

CS2100

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COMPUTER ORGANISATION

Lecture #7

MIPS

Part I: Introduction



NUS
National University
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School of
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Lecture #7: MIPS Part 1: Introduction (1/2)

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2. Machine Code vs Assembly Language
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Lecture #7: MIPS Part 1: Introduction (1/2)

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- 7. MIPS Basic Instructions Checklist

Recap

High-level
language
program
(in C)

```
swap(int v[], int k)
{int temp;
  temp = v[k];
  v[k] = v[k+1];
  v[k+1] = temp;
}
```

Compiler

Assembly
language
program
(for MIPS)

```
swap:
  muli $2, $5, 4
  add  $2, $4, $2
  lw   $15, 0($2)
  lw   $16, 4($2)
  sw   $16, 0($2)
  sw   $15, 4($2)
  jr   $31
```

Assembler

Binary machine
language
program
(for MIPS)

```
000000001010000100000000000011000
000000000000110000001100000100001
100011000110001000000000000000000
100011001111001000000000000000100
101011001111001000000000000000000
101011000110001000000000000000100
00000011111000000000000000001000
```

- You write programs in high level programming languages, e.g., C/C++, Java:

A + B

- Compiler** translates this into **assembly language** statement:

add A, B

- Assembler** translates this statement into **machine language instructions** that the processor can execute:

1000 1100 1010 0000

1. Instruction Set Architecture (1/2)

- **Instruction Set Architecture (ISA):**
 - An abstraction on the **interface** between the hardware and the low-level software.

Software

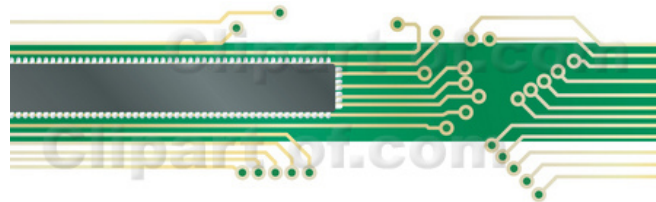
(to be translated to the instruction set)



Instruction Set Architecture

Hardware

(implementing the instruction set)



1. Instruction Set Architecture (2/2)

- Instruction Set Architecture
 - Includes everything programmers need to know to make the machine code work correctly
 - Allows computer designers to talk about functions independently from the hardware that performs them
- This abstraction allows many implementations of varying cost and performance to run identical software.
 - Example: Intel x86/IA-32 ISA has been implemented by a range of processors starting from 80386 (1985) to Pentium 4 (2005)
 - Other companies such as AMD and Transmeta have implemented IA-32 ISA as well
 - A program compiled for IA-32 ISA can execute on any of these implementations

2. Machine Code vs Assembly Language

■ **Machine code**

- Instructions are represented in binary
- `1000110010100000` is an instruction that tells one computer to add two numbers
- Hard and tedious for programmer

■ **Assembly language**

- Symbolic version of machine code
- Human readable
- `add A, B` is equivalent to `1000110010100000`
- **Assembler** translates from assembly language to machine code
- Assembly can provide '**pseudo-instructions**' as syntactic sugar
- When considering performance, only real instructions are counted.

MIPS instruction

sll \$8, \$9, 4

Note the placement of the source register

Field representation in decimal:

opcode	rs	rt	rd	shamt	funct
0	0	9	8	4	0

Field representation in binary:

000000	00000	01001	01000	00100	000000
--------	-------	-------	-------	-------	--------

Split into 4-bit groups for hexadecimal conversion:

0000	0000	0000	1001	0100	0001	0000	0000
0 ₁₆	0 ₁₆	0 ₁₆	9 ₁₆	4 ₁₆	1 ₁₆	0 ₁₆	0 ₁₆

3. Walkthrough: An Example Code (1/15)

- Let us take a journey with the execution of a simple code:
 - Discover the components in a typical computer
 - Learn the type of instructions required to control the processor
 - Simplified to highlight the important concepts ☺

```
// assume res is 0 initially  
for (i=1; i<10; i++) {  
    res = res + i;  
}
```

C-like code
fragment

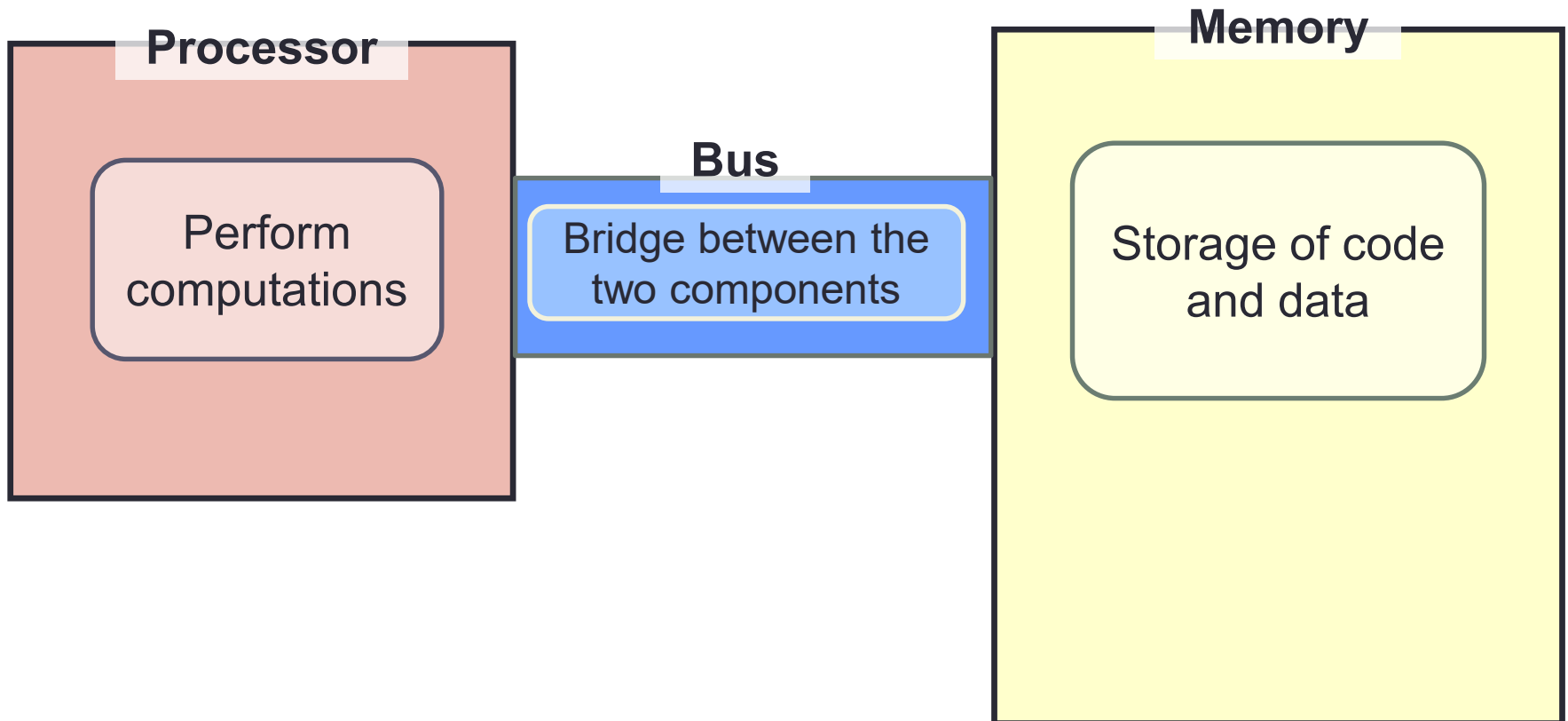


```
res ← res + i  
i ← i + 1  
if i < 10, repeat
```

