

GALOIS THEORY

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These notes are for use in the first half of the Masters course on Galois Theory, Bristol, 2009. The treatment is based on notes by Andrey Lazarev and Trevor Woolley (both Bristol) and Rick McFeat (Western Australia). Two texts for the course are recommended - see the bibliography. The official book for this course is [Gar86]; the secondary book, [Ste03], is more informal in tone.

Most of the exercises in the following set of notes are very easy. So make sure you do them! Some proofs are not included, but citations are provided. You should regard these as hard exercises - try and prove them yourself. If repeated attempts don't yield success, then refer to the relevant text.

The website for (my part of) the course is:

<http://www.maths.bris.ac.uk/~manpg/teaching.html>

1. MOTIVATION

Galois theory started with... wait for it... a man called Evariste Galois. It straddles, and informs, several areas of mathematics, most notably group theory and number theory. The beauty of Galois theory, and its importance in many areas of pure mathematics, will hopefully become very obvious as the course proceeds.

1.1. Some history. Right from its first conception, Galois theory provided answers to some of the most long-standing questions in mathematics.

1.1.1. Solutions to polynomial equations. We all learn at school how to solve the quadratic equation,

$$ax^2 + bx + c = 0,$$

where x is a variable and a, b, c are fixed real or complex numbers. Our method is to *complete the square* and then solve for x ; this yields the well-known quadratic formula:

$$x = \frac{-b \pm \sqrt{b^2 - 4ac}}{2a}.$$

This stuff aint new: the Babylonians could solve quadratic equations using algebraic methods in 1600BC. Some time later the Greeks came up with geometric methods that did the job. Indian mathematicians also knew the answer to this a very long time ago.

Cubic equations proved much harder. Take the polynomial equation,

$$a_n x^n + a_{n-1} x^{n-1} + \cdots + a_2 x^2 + a_1 x + a_0 = 0.$$

Now, we can reduce the problem of solving this equation to that of solving a *monic polynomial*:

$$x^n + a_{n-1} x^{n-1} + \cdots + a_2 x^2 + a_1 x + a_0 = 0.$$

(*Exercise*) In the cubic case it is easy to show that this general monic case can be reduced to the situation

$$x^3 + a_1x + a_0 = 0.$$

(*Exercise*) A bunch of renaissance mathematicians in Bologna came up with a general solution to this last equation. There was a vast amount of intrigue and murkiness over who deserved the credit for this discovery. Niccolo “Tartaglia” Fontana demonstrated his methods in public around 1535.

For the equation $x^3 + px = q$, we have the solution

$$x = \sqrt[3]{\frac{q}{2} + \sqrt{\frac{p^3}{27} + \frac{q^2}{4}}} + \sqrt[3]{\frac{q}{2} - \sqrt{\frac{p^3}{27} + \frac{q^2}{4}}}.$$

Notice that the expression for this general solution depends on the coefficients via a finite number of operations involving $+$, $-$, \times , \div and extraction of roots, $\sqrt[n]{}$. We refer to such an expression as a *radical* expression. This will be a very important concept in what follows.

Quartic equations were solved around the same time. Ludovico Ferrari came up with a method for solving the quartic by reducing it to a cubic. The method also expressed the solutions of a quartic as a radical expression.

Quintic equations were the natural next equation to be solved.... But nobody could do it. Various big names - Euler, Lagrange - proved results related to this question, some of which made people start to think that perhaps no such equation existed.

In 1824, Abel proved that there is no general formula for the equation

$$x^5 + bx^4 + cx^3 + dx^2 + ex + f = 0,$$

using $+$, $-$, \times , \div , and $\sqrt[n]{}$. In other words the quintic is *insoluble by radicals*.

So now, supposing we have a particular quintic equation, how do we know if it can be solved by radicals? Abel was working on this question in the years after 1824. Unfortunately tuberculosis brought an end to his mathematical research and, indeed, his life in 1829, at the age of 26.

Our man Evariste Galois appeared on the scene around now. He had a whole heap of trouble with the mathematical establishment, with various political groups, and - most fatal of all - with the ladies. It would seem likely that it was an affair of the heart that led to his accepting a challenge to a duel one day in 1832. And so, at the age of 20, he died.

The day before this duel Galois wrote to his friend Auguste Chevalier outlining his mathematical discoveries; hoping that Chevalier could make these discoveries known. Perhaps the most significant of these discoveries can be summed up in the statement that *an equation is soluble by radicals provided its group is soluble*. What this mysterious sentence really means will be the subject of a large part of this course.

We won't dwell on the historical events surrounding the lives of both Galois, and Abel. Suffice to say that they are both fascinating figures, and you could do worse than look them up on wikipedia.

1.1.2. *Construction problems.* We should also briefly mention an application of Galois theory that was of great historical importance.

Suppose that one is given two instruments - a ruler (an unmarked straight edge), and a compass.

Now suppose I draw an angle, α , from two straight lines. You probably know from highschool that, using only a ruler and compass, it is possible to draw the angle $\frac{\alpha}{2}$. In other words, I can *bisect the angle*.

The Greeks loved this sort of game, and could do many nice constructions of the same ilk. But there were some constructions that they couldn't do:

- *duplication of the cube*: given a cube C , construct a cube D of twice the volume;
- *trisection of the angle*;
- *quadrature of the circle*: given a circle C , construct a square S of the same area.

It turns out that all of these problems are impossible, for reasons that Galois theory makes clear.

There is a related problem that I should mention. We would like to draw a regular polygon using only a ruler and compass. It turns out that some regular polygons (e.g. a pentagon) are easy to construct with ruler and compass; others are not. This led to the question: Is it possible to construct all regular polygons with ruler and compass?

Carl Friedrich Gauss in 1796 showed that a regular n -sided polygon can be constructed with ruler and compass if the odd prime factors of n are distinct Fermat primes. Gauss conjectured that this condition was also necessary, but he offered no proof of this fact, which was proven by Pierre Wantzel in 1837... using Galois theory.

1.2. The “group of an equation”. Before we get to some proper mathematics, I want to have an (arm-waving) look at that mysterious sentence from earlier, “an equation is soluble by radicals provided its group is soluble”. Specifically, I have this question: *what is the “group of an equation”?*

We proceed by example Consider the polynomial equation

$$x^4 - 2 = 0,$$

with integer coefficients. We have four solutions:

$$\alpha = \sqrt[4]{2}, \beta = i\sqrt[4]{2}, \gamma = -\sqrt[4]{2}, \delta = -i\sqrt[4]{2}.$$

Note that two of these solutions are not real. Now let us list some equations that these four roots satisfy:

$$\alpha + \gamma = 0, \alpha\beta\gamma\delta = -2, \alpha\beta - \gamma\delta = 0, \dots$$

What happens if we swap α and γ in this list? We get

$$\gamma + \alpha = 0, \gamma\beta\alpha\delta = -2, \gamma\beta - \alpha\delta = 0, \dots$$

which are all still valid. Likewise we could permute the variables as follows

$$\alpha \mapsto \beta \mapsto \gamma \mapsto \delta.$$

You can check that the equations above remain valid. So, does any such permutation work? For any such equation? The answer is “no”: try, for instance, swapping β and γ , and you find that the first of the equations above becomes false.

The set of permutations of the set $\{\alpha, \beta, \gamma, \delta\}$ that preserve the validity of all polynomial equations (with coefficients in \mathbb{Q}) in these variables is called the *Galois group* of the equation $x^4 = 2$. This group is the symmetric group of the square (which makes sense if you look at the location of α, β, γ , and δ in the Complex plane. It is also known as D_8 , the dihedral group of order 8.

In fact this Galois group can be defined more generally. Take the equation

$$x^2 + 1 = 0$$

which has roots $\alpha = i$, and $\beta = -i$. Now the conjugation map

$$\bar{} : \mathbb{C} \rightarrow \mathbb{C}, z \mapsto \bar{z},$$

acts as a permutation on the set of roots $\{\alpha, \beta\}$. It's not hard to prove that this permutation preserves the validity of all polynomial equations in α and β with real coefficients.

But now the conjugation map is, of course, a map of the whole field. This phenomenon is general: those permutations relating to the equation $x^4 - 2$ that we discussed above can be extended to yield maps $\mathbb{C} \rightarrow \mathbb{C}$. So, the Galois group is a set of maps between fields.

We're going to need to move from arm-waving to mathematics. Let's highlight some issues:

- What is the relation between \mathbb{R} and \mathbb{C} for the equation $x^2 + 1$? Or between \mathbb{Q} and \mathbb{C} for the equation $x^4 - 2$? More generally, what is a *field extension*?
- What about other fields?
- What do we mean by a “root of a polynomial”, when there is no root in that field?

2. SOME ALGEBRA

We assume a number of basic facts about the integers \mathbb{Z} , which I will not review here. Students should be familiar with unique factorization, the division algorithm, Euler's theorem, the Euclidean algorithm to find the highest common factor of two integers, and the Chinese remainder theorem.

