

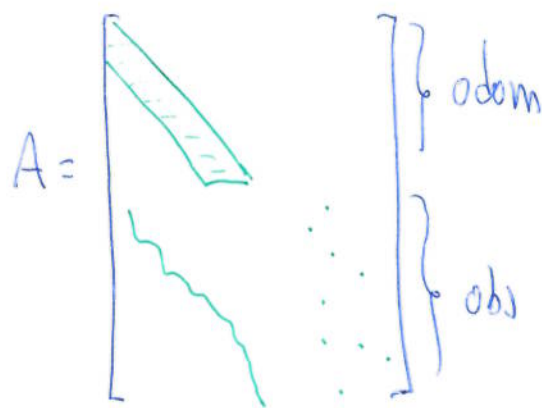
## L13: Square root SAM ( $\sqrt{\text{SAM}}$ )

\* Summary from L12: smoothing and mapping.

SAM as a factor graph (prev. Bayes network)

$$\theta^* = \underset{\theta}{\operatorname{argmax}} p(X, M | Z, U) = \underset{\theta}{\operatorname{argmin}} \left\{ \underbrace{\sum_{i=1}^M \| \cdot \|_{\xi_i}^2}_{\text{odom}} + \underbrace{\sum_{j=1}^K \| \cdot \|_{\xi_{jk}}^2}_{\text{obs}} \right\}$$

$$f^* = \underset{f}{\operatorname{argmin}} \|As - b\|_2^2 \quad (\text{LLSQ})$$



Adjacency matrix

$$A^T A = \Lambda \quad \text{information matrix}$$

Solve:  $As = b$  today's topic.

\* Normal equation

$$A^T A s = A^T b$$

$$s = (A^T A)^{-1} A^T b \quad \text{inversion } O(n^3)$$

- $A$  is sparse, we want to exploit sparsity
  - exploit problem knowledge
- { State of the art linear algebra methods

## → Cholesky factorization

$$\Lambda = A^T A = L \cdot L^T = R^T R$$

(information matrix)  $\Delta \nabla \Delta \nabla$

$$\text{chol} \begin{cases} A^T A \delta = A^T b \\ R^T R \delta = A^T b \end{cases}$$

Hence, square root methods

$$\begin{cases} R^T \cdot y = A^T b \\ R \cdot \delta = y \end{cases}$$

Solved efficiently by back-substitution

$$\begin{array}{|c|} \hline \text{ } \\ \hline \end{array} = \begin{array}{|c|} \hline \text{ } \\ \hline \end{array} \begin{array}{l} \text{back} \\ \text{var.} \rightarrow \text{easy} \end{array}$$

Ex:

$$\begin{bmatrix} 1 & 0 & 0 \\ 5 & 7 & 0 \\ 6 & -7 & 3 \end{bmatrix} \begin{bmatrix} y_1 \\ y_2 \\ y_3 \end{bmatrix} = \begin{bmatrix} 2 \\ 5 \\ 5 \end{bmatrix}$$

$$1) y_1 = 2$$

$$2) 5 \cdot 2 + 7 \cdot y_2 = 5 \\ y_2 = -\frac{5}{7}$$

$$3) 6y_1 - 7y_2 + 3y_3 = 5$$

$$6 \cdot 2 - 7 \cdot \left(-\frac{5}{7}\right) + 3y_3 = 5 \Rightarrow y_3 = \frac{5 - 12 - 5}{3} = -4$$

Cholesky factorization requires to solve 2 systems by back-substitution

QR factorization

$$Q^T A = \begin{bmatrix} R \\ 0 \end{bmatrix}$$

Q: orthonormal matrix (square)

R: Upper triangular matrix.

$$Q^T b = \begin{bmatrix} d \\ e \end{bmatrix}$$

$$\|As - b\|_2^2 = \|QAs - Q^T b\|_2^2 =$$

$$= \left\| \begin{bmatrix} R \\ 0 \end{bmatrix} s - \begin{bmatrix} d \\ e \end{bmatrix} \right\|_2^2 = \|Rs - d\|_2^2 + \|e\|_2^2$$

$$\boxed{Rs = d}$$

1 backsubstitution.

