

## 4 Integer arithmetic

(4.1) In C,

- Integer data is stored in blocks of
  - 2 bytes, ‘short’ or ‘short int’
  - 4 bytes, ‘int’,
  - 4 bytes, ‘long’ or ‘long int’ on a 32-bit machine,
  - 8 bytes, ‘long’ or ‘long int’ on a 64-bit machine.
- Short integers are never used except where memory is scarce.
- At *face value*, these integers have ranges

$$0 \dots 2^{16} - 1, \quad 0 \dots 2^{32} - 1, \quad 0 \dots 2^{64} - 1,$$

respectively. But *negative numbers* must be allowed for.

- Where the *high-order bit* is 1, a *negative number* is represented.

**Negative integers.** A short integer with high-order bit 1 represents the negative integer

$$v - 2^{16}$$

where  $v$  is its face value (from 32768 to 65536). Where the high-order bit is 0, it represents the number given by its face value.

Conversely, a negative number  $s$  is represented as a short integer as

$$s + 2^{16}.$$

This representation is called *2s complement*.

Thus if the face value is 32000 then this represents a positive number, but 33000 represents  $33000 - 2^{32} = -32536$ .

Short integers have values in the range

$$-2^{15} \dots 2^{15} - 1, \quad -32768 \dots 32767, \quad -(10000000000000000)_2 \dots (-80)_{16} \dots (ff)_{16}.$$

This is called *short integer range*.

In general, an  $N$ -bit integer has range

$$-2^{N-1} \dots 2^{N-1} - 1$$

The range of these kinds of integer are roughly

- 16 bit short int:  $\pm 32,000$ .
- 32 bit int;  $\pm 2$  billion.

- 64 bit long int:  $\pm 9,223,372,036,854,775,808$ .  
 $\pm 9 \times 10^{18}$

respectively.

**Converting a number  $s$ , in short integer range, to a short int.**

- If  $s$  is nonnegative, just convert it to 4 hex digits. Otherwise,
- Suppose  $s = -t$ .
- Convert  $t$  to 4 hexadecimal digits.
- Subtract from  $(fff)_{16}$ .
- Add 1.

**Example.** Convert 3276 to short int.

	204		12
16 )	3276	16 )	204
	32		16
	07		44
	0		32
	76	rmndr	12
	64		
rmndr	12		answer 0ccc

**Little endian.** The maths machines are all Dell computers using Intel processors, which store numerical data ‘little endian.’ That is, the *bytes* are stored low-order byte first, but within the byte face value is observed. The integer 3276 is stored as *cc 0c*.

**Example.** Convert  $-2768$  to short int.

	173		10		ffff
16 )	2768	16 )	173		- 0ad0
	16		16		f52f
	116		13		+ 1
	112		0		answer f530
	48	rmndr	13		
	48				
rmndr	0		2768 =	0a d0	

The answer is *f5 30*, or *30 f5* little-endian.

**Addition of ints.** The computer effectively adds  $N$ -bit integers modulo  $2^N$ , where  $N = 16, 32, 64$  for 16-, 32-, and 64-bit respectively.

Note that if  $x = 32,000$  then it is in short integer range, but  $x + x$  corresponds to a negative number. However,

**(4.2) Proposition** *If  $x$  and  $y$  are in short integer range, and  $x + y$  is in short integer range, then  $x + y$  will be computed correctly as short (16-bit) integers.*

*The same goes for 32-bit and 64-bit integers. (No proof.)* ■

**Example.** Convert 3276 and  $-2768$  to short integers, add them (big-endian), and convert the result to decimal.

```
3276    is    0ccc
-2768   is    f530
          01fc
```

One can check this by converting to decimal.

$01fc = (16 \times 1 + 15) \times 16 + 12 = 508$

```
3276
-2768
508 ----- results agree.
```

It is almost as easy to calculate 32-bit integer representations. For example, Convert 31415 and 31415 to int (big endian), and add

```
byte to hex 183 is b 7
byte to hex 122 is 7 a
byte to hex 0 is 0 0
byte to hex 0 is 0 0
31415: 00 00 7a b7
int 31415 hex  b7 7a 00 00
```

```
byte to hex 183 is b 7
byte to hex 122 is 7 a
byte to hex 0 is 0 0
byte to hex 0 is 0 0
31415: 00 00 7a b7
int 31415 hex  b7 7a 00 00
```

```
byte to hex 110 is 6 e
byte to hex 245 is f 5
byte to hex 0 is 0 0
byte to hex 0 is 0 0
The sum 62830: 00 00 f5 6e
int 62830 hex  6e f5 00 00
```