2.1. A review of rings.

Definition 1. A **commutative ring with 1** is a set R equipped with two binary operations $+$ and \cdot such that

- $(R, +)$ is an abelian group;
- \cdot is associative and commutative;
- distributivity holds: for all $a, b, c \in R$

$$(a + b) \cdot c = (a \cdot c) + (b \cdot c), \text{ and}$$

$$a \cdot (b + c) = (a \cdot b) + (a \cdot c).$$

- there exists an element $1 \in R$ such that $a1 = 1a = a$ for all $a \in R$.

We will write ab for $a \cdot b$, and label the additive identity by 0. There are more general definitions of rings, but we will not need them. Thus whenever we refer to a **ring** we mean a commutative ring with 1.

Definition 2. We say that a set R is an **integral domain** if it is a ring such that

- If $ab = 0$ for $a, b \in R$, then either $a = 0$ or $b = 0$.

A set R is a **field** if it is a ring such that

- $F \setminus \{0\}$ is an abelian group under \cdot .

Note that then 1 is the identity of $F \setminus \{0\}$; what is more we often write F^* for $F \setminus \{0\}$.

We recall the definitions of *subring*, *subfield*, and *ideal*. If a ring R has an ideal I , then we can form the *quotient ring* R/I .

Definition 3. Let R and S be two rings. A homomorphism ϕ between R and S is a map $R \rightarrow S$ such that, for all $a, b \in R$,

$$\begin{aligned}\phi(a + b) &= \phi(a) + \phi(b), \text{ and} \\ \phi(ab) &= \phi(a)\phi(b).\end{aligned}$$

We have the consequent definitions of *epimorphism*, *monomorphism* and *isomorphism*. In addition, if we take $R = S$, then a homomorphism ϕ is called an *endomorphism*, while an isomorphism ϕ is called an *automorphism*.

Obviously all of the above definitions apply if R is not just a ring, but a field. *Field automorphisms* will be central objects in this course. (See the earlier arm-waving!)

Exercise 2.1. A field is an integral domain.

Exercise 2.2. $\mathbb{Q}, \mathbb{R}, \mathbb{C}$ are fields; \mathbb{Z} is an integral domain.

Exercise 2.3. Let $R = \mathbb{Z}$, and consider the ideal $I = n\mathbb{Z} = \{nz \mid z \in \mathbb{Z}\}$. Then R/I is the ring $\mathbb{Z}/n\mathbb{Z}$, the ring of integers modulo n . We can think of elements of $\mathbb{Z}/n\mathbb{Z}$ as being elements in the set

$$\{0, 1, \dots, n-1\},$$

with the ring operations being defined by

$$ab = ab \pmod{n}, \quad a + b = (a + b) \pmod{n}.$$

Then $\mathbb{Z}/n\mathbb{Z}$ is a field if and only if n is a prime number. For n prime we write \mathbb{F}_n for $\mathbb{Z}/n\mathbb{Z}$.

Exercise 2.4. Let $F = \{a + b\sqrt{2} \mid a, b \in \mathbb{Q}\}$. Then F is a field.

Exercise 2.5. Let $\phi : R \rightarrow S$ be a ring homomorphism. Then $\ker(\phi)$ is an ideal of R , $\text{im}(S)$ is a subring of S , and, moreover,

$$R/\ker(\phi) \cong \text{im}(S).$$

2.2. The characteristic.

Definition 4. Let K be a field. The **prime subfield** of K is the intersection of all subfields of K .

Note that the prime subfield must contain the elements 0 and 1.

Theorem 2.6. [Ste03, p.3] Every prime subfield is isomorphic to \mathbb{Q} or to the field \mathbb{F}_p , for p some prime.

If K has prime subfield isomorphic to \mathbb{Q} , then we say K has *characteristic 0*. If K has prime subfield isomorphic to \mathbb{F}_p , then we say that K has *characteristic p* .

Exercise 2.7. If K is a subfield of L , then K and L have the same characteristic.

If n is an integer, and k an element of the field K , then we write nk to mean

$$\underbrace{k + \dots + k}_n = \underbrace{(1 + \dots + 1)}_n k.$$

Exercise 2.8. If K is a non-zero element of the field K , and if n is an integer such that $nk = 0$, then n is a multiple of the characteristic of K .

2.3. Fields of fractions.

Definition 5. A **field of fractions** for a ring R is a field K containing a subring R' isomorphic to R , such that every element of K can be expressed in the form $\frac{r}{s}$ for $r, s \in R'$, where $s \neq 0$.

Note that this definition generalizes the relationship of the field \mathbb{Q} to the ring \mathbb{Z} . Just as in that example, one shouldn't assume that an element has a unique expression $\frac{r}{s}$ (for instance $\frac{1}{2} = \frac{2}{4}$ in \mathbb{Q}).

Theorem 2.9. [Ste03, p.5] Every integral domain possesses a field of fractions.

3. POLYNOMIALS

Definition 6. Let R be a ring. A **polynomial over R in the indeterminate t** is an expression

$$(3.1) \quad a_n t^n + a_{n-1} t^{n-1} + \cdots + a_1 t + a_0,$$

in which $n \in \mathbb{Z}^+$, and $a_0, \dots, a_n \in R$. The quantity a_i is called the **coefficient of t^i** ; note that if $a_i = 0$ then we often omit it.

If $i > n$, then we sometimes like to think that a_i exists, and is equal to 0. This allows us to write polynomials as $\sum a_i t^i$ where the sum is considered over all integers $i \geq 0$.

Given two polynomials $\sum a_i t^i$ and $\sum b_i t^i$, we define

$$\begin{aligned} \sum a_i t^i + \sum b_i t^i &= \sum (a_i + b_i) t^i, \text{ and} \\ \sum a_i t^i \cdot \sum b_i t^i &= \sum c_i t^i, \end{aligned}$$

where $c_j = \sum_{h+i=j} a_h b_i$. The set of all polynomials over R in indeterminate t is a ring, $R[t]$, the *ring of polynomials over R in the indeterminate t* .

Exercise 3.1. Let $R = \mathbb{Z}$. Define $f = 3t + 1$, $g = 4t^3 + 2t - 7$ calculate $f + g$, and fg .

Now take $f \in R[t]$, $f \neq 0$; it can be written in the form (3.1), with $a_n \neq 0$. Then we say that f has *degree n* , write $\deg f = n$. We call a_n the *leading coefficient* of f ; if $a_n = 1$, then we call f *monic*. (If $f = 0$ then, by convention, we say that $\deg f = -\infty$.)

Note that there is a natural ring monomorphism

$$R \rightarrow R[t], \lambda \mapsto 0t + \lambda.$$

This embeds R into $R[t]$ and the image of R is just the set of *constant polynomials*.

Lemma 3.2. Let R be an integral domain. Then $R[t]$ is an integral domain.

Proof. Take $f, g \in R[t]$ non-zero. Then they have nonzero leading coefficients, call these a and b . The product fg in $R[t]$ has leading coefficient ab . Since R is an integral domain we know that $ab \neq 0$; we conclude therefore that $fg \neq 0$, as required. \square

3.1. The Euclidean algorithm. Throughout this section K is a field. In this section we establish a number of facts about $K[t]$ that mirror properties of \mathbb{Z} with which we are (hopefully!) already familiar. We start with a simple observation:

Exercise 3.3. *If R is an integral domain and $f, g \in R[t]$, then*

$$\begin{aligned}\deg(f + g) &\leq \max\{\deg f, \deg g\}; \\ \deg(fg) &= \deg f + \deg g.\end{aligned}$$

(Why the inequality in the first line?)

Proposition 3.4. (Division algorithm) *Let f and g be polynomials over a field K , and suppose that f is non-zero. Then there exist unique polynomials q and r over K such that*

$$g = qf + r$$

and $\deg r < \deg f$.

Proof. We start by showing that q and r exist; we proceed by induction on $\deg g$. If $\deg g < \deg f$ then $g = 0f + g$, and we are done. If $\deg f = 0$ then $f = j \in K$ and we may take $q = k/j$ and $r = 0$. So assume that $0 < \deg f \leq \deg g$.

Thus, we take $m \leq n$, and

$$\begin{aligned}f &= a_m t^m + \cdots + a_0, \\ g &= b_n t^n + \cdots + b_0,\end{aligned}$$

where $a_m \neq 0 \neq b_n$. Let

$$g_1 = g - \frac{b_n}{a_m} t^{n-m} f.$$

Then $\deg g_1 < \deg g$ since leading terms cancel. Thus, by the inductive hypothesis,

$$g_1 = q_1 f + r_1,$$

with $\deg r_1 < \deg f$. But then

$$g = g_1 + b_n a_m^{-1} t^{n-m} f = (q_1 + b_n a_m^{-1} t^{n-m}) f + r_1,$$

and, since $\deg r_1 < \deg f$, we have the desired representation.

We must now prove uniqueness. Suppose that

$$g = f q_1 + r_1 = f q_2 + r_2, \text{ where } \deg r_1, \deg r_2 < \deg f.$$

Then

$$f(q_1 - q_2) = r_2 - r_1.$$

Now we use Ex. 3.3. If $q_1 \neq q_2$, then the left-hand side of this equation is a non-zero polynomial with degree at least $\deg f$. The right-hand side has degree strictly less than $\deg f$, thus we have a contradiction.