No need to calculate  $\Lambda = A^T A$

Computationally equivalent to Cholesky (for dense matrices)

Is R the same as for Cholesky?

$$\Lambda = A^T A = (QR)^T QR = R^T \underbrace{Q^T Q}_I R = R^T R$$

Cholesky with positive diagonal terms is unique.

\* Schur - complement (g2o, Kummerle' 2011)

$$\Lambda = A^T A = \begin{bmatrix} \Lambda_x & \Lambda_{xm} \\ \Lambda_{mx} & \boxed{\Lambda_m} \end{bmatrix}, \quad \delta = \begin{bmatrix} \delta_x \\ \delta_m \end{bmatrix}, \quad A^T b = \begin{bmatrix} b_x \\ b_m \end{bmatrix}$$

diagonal = easy to invert!

$$A^T A \delta = A^T b$$

$$\begin{cases} \Lambda_x \delta_x + \Lambda_{xm} \delta_m = b_x \\ \Lambda_{mx} \delta_x + \Lambda_m \delta_m = b_m \end{cases} \quad (\times \Lambda_{xm} \Lambda_m^{-1})$$

$$\begin{cases} \Lambda_x \delta_x + \Lambda_{xm} \delta_m = b_x \\ \Lambda_{xm} \Lambda_m^{-1} \Lambda_{mx} \delta_x + \cancel{\Lambda_{xm} \Lambda_m^{-1} \Lambda_m} \delta_m = \Lambda_{xm} \Lambda_m^{-1} b_m \end{cases}$$

$$(\Lambda_x - \Lambda_{xm} \Lambda_m^{-1} \Lambda_{mx}) \delta_x + 0 = b_x - \Lambda_{xm} \Lambda_m^{-1} b_m$$

Schur complement

1) solve  $S \boxed{\delta_x} = b_s$  using Cholesky for instance.

$$2) \Lambda_{mx} \delta_x + \Lambda_m \delta_m = b_m$$

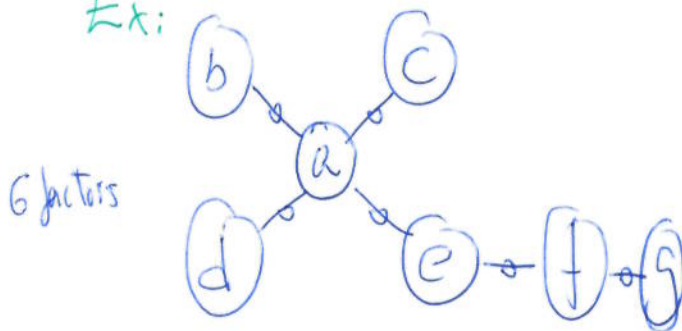
$$\delta_m = \Lambda_m^{-1} (b_m - \Lambda_{mx} \delta_x)$$

## \* Ordering of nodes to enhance solutions

Every graph has an optimal ordering of nodes:

- fewer edges when eliminating nodes  
(equivalent to backsubstitution in linear algebra)
- fewer fill-ins in the square root factorization.

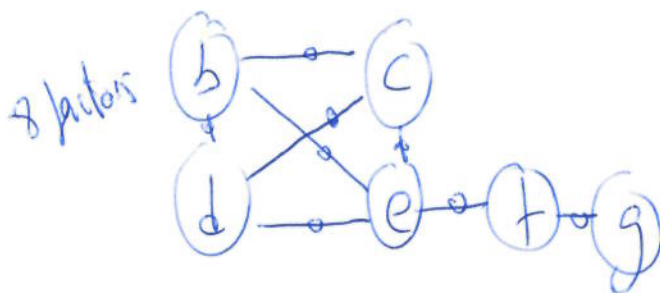
Ex:



$$\Lambda = A^T A$$

|   | a | b | c | d | e | f | g |
|---|---|---|---|---|---|---|---|
| a | x | x | x | x | x |   |   |
| b | x | x |   |   |   |   |   |
| c | x |   | x |   |   |   |   |
| d | x |   |   | x |   |   |   |
| e | x |   |   |   | x | x |   |
| f |   |   |   |   | x | x | x |
| g |   |   |   |   |   | x | x |

eliminate (a)



$$\Lambda'$$

|   | b | c | d | e | f | g |
|---|---|---|---|---|---|---|
| b | x | x | x | x |   |   |
| c | x | x | x | x |   |   |
| d | x | x | x | x |   |   |
| e | x | x | x | x | x | x |
| f |   |   |   | x | x | x |
| g |   |   |   |   | x | x |

New information matrix is denser than expected.

