

Introduction to Homotopy Type Theory

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Ways to introduce (homotopy) type theory

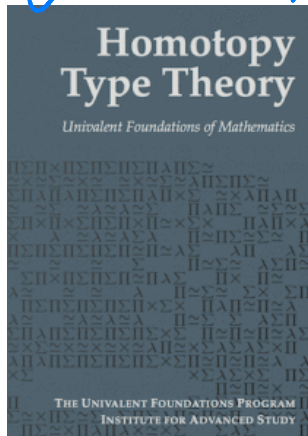
↪ Agda

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x·-mono : ∀ {x y z : Brw} → y ≤ z → x · y ≤ x · z
x·-mono ≤-zero = ≤-zero
x·-mono (≤-trans ysw wsz) = ≤-trans (x·-mono ysw) (x·-mono wsz)
x·-mono {x} (≤-succ-mono {y} {z} y≤z) = +x-mono x (x·-mono {x} y≤z)
x·-mono {x} {y} (≤-cocone f {k = k} y≤z) with decZero x
... | yes x≡zero = subst (λ z → z ≤ zero) (sym (zero·y≡zero y)) • cong
... | no x≠zero = ≤-cocone (λ n → x · f n) {k = k} (x·-mono y≤z)
x·-mono {x} (≤-limiting f f≤z) with decZero x
... | yes x≡zero = ≤-zero
... | no x≠zero = ≤-limiting (λ n → x · f n) λ k → x·-mono (f≤z k)
x·-mono (≤-trunc p q i) = ≤-trunc (x·-mono p) (x·-mono q) i
```

↪ inference rules

$$\frac{\Gamma, x:1 \vdash C : \mathcal{U}_i \quad \Gamma \vdash c : C[\star/x] \quad \Gamma \vdash a : 1}{\Gamma \vdash \text{ind}_1(x.C, c, a) : C[a/x]} \text{1-ELIM}$$

↪ natural
math language



Prerequisites

- Basic examples of types

Nat Bool Unit Empty

- Universe(s)

Type

$A, B : \text{Type}$

- Function types

$A \rightarrow B$

- Dependent function types

if $C : A \rightarrow \text{Type}$, then

- Binary products (pairs) and coproducts (sums)

$A \times B$

$A \cup B$

- Dependent pairs

$\sum (a:A). C a$

$(a:A) \rightarrow C a$
 $[\lambda a a \quad \prod (a:A). C a]$

- Inductive types

see above

Program = Proof

Exercise: Formulate a type which says that there are infinitely many primes
(or simply: arbitrarily large primes)!

isPrime : $\text{Nat} \rightarrow \text{Type}$

isPrime(n) \equiv $\left((m\ k : \text{Nat}) \rightarrow (m \cdot k = n) \rightarrow (m=1) \vee (k=n) \right)$
 $\times (n > 1)$

bigPrimes : $(n : \text{Nat}) \rightarrow \sum (p : \text{Nat}). \text{isPrime}(p)$

bigPrimes \equiv (exercise) $\times p > n$

Caveat: judgments (meta-theoretic)

- A type "A is a valid type"
- $a : A$ "a is term of type A"
- $a \equiv b$ "a and b are the same" $(\lambda x. fx)y$
- $a : \equiv b$ same, by definition $\equiv fy$