We conclude that $q_1 = q_2$ and, therefore, that $r_1 = r_2$. Thus q and r are unique. \square

Proposition 3.5. *If K is a field then $K[t]$ is a principal ideal domain.*

Proof. Let I be a non-trivial ideal of $K[t]$; we need to prove that $I = (f) = \{fg \mid g \in K[t]\}$ for some $f \in K[t]$. Let f be a non-zero polynomial of minimal degree in I . If $g \in I$ then the Division Algorithm implies that there exists $q, r \in K[t]$ such that $g = qf + r$ with $\deg(r) < \deg(f)$. Now $r = g - qf \in I$ and, since f is of minimal degree, we conclude that $r = 0$. In other words $f \mid g$ and $g \in (f)$. Thus $I \subseteq (f)$; since $(f) \subseteq I$ we conclude that $I = (f)$ as required. \square

Take $f, g \in K[t]$. We say that f divides g and write $f \mid g$, when there exists $h \in K[t]$ such that $g = fh$. If f does not divide g then we write $f \nmid g$.

Definition 7. A polynomial $d \in K[t]$ is a **highest common factor (h.c.f.)** of f and g if $d \mid f$ and $d \mid g$ and further, whenever $e \mid f$ and $e \mid g$ we have $e \mid d$. We often write (f, g) for a highest common factor of f and g . If 1 is a highest common factor of f and g then we call f and g **relatively prime**.

Exercise 3.6. Take $d, e, f, g \in K[t]$. If d is an h.c.f. of f and g , and $k \in K^*$, then kd is also a h.c.f. of f and g .

If d and e are two h.c.f.'s of f and g , then there exists $k \in K^*$ such that $e = kd$.

The polynomial d is a common factor of f and g of largest degree if and only if d is a h.c.f. of f and g .

Theorem 3.7. (Euclidean algorithm for polynomials) Let $f, g \in K[t]$. Write $r_{-1} = g, r_0 = f$ and define q_i and r_i for $i \geq 1$ via the relation

$$(3.2) \quad r_{i-2} = q_i r_{i-1} + r_i,$$

in which $\deg r_i < \deg r_{i-1}$. Then, for some non-negative integer I , one has $r_I = 0$ and then a highest common factor of f and g is r_{I-1} .

Proof. Clearly the algorithm is well-defined: the sequence $(\deg r_i)_{i \geq 0}$ is decreasing and bounded below by zero; thus the algorithm terminates with an integer I such that $r_I = 0$.

Next we prove that r_{I-1} divides both f and g by examining equation (3.2) for different values of i . Setting $i = I$, we observe that, since $r_I = 0$, r_{I-1} divides r_{I-2} . Setting $i = I - 1$, this in turn implies that r_{I-1} divides r_{I-3} . Continuing in this way we find that r_{I-1} divides $r_0 = g$ and $r_{-1} = f$.

Now suppose that $e \mid f$ and $e \mid g$. Then (3.2) implies that e divides r_1, r_2, \dots, r_{I-1} . Thus r_{I-1} is an h.c.f. of f and g . \square

This result is sometimes expressed in the literature as: $K[t]$ is a *Euclidean Domain*. Let us work through an example: we find a highest common factor of $g = t^5 - 3t + 2$ and $f = t^2 - 2t + 1$ in $\mathbb{Q}[t]$:

$$\begin{aligned} t^5 - 3t + 2 &= (t^3 + 2t^2 + 3t - 4)(t^2 - 2t + 1) + (2t + 2) \\ t^2 - 2t + 1 &= \left(\frac{1}{2}t - \frac{1}{2}\right)(2t - 2) \end{aligned}$$

Thus $2t - 2$ is a h.c.f. of f and g in $\mathbb{Q}[t]$. In fact $t - 1$ is also a h.c.f. of f and g .

Exercise 3.8. Take f and g in the example but consider them as polynomials in $\mathbb{F}_2[t]$. Now find a h.c.f.

Theorem 3.9. *Let d be a h.c.f. of f and g in $K[t]$. Then there exist polynomials $a, b \in K[t]$ such that*

$$d = af + bg.$$

Proof. We apply the Euclidean algorithm in reverse: as before write $r_{-1} = g, r_0 = f$ and define q_i and r_i for $i \geq 1$ via the relation

$$(3.3) \quad r_{i-2} = q_i r_{i-1} + r_i,$$

in which $\deg r_i < \deg r_{i-1}$. Let I be the smallest non-negative integer such that $r_I = 0$; then r_{I-1} is a highest common factor of f and g . Since h.c.f.'s are unique up to constant factors we may assume that $d = r_{I-1}$.

Now by (3.3) with $i = 1$, r_1 is a linear combination of r_{-1} and r_0 . Indeed for $i > 1$, r_i is a linear combination of r_{i-1} and r_{i-2} . Thus, by induction, every r_i is a linear combination of r_{-1} and r_0 .

In particular this is true for $i = I - 1$. Thus there are polynomials $a, b \in K[t]$ such that

$$r_{I-1} = af + bg.$$

□

Note that, for given f and g , this proof gives a method for calculating a and b .

3.2. Factorization of polynomials. Throughout this section R is a ring, K a field.

Definition 8. *An irreducible polynomial in $R[t]$ is a non-constant polynomial which cannot be written as a product of two polynomials of smaller degree. A polynomial which is not irreducible is reducible.*

Note that all polynomials of degree 0 are reducible, while those of degree 1 are irreducible. Beware: some texts define degree 0 polynomials to be irreducible.

The irreducible polynomials in $R[x]$ are analogous to the prime numbers in \mathbb{Z} . This should become clear as we prove various prime-number type properties...

Exercise 3.10. *Is the polynomial $f = t^2 - 2 \in \mathbb{Q}[t]$ reducible? Is the polynomial $f = t^2 - 2 \in \mathbb{R}[t]$ reducible?*

Lemma 3.11. *Suppose that $f, g, h \in K[t]$ and that f is irreducible. If $f \mid gh$ then either $f \mid g$ or $f \mid h$.*

Proof. Suppose that $f \nmid g$. A h.c.f. of f and g divides f which is irreducible. Thus such an h.c.f. must equal λ for some $\lambda \in K^*$. But then there exist $a, b \in K[t]$ such that

$$\lambda = af + bg.$$

Then

$$h = \lambda^{-1}(af + bg)h = (\lambda^{-1}ah)f + \lambda^{-1}b(gh).$$

Now, since f divides gh , we conclude that f divides the term on the right-hand side of the equation. Thus $f \mid h$. □

Theorem 3.12. (Unique factorization theorem) *If $f \in K[t]$, and $\deg f \geq 1$, then f can be written as a product of irreducible polynomials in $K[t]$, and this factorization is unique apart from the order of the irreducible factors and the presence of constant factors.*

Proof. Any polynomial $f \in K[t]$ can be written as a product of irreducible polynomials. For, if it is not irreducible, it is a product of polynomials of positive degree, and so the conclusion follows by induction.

Now we need to prove that this product is unique up to ordering and constant factors. Suppose that $f = g_1 \cdots g_r = h_1 \cdots h_s$, with each of g_i, h_j irreducible. Lem. 3.11 implies that $g_1 \mid h_1$ for some $1 \leq i \leq s$. Without loss of generality assume that $g_1 = h_1$, so $g_1 = \lambda_1 g_1$ for some $\lambda_1 \in K^*$. Then

$$\begin{aligned} (\lambda_1 g_1) h_2 \cdots h_s &= g_1 \cdots g_r \\ \implies g_1 (\lambda_1 h_2 \cdots h_s - g_2 \cdots g_r) &= 0. \end{aligned}$$

Since $K[t]$ is an integral domain we conclude that

$$g_2 \cdots g_r = (\lambda_1 h_2) h_3 \cdots h_s.$$

An inductive argument implies that we can choose an ordering such that $h_i = \lambda_i g_i$ for $1 \leq i \leq r$. We then obtain

$$h_{r+1} \cdots h_s = \lambda,$$

a constant. Thus $r = s$, and our factorization is unique up to constant factors. \square

This result is sometimes expressed in the literature as: $K[t]$ is a *Unique Factorization Domain (UFD)*. This is a broader class of objects than the Euclidean domains mentioned earlier.

Corollary 3.13. *If $f \in K[t]$ is monic, and $\deg f \geq 1$, then f can be written as a product of irreducible monic polynomials in $K[t]$, and this factorization is unique apart from the order of the irreducible factors.*

Exercise 3.14. *Prove this corollary!*

3.3. Zeros of polynomials.

Definition 9. Let R be a ring, $f \in R[t]$. We say that $\alpha \in R$ is a **zero** or a **root** of $f \in R[t]$ if $f(\alpha) = 0$.

Lemma 3.15. Let K be a field and suppose that $\alpha \in K$ and $f \in K[t]$. Then $f(\alpha) = 0$ if and only if $(t - \alpha) \mid f$.

Proof. Suppose, first of all, that $(t - \alpha) \mid f$; then $f = (t - \alpha)g$ for some $g \in K[t]$. But then $f(\alpha) = (\alpha - \alpha)g(\alpha) = 0$.