Adjacency matrix has more (factors) rows  $A'_{8 \times 6}$   
while before  $A_{6 \times 7}$



### \* Minimum order degree

Intuition: The number of edges captures how connected the node is and we want to minimize that.

Permute the nodes on a non-decreasing order (heuristic)  
(Solving the optimal ordering is NP-hard)

Ex:

|   |   |   |   |   |   |   |   |
|---|---|---|---|---|---|---|---|
|   | b | c | d | g | e | f | a |
| b | x |   |   |   |   |   | x |
| c |   | x |   |   |   |   | x |
| d |   |   | x |   |   |   | x |
| g |   |   |   | x |   |   |   |
| e |   |   |   |   | x | x | x |
| f |   |   |   | x | x | x |   |
| a | x | x | x |   | x | x |   |

This ordered (permuted)  
 $\Lambda_P = \Lambda(p, p)$  provides  
less fill-ins during  
factorization

col perm  $(A^T A)$  Cholesky

col perm  $(A)$  QR

COLAMD heuristic for ordering nodes (Davis' 2001)

As reported in Dellavent' 2006 they performed better using Cholesky and colamd. But it is problem dependent.

# \* Incremental square root factorization (iSAM Kess' 2008)

As new observations are available, update  $A, R$  without recalculating everything.

## QR factorization incrementally.

$$A = QR, \quad R = Q^T A$$

New obs:

$$\begin{bmatrix} Q^T \\ 1 \end{bmatrix} = \begin{bmatrix} A \\ W^T \end{bmatrix} = \begin{bmatrix} R \\ W^T \end{bmatrix}$$

$$b = \begin{bmatrix} d \\ e \end{bmatrix} \quad \text{accordingly update the vector } d.$$

$$\text{odom: } W^T = \begin{bmatrix} G_i^{i-1} & -I \end{bmatrix} \quad (\text{sparse})$$

$$\text{obs: } W^T = \begin{bmatrix} H^i & J^i \end{bmatrix}$$

$$R' = \begin{bmatrix} R & \\ & W^T \end{bmatrix} \quad \text{if new odometry we increase the dimensions}$$

## Givens rotations

$$\Phi = \begin{bmatrix} \cos \phi & \sin \phi \\ -\sin \phi & \cos \phi \end{bmatrix}$$

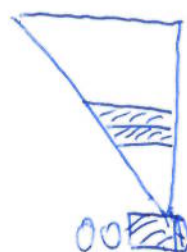
$$\text{such as } \Phi \begin{bmatrix} a & c \\ b & d \end{bmatrix} = \begin{bmatrix} r & s \\ 0 & t \end{bmatrix}$$

A sequence of given rotations produce the QR decomposition

$$I_i \begin{bmatrix} 1 & & & 0 \\ & \ddots & & \\ & & c & s \\ 0 & & -s & c \end{bmatrix} \begin{array}{c} \text{Diagram 1: A square with a diagonal line from top-left to bottom-right. A small shaded rectangle is at the bottom-left corner, representing a rotation in the first two dimensions.} \end{array} = \begin{array}{c} \text{Diagram 2: A square with a diagonal line from top-left to bottom-right. A shaded rectangle is at the bottom-left corner, and another shaded rectangle is at the top-right corner, representing a rotation in the last two dimensions.} \end{array}$$

$\Phi$

We keep applying Givens below the diagonal until we get an upper triangular matrix



iSAM algorithm:

- 1: New information  $w^T$ , update  $\begin{bmatrix} R \\ w^T \end{bmatrix}$
- 2: Givens rotations until  $R'$  upper triangular. Update  $d'$
- 3: Solve  $R'S = d'$

Discussion

- It is only possible for some time, eventually we need recalculate RA
- The ordering is important for next incremental poses. May produce fill-ins  $\rightarrow$  iSAM2 (Kaess 2011) and the Bayes tree graph.