

$$A \times \coprod_{G(n)} P(n) \times B^n \cong \coprod_{G(n)} A \times (P(n) \times B^n)$$

$$\rightarrow \coprod_{G(n)} A^n \times (P(n) \times B^n)$$

$$\rightarrow \coprod_{G(n)} P(n) \times (A \times B)^n \cong \underline{P}(A \times B)$$

$$t: A \times \coprod_{G(n)} P(n) \times B^n \rightarrow A^n \times \coprod_{G(n)} P(n) \times B^n \rightarrow \coprod_{G(n)} P(n) \times (A \times B)^n$$

$$t^*: \coprod_{G(n)} P(n) \times A^n \times B \rightarrow \coprod_{G(n)} P(n) \times A^n \times B^n \rightarrow \coprod_{G(n)} P(n) \times (A \times B)^n$$

$$(a_1, b_1) (a_1, b_2) \dots (a_1, b_n) (a_2, b_1) \dots (a_2, b_n) \dots (a_m, b_n)$$

$$(a_1, b_1) (a_2, b_1) \dots (a_m, b_1) (a_1, b_2) \dots (a_m, b_2) \dots (a_m, b_n)$$

$$\begin{aligned} p &\in P(n) & q &\in P(n) \\ \mu(p, q, \dots, q) &\in P(nm) \\ \mu(q, p, \dots, p) &\in P(nm) \end{aligned}$$

$$\text{Iso in } \coprod_{G(n)} P(n) \times (A \times B)^n$$

The operad for braided strict mon. cats is

$$\begin{aligned} \bullet \quad G(n) &= Br_n, \quad \pi: G(n) \rightarrow \Sigma_n \\ \circ \quad P(n) &= E Br_n \end{aligned}$$

$$A \times B \xrightarrow{\cong} B \times A \begin{cases} \rightarrow B \\ \rightarrow A \end{cases}$$

3.2

Why is the free strict monoidal cat 2-monod not ps-commutative.

t, t^* both given by diagonals

Gives a strength for any operad
 $\gamma_{A,B}$ is just the permutation $(A \times B)^{nn} \xrightarrow{\cong} (A \times B)^{nn}$

We think it is ps-commutative.

Thm #1: $P(n) = E G(n)$, P is ps-comm.

$$[p, f] = [q, g]$$

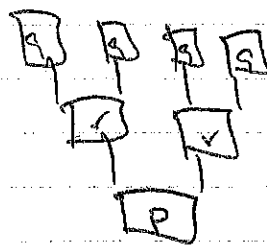
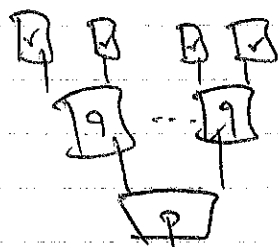
PF: Read if Land-Power better.

Cor: Free strict monoidal cat 2-monod
 & Free braided " 2-monod
 are ps-comm.

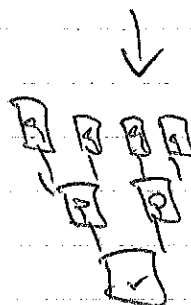
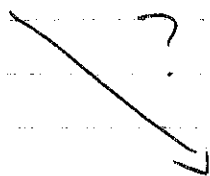
Aside: $A_0 = 1$ $G(n) = A_n$
 $A_1 = 1$ $P(n) = A_n$
 $A_2 = 1$ $P \text{ acts } X \mapsto X/A_n \rightarrow X$
 $A_3 = \mathbb{Z}/3$ $abc = bac = cab$
 $A_4 = \text{semidirect of } \mathbb{Z}/3 \text{ and } \mathbb{Z}/2 \times \mathbb{Z}/2?$
 $abcd = bacd = badc$

Thm #2: $P(n)$ contractible $\Rightarrow P$ ps-comm.
 This implies Thm #1 (obviously).

3.3



(1) - (3)



= id

(4) - (5)

Strength!