On the other hand suppose that $f(\alpha) = 0$. By the division algorithm there exist $q, r \in K[t]$ such that $f = q(t - \alpha) + r$ where $\deg r < 1$, i.e. $r = \lambda$ for some $\lambda \in K$. Then

$$0 = f(\alpha) = q(\alpha)(t - \alpha) + \lambda = \lambda.$$

Thus $f = q(t - \alpha)$ and $(t - \alpha) \mid f$. \square

We will often be concerned with the question of when a given polynomial f is irreducible. Lem. 3.15 gives us one useful test.

Exercise 3.16. *Show that the polynomial $f(t) = t^3 + t + 2$ is irreducible in $\mathbb{F}_3[t]$.*

In general it is very difficult to tell whether or not a given polynomial is irreducible. There are a myriad of results like Lem. 3.15 that give partial answers to this question in particular situations. See [Ste03, §2.2] for a number of such results; we note one in particular:

Proposition 3.17. [Ste03, Prop. 2.4] *Let f be a polynomial in $\mathbb{Z}[t]$ which is irreducible over \mathbb{Z} . Then, f considered as a polynomial in $\mathbb{Q}[t]$ is irreducible over \mathbb{Q} .*

4. FIELD EXTENSIONS

We begin with a construction that will be very important throughout the rest of the course. Take a polynomial $f \in K[t]$, where K is some field. Consider the quotient ring $K[t]/(f)$ (by (f) we mean the ideal $\{fg \mid g \in K[t]\}$).

Lemma 4.1. *If f is irreducible then $K[t]/(f)$ is a field.*

Proof. We need to check the field axioms. The only one that presents a serious challenge is to show that every non-zero element of $K[t]/(f)$ has an inverse.

Take $g \in K[t]$. Then g is in a non-zero residue class of $K[t]/(f)$ if and only if f does not divide g . But now, since f is irreducible, this implies that any highest common factor of f and g will be a constant.

Then Theorem 3.9 implies that there exist polynomials a and b such that

$$d = af + gb,$$

where d is a constant. Then

$$\begin{aligned} 1 &= (d^{-1}a)f + (d^{-1}b)g \\ \implies (d^{-1}a)f &= 1 \pmod{f}. \end{aligned}$$

□

What does $K[t]/(f)$ look like? Assume that f is irreducible of degree d . The Division Algorithm implies that every residue class in $K[t]/(f)$ contains **precisely one** polynomial of degree at most $d - 1$. Thus the set of residues is in 1-1 correspondence with the set

$$\{a_0 + a_1t + \cdots + a_{d-1}t^{d-1} \mid a_i \in K\}.$$

Exercise 4.2. *Prove that, as an additive group, $K[t]/(f)$ is isomorphic to K^d .*

Note that the set of special residue classes which contain the constant polynomials a_0 (for some $a_0 \in K$) form a field isomorphic to K . In other words the field K lies inside the field $K[t]/(f)$. This motivates our next definition.

Definition 10. *A field extension of a field K is a triple (i, K, L) where L is a field, and i is a ring monomorphism of K into L . We identify the image $i(K)$ with K , and think of K lying inside L ; we then denote the field extension $L : K$, or L/K (omitting mention of i).*

Some examples:

- The field extension $\mathbb{C} : \mathbb{R}$ is formally given by

$$i : \mathbb{R} \rightarrow \mathbb{C}, x \mapsto x + 0i.$$

- Likewise with $\mathbb{C} : \mathbb{Q}$, or $\mathbb{R} : \mathbb{Q}$.
- The construction we considered above yielded a field extension $K[t]/(f) : K$.

4.1. Degree.

Theorem 4.3. *Let $L : K$ be a field extension. Under the operations*

$$\begin{aligned} L \times L &\rightarrow L, (l_1, l_2) \mapsto l_1 + l_2, \text{ and} \\ K \times L &\rightarrow L, (k, l) \mapsto kl, \end{aligned}$$

L forms a vector space over K .

The first operation is *vector addition*, the second *scalar multiplication*. So in this vector space the scalars are elements of K , the vectors are elements of L .

Proof. Check that the axioms for a vector space are satisfied by $L : K$. □

So in all the examples of field extensions that we gave earlier, we were also listing vector spaces!

Definition 11. *The **degree of an extension** $L : K$ is the dimension of L as a vector space over K . We write $[L : K]$ for the degree of $L : K$. We call the extension **finite** if $[L : K] < \infty$, and **infinite** if $[L : K] = \infty$.*

Exercise 4.4. *Calculate the degree of the extensions that were listed above. The only (slightly) hard one is the degree of the extension $K[t]/(f)$.*

Theorem 4.5. *Suppose that $M : L$ and $L : K$ are field extensions. Then $M : K$ is a field extension, and*

$$(4.1) \quad [M : K] = [M : L][L : K].$$

Proof. That $M : K$ is a field extension is immediate (compose the inclusion maps given by $M : L$ and $L : K$). We need to prove the identity (4.1).

Suppose first that the RHS is finite. Thus $[M : L] = m \in \mathbb{Z}^+$ and $[L : K] = n \in \mathbb{Z}^+$. Let $\{e_1, \dots, e_m\}$ be a basis for M over L , and let $\{f_1, \dots, f_n\}$ be a basis for L over K . We claim that the lm elements $e_i f_j, 1 \leq i \leq m, 1 \leq j \leq n$, form a basis for M over K . This claim implies immediately that

$$[M : K] = lm = [L : K][K : M],$$

as required.

We prove first that

$$M \subseteq \text{span}\{e_i f_j \mid 1 \leq i \leq m, 1 \leq j \leq n\}.$$

Let z be an element of M . Then, since $\{e_1, \dots, e_m\}$ is a basis for M over L ,

$$z = a_1 e_1 + \dots + a_m e_m,$$

for elements a_1, \dots, a_m in L . Now, since $\{f_1, \dots, f_n\}$ is a basis for L over K , we have for each $i = 1, \dots, m$,

$$a_i = b_{i1} f_1 + \dots + b_{in} f_n,$$

for elements $b_i 1, \dots, b_i n$ in K . This implies that

$$z = \sum_{i=1}^m \sum_{j=1}^n b_{ij} e_i f_j,$$

as required.

Next we prove linear independence. Suppose that

$$\sum_{i=1}^m \sum_{j=1}^n c_{ij} e_i f_j = 0$$

for some elements c_{ij} , $1 \leq i \leq m$, $1 \leq j \leq n$. This implies that

$$\left(\sum_{i=1}^m c_{i1} e_i \right) f_1 + \dots + \left(\sum_{i=1}^m c_{in} e_i \right) f_n = 0.$$

Since $\{f_1, \dots, f_n\}$ is a basis for L over K this implies that $\sum_{i=1}^m c_{ij} e_i = 0$ for each $j = 1, \dots, n$. But now, since $\{e_1, \dots, e_m\}$ is a basis for M over L , this implies that $c_{ij} = 0$ for all i, j . Linear independence is proved.

Suppose next that the RHS is not finite; we need to show that the LHS is not finite. Suppose that, instead, $[M : K] = n < \infty$. Then we can find a basis $\{g_1, \dots, g_n\}$ for M over K . But $\{g_1, \dots, g_n\}$ spans M over K , so it certainly spans M over L . Thus $[M : L] < \infty$. Moreover L is a K -linear subspace of M so $[L : K] < \infty$. Then $[M : L][L : K] < \infty$ which is a contradiction, as required. \square

Definition 12. A sequence $K_n : K_{n-1}, K_{n-1} : K_{n-2}, \dots, K_1 : K_0$ is called a **tower** of field extensions, and we write it as $K_n : K_{n-1} : \dots : K_1 : K_0$.

Exercise 4.6. Generalize the previous theorem to towers of arbitrary length. Thus, for a tower, $K_n : K_{n-1} : \dots : K_1 : K_0$, prove that

$$[K_n : K_{n-1} : \dots : K_1 : K_0] = [K_n : K_{n-1}][K_{n-1} : K_{n-2}] \cdots [K_1 : K_0].$$

Exercise 4.7. Suppose that $L : K$ is a field extension with $[L : K]$ prime. Then there exists no subfield M of L containing K other than K itself.

4.2. Simple extensions.

Definition 13. Suppose that $L : K$ is a field extension and $A \subset L$. Write $K(A)$ for the intersection of all subfields of L that contain K and A . We call $K(A)$ the extension of K generated by A . If $A = \{a_1, \dots, a_n\}$ then write $K(a_1, \dots, a_n)$ for $K(A)$. We call $L : K$ a **simple extension** if $L = K(\alpha)$ for some $\alpha \in L$; we say L is **finitely generated over K** if $L = K(A)$ for a finite set A .

Note that $K(A)$ is a field, in fact the smallest subfield of L containing both K and A . Some examples:

- $\mathbb{C} = \mathbb{R}(i)$, a simple extension.
- The set $\{a + b\sqrt{2} \mid a, b \in \mathbb{Q}\}$ which we encountered earlier is a simple field extension of \mathbb{Q} ; it is $\mathbb{Q}(\sqrt{2})$.
- $\mathbb{C} : \mathbb{Q}$ is not simple (since any simple extension of a countable field is countable).