We cannot ask or prove these statements *in* the language.

Exercise: Which of the following can we ask internally?

- (1) Is 8 a natural number? no, b/c $8 : \text{Nat}$ is judgment
- (2) Is 8 a prime number? yes, b/c $\text{isPrime}(8)$
- (3) Is "hello" a prime number? no, $\text{isPrime}(\text{"hello"})$
is not a type

(Identity a.k.a. equality a.k.a. identification a.k.a. path) types

- if $a, b : A$ then $a = b$ type "formation"
- given $a : A$ we have $\text{refl} : a = a$ "introduction"
- given $C : (a b : A) \rightarrow (a = b) \rightarrow \text{Type}$ "elimination"
 and $C a a \text{ refl}$ (for all $a : A$)
 we get $C a b p$ (for all $a, b, p : a = b$) } aka path induction
- applying this rule and asking for the refl case gives back the assumption "computation"

$\text{sym} : (a b : A) \rightarrow a = b \rightarrow b = a$

$\text{sym } a = a$
 $\text{refl} := \text{refl}$

here:

$C a b p := (b = a)$

computation: $\text{sym } a a \text{ refl} \equiv \text{refl}$

Model: types as spaces

The homotopical interpretation of dependent type theory works roughly like this:

judgment

interpretation

A type

" A is a space"

$a : A$

" a is a point in space A "

$c \equiv d$

path
 \hookrightarrow in A

" c and d are the same point"

$p : a = b$

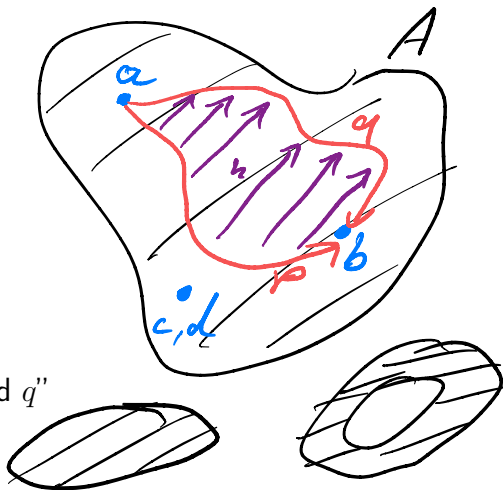
" p is a path between a and b "

$h : p = q$

if p and q are equalities $a = b$:

" h is a homotopy between p and q "

\uparrow
path in
 $a = b$

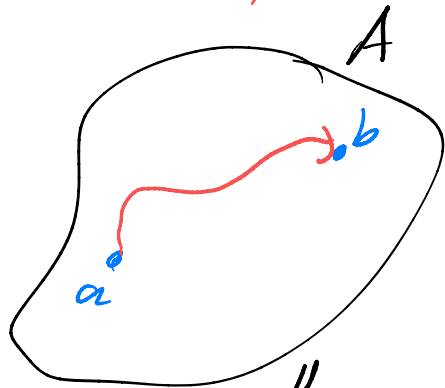


Types as spaces, examples: symmetry, transitivity

$\text{sym} : (a\ b : A) \rightarrow a = b \rightarrow b = a$

$\text{sym } a\ ~~b~~$
 a

~~p~~ $\text{refl} \equiv \text{refl}$

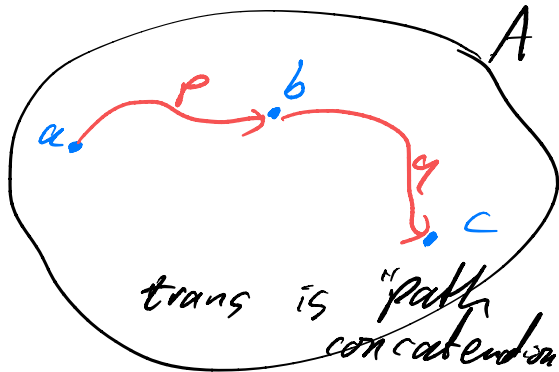


sym is "path reversal"

$\text{trans} : (a\ b\ c : A) \rightarrow$

$a = b \rightarrow b = c \rightarrow a = c$

$\text{trans } a\ ~~b~~\ c$
 a ~~p~~ $q \equiv q$
 refl



trans is "path concatenation"

Types as spaces, example: uniqueness of identity proofs

Can we prove this for all types A ?

VIP

$uip : (a b : A) \rightarrow (p q : a = b) \rightarrow p = q$

$uip\ a\ ~~b~~\ ~~p~~\ q$
 $\quad\quad\quad a\ refl$
 $\quad\quad\quad : \equiv ?$

