

8 2019 04 02 Multi-stage games

Definition 42. A multi-stage game consists of the following components:

- A sequence $(0, \dots, K)$ represents the stage of the game;
- there are N players who play repeatedly and chose their actions simultaneously at each stage of the game;
- Each player observes the actions of all N players in all previous stages;
- For each stage $k \in \{0, \dots, K\}$, the action profile of players is given by

$$\underline{\alpha}_k = (\alpha_k^1, \dots, \alpha_k^N)$$

where α_k^i is the action chosen by player i at stage k ;

- The available information prior to stage k , called *history*, is defined as

$$h^k = (\underline{\alpha}_0, \dots, \underline{\alpha}_{k-1});$$

- For every player i at stage k , the feasible set of actions is denoted by $A^i(h^k)$, which depends on the information of history prior to stage k . For example, the resources that have been used in the past can affect the current actions of players. So, at stage k , the action chosen by player i can be viewed as a function of history h^k ;
- We would like to include “no action” as part of $A^i(h^k)$ for every $i = 1, \dots, N$ and for every $k = 0, \dots, K$. When player i choses “no action”, she is inactive at this stage.

Here are some examples for multi-stage games:

1. $K = 0$, Cournot competition.
2. $K = 1$, Stackelberg game. For example $N = 2$, one player is called the leader, and the other called the follower. Then
 - at stage $k = 0$, leader chooses an action $\alpha_0 \in A^{leader}(\emptyset)$. The follower is inactive, namely $A^{follower}(\emptyset) = \{\text{“no action”}\}$.
 - at stage $k = 1$, leader is inactive $A^{leader}(h^1) = \{\text{“no action”}\}$ and the follower chooses an action $\alpha^{follower} \in A^{follower}(h^1)$ where $h^1 = ((\alpha_0, \text{“no action”}))$.
3. $K = 2$, Entry-deterrence game.
 - at stage $k = 0$, incumbent raises money α_0 , entrant “no action”;
 - at stage $k = 1$, incumbent “no action”, entrant choses between “enter” or “not enter”;

- at stage $k = 2$, Cournot competition.

Definition 43. A pure strategy for player $i \in \{1, \dots, N\}$ at stage $k \in \{0, \dots, K\}$ is denoted by $\hat{\alpha}_k^i$.

A pure strategy profile, denoted by $\hat{\alpha}$, is the collection of pure strategies of all players at all stages except the last one $k = K$, namely

$$\hat{\alpha} = (\hat{\alpha}_k^i)_{i=1, \dots, N, k=0, \dots, K-1}.$$

It is said to be an admissible pure strategy profile if $\hat{\alpha}_k^i \in A^i(h^k)$ for all $k = 0, \dots, K-1$ and $i = 1, \dots, N$.

Remark 30. A stage $k = K$, we collect everything and decide the reward or cost for players.

Definition 44. The set of histories up to stage $k \in \{0, \dots, K\}$ is denoted by $H(k)$. So $H(K)$ represents the set of all histories. For each player $i \in \{1, \dots, N\}$, we associate her with a cost function $J^i : H(K) \rightarrow \mathbb{R}$.

Very often, the cost function J^i is additive, i.e.

$$J^i(h^K) = \sum_{k=0}^{K-1} c_k f^i(\alpha_k^i)$$

The coefficients $(c_k)_{k=0, \dots, K-1}$ can be discounted factors, for example, $c_k = c_K \delta^k$ for some $c_K \in \mathbb{R}$ and $\delta \in [0, 1]$.

When $K = \infty$, we need to make sure $\delta < 1$.

Definition 45. A pure strategy profile $\hat{\alpha}$ is said to be a pure strategy Nash equilibrium if for every $i \in \{1, \dots, N\}$, for every admissible sequence of actions $\alpha = (\alpha_0, \dots, \alpha_{K-1})$,

$$J^i(\hat{\alpha}_0, \dots, \hat{\alpha}_{K-1}) \leq J^i((\alpha_0, \hat{\alpha}_0^{-i}), \dots, (\alpha_{K-1}, \hat{\alpha}_{K-1}^{-i}))$$

where $(\alpha_k, \hat{\alpha}_k^{-i}) = (\hat{\alpha}_k^1, \dots, \hat{\alpha}_k^{i-1}, \alpha_k, \hat{\alpha}_k^{i+1}, \dots, \hat{\alpha}_k^N)$ is the action profile at stage k for $k \in \{0, \dots, K-1\}$.