- The field extension $\mathbb{Q}(\sqrt{2}, \sqrt{3})$ is simple even though it does not appear to be! To see this, define $\alpha = \sqrt{2} + \sqrt{3}$. Then

$$\alpha^2 = 5 + 2\sqrt{6},$$

and, in particular, $\sqrt{6} = \frac{1}{2}(\alpha^2 - 5) \in \mathbb{Q}(\alpha)$. This implies that

$$(\sqrt{6} - 2)\alpha = \sqrt{2} \in \mathbb{Q}(\alpha);$$

$$(3 - \sqrt{6})\alpha = \sqrt{3} \in \mathbb{Q}(\alpha).$$

Thus $\mathbb{Q}(\sqrt{2}, \sqrt{3}) \subseteq \mathbb{Q}(\alpha) \subseteq \mathbb{Q}(\sqrt{2}, \sqrt{3})$; in other words $\mathbb{Q}(\sqrt{2}, \sqrt{3}) = \mathbb{Q}(\alpha)$, a simple extension of \mathbb{Q} .

- It is useful to imagine what a field extension looks like in general. Consider the previous example, $\mathbb{Q}(\sqrt{2}, \sqrt{3})$. Observe that this must contain the set

$$S = \{ a + b\sqrt{2} + c\sqrt{3} + d\sqrt{6} \mid a, b, c, d \in \mathbb{Q} \}.$$

If we can prove that S is a field then, by definition, $S = \mathbb{Q}(\sqrt{2}, \sqrt{3})$. In fact, proving this is easy; an immediate consequence is that $[\mathbb{Q}(\sqrt{2}, \sqrt{3}) : \mathbb{Q}] = 4$.

You might also like to consider the example $\mathbb{Q}(\sqrt[3]{2})$. Can you find a “natural” basis for this extension? What is its degree? We will give a result later that answers this sort of question for general field extensions $K(A)$.

4.3. Algebraic extensions.

Definition 14. Suppose that $L : K$ is a field extension, and that $\alpha \in L$. If there is a polynomial $f = a_n t^n + \cdots a_1 t + a_0 \in K[t]$ for which $f(\alpha) = 0$, then we say that α is **algebraic** over K . If no such polynomial exists, then we say that α is **transcendental** over K .

Observe that, since L is a vector space over K , the element α is algebraic if the set $\{1, \alpha, \dots, \alpha^n\}$ is linearly independent for some natural number n . Conversely α is transcendental if $\{1, \alpha, \dots, \alpha^n\}$ is linearly independent for each $n \in \mathbb{N}$.

It is a fact (not easily proved) that π , e and $2^{\sqrt{2}}$ are transcendental over \mathbb{Q} . In fact the set of transcendental elements in \mathbb{R} has measure 1. (When applied to real numbers, the adjective *transcendental* is always relative to \mathbb{Q} unless otherwise stated.)

Definition 15. Suppose that $L : K$ is a field extension and $\alpha \in L$. We define the **evaluation map**,

$$E_\alpha : K[t] \rightarrow L, f \mapsto f(\alpha).$$

Observe that E_α is a ring homomorphism. What is more

- α is transcendental if and only if $\ker(E_\alpha) = \{0\}$;
- α is algebraic if and only if $\ker(E_\alpha) \neq \{0\}$.

Since $K[t]$ is a principal ideal domain we know that $\ker(E_\alpha) = (f)$, for some polynomial f . If α is algebraic, then f is non-zero; in fact we may take f to be monic, and this defines f uniquely (see the proof of Prop. 3.5.)

Definition 16. The **minimal polynomial** of the element $\alpha \in L$ from a field extension $L : K$ is the monic polynomial m_α such that $\ker(E_\alpha) = (m_\alpha)$.

The next theorem connects the two primary constructions of field extensions that we have seen: $K[t]/(f)$ for irreducible f , and $K(\alpha)$. First of all we need to formally define what it means for two field extensions to be the same.

Definition 17. Let (i_1, K_1, L_1) and (i_2, K_2, L_2) be two field extensions. An **isomorphism of field extensions** is a pair (λ, μ) of field isomorphism $\lambda : K_1 \rightarrow K_2$, $\mu : L_1 \rightarrow L_2$ such that the following diagram commutes:

$$\begin{array}{ccc} K_1 & \xrightarrow{i_1} & L_1 \\ f \downarrow & & \downarrow g \\ K_2 & \xrightarrow{i_2} & L_2. \end{array}$$

Of course, most of the time we do not think of field extensions in terms of monomorphisms; instead we identify $i_j(K_j)$ and K_j , so that K_j is just a *subfield* of L_j ($j = 1, 2$). In this formulation, an isomorphism of field extensions is a field isomorphism $\mu : L_1 \rightarrow L_2$ such that

$$\mu|_{K_1} = K_2.$$

Theorem 4.8. Suppose that $L : K$ is a field extension, and that $\alpha \in L$ is algebraic. Then

- (a) the minimal polynomial m_α of α is irreducible in $K[t]$;
- (b) the image $E_\alpha(K[t])$ of the polynomial ring $K[t]$ is the subfield $K(\alpha)$ of L ;
- (c) the evaluation map E_α factorizes as $i\tilde{E}_\alpha q$ in the following way:

$$\begin{array}{ccc} K[t] & \xrightarrow{E_\alpha} & L \\ q \downarrow & & \uparrow i \\ K[t]/(m_\alpha) & \xrightarrow{\tilde{E}_\alpha} & K(\alpha). \end{array}$$

Here q is the quotient map, i is the inclusion map, and \tilde{E}_α is an isomorphism of field extensions.

Proof. (a) Suppose that $m_\alpha = fg$. Then applying the evaluation map we obtain

$$0 = E_\alpha(m_\alpha) = E_\alpha(f)E_\alpha(g) = f(\alpha)g(\alpha).$$

Then, without loss of generality, $f(\alpha) = 0$. Thus $f \in (m_\alpha)$ and so $m_\alpha \mid f$. In particular $\deg(m_\alpha) = \deg(f)$, and m_α is irreducible.

- (b) The field $E_\alpha(K[t])$ is a subfield of L . Furthermore $E_\alpha(k) = k$ for $k \in K$, and $E_\alpha(t) = \alpha$. Thus $K \cup \{\alpha\} \subseteq \text{im}(E_\alpha)$ and so $K(\alpha) \subseteq \text{im}(E_\alpha)$. On the other hand it is obvious that $\text{im}(E_\alpha) \subseteq K(\alpha)$, so we are done.
- (c) Since m_α is irreducible, Lem. 4.1 implies that $K[t]/(m_\alpha)$ is a field. The ring isomorphism theorem implies that one has the following factorization:

$$\begin{array}{ccc} K[t] & \xrightarrow{E_\alpha} & L \\ q \downarrow & & \uparrow i \\ K[t]/\ker(E_\alpha) & \xrightarrow{\cong} & \text{im}(E_\alpha), \end{array}$$

which yields the result. □

Recall that, given an irreducible polynomial $f \in K[t]$, we have a good understanding of what the field $K[t]/(f)$ looks like: its elements are in 1-1 correspondence with polynomials in t of degree strictly smaller than f . This information can now be used to understand the structure of field extensions of form $K(A)$, especially $K(\alpha)$: in particular, observe that the ring homomorphism $E_\alpha : K[t] \rightarrow L$ maps t to α . Following the proof above, we conclude that, if $\deg(m_\alpha) = n$, the elements of the field $K(\alpha)$ are in 1-1 correspondence with the set

$$\{a_{n-1}\alpha^{n-1} + \cdots + a_1\alpha + a_0 \mid a_0, \dots, a_{n-1} \in K\}.$$

For instance, recall the example of a field

$$F = \{a + b\sqrt{2} \mid a, b \in \mathbb{Q}\}.$$

We have already seen that $F = \mathbb{Q}(\sqrt{2})$. One can easily check that, for $\alpha = \sqrt{2}$, the minimal polynomial over \mathbb{Q} is $m_\alpha = t^2 - 2$. Thus $F \cong \mathbb{Q}[t]/(t^2 - 2)$ and we also have a 1-1 correspondence between F and the set

$$\{a + bt \mid a, b \in \mathbb{Q}\}.$$

These different representations of the same field have their advantages, mostly when it comes to multiplying elements.

Theorem 4.9. *Suppose that $L : K$ is a field extension, and that $\alpha \in L$. Then α is algebraic over K if and only if $[K(\alpha) : K] < \infty$, in which case $[K(\alpha) : K] = \deg(m_\alpha)$.*

Proof. Suppose that $[K(\alpha) : K] = n < \infty$. Now $K(\alpha)$ is a vector space over K of dimension n and so the elements $1, \alpha, \alpha^2, \dots, \alpha^n$ are linearly independent in $K(\alpha) : K$. Thus, for some $k_0, \dots, k_n \in K$ not all zero, we have

$$k_n\alpha^n + \cdots k_1\alpha + k_0 = 0.$$

Then the polynomial $k_nt^n + \cdots k_1t + k_0 \in \ker(E_\alpha)$. Thus $\ker(E_\alpha) \neq \{0\}$, and so α is algebraic over K .