“Assembly”
Code

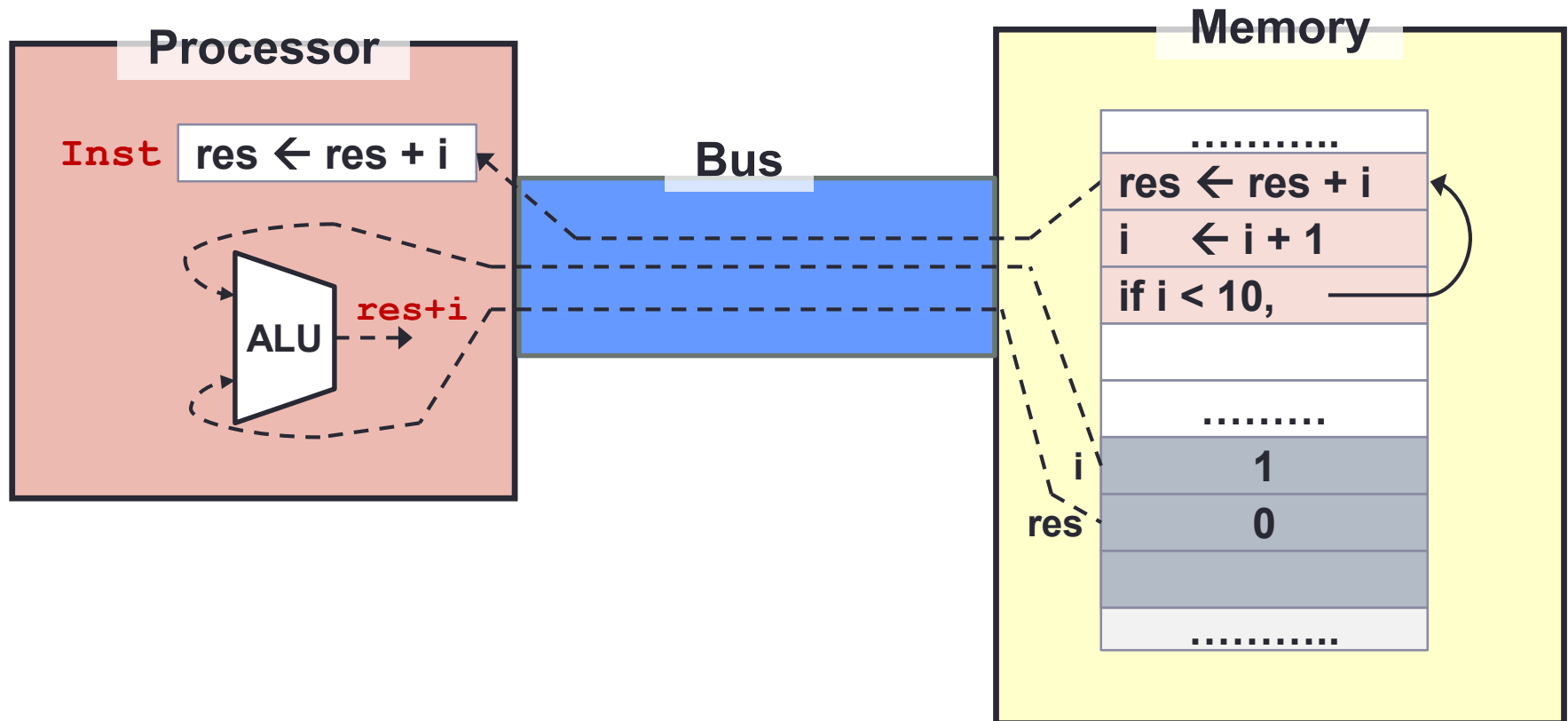
3. Walkthrough: The Components (2/15)

- The two major components in a computer
 - **Processor** and **Memory**
 - Input/Output devices omitted in this example



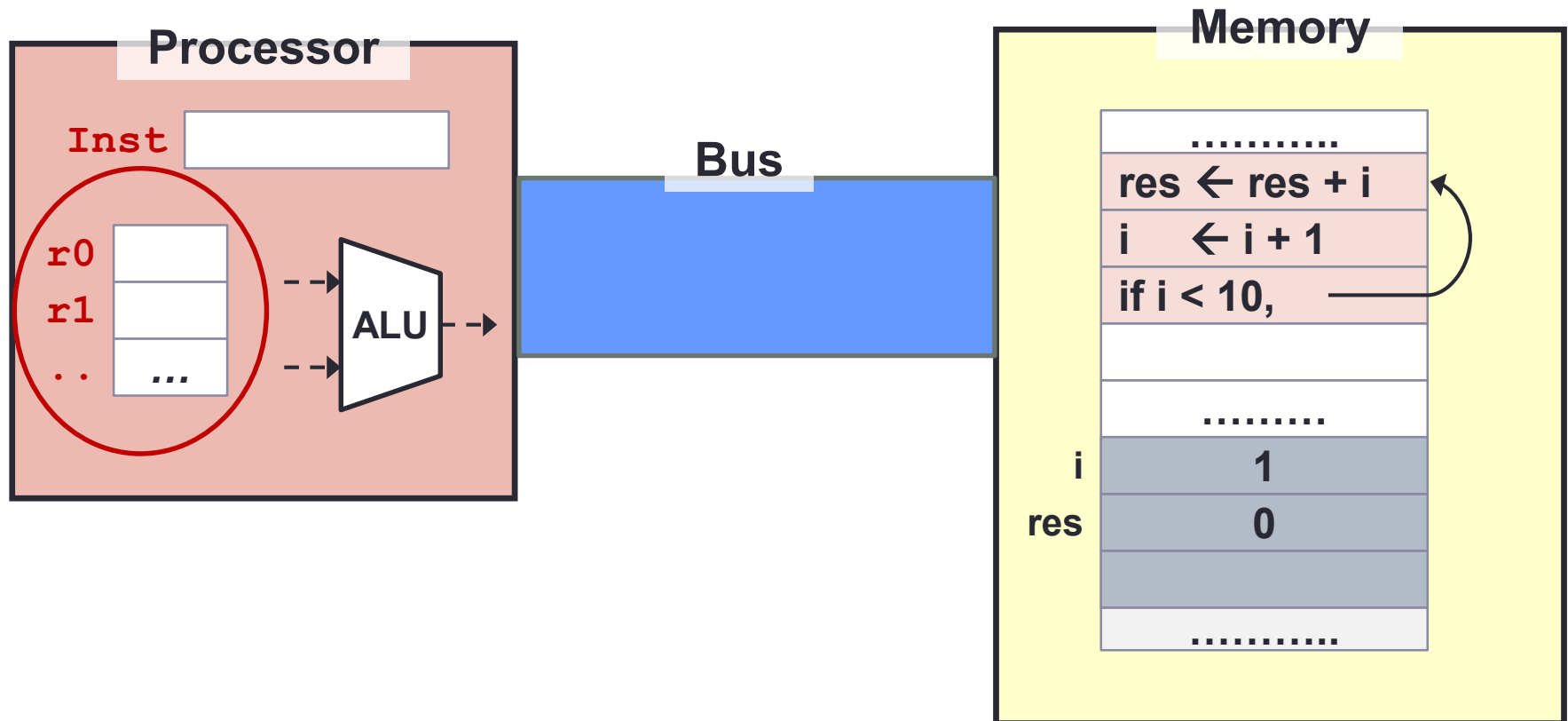
3. Walkthrough: The Code in Action (3/15)

- The code and data reside in memory
 - Transferred into the processor during execution



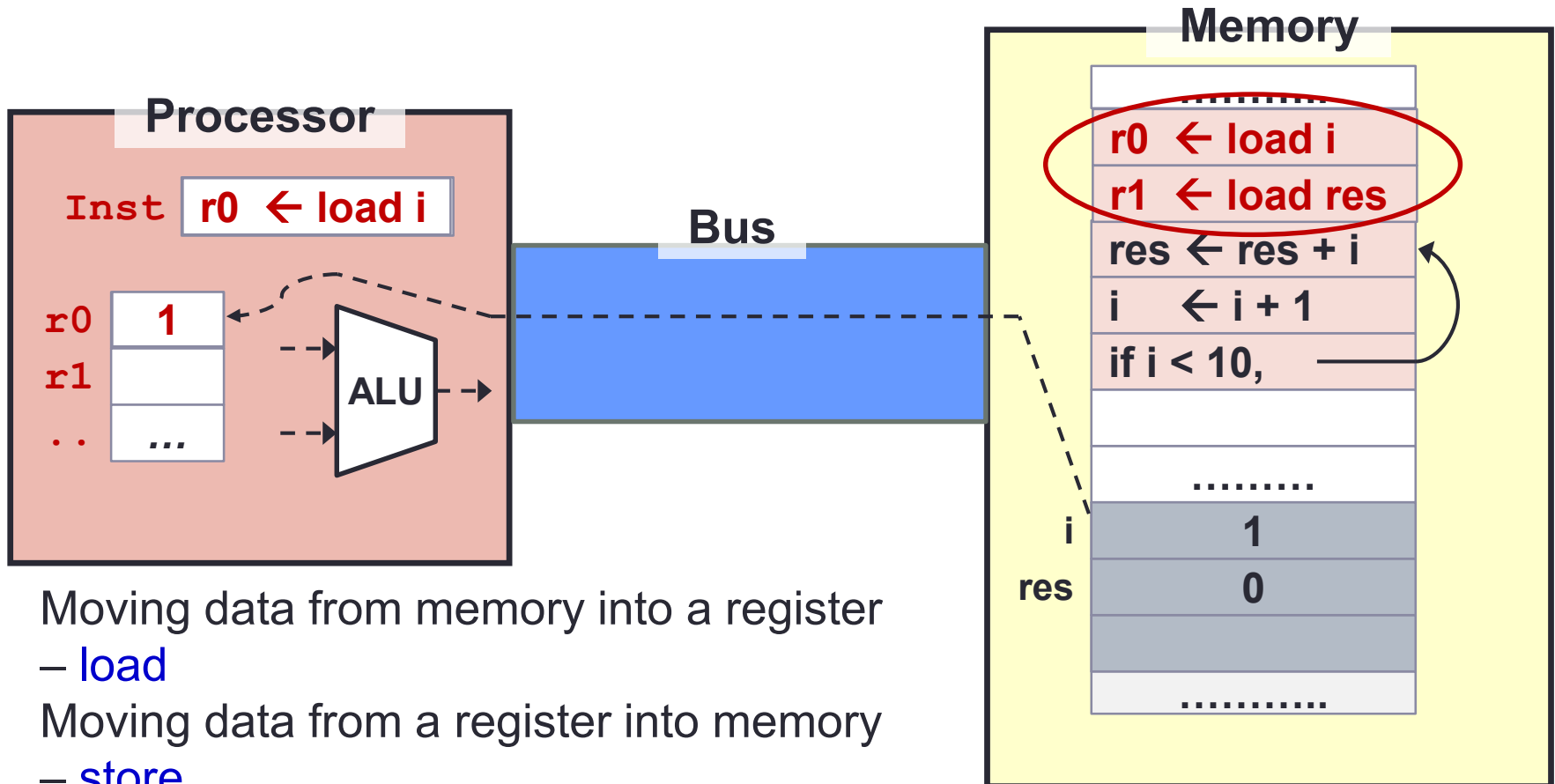
3. Walkthrough: Memory Access is Slow! (4/15)

