[k];
    return s;
}

// Driver Code
int main()
{
    int keys[] = {10, 20, 30, 40};
    int freq[] = {6, 4, 2, 3};
    int n = sizeof(keys) / sizeof(keys[0]);
    cout << "Cost of Optimal BST is "
          << optimalSearchTree(keys, freq, n);
    return 0;
}

```

### OUTPUT:

```

[Running] cd "c:\Users\Nick\Desktop\Algorithm-Engineering-Work\Homework 4\" &&
g++ OBST.cpp -o OBST && "c:\Users\Nick\Desktop\Algorithm-Engineering-
Work\Homework 4\OBST
Cost of Optimal BST is 28
[Done] exited with code=0 in 0.994 seconds

```



# Written Work:

Items	1	2	3	4
Data	10	20	30	40
cost (\$)	6	4	2	3

14 trees based on Catalan Numbers 3

1) 
$$\begin{array}{c} 10 \\ / \quad \backslash \\ 1 \quad 20 \\ \quad / \quad \backslash \\ \quad 2 \quad 30 \\ \quad \quad / \quad \backslash \\ \quad \quad 3 \quad 40 \\ \quad \quad \quad / \quad \backslash \\ \quad \quad \quad 4 \end{array}$$
 
$$C[0,0] = 0 \quad C[1,1] = 6 \quad C[2,2] = 8$$
 
$$1 \times 6 + 2 \times 4 + 3 \times 2 + 4 \times 3 = 38$$

2) 
$$\begin{array}{c} 10 \\ / \quad \backslash \\ 1 \quad 20 \\ \quad / \quad \backslash \\ \quad 2 \quad 30 \\ \quad \quad / \quad \backslash \\ \quad \quad 3 \quad 40 \\ \quad \quad \quad / \quad \backslash \\ \quad \quad \quad 4 \end{array}$$
 
$$1 \times 6 + 2 \times 4 + 3 \times 3 + 4 \times 2 = 31$$

3) 
$$\begin{array}{c} 10 \\ / \quad \backslash \\ 1 \quad 20 \\ \quad / \quad \backslash \\ \quad 2 \quad 30 \\ \quad \quad / \quad \backslash \\ \quad \quad 3 \quad 40 \\ \quad \quad \quad / \quad \backslash \\ \quad \quad \quad 4 \end{array}$$
 
$$1 \times 6 + 2 \times 2 + 3 \times 4 + 4 \times 3 = 34$$

4) 
$$\begin{array}{c} 10 \\ / \quad \backslash \\ 1 \quad 20 \\ \quad / \quad \backslash \\ \quad 2 \quad 30 \\ \quad \quad / \quad \backslash \\ \quad \quad 3 \quad 40 \\ \quad \quad \quad / \quad \backslash \\ \quad \quad \quad 4 \end{array}$$
 
$$1 \times 6 + 2 \times 2 + 3 \times 3 + 4 \times 4 = 35$$

5) 
$$\begin{array}{c} 10 \\ / \quad \backslash \\ 1 \quad 20 \\ \quad / \quad \backslash \\ \quad 2 \quad 30 \\ \quad \quad / \quad \backslash \\ \quad \quad 3 \quad 40 \\ \quad \quad \quad / \quad \backslash \\ \quad \quad \quad 4 \end{array}$$
 
$$1 \times 6 + 2 \times 3 + 3 \times 4 + 4 \times 2 = 32$$

6) 
$$\begin{array}{c} 10 \\ / \quad \backslash \\ 1 \quad 20 \\ \quad / \quad \backslash \\ \quad 2 \quad 30 \\ \quad \quad / \quad \backslash \\ \quad \quad 3 \quad 40 \\ \quad \quad \quad / \quad \backslash \\ \quad \quad \quad 4 \end{array}$$
 
$$1 \times 6 + 2 \times 3 + 3 \times 2 + 4 \times 3 = 34$$

7) 