Now for the converse: suppose that α is algebraic over K , and m_α is its minimal polynomial. Let $n = \deg(m_\alpha)$ and suppose that $1, \alpha, \dots, \alpha^{n-1}$ are linearly dependent over K . Then there exist $k_0, \dots, k_{n-1} \in K$, not all zero, satisfying

$$k_{n-1}\alpha^{n-1} + \cdots + k_1\alpha + k_0 = 0.$$

This implies that $\deg(m_\alpha) < n$ which is a contradiction. We conclude that $1, \alpha, \dots, \alpha^{n-1}$ are linearly independent over K .

The result follows if we can show that the set $\{1, \alpha, \dots, \alpha^{n-1}\}$ spans $K(\alpha)$, since this set will then be a basis of $K(\alpha)$. So take $\beta \in K(\alpha)$ and observe that $\beta = E_\alpha(f)$ for some $f \in K[t]$. We can write $f = m_\alpha q + r$ where $\deg r < n$. Then

$$\beta = E_\alpha(f) = E_\alpha(m_\alpha)E_\alpha(q) + E_\alpha(r) = E_\alpha(r).$$

Thus if $r = k_{n-1}t^{n-1} + \cdots + k_1t + k_0$, then

$$\beta = k_{n-1}\alpha^{n-1} + \cdots + k_1\alpha + k_0 \in \text{span}\{1, \alpha, \dots, \alpha^{n-1}\}$$

as required. □

Definition 18. For $L : K$ a field extension define the set L^{alg} to be those elements of L which are algebraic over K .

Theorem 4.10. If $L : K$ is a field extension then L^{alg} is a subfield of L .

Proof. We need to check that $(L^{alg}, +)$ and $(L^{alg} \setminus \{0\}, \cdot)$ are abelian groups. Since they are sets inside larger abelian groups we just need to check closure under composition and inversion in each case.

Suppose that α and β are algebraic over K . Then $m_\beta \in K[t] \subseteq K(\alpha)[t]$. Thus β is algebraic over $K(\alpha)$. Then Thm. 4.9 implies that

$$[K(\alpha, \beta) : K(\alpha)] = [K(\alpha)(\beta) : K(\alpha)] < \infty.$$

But then

$$[K(\alpha, \beta) : K] = [K(\alpha, \beta) : K(\alpha)][K(\alpha) : K] < \infty.$$

Now observe that $K(\alpha + \beta)$, $K(-\alpha)$, and $K(\alpha\beta)$ are all subfields of $K(\alpha, \beta)$. Furthermore, if $\alpha \neq 0$, then $K(\alpha^{-1})$ is a subfield of $K(\alpha, \beta)$. Thus they are all finite extensions of K , and so $\alpha + \beta$, $-\alpha$, $\alpha\beta$, and α^{-1} are all algebraic as required. \square

Definition 19. A field extension $L : K$ is **algebraic** if every element of L is algebraic over K . i.e. $L = L^{alg}$.

Theorem 4.11. Suppose that $L : K$ is a field extension. The following are equivalent:

- (a) $[L : K] < \infty$;
- (b) $L : K$ is algebraic, and L is finitely generated over K ;
- (c) L is finitely generated over K

Proof. First we prove (a) \implies (b). Suppose that $[L : K] = s < \infty$. If $\alpha \in L$, then

$$[K(\alpha) : K] \leq [L : K] < \infty,$$

so α is algebraic over K . Thus $L : K$ is algebraic. Moreover if $\{\gamma_1, \dots, \gamma_s\}$ is a basis for L over K , then $L = K(\gamma_1, \dots, \gamma_s)$, and L is finitely generated over K .

(b) \implies (c) is trivial. So now prove (c) \implies (a). Suppose that $L = K(\alpha_1, \dots, \alpha_n)$. Define

$$K_0 = K, K_1 = K_0(\alpha_1), K_2 = K_1(\alpha_2), \dots, K_n = K_{n-1}(\alpha_n) = L.$$

Then $[K_j : K_{j-1}]$ is finite for $1 \leq j \leq n$, since α_j is algebraic in each case. Thus

$$[L : K] = [K_n : K_{n-1}] \cdots [K_2 : K_1][K_1 : K_0] < \infty,$$

as required. \square

Let $L : K$ be a field extension, $\alpha \in L$. We have a number of equivalent statements:

- (a) α is algebraic;
- (b) $f(\alpha) = 0$ for some non-zero $f \in K[t]$;
- (c) $\ker(E_\alpha) \neq \{0\}$;
- (d) α has a well-defined non-zero minimal polynomial;
- (e) $K(\alpha) : K$ is algebraic.

Of course one has a load of similar statements if one starts instead with α is transcendental.

Exercise 4.12. Suppose that $M : L$ and $L : K$ are algebraic extension. Then $M : K$ is algebraic.

We have proved a lot of results about algebraic extensions, and almost entirely neglected transcendental extensions. However that's because, when the extension is simple, there is only one transcendental extension. So everything's (kind of) easy!

Exercise 4.13. *Every simple transcendental extension $K(\alpha) : K$ is isomorphic to the extension $K(t) : K$ of rational expressions in an indeterminate t over K . The isomorphism can be chosen to carry t into α .*

5. RULER AND COMPASS CONSTRUCTIONS

This section can be thought of as a bit of a diversion from the main mathematical train of thought. But the results we discuss here are not only beautiful, they are also historically very important. Some of the problems we solve in this section were open for upwards of 2000 years.

Here is the set-up: Assume two points P_0 and P_1 are given in \mathbb{R}^2 . By rescaling and rotation, there is no loss of generality in supposing these points to be $P_0 = (0, 0)$ and $P_1 = (1, 0)$. We can then make the following definition:

Definition 20. *A point $P \in \mathbb{R}^2$ is **constructible** if there exists a sequence of points $P_0, P_1, \dots, P_n = P$ with the following property. Write $A_j = \{P_0, \dots, P_j\}$ for $2 \leq j \leq n$. Then for each j with $2 \leq j \leq n$, the point P_j is one of:*

- (a) *the intersection of two distinct lines, each joining two points in A_{j-1} ;*
- (b) *the intersection of a line L and a circle C where*
 - *L joins two points in A_{j-1} ;*
 - *the centre of C is a point of A_{j-1} and the radius of C is equal to the distance between two points of A_{j-1} ;*
- (c) *the intersection of two circles with distinct centres in A_{j-1} , and so that both radii are equal to the respective distances of two pairs of points of A_{j-1} .*

Now let us see how this relates to Galois theory:

Theorem 5.1. *If $P = (x, y)$ is a constructible point, then the field extension $\mathbb{Q}(x, y) : \mathbb{Q}$ is finite and satisfies $[\mathbb{Q}(x, y) : \mathbb{Q}] = 2^r$ for some natural number r .*

Proof. Since P is constructible, there exists a finite sequence, $P_0, P_1, \dots, P_n = P$ generated in the process describe above. For $j = 0, \dots, n$, write $P_j = (x_j, y_j)$ and

$$K_j = \mathbb{Q}(x_0, y_0, x_1, y_1, \dots, x_j, y_j).$$

Observe that $K_j = K_{j-1}(x_j, y_j)$. We make the following claim:

Claim: For $j = 2, \dots, n$, $[K_j : K_{j-1}] \in \{1, 2\}$.

Proof that the claim implies the theorem: Observe that $\mathbb{Q}(x, y)$ is a subfield of K_n . The claim implies that

$$[K_n : \mathbb{Q}] = [K_n : K_{n-1}][K_{n-1} : K_{n-2}] \cdots [K_1 : K_0] = 2^s,$$

for some natural number s . (Note that $K_0 = \mathbb{Q}$.) Thus $[\mathbb{Q}(x, y) : \mathbb{Q}] = 2^r$ for some natural number r , as required.

Proof of claim: The point P_j can be constructed in one of three different ways; we need to check that the claim holds for each of these methods.

- $P = L_1 \cap L_2$ for two lines L_1 and L_2 . Since L_i connects two points in K_{j-1} we conclude (check!) that the equation for L_i has form $\lambda_i x + \mu_i y + \nu_i = 0$ for some $\lambda_i, \mu_i, \nu_i \in K_j$, $i = 1, 2$.

Now the intersection of two such lines is a point with coordinates in K_j (check!), and so $[K_{j+1} : K_j] = 1$ in this case.

- $P = C \cap L$ for a circle C and line L . Now C has centre (a_1, b_1) and radius equal to the distance between the points (a_2, b_2) and (a_3, b_3) ; then the equation for C is

$$(x - a_1)^2 + (y - b_1)^2 = (a_3 - a_2)^2 + (b_3 - b_2)^2.$$

Since the three points (a_i, b_i) , $i = 1, 2, 3$, are in A_j we conclude that C has equation

$$x^2 + y^2 + \lambda x + \mu y + \nu = 0$$

for some $\lambda, \mu, \nu \in K_{j-1}$. Now, as we have seen, L has equation

$$\rho x + \sigma y + \tau = 0$$

for some $\rho, \sigma, \tau \in K_{j-1}$. Substitution yields an equation

$$x^2 + \alpha x + \beta = 0$$

for some $\alpha, \beta \in K_{j-1}$. If the quadratic on the LHS of this equation is reducible then x lies in K_{j-1} (it is a solution to a linear equation) and $[K_j : K_{j-1}] = 1$. If it is irreducible then $[K_j : K_{j-1}] = 2$ and we are done.