17/04/23

$$t_{AB} : A \times P B \longrightarrow P(A \times B)$$

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$$A \times \coprod_{n \geq 0} P(n) \times_{\mathcal{A}(n)} B^n \longrightarrow \coprod_{n \geq 0} P(n) \times_{\mathcal{A}(n)} (A \times B)^n$$

$$(a, [p, \underline{b}])$$

$$[p, (a, \underline{b})]$$

$$(a, [p, \underline{b}]) \longmapsto [p, (a, \underline{b})]$$

$$(a, [p, b_1, \dots, b_n]) \longmapsto [p, (a, b_1), \dots, (a, b_n)]$$

$$(a, [p \cdot g, \underline{b}]) \longmapsto [p \cdot g, (a, \underline{b})]$$

$$= [p, \pi_g(a, \underline{b})] = [p, (a, \pi_g \underline{b})]$$

||

$$(a, [p, \pi_g \underline{b}]) \longmapsto [p, (a, \pi_g \underline{b})]$$

well-defined.

$$\begin{array}{ccc} I \times PA & \xrightarrow{t_{IA}} & P(I \times A) \\ & \searrow \lambda_{PA} & \downarrow P(\lambda_A) \\ & & PA \end{array}$$

$$\begin{array}{ccc} I \times \coprod_{n \geq 0} P(n) \times_{\mathcal{A}(n)} A^n & \longrightarrow & \coprod_{n \geq 0} P(n) \times_{\mathcal{A}(n)} (I \times A)^n \\ & \searrow & \downarrow \\ & & \coprod_{n \geq 0} P(n) \times_{\mathcal{A}(n)} A^n \end{array}$$

$$\begin{array}{ccc} (*, [p, \underline{a}]) & \longmapsto & [p, (*, \underline{a})] \\ & \searrow & \downarrow \\ & & [p, \underline{a}] \end{array}$$

$$(a, [p, b]) \mapsto [p, (a, b)]$$

$$\begin{array}{ccc} A \times B & \xrightarrow{1 \times \iota_B} & A \times \mathbb{P}B \\ & \searrow \pi_{A \times B} & \downarrow \iota_{A \times B} \\ & & \mathbb{P}(A \times B) \end{array}$$

$$\begin{array}{ccc} A \times B & \xrightarrow{1 \times \iota_B} & A \times \mathbb{P}(\mathbb{P}(\mathbb{P}(A) \times_{\mathbb{P}(A)} B)) \\ & \searrow \pi_{A \times B} & \downarrow \iota_{A \times B} \\ & & \mathbb{P}(\mathbb{P}(\mathbb{P}(A) \times_{\mathbb{P}(A)} (A \times B))) \end{array}$$

$$\begin{array}{ccc} (a, b) & \xrightarrow{\quad} & (a, [\mathbb{I}, b]) \\ & \searrow & \downarrow \\ & & [\mathbb{I}, (a, b)] \end{array}$$

$$\begin{array}{ccc} \cancel{(A \times B) \times \mathbb{P}C} & \xrightarrow{\quad} & \cancel{A \times \mathbb{P}((A \times B) \times C)} \\ & \searrow & \downarrow \\ & & (A \times B) \times \mathbb{P}C \xrightarrow{\iota_{A \times B, C}} \mathbb{P}((A \times B) \times C) \\ & & \parallel \\ & & A \times (B \times \mathbb{P}C) \xrightarrow{1 \times \iota_{B \times C}} A \times \mathbb{P}(B \times C) \xrightarrow{\iota_{A \times B, C}} \mathbb{P}(A \times (B \times C)) \end{array}$$

$$(a, b), [p, s] \mapsto [p, (a, b), c]$$

$$\begin{array}{ccc} \downarrow & & \downarrow \\ (a, (b, [p, c])) & \mapsto & (a, [p, (b, s)]) \mapsto [p, (a, (b, c))] \end{array}$$

$$\begin{array}{ccccc}
 A \times P^2 B & \xrightarrow{t_{AB}} & P(A \times P B) & \xrightarrow{P_{t_{AB}}} & P^2(A \times B) \\
 \downarrow \mu_B & & & & \downarrow \mu_{A \times B} \\
 A \times P B & \xrightarrow[t_{AB}]{} & P(A) & \xrightarrow{} & P(A \times B)
 \end{array}$$