$$\begin{array}{c} 10 \\ / \quad \backslash \\ 1 \quad 20 \\ \quad / \quad \backslash \\ \quad 2 \quad 30 \\ \quad \quad / \quad \backslash \\ \quad \quad 3 \quad 40 \\ \quad \quad \quad / \quad \backslash \\ \quad \quad \quad 4 \end{array}$$
 
$$1 \times 4 + 2 \times 6 + 3 \times 2 + 4 \times 3 = 51$$

14) 
$$\begin{array}{c} 10 \\ / \quad \backslash \\ 1 \quad 20 \\ \quad / \quad \backslash \\ \quad 2 \quad 30 \\ \quad \quad / \quad \backslash \\ \quad \quad 3 \quad 40 \\ \quad \quad \quad / \quad \backslash \\ \quad \quad \quad 4 \end{array}$$
 
$$1 \times 4 + 2 \times 6 + 2 \times 3 + 3 \times 2 = 28$$

4. (10 points) Use the following table to group each chain of matrices to minimize the number of multiplications

D	A1	A2	A3	A4	A5	A6
A1	0	d1	d2	d1	d1	d2
A2		0	d2	d1	d2	d3
A3			0	d1	d2	d3
A4				0	d2	d1
A5					0	d2
A6						0

Complete the following table

Chain	Grouping with min multiplications (just write the grouping order)
A1.A2.A3.A4.A5.A6	
A2.A3.A4.A5.A6.A7	
A3.A4.A5.A6	
A4.A5.A6	



4. D

	A1	A2	A3	A4	A5	A6
A1	0	d1	d2	d1	d1	d2
A2		0	d2	d1	d2	d3
A3			0	d1	d2	d3
A4				0	d2	d1
A5					0	d2
A6						0

Chain

A1 A2 A3 A4 A5 A6

A2 A3 A4 A5 A6  
(A1 DNE)

A3 A4 A5 A6

A4 A5 A6

Group

$$= (A1 A2) (A3 A4 A5 A6), D(A1, A6) = d2$$

$$= (A1 A2) ((A3 A4 A5) A6), D(A3, A6) = d3$$

$$= \boxed{(A1 A2) ((A3 A4) (A5) \cdot A6)}, D(A3, A5) = d2$$

$$= (A2 A3 A4) (A5 A6), D(A2, A6) = d3$$

$$= \boxed{(A2) \cdot (A3 A4) (A5 A6)}, D(A2, A4) = d1$$

$$= (A3 A4 A5) (A6), D(A3, A6) = d3$$

$$= \boxed{((A3 A4) \cdot (A5)) (A6)}, D(A3, A5) = d2$$

$$= \boxed{(A4) \cdot (A5 \cdot A6)}, D(A4, A6) = d1$$



5. Write a program to display the first 10 Fibonacci numbers using the Cassini's identity  
 $F(n-1)F(n-2) - [F(n)]^2 = (-1)^n, n \geq 1$

```
#include <iostream>

using namespace std;

size_t s[100] = { 0 };

int cassini(int n) { return (n & 1) ? -1 : 1; }

size_t fib(const int& n) {
    if (s[n]) return s[n];
    if (n <= 2) return s[n] = 1;
    else
    {
        s[n] = fib(n-1) + fib(n-2);
        return s[n];
    }
}

int main(int argc, char** argv) {

    for(int i = 0; i < 11; i++){
        cout << fib(i) << "          " << cassini(i) << endl;
    }
    return 0;
}
```

**OUTPUT:**

```
[Running] cd "c:\Users\Nick\Desktop\Algorithm-Engineering-Work\Homework 4\" &&
g++ fib.cpp -o fib && "c:\Users\Nick\Desktop\Algorithm-Engineering-Work\Homework
4\"fib
1          1
1          -1
1          1
2          -1
3          1
5          -1
8          1
13         -1
21         1
34         -1
55         1

[Done] exited with code=0 in 0.421 seconds
```