- To avoid frequent access of memory
 - Provide temporary storage for values in the processor (known as **registers**)



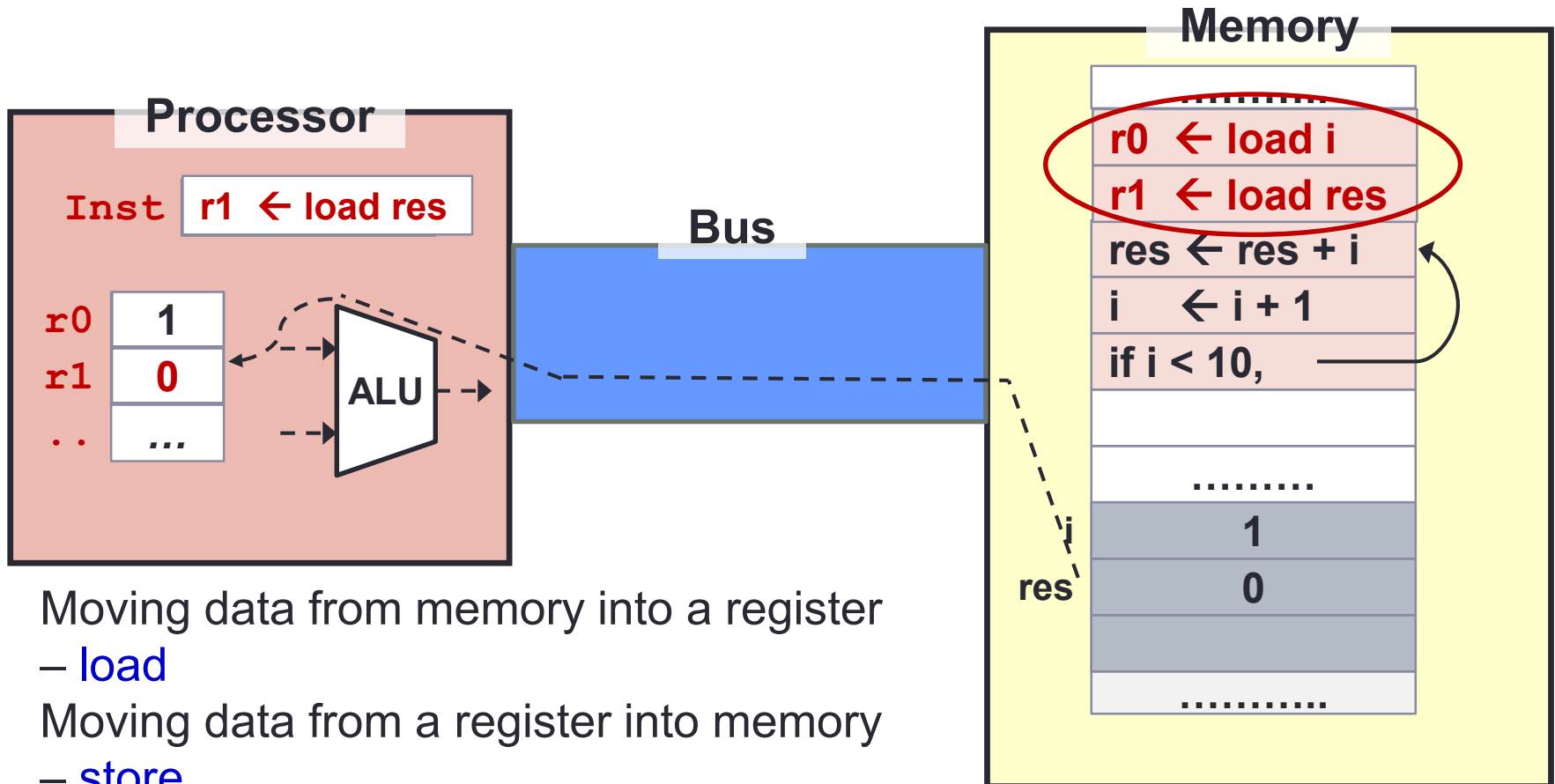
3. Walkthrough: Memory Instruction (5/15)

- Need instruction to move data into registers
 - Also to move data from registers to memory later



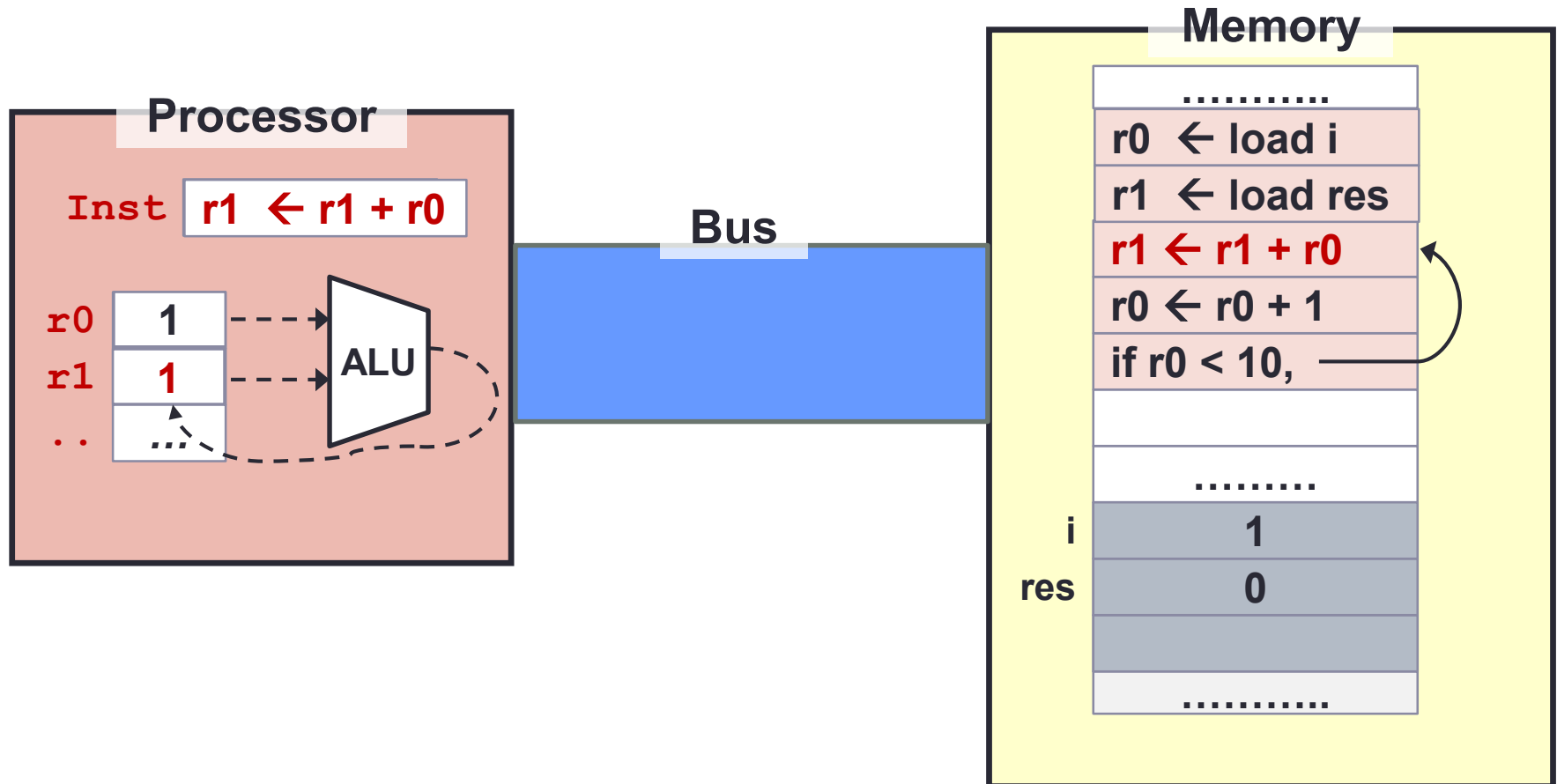
3. Walkthrough: Memory Instruction (6/15)

- Need instruction to move data into registers
 - Also to move data from registers to memory later