- $P = C_1 \cap C_2$ for two circles C_1 and C_2 . Then, just as in the previous case, for $i = 1, 2$, we have equations

$$x^2 + y^2 + \lambda_i x + \mu_i y + \nu_i = 0.$$

By elimination these intersect at the common solution of the equations

$$x^2 + y^2 + \lambda_1 x + \mu_1 y + \nu_1 = 0;$$

$$(\lambda_2 - \lambda_1)x + (\mu_2 - \mu_1)y + (\nu_2 - \nu_1) = 0.$$

The second equation defines a line with coefficients in K_{j-1} . Hence we are back in case (2) where, as we have seen, the claim holds.

□

We are now in a position to prove a number of “impossibility” theorems. Each of the next three results relates to a problem that was open for 2000 odd years! We begin with an exercise.

Exercise 5.2. Call a line “constructible” if it joins two constructible points. Suppose that two constructible lines, L_1 and L_2 , intersect at angle θ . Let Q be a constructible point, and M a constructible line through Q . Then a line N can be constructed that intersects M at Q , forming the angle θ .

Theorem 5.3. There are angles which cannot be trisected, such as $\frac{\pi}{3}$.

Proof. Let $\theta = \frac{\pi}{9}$. Suppose that $\frac{\pi}{3}$ can be trisected, then (by the exercise) there is a constructible point P other than $(0, 0)$ on the line L defined by,

$$y = (\tan \theta)x.$$

Now we can intersect L with the circle C of radius 1 and centre $(0, 0)$, thereby constructing the point $(\cos \theta, \sin \theta)$. Now there exists a natural number r such that

$$2^r = [\mathbb{Q}(\cos \theta, \sin \theta) : \mathbb{Q}] = [\mathbb{Q}(\cos \theta, \sin \theta) : \mathbb{Q}(\cos \theta)][\mathbb{Q}(\cos \theta) : \mathbb{Q}].$$

Thus, in particular, $[\mathbb{Q}(\cos \theta) : \mathbb{Q}] = 2^s$ for some natural number s .

But now observe that

$$\begin{aligned} \cos(3\theta) &= 4(\cos \theta)^3 - 3\cos \theta \\ \implies (2\cos \theta)^3 - 2(2\cos \theta) &= 2\cos(3\theta) = 2\cos\left(\frac{\pi}{3}\right) = 1. \end{aligned}$$

Thus $z = 2\cos \theta$ satisfies $z^3 - 3z - 1$. **Claim:** the polynomial $t^3 - 3t - 1$ is irreducible.

Proof that claim implies theorem. The claim implies that

$$[\mathbb{Q}(\cos \theta) : \mathbb{Q}] = [\mathbb{Q}(2\cos \theta) : \mathbb{Q}] = 3.$$

Thus, in particular, $[\mathbb{Q}(\cos \theta) : \mathbb{Q}]$ does not divide 2^s for $s > 1$ and we are done.

Proof of claim: If the polynomial $t^3 - 3t - 1$ is reducible, then it has a rational zero. Let $z = \frac{p}{q}$ be such a solution, expressed in lowest terms. Then $p^3 - 3q^2p - q^3 = 0 \implies p \mid q^3$, which implies that $p \mid q^3$. This implies that $p = 1$. Then we have $q \mid p^3$, and so $q = 1$. Thus $z = 1$; but 1 is not a zero of the given polynomial so we have a contradiction as required. \square

Note that this result is a proof that the angle $\frac{\pi}{3}$ cannot be trisected. On the other hand the angle π **can** be trisected!

Exercise 5.4. Describe the procedure for trisecting π with ruler and compass.

Theorem 5.5. The cube cannot be duplicated by ruler and compass.

Proof. “Duplicating the cube” means constructing a cube Q of volume 2. Suppose that we can do this: then Q has side length $\sqrt[3]{2}$, and so we are able to construct the point $P = (\sqrt[3]{2}, 0)$. Now $\sqrt[3]{2}$ is a root of the polynomial $f = t^3 - 2$. **Claim:** f is irreducible.

Proof that claim implies theorem: Since f is irreducible, $[\mathbb{Q}(\sqrt[3]{2}) : \mathbb{Q}] = 3$. Since 3 does not divide 2^s for any natural number s we conclude that P is not constructible.

Proof of claim: As before it is enough to show that f has no rational roots. It is easy enough to verify that $\sqrt[3]{2}$ is not rational, and we are done. \square

Theorem 5.6. We cannot “square the circle” using a ruler and compass.

Proof. “Squaring the circle” means drawing a square of area equal to π . Suppose we can do this; then we are able to construct two points of distance $\sqrt{\pi}$ apart; this in turn means that we can construct the point $(\sqrt{\pi}, 0)$. This implies that $[\mathbb{Q}(\sqrt{\pi}) : \mathbb{Q}] = 2^s$ for some natural number s . Since $\mathbb{Q}(\pi)$ is a subfield of $\mathbb{Q}(\sqrt{\pi})$ we conclude that $[\mathbb{Q}(\pi) : \mathbb{Q}] = 2^r$ for some natural number r .

But now π is transcendental (Lindemann, 1882) and we have a contradiction. \square

There is nothing particularly special about the rules of construction corresponding to a ruler and compass. Given any set of rules for constructing points, one can (try to) prove results similar to those above. For instance one can study constructibility problems corresponding to origami-type procedures, or to cord-and-nail procedures.

6. SOME RESULTS ON IRREDUCIBILITY

The previous section highlighted how important it is to know when a given polynomial $f \in R[t]$ is irreducible. In this section we outline a number of results that tell us when a polynomial is irreducible in $\mathbb{Z}[t]$. In fact analogues of (most of) these results can be proved in much greater generality.

Definition 21. The **content** of a polynomial $f = a_n t^n + \cdots + a_1 t + a_0 \in \mathbb{Z}[t]$ is defined to be the greatest common divisor of the coefficients: (a_n, \dots, a_1, a_0) . If the content of f is equal to 1 then these coefficients are **relatively prime** and we call f **primitive**.

Lemma 6.1. If f and g are primitive elements of $\mathbb{Z}[t]$, then so is fg .

Proof. Write

$$\begin{aligned} f &= f_n t^n + \cdots + f_1 t + f_0; \\ g &= g_n t^n + \cdots + g_1 t + g_0. \end{aligned}$$

Let d be the content of fg , and suppose that $d > 1$. Let p be a prime dividing d . Since f and g are primitive we know that p does not divide any f_i nor any g_i . Let j (resp. k) be the minimum value such that p does not divide f_j (resp. g_k).

Let c_{j+k} be the coefficient of t^{j+k} in fg . Then

$$c_{j+k} = \sum_{l+m=j+k} f_l g_m = \sum_{l < j} f_l g_{j+k-l} + \sum_{m < k} f_{j+k-m} g_m + f_j g_k.$$

Now p divides the left hand side. On the right hand side, p divides all terms except for the last. This is a contradiction. \square

Theorem 6.2. (Gauss' Lemma) Let $f \in \mathbb{Z}[t]$. Then f is irreducible in $\mathbb{Z}[t]$ if and only if f is irreducible in $\mathbb{Q}[t]$.

Proof. Suppose that f is irreducible in $\mathbb{Z}[t]$, and $f = gh$ with $g, h \in \mathbb{Q}[t]$. We need to show that $\deg g = \deg f$ or $\deg h = \deg f$. We can assume that f is primitive, since $\frac{1}{d}f = (\frac{1}{d}g)h$ where d is the content of f .

Let d_g (resp. d_h) be the lowest common multiple of the set of denominators in coefficients of g (resp. h). Then $d_g g$ and $d_h h$ are both elements of $\mathbb{Z}[t]$. Now let n_g (resp. n_h) be the content of $d_g g$ (resp. $d_h h$). Set $\lambda_g = \frac{n_g}{d_g}$ and $\lambda_h = \frac{n_h}{d_h}$. Then $\lambda_g g$ and $\lambda_h h$ are primitive elements of $\mathbb{Z}[t]$.

Now observe that

$$\lambda_g \lambda_h f = (\lambda_g g)(\lambda_h h).$$

The RHS is a product of polynomials in $\mathbb{Z}[t]$. Since f is primitive in $\mathbb{Z}[t]$ we conclude that $\lambda_g \lambda_h \in \mathbb{Z}$. Furthermore the RHS is a product of primitive polynomials and so Lem. 6.1 implies that $\lambda_g \lambda_h = \pm 1$. Thus we obtain that

$$\pm f = (\lambda_g g)(\lambda_h h).$$

Since f is irreducible in $\mathbb{Z}[t]$, we conclude that either $\deg g = \deg f$ or $\deg h = \deg f$, as required. \square

Exercise 6.3. Formulate and prove the two previous results for any unique factorization domain, R . (Instead of \mathbb{Q} , one needs to use the field of fractions for R .)