Remark 31. Given a strategy profile $\hat{\alpha}$ and a player i , a sequence of actions $\alpha = (\alpha_0, \dots, \alpha_{K-1})$ is said to be admissible if for every $k = 1, \dots, K-1$,

$$\alpha_k \in A^i(h_{\alpha}^k)$$

where $h_{\alpha}^0 = \emptyset$ and $h_{\alpha}^k = ((\alpha_0, \hat{\alpha}_0^{-i}), \dots, (\alpha_{k-1}, \hat{\alpha}_{k-1}^{-i}))$

Definition 46. For every stage $k = 0, \dots, K$, let $G(h^k)$ be a new game depending on the history $h^k = (\alpha_0, \dots, \alpha_{k-1})$.

- The new strategy profile for N players in game $G(h^k)$ is denoted by $\alpha_{h^k} = (\alpha_l^i)_{i=1, \dots, N, l=k, k+1, \dots, K-1}$ (similarly, $\hat{\alpha}_{h^k}$ for pure strategy profile);

- The new history is denoted by $(h^k, \underline{\alpha}_k, \dots, \underline{\alpha}_{K-1})$;
- The new cost function is a mapping $(\underline{\alpha}_k, \dots, \underline{\alpha}_{K-1}) \mapsto J^i(h^k, \underline{\alpha}_k, \dots, \underline{\alpha}_{K-1})$

A strategy profile $\hat{\alpha}$ is said to be a *sub-game perfect Nash equilibrium* if for every $k \in \{0, \dots, K\}$, for every $h^k \in H(k)$, the strategy profile $(\hat{\alpha}_l^i)_{i=1, \dots, N, l=k, \dots, K-1}$ is a Nash equilibrium for the sub game $G(h^k)$.

In words, at every stage k , whatever the history is, if players play according to the strategy profile $(\hat{\alpha}_l^i)_{i=1, \dots, N, l=k, \dots, K-1}$ afterwards, then they are in a Nash equilibrium associated to the new game.

Remark 32. We need to clarify the notion of closed loop strategy and open loop strategy. Let α_k^i be the action taken by player i at stage k . For the sake of simplicity, we assume that $A^i(h^k) = \mathbb{R}$ for every $h^k \in H(k)$ for every $k = 0, \dots, K$.

- α_k^i is said to be an open loop strategy if there exists a function $\varphi^i : \{0, \dots, K\} \rightarrow \mathbb{R}$ such that $\alpha_k^i = \varphi^i(k)$ for every $k = 0, \dots, K$. The function φ^i is decided before the game.
- α_k^i is said to be a closed loop strategy if there exists a function $\varphi_k^i : H(k) \rightarrow \mathbb{R}$ such that $\alpha_k^i = \varphi_k^i(h^k)$ for some $h^k \in H(k)$.

It will be rare for open loop strategy profile to have a sub-game perfect Nash equilibrium.

8.1 Prisoner's dilemma

Assume that there are only two players $N = 2$, and the feasible set of actions are $A_1 = A_2 = \{C, D\}$ for every stage k of the game. Here C stands for cooperation and D stands for defection.

Let $T < R < P < S$ where $T, R, P, S \in \mathbb{R}$ and T stands for *temptation*, R stands for *negative rewards* or *cost*, P stands for *punishment*, and S stands for *sucker*. Assume that the cost functions are defined as follow (see also the table 1):

$$J^1(\alpha_1, \alpha_2) = \mathbb{1}_{\alpha_1=C}(R\mathbb{1}_{\alpha_2=C} + T\mathbb{1}_{\alpha_2=D}) + \mathbb{1}_{\alpha_1=D}(S\mathbb{1}_{\alpha_2=C} + P\mathbb{1}_{\alpha_2=D})$$

and

$$J^2(\alpha_1, \alpha_2) = \mathbb{1}_{\alpha_2=C}(R\mathbb{1}_{\alpha_1=C} + T\mathbb{1}_{\alpha_1=D}) + \mathbb{1}_{\alpha_2=D}(S\mathbb{1}_{\alpha_1=C} + P\mathbb{1}_{\alpha_1=D}).$$

The objective of players are to minimize their own cost function. We observe that the choice of relationships $T < R$ and $P < S$ imply that the defection is a preferable strategy for both player 1 and 2.

It can be verified that the strategy (D, D) is a Nash equilibrium.

Lemma 14. The strategy profile $((D, D), \dots, (D, D))$ is a sub-game perfect Nash equilibrium.

player 2 \ palyer 1	C	D
C	(R, R)	(T, S)
D	(S, T)	(P, P)

Table 1: the cost function table. The tuple $(J^1(\cdot), J^2(\cdot))$ represents the costs of player 1 and player 2 in different situation. The column names represent the strategies chosen by player 1 and the row names are for player 2.

Proof. Step 1: If I am looking for a sub-game perfect Nash equilibrium, the last choice should be (D, D) whatever the history is.

Step 2: One step perturbation principle (for sub-game perfect Nash equilibrium) for the induction part of the proof. \square

Remark 33. The result may be different if the number of stages equals to infinity, i.e. $K = \infty$.