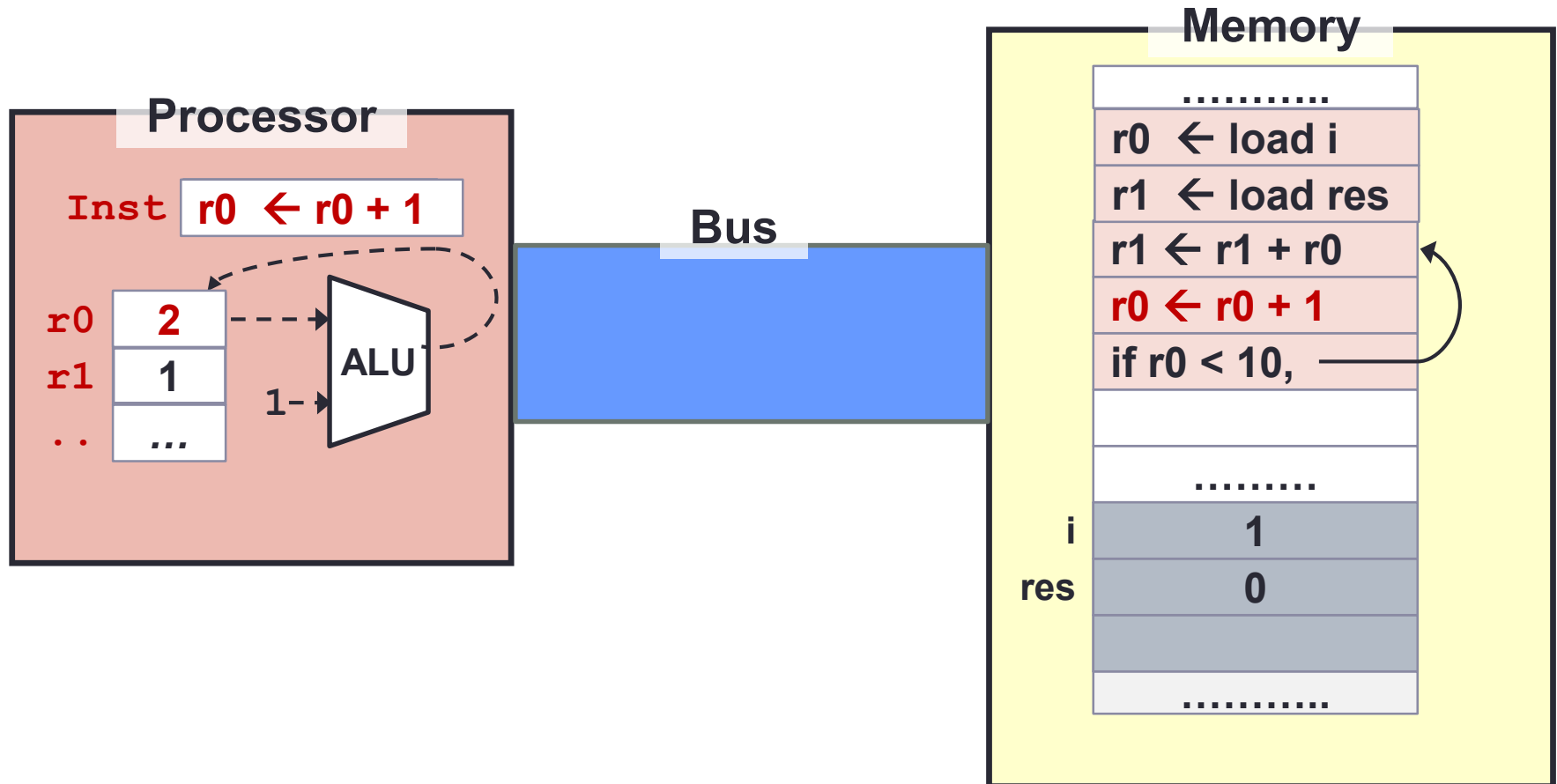
3. Walkthrough: Reg-to-Reg Arithmetic (7/15)

- Arithmetic operations can now work directly on registers only (much faster!)



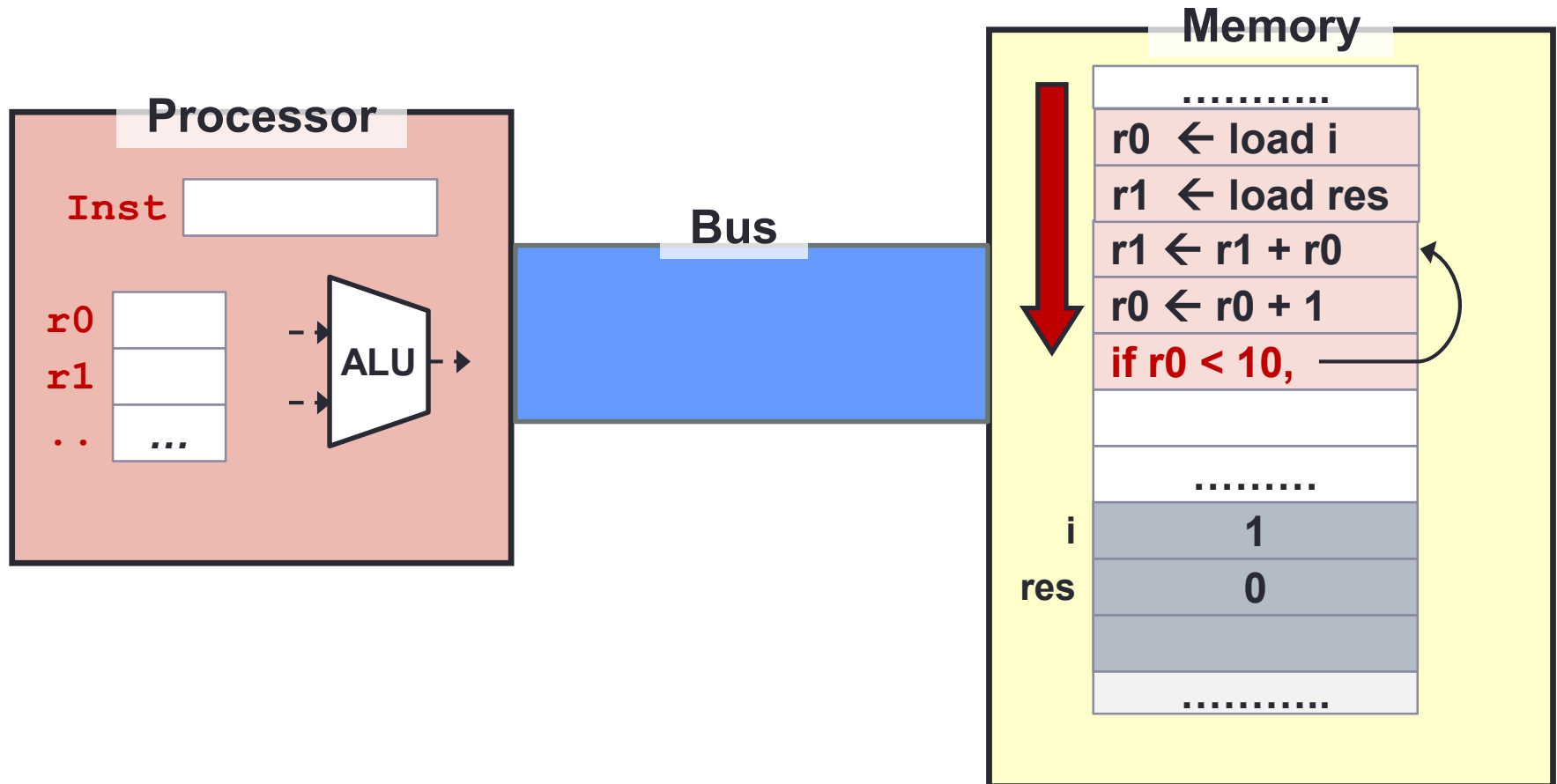
3. Walkthrough: Reg-to-Reg Arithmetic (8/15)

- Sometimes, arithmetic operation uses a **constant** value instead of register value



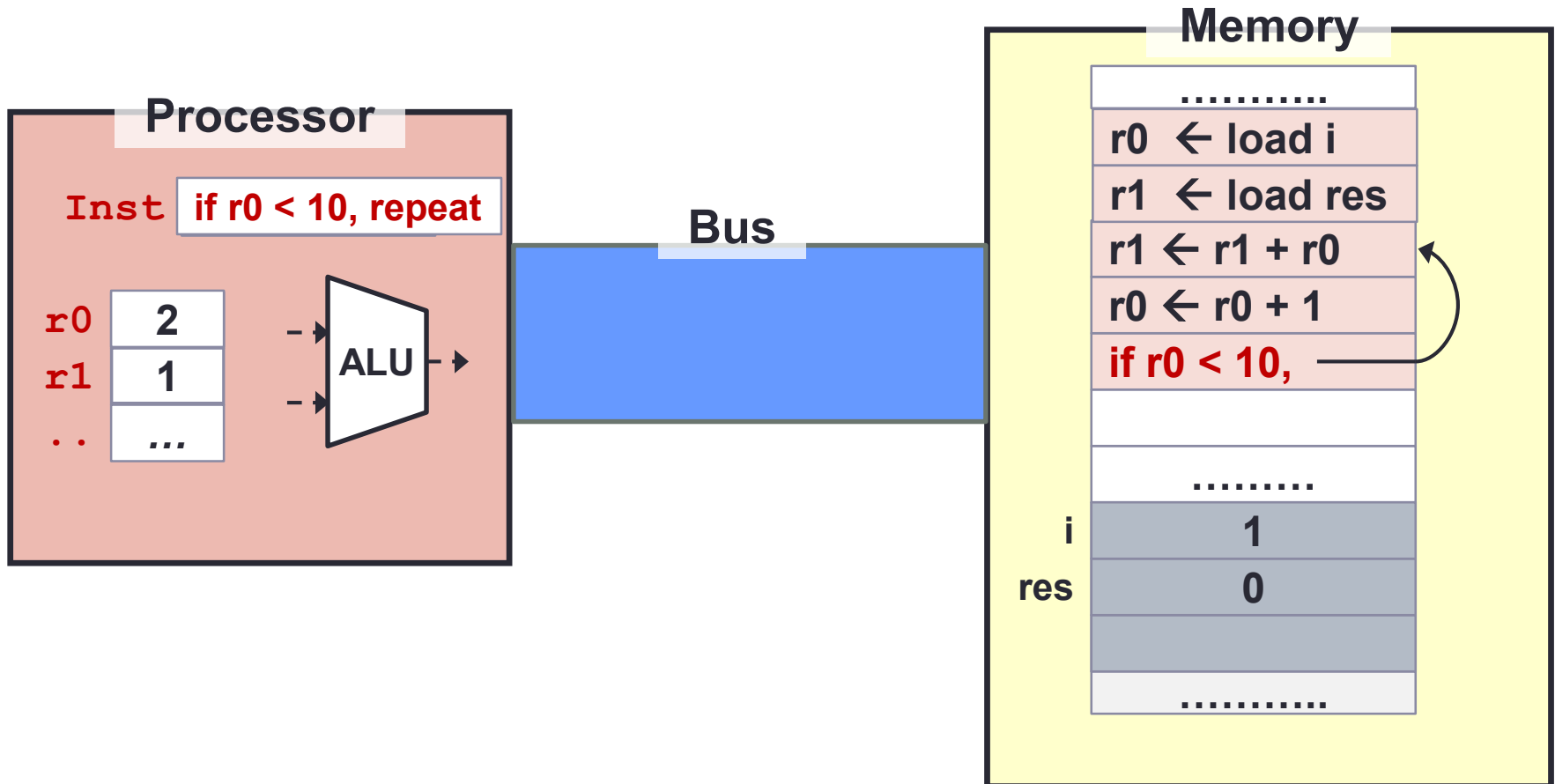
3. Walkthrough: Execution Sequence (9/15)