We need some notation for the next result. For $f \in \mathbb{Z}[t]$ and $m \in \mathbb{Z}^+$, we write $[f]_m$ for the image of f under the canonical map $\mathbb{Z}[t] \rightarrow \mathbb{Z}/m\mathbb{Z}[t]$ (where we reduce every coefficient modulo m).

Theorem 6.4. (Localization principle) *Suppose that $f, g, h \in \mathbb{Z}[t]$. If $f = gh$ then, for all $m \in \mathbb{Z}^+$,*

$$[f]_m = [g]_m [h]_m.$$

Furthermore if m does not divide the leading coefficient of f , then $[f]_m$ irreducible implies that f is irreducible.

Proof. That $[f]_m = [g]_m [h]_m$ is trivial. Now suppose that $f = gh$ where $\deg g, \deg h < \deg f$, and that m does not divide the leading coefficient of f . Then $\deg [g]_m \deg [h]_m < \deg [f]_m$, and the second part follows. \square

Corollary 6.5. (Eisenstein's criterion) *Suppose that*

$$f = f_n t^n + \cdots + f_1 t + f_0$$

is a polynomial in $\mathbb{Z}[t]$. suppose that p is a prime such that $p \mid f_i$ for $i = 0, \dots, n-1$, $p \nmid f_n$, and $p^2 \nmid f_0$. Then f is irreducible in $\mathbb{Z}[t]$.

Proof. Suppose f is reducible of degree n , and localize at p :

$$\begin{aligned} f = gh &\implies [f]_p = [g]_p [h]_p \\ &\implies x^n = [g]_p [h]_p. \end{aligned}$$

Now the final equation is in $\mathbb{Z}/p\mathbb{Z}[t]$, which is a unique factorization domain. It follows that $[g]_p = t^r$ and $[h]_p = t^{n-r}$ where r is some positive integer. Then p must divide all non-leading coefficients of g and h ; in particular p divides the constant term in both g and h . But this implies that p^2 divides f_0 , which is a contradiction. \square

Exercise 6.6. *Formulate and prove Eisenstein's criterion for an arbitrary integral domain R . Note that in this case one must be careful to choose p to be a **prime** element, rather than an **irreducible** element of R . In arbitrary integral domains, these two concepts are distinct.*

We should note an important variation in definitions: Garling [Gar86], and others, require a polynomial in $\mathbb{Z}[t]$ to have content equal to 1 if it is to be irreducible. This makes some sense: since \mathbb{Z} is not a field, not all polynomials will be divisible by a constant polynomial other than ± 1 . However we will not use this definition of irreducibility in this course.

7. SPLITTING FIELDS

We return to the main stream of the course. We want to examine field extensions, and their connection to the factorization of polynomials. If we have a polynomial $f \in \mathbb{Z}[t]$, then we know that there exists a field extension - \mathbb{C} - within which f factorizes into linear factors. But is this true for a general field K ?

Definition 22. Take $f \in K[t]$ of degree n , with $L : K$ a field extension. We say that f **splits** over L if one has

$$f(t) = \lambda(t - \alpha_1) \cdots (t - \alpha_n),$$

for some $\lambda \in K$, and $\alpha_1, \dots, \alpha_n \in L$.

We say, furthermore, that $L : K$ is a **splitting field extension** for f over K (or simply that L is a **splitting field** for f) if

- (a) f splits over L , and
- (b) there exists no proper subfield M of L containing K such that f splits over M .

Observe that if f splits over some field L_0 , then a splitting field for f must exist.

Exercise 7.1. Let $f = t^2 + 1$. Show that \mathbb{C} is a splitting field for f over \mathbb{R} , but is not a splitting field for f over \mathbb{Q} .

Theorem 7.2. Suppose that $L : K$ is a field extension, and that $f \in K[t]$. Suppose that f splits over L as

$$f(t) = \lambda(t - \alpha_1) \cdots (t - \alpha_n).$$

Then $K(\alpha_1, \dots, \alpha_n)$ is a splitting field for f over K .

Proof. It is plain that f splits over $K(\alpha_1, \dots, \alpha_n)$. Suppose now that M is a field with $K \subseteq M \subseteq K(\alpha_1, \dots, \alpha_n)$, and that f splits over M . Then

$$f = \mu(t - \beta_1) \cdots (t - \beta_n),$$

for some $\mu \in K$ and $\beta_1, \dots, \beta_n \in M$. Now $K[t]$ is a unique factorization domain; thus, by reordering, we have $\beta_i = \alpha_i$ for $i = 1, \dots, n$. Thus $\alpha_1, \dots, \alpha_n \in M$ and so $K(\alpha_1, \dots, \alpha_n) \subseteq M$. We are done. \square

Corollary 7.3. If $L : K$ is a splitting field extension for $f \in K[t]$, then $L : K$ is an algebraic extension.

The problem with the previous result is that we had to assume that K lay inside a field L over which the polynomial f split. But how do we know that such a field L exists for arbitrary K and f ? We need an “intrinsic” construction of a splitting field...

Theorem 7.4. Suppose that $f \in K[t]$ is irreducible of degree n . Then there is a simple algebraic extension $K(\alpha) : K$ such that $[K(\alpha) : K] = n$ and $f(\alpha) = 0$.

This result is very much reminiscent of Thm. 4.8. Crucially, though, in this instance we don’t require the existence of an ambient field L .

Proof. Let $\phi : K \rightarrow K[t]$ be the canonical monomorphism embedding K in $K[t]$ as the constant polynomials. Write $L = K[t]/(f)$, and let $q : K[t] \rightarrow L$ be the induced quotient map. Since f is irreducible, Lem. 4.1 implies that L is a field.

Let $i = q\phi$; then i is a monomorphism of K into L so that (i, K, L) is a field extension. Write $\alpha = q(t) = t + (f)$. Since t generates $K[t]$ over K , we know that $L = K(\alpha)$. Furthermore, since q is a ring homomorphism,

$$f(\alpha) = f(q(t)) = q(f) = q(0 + (f)) = 0.$$

Thus α is algebraic over K . Furthermore, since f is irreducible and $f(\alpha) = 0$, we know that f is a scalar multiple of the minimal polynomial m_α of α over K . Then $[L : K] = \deg m_\alpha = n$, as required. \square

Theorem 7.5. *Suppose that $f \in K[t]$. Then there exists a splitting field extension $L : K$ for f with $[L : K] \leq n!$.*

Proof. We prove this by induction on the degree of f . If $\deg f \leq 1$, then the conclusion is trivial. So suppose that $\deg f > 1$.

Suppose that f is irreducible over K . By Thm. 7.4, there exists a simple algebraic extension $K(\alpha) : K$ with $[K(\alpha) : K] = n$, and such that

$$f(t) = (t - \alpha)g(t),$$

with $g \in K(\alpha)[t]$, and $\deg g = n - 1$. By the inductive hypothesis, there exists a splitting field for g , say $L : K(\alpha)$, with $[L : K(\alpha)] \leq (n - 1)!$. Write

$$g = \lambda(t - \beta_1) \cdots (t - \beta_{n-1}),$$

where $\lambda \in K(\alpha)$ is the leading coefficient of g . But then

$$f = \lambda(t - \alpha)(t - \beta_1) \cdots (t - \beta_{n-1}),$$

and so λ is the leading coefficient of f and lies in K . Thus f splits over L . Then $L = K(\alpha, \beta_1, \dots, \beta_{n-1})$ is a splitting field extension of f over K . What is more

$$[L : K] = [L : K(\alpha)][K(\alpha) : K] \leq (n - 1)!n = n!.$$

Suppose that f is not irreducible over K . Write $f = gh$ with $\deg g = k$, $\deg h = l$, $\deg f = k + l$, and both k and l at least 1. By the inductive hypothesis, there is a splitting field extension $L : K$ for g with $[L : K] \leq k!$. Thus we may write

$$g = \lambda(t - \alpha_1) \cdots (t - \alpha_k),$$

with $\lambda \in K$ and $\alpha_1, \dots, \alpha_k \in L$; then $L = K(\alpha_1, \dots, \alpha_k)$.

Now we consider h as an element of $L[t]$. By the inductive hypothesis, there is a splitting field $M : L$ for h with $[M : L] \leq l!$, and we can write

$$h = \mu(t - \beta_1) \cdots (t - \beta_l),$$

where $\mu \in L$ and $\beta_1, \dots, \beta_l \in M$. Clearly then

$$f = gh = \lambda\mu(t - \alpha_1) \cdots (t - \alpha_k)(t - \beta_1) \cdots (t - \beta_l),$$

where the leading coefficient of f , namely $\lambda\mu$, must lie in K . Then

$$M = L(\beta_1, \dots, \beta_l) = K(\alpha_1, \dots, \alpha_k, \beta_1, \dots, \beta_l),$$

so $M : K$ is a splitting field extension for f over K . Moreover by the tower law

$$[M : K] = [M : L][L : K] \leq l!k! \leq (l + k)! = n!.$$

□

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