

## **Basic blocks and Traces**

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## **After translation**

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To simplify translation there are mismatches between tree code and actual machine instructions:

1. CJUMP to two labels; machine conditionals fall through on false
2. ESEQ and CALL order evaluation of subtrees for side-effects – constrains optimization
3. CALL as argument to another CALL causes interference between register arguments

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## **Canonical trees**

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Rewrite into an equivalent canonical form:

- SEQ can only be subtree of another SEQ
- SEQs clustered at top of tree
- might as well turn into simple linear list of statements

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## **Basic blocks and traces**

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3-stage transformation:

1. to linear list of *canonical trees* without SEQ/ESEQ
2. to *basic blocks* with no internal jumps or labels
3. to *traces* with every CJUMP immediately followed by false target

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## Linear trees

1. No SEQ or ESEQ nodes
2. A CALL can only be a subtree of an EXP(...) or a MOVE(TEMP t,...)

Transformations:

- lift ESEQs up tree until they can become SEQs
- turn SEQs into linear list

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## Taming conditional branches

1. Form *basic blocks*: sequence of statements always entered at the beginning and exited at the end:
  - first statement is a LABEL
  - last statement is a JUMP or CJUMP
  - contains no other LABELs, JUMPS or CJUMPS
2. Order blocks into *trace*:
  - every CJUMP followed by false target
  - JUMPs followed by target, if possible, to eliminate JUMP

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## Linearizing trees

$$\begin{aligned}
 & \text{ESEQ}(s_1, \text{ESEQ}(s_2, e)) &&= \text{ESEQ}(\text{SEQ}(s_1, s_2), e) \\
 & \text{BINOP}(op, \text{ESEQ}(s, e_1), e_2) &&= \text{ESEQ}(s, \text{BINOP}(op, e_1, e_2)) \\
 & \text{MEM}(\text{ESEQ}(s, e_1)) &&= \text{ESEQ}(s, \text{MEM}(e_1)) \\
 & \text{JUMP}(\text{ESEQ}(s, e_1)) &&= \text{SEQ}(s, \text{JUMP}(e_1)) \\
 & \text{CJUMP}(op, \text{ESEQ}(s, e_1), e_2, l_1, l_2) &&= \text{SEQ}(s, \text{CJUMP}(op, e_1, e_2, l_1, l_2)) \\
 & \text{BINOP}(op, e_1, \text{ESEQ}(s, e_2)) &&= \text{ESEQ}(\text{MOVE}(\text{TEMP } t, e_1), \text{ESEQ}(s, \text{BINOP}(op, \text{TEMP } t, e_2))) \\
 & \text{CJUMP}(op, e_1, \text{ESEQ}(s, e_2), l_1, l_2) &&= \text{SEQ}(\text{MOVE}(\text{TEMP } t, e_1), \text{SEQ}(s, \text{CJUMP}(op, \text{TEMP } t, e_2, l_1, l_2))) \\
 & \text{MOVE}(\text{ESEQ}(s, e_1), e_2) &&= \text{SEQ}(s, \text{MOVE}(e_1, e_2)) \\
 & \text{CALL}(f, a) &&= \text{ESEQ}(\text{MOVE}(\text{TEMP } t, \text{CALL}(f, a)), \text{TEMP}(t))
 \end{aligned}$$

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## Basic blocks

*Control flow analysis* discovers basic blocks and control flow between them:

1. scan from beginning to end:
  - LABEL *l* starts a new block and previous block ends (append JUMP *l* if necessary)
  - JUMP or CJUMP ends a block and starts next block (prepend new LABEL if necessary)
2. prepend new LABELs to blocks with non-LABEL at beginning
3. append JUMP(NAME done) to last block

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## Traces

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1. Pick an untraced block, the start of some trace
2. Follow a *possible* execution path, choosing false targets first
3. Repeat until all blocks are traced

*Cleaning up:*

- CJUMP followed by true target: switch targets, negate condition
- CJUMP( $o, a, b, l_t, l_f$ ) followed by neither  $l_t$  nor  $l_f$ :
  1. create new  $l'_f$
  2. rewrite as CJUMP( $o, a, b, l_t, l'_f$ ), LABEL  $l'_f$ , JUMP  $l_f$
- JUMP  $l$ , LABEL  $l \rightarrow$  LABEL  $l$