- Instruction is executed sequentially by default
 - How do we “repeat” or “make a choice”?



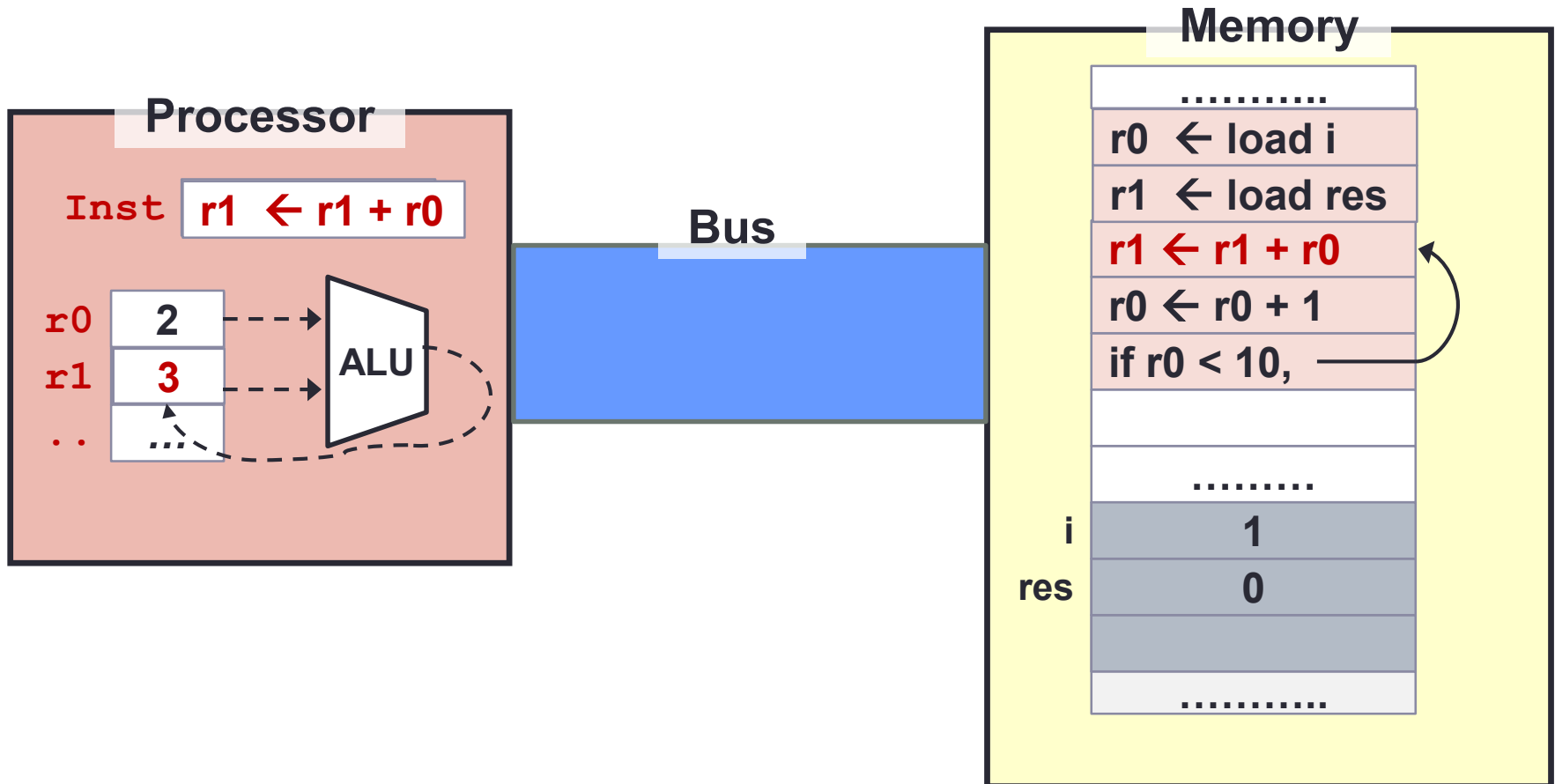
3. Walkthrough: Control Flow (10/15)

- We need instructions to **change** the control flow based on **condition**:
 - Repetition (loop) and Selection (if-else) can both be supported



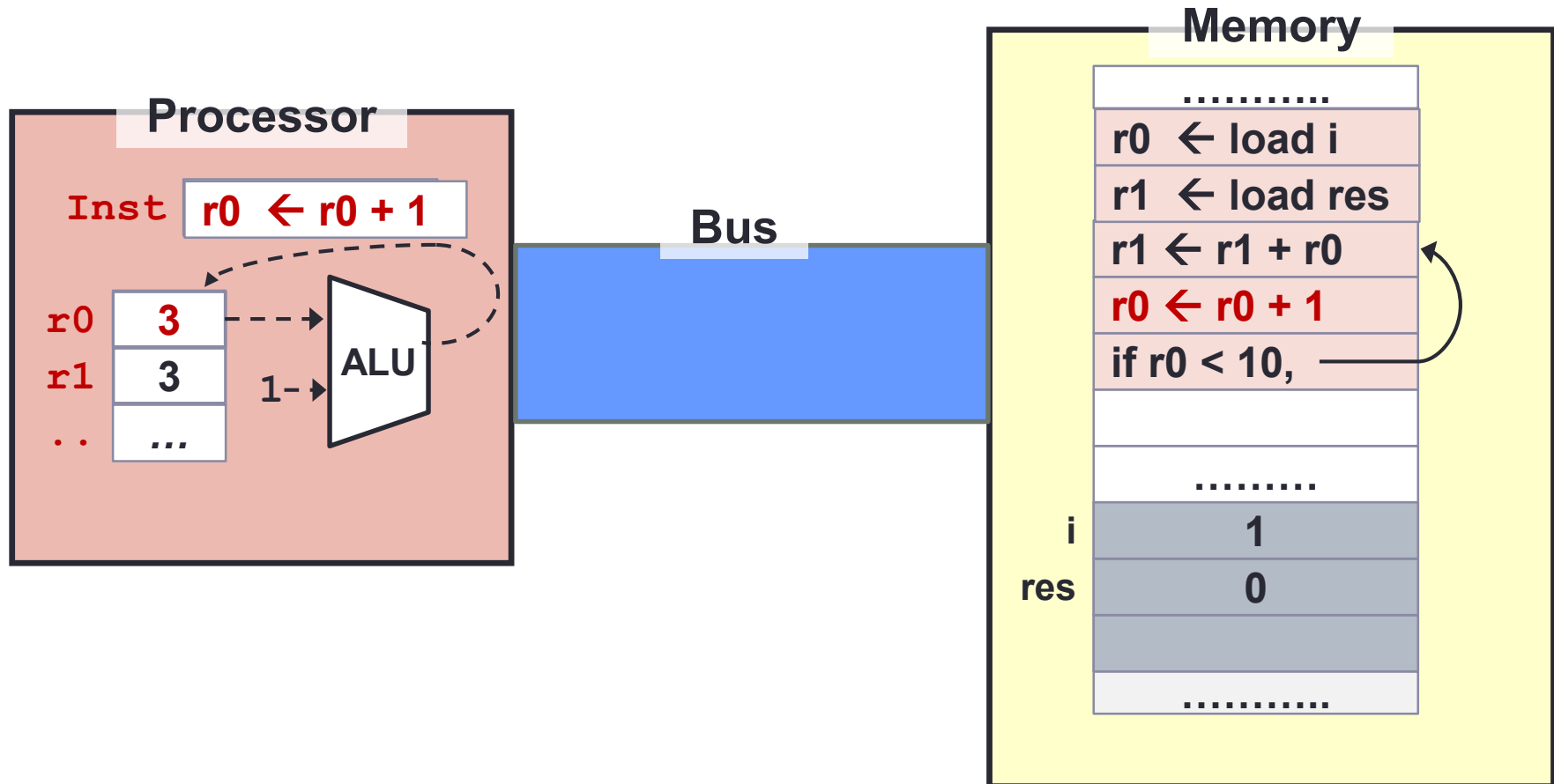
3. Walkthrough: Looping! (11/15)