$$\begin{array}{ccc}
 (a, [q, [v_1, b], \dots, [v_n, b_n]]) & \mapsto & [q, (a, [v_1, b])] \mapsto [q, [v_1, (a, b)]] \\
 \downarrow & & \downarrow \\
 (a, [q(v_1, \dots, v_n), b]) & \xrightarrow{} & [q(v_1, \dots, v_n), (a, b)] \\
 & & \parallel
 \end{array}$$

Coherence:

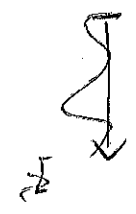
$$\begin{array}{l}
 P A \times B \xrightarrow{t_{AB}^*} P(A \times B) \\
 (\perp P(A) \times_{\Sigma(A)} A^*) \times B \\
 ([p, a], b) \mapsto [p, (a, b)]
 \end{array}$$

$$([p, a], b) \mapsto [p, (a, b)] = [p, \pi_2(a, b)] = [p, (\pi_2 a, b)] = t_{AB}^*([p, \pi_2 a], b)$$

Pseudo-compatibility:

$$\begin{array}{ccc}
 t_{AB} \circ (\perp P(A) \times_{\Sigma(A)} A^*) \times (\perp P(A) \times_{\Sigma(A)} B^*) & \longrightarrow & \perp P(A) \times_{\Sigma(A)} (A \times (\perp P(A) \times B^*))^* \\
 \downarrow & & \searrow \\
 \perp P(A) \times_{\Sigma(A)} ((\perp P(A) \times A^*) \times B)^* & \longrightarrow & \perp P(A) \times_{\Sigma(A)} P(A) \times_{\Sigma(A)} (A \times B)^* \times \dots \\
 & & \downarrow \\
 \perp P(A) \times_{\Sigma(A)} ((\perp P(A) \times A^*) \times B)^* & \longrightarrow & \perp P(A) \times_{\Sigma(A)} (A \times B)^*
 \end{array}$$

$$[p, a], [q, b] \mapsto [p, (a, [q, b])] \mapsto [p, [q, (a, b)]]$$



$$([p, a], [q, b]) \mapsto$$

$P A \times P B$



$$[p, [q, ([p, a], b)]] \mapsto [q, [p, (a, b)]]$$

$$([p, a], [q, b]) \mapsto [p, (a_1, [q, b]), \dots, (a_n, [q, b])] \mapsto [p, [q, (a_1, b)], \dots, [q, (a_n, b)]]$$



$p \in P(A), \quad q \in P(B),$
 $a = a_1, \dots, a_n, \quad b = b_1, \dots, b_m.$



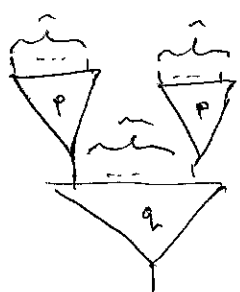
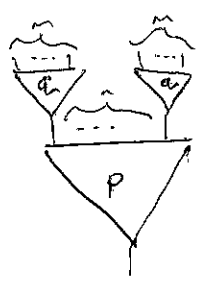
$$[p(q_1, \dots, q_n), (a_1, b)]$$

$$[q, ([p, a], b_1), \dots, ([p, a], b_m)] \mapsto [q, [p, (a, b_1)], \dots, [p, (a, b_m)]] \mapsto [q(p_1, \dots, p_n), (a, b)]$$

Need isomorphisms

$$[q(p_1, \dots, p_n), (a, b_1), \dots, (a, b_m)] \cong [p(q_1, \dots, q_n), (a_1, b), \dots, (a_n, b)]$$

Do $p(q_1, \dots, q_n)$ and $q(q_1, \dots, p_n)$ make sense? Yes, 'n' q's from $[q, (a_1, b)]$'s.
 ; vice versa.



$P(A \times B)$