- Since the condition succeeded, execution will repeat from the indicated position



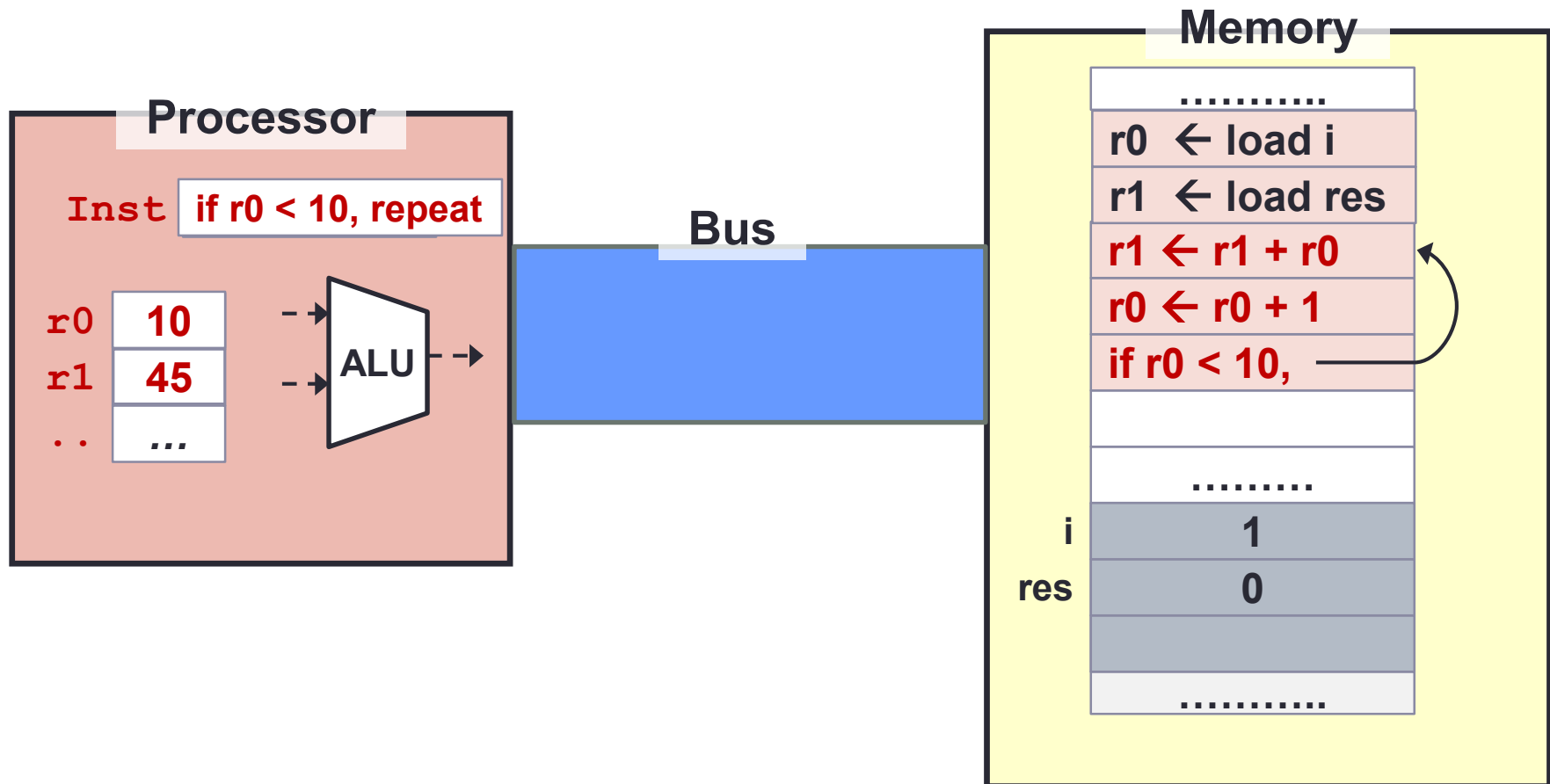
3. Walkthrough: Looping! (12/15)

- Execution will continue sequentially
 - Until we see another control flow instruction



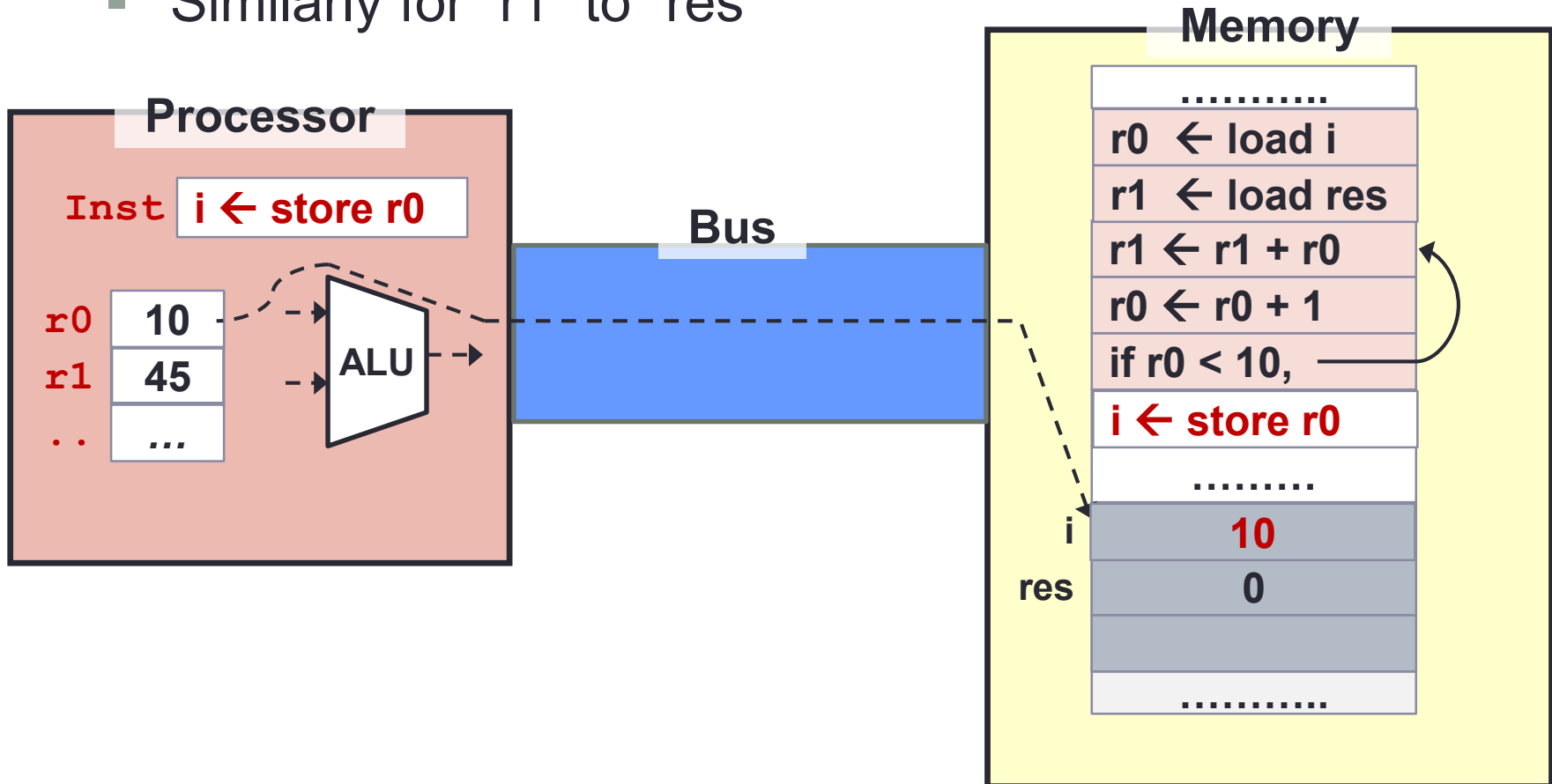
3. Walkthrough: Looping! (13/15)

- The three instructions will be repeated until the condition fails



3. Walkthrough: Memory Instruction (14/15)

- We can now move back the values from register to their “home” in memory
 - Similarly for “r1” to “res”



3. Walkthrough: Summary (15/15)

- The stored-memory concept:
 - Both **instruction** and **data** are stored in memory
- The load-store model:
 - Limit memory operations and relies on registers for storage during execution
- The major types of assembly instruction:
 - **Memory:** Move values between memory and registers
 - **Calculation:** Arithmetic and other operations
 - **Control flow:** Change the sequential execution

4. General Purpose Registers (1/2)

- Fast memories in the processor:
 - Data are transferred from memory to registers for faster processing
- Limited in number:
 - A typical architecture has 16 to 32 registers
 - Compiler associates variables in program with registers
- Registers have **no data type**
 - Unlike program variables!
 - Machine/Assembly instruction assumes the data stored in the register is of the correct type

4. General Purpose Registers (2/2)

- There are **32 registers** in **MIPS** assembly language:
 - Can be referred by a number (\$0, \$1, ..., \$31) OR
 - Referred by a name (eg: \$a0, \$t1)

Name	Register number	Usage
\$zero	0	Constant value 0
\$v0-\$v1	2-3	Values for results and expression evaluation
\$a0-\$a3	4-7	Arguments
\$t0-\$t7	8-15	Temporaries
\$s0-\$s7	16-23	Program variables

Name	Register number	Usage
\$t8-\$t9	24-25	More temporaries
\$gp	28	Global pointer
\$sp	29	Stack pointer
\$fp	30	Frame pointer
\$ra	31	Return address

\$at (register 1) is reserved for the assembler.

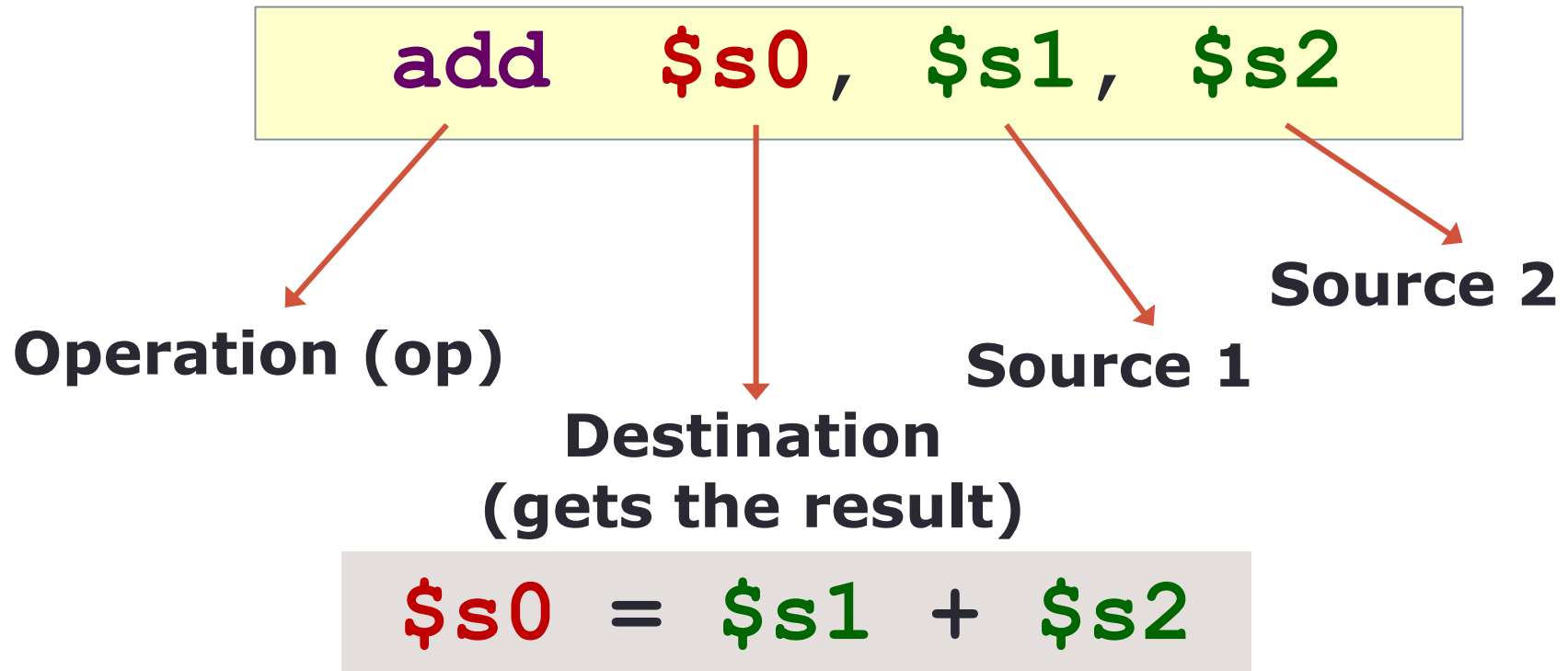
\$k0-\$k1 (registers 26-27) are reserved for the operation system.

5. MIPS Assembly Language

- Each instruction executes a simple command
 - Usually has a counterpart in high level programming languages like C/C++, Java etc
- Each line of assembly code contains at most 1 instruction
- # (hex-sign) is used for comments
 - Anything from # to end of line is a comment and will be ignored by the assembler

```
add $t0, $s1, $s2 # $t0 ← $s1 + $s2  
sub $s0, $t0, $s3 # $s0 ← $t0 - $s3
```

5.1 General Instruction Syntax



Naturally, most of the MIPS arithmetic/logic operations have three operands: **2 sources** and **1 destination**

5.2 Arithmetic Operation: Addition

C Statement	MIPS Assembly Code
<code>a = b + c;</code>	<code>add \$s0, \$s1, \$s2</code>

- We assume the values of "**a**", "**b**" and "**c**" are loaded into registers "**\$s0**", "**\$s1**" and "**\$s2**"
 - Known as **variable mapping**
 - Actual code to perform the loading will be shown later in ***memory instruction***
- Important concept:
 - MIPS arithmetic operations are mainly **register-to-register**

5.3 Arithmetic Operation: Subtraction

C Statement	MIPS Assembly Code
<code>a = b - c;</code>	<code>sub \$s0, \$s1, \$s2</code> \$s0 → variable a \$s1 → variable b \$s2 → variable c

- Positions of `$s1` and `$s2` (i.e., source1 and source2) are important for subtraction

5.4 Complex Expression (1/3)

C Statement	MIPS Assembly Code
<code>a = b + c - d;</code>	<p>??? ??? ???</p> <p>\$s0 → variable a \$s1 → variable b \$s2 → variable c \$s3 → variable d</p>

- A single MIPS instruction can handle at most two source operands

➔ **Need to break a complex statement into multiple MIPS instructions**

MIPS Assembly Code			
add	\$t0	\$s1, \$s2	# tmp = b + c
sub	\$s0	\$t0, \$s3	# a = tmp - d

Use temporary registers **\$t0** to **\$t7** for intermediate results

5.4 Complex Expression: Example (2/3)

C Statement	Variable Mappings
<code>f = (g + h) - (i + j);</code>	<code>\$s0</code> → variable <code>f</code> <code>\$s1</code> → variable <code>g</code> <code>\$s2</code> → variable <code>h</code> <code>\$s3</code> → variable <code>i</code> <code>\$s4</code> → variable <code>j</code>

- Break it up into multiple instructions
 - Use two temporary registers `$t0`, `$t1`

```
add $t0, $s1, $s2    # tmp0 = g + h
add $t1, $s3, $s4    # tmp1 = i + j
sub $s0, $t0, $t1    # f = tmp0 - tmp1
```

5.4 Complex Expression: Exercise (3/3)

C Statement	Variable Mappings
$z = a + b + c + d;$ <div> add $\\$s4$, $\\$s0$, $\\$s1$ add $\\$s4$, $\\$s4$, $\\$s2$ add $\\$s4$, $\\$s4$, $\\$s3$ </div>	$\$s0 \rightarrow$ variable a $\$s1 \rightarrow$ variable b $\$s2 \rightarrow$ variable c $\$s3 \rightarrow$ variable d $\$s4 \rightarrow$ variable z

C Statement	Variable Mappings
$z = (a - b) + c;$ <div> sub $\\$s3$, $\\$s0$, $\\$s1$ add $\\$s3$, $\\$s3$, $\\$s2$ </div>	$\$s0 \rightarrow$ variable a $\$s1 \rightarrow$ variable b $\$s2 \rightarrow$ variable c $\$s3 \rightarrow$ variable z

5.5 Constant/Immediate Operands

C Statement	MIPS Assembly Code
<code>a = a + 4;</code>	<code>addi \$s0, \$s0, 4</code>

- **Immediate** values are numerical constants
 - Frequently used in operations
 - MIPS supplies a set of operations specially for them
- “Add immediate” (**addi**)
 - Syntax is similar to **add** instruction; but source2 is a constant instead of register
 - **The constant ranges from $[-2^{15}$ to $2^{15}-1$]**

Can you guess what number system is used?

Answer: 16-bit 2s complement number system

There's no **subi**. Why?

Answer: Use **addi** with negative constant

5.6 Register Zero (\$0 or \$zero)

- The number zero (0), appears very often in code
 - Provide register zero (**\$0** or **\$zero**) which always have the **value 0**

C Statement	MIPS Assembly Code
<code>f = g;</code>	<code>add \$s0, \$s1, \$zero</code> \$s0 → variable f \$s1 → variable g

- The above assignment is so common that MIPS has an equivalent **pseudo instruction** (**move**):

MIPS Assembly Code
<code>move \$s0, \$s1</code>

Pseudo-Instruction

"Fake" instruction that gets translated to corresponding MIPS instruction(s).
Provided for convenience in coding only.

5.7 Logical Operations: Overview (1/2)

- Arithmetic instructions view the content of a register as a single quantity (signed or unsigned integer)
- New perspective:**
 - View register as 32 raw bits rather than as a single 32-bit number
- Possible to operate on individual bits or bytes within a word

Logical operation	C operator	Java operator	MIPS instruction
Shift Left	<<	<<	sll
Shift right	>>	>>, >>>	srl
Bitwise AND	&	&	and, andi
Bitwise OR			or, ori
Bitwise NOT*	~	~	nor
Bitwise XOR	^	^	xor, xori

*with some tricks

5.7 Logical Operations: Overview (2/2)

- Truth tables of logical operations
 - 0 represents false; 1 represents true

AND

a	b	a AND b
0	0	0
0	1	0
1	0	0
1	1	1

OR

a	b	a OR b
0	0	0
0	1	1
1	0	1
1	1	1

NOR

a	b	a NOR b
0	0	1
0	1	0
1	0	0
1	1	0

XOR

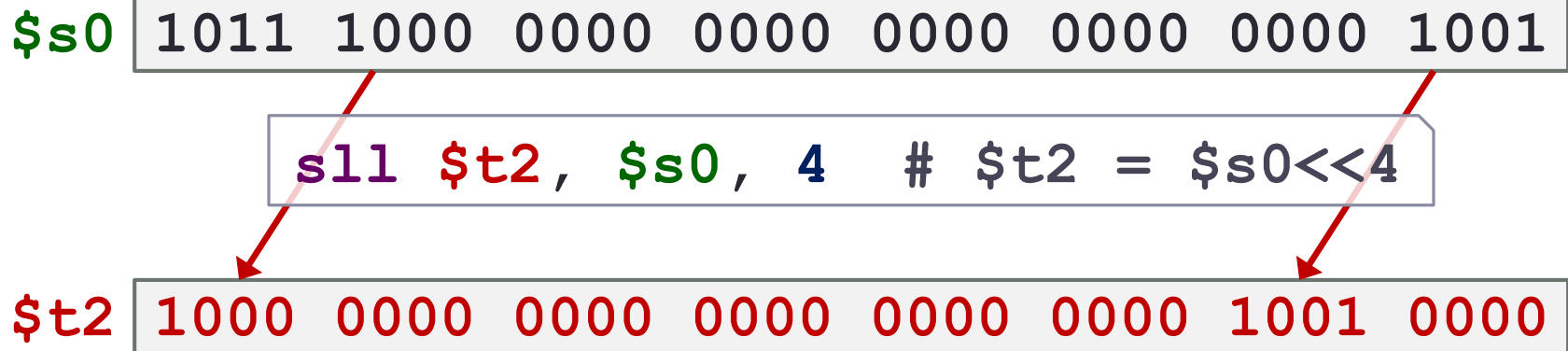
a	b	a XOR b
0	0	0
0	1	1
1	0	1
1	1	0

5.8 Logical Operations: **Shifting** (1/2)

Opcode: **s11** (s**h**ift **l**eft **l**ogical)

Move all the bits in a word to the left by a number of positions; fill the emptied positions with zeroes.

- E.g. Shift bits in **\$s0** to the left by 4 positions



5.8 Logical Operations: **Shifting** (2/2)

Opcode: **sr1** (shift right logical)

Shifts right and fills emptied positions with zeroes.

- What is the equivalent math operations for shifting **left/right** n bits? Answer: **Multiply/divide** by 2^n
- Shifting is faster than multiplication/division
 - Good compiler translates such multiplication/division into shift instructions

C Statement	MIPS Assembly Code
a = a * 8;	sll \$s0, \$s0, 3

5.9 Logical Operations: Bitwise AND

Opcode: `and` (bitwise **AND**)

Bitwise operation that leaves a 1 only if both the bits of the operands are 1

- E.g.: `and $t0, $t1, $t2`

	<code>\$t1</code>	0110	0011	0010	1111	0000	1101	0101	1001
mask	<code>\$t2</code>	0000	0000	0000	0000	0011	1100	0000	0000
	<code>\$t0</code>	0000	0000	0000	0000	0000	1100	0000	0000

- `and` can be used for **masking** operation:
 - Place **0s** into the positions to be ignored → bits will turn into 0s
 - Place **1s** for interested positions → bits will remain the same as the original.

5.9 Exercise: Bitwise AND

- We are interested in the last 12 bits of the word in register `$t1`. Result to be stored in `$t0`.
 - Q: What's the mask to use?

<code>\$t1</code>	0000	1001	1100	0011	0101	1101	1001	1100
mask	0000	0000	0000	0000	0000	1111	1111	1111
<code>\$t0</code>	0000	0000	0000	0000	0000	1101	1001	1100

Notes:

The **and** instruction has an immediate version, **andi**

5.10 Logical Operations: Bitwise OR

Opcode: `or` (bitwise `OR`)

Bitwise operation that places a 1 in the result if either operand bit is 1

Example: `or $t0, $t1, $t2`

- The `or` instruction has an immediate version `ori`
- Can be used to force certain bits to 1s
- E.g.: `ori $t0, $t1, 0xFFF`

<code>\$t1</code>	0000	1001	1100	0011	0101	1101	1001	1100
<code>0xFFF</code>	0000	0000	0000	0000	0000	1111	1111	1111
<code>\$t0</code>	0000	1001	1100	0011	0101	1111	1111	1111

5.11 Logical Operations: Bitwise NOR

- Strange fact 1:
 - There is no **NOT** instruction in MIPS to toggle the bits ($1 \rightarrow 0, 0 \rightarrow 1$)
 - However, a **NOR** instruction is provided:

Opcode: **nor** (bitwise **NOR**)

Example: **nor** \$t0, \$t1, \$t2

- Question: How do we get a NOT operation?

nor \$t0, \$t0, \$zero

- Question: Why do you think is the reason for not providing a NOT instruction?

One of design principles: **Keep the instruction set small.**

5.12 Logical Operations: Bitwise XOR

Opcode: `xor` (bitwise **XOR**)

Example: `xor $t0, $t1, $t2`

- Question: Can we also get **NOT** operation from **XOR**?

Yes, let `$t2` contain all 1s.

`xor $t0, $t0, $t2`

- Strange Fact 2:
 - There is no **NORI**, but there is **XORI** in MIPS
 - Why?

6. Large Constant: Case Study

- Question: How to load a 32-bit constant into a register? e.g **10101010 10101010 11110000 11110000**

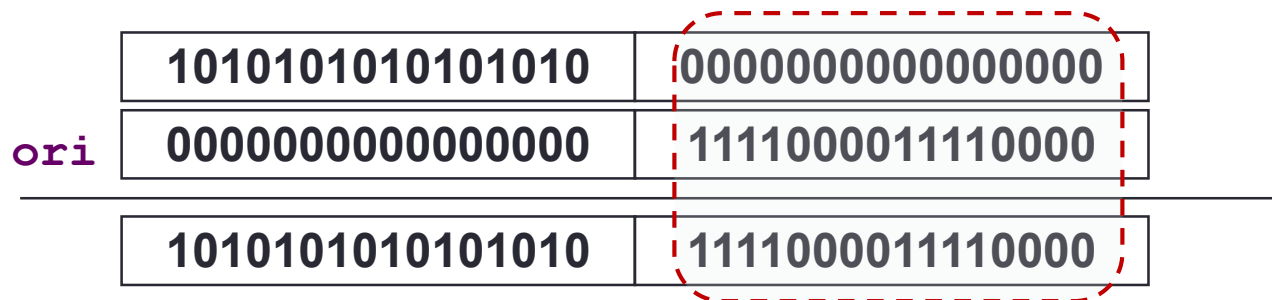
- Use “load upper immediate” (**lui**) to set the upper 16-bit:

```
lui    $t0, 0xAAAA    #1010101010101010
```



- Use “or immediate” (**ori**) to set the lower-order bits:

```
ori    $t0, $t0, 0xF0F0 #1111000011110000
```



7. MIPS Basic Instructions Checklist

Operation	Opcode in MIPS	Meaning
Addition	<code>add \$s0, \$s1, \$s2</code>	$\$s0 = \$s1 + \$s2$
	<code>addi \$s0, \$s1, C16_{2s}</code>	$\$s0 = \$s1 + C16_{2s}$
Subtraction	<code>sub \$s0, \$s1, \$s2</code>	$\$s0 = \$s1 - \$s2$
Shift left logical	<code>sll \$s0, \$s1, C5</code>	$\$s0 = \$s1 \ll C5$
Shift right logical	<code>srl \$s0, \$s1, C5</code>	$\$s0 = \$s1 \gg C5$
AND bitwise	<code>and \$s0, \$s1, \$s2</code>	$\$s0 = \$s1 \& \$s2$
	<code>andi \$s0, \$s1, C16</code>	$\$s0 = \$s1 \& C16$
OR bitwise	<code>or \$s0, \$s1, \$s2</code>	$\$s0 = \$s1 \$s2$
	<code>ori \$s0, \$s1, C16</code>	$\$s0 = \$s1 C16$
NOR bitwise	<code>nor \$s0, \$s1, \$s2</code>	$\$s0 = \$s1 \downarrow \$s2$
XOR bitwise	<code>xor \$s0, \$s1, \$s2</code>	$\$s0 = \$s1 \oplus \$s2$
	<code>xori \$s0, \$s1, C16</code>	$\$s0 = \$s1 \oplus C16$

C5 is $[0 \text{ to } 2^5-1]$

C16_{2s} is $[-2^{15} \text{ to } 2^{15}-1]$

C16 is a 16-bit pattern

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