Team notebook

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I BitTricks

```
y = x & (x-1) // Turn off rightmost 1bit
y = x & (-x) // Isolate rightmost 1bit
y = x | (x-1) // Right propagate rightmost 1bit(fill in 1s)
y = x | (x+1) // Turn on rightmost 0bit
y = ~x & (x+1) // Isolate rightmost 0bit
```

```
// If x is of long type, use __builtin_popcountl(x)
// If x is of long long type, use __builtin_popcountll(x)
// 1. Counts the number of ones(set bits) in an integer.
__builtin_popcount(x)
// 2. Checks the Parity of a number. Returns true(1) if the
// number has odd number of set bits, else it returns
// false(0) for even number of set bits.
 __builtin_parity(x)
// 3. Counts the leading number of zeros of the integer.
__builtin_clz(x)
// 4. Counts the trailing number of zeros of the integer.
__builtin_ctz(x)
// 5. Returns 1 + the index of the least significant 1-bit.
__builtin_ffs(x) // If x == 0, returns 0.
// Iterate over non empty subsets of bitmask
for(int s=m;s;s=(s-1)&m) // Decreasing order
for(int s=0;s=s-m&m;) // Increasing order
```

2 Catalan

```
/*Catalan, counts the number of ways of:
( A ) B, where |A|+|B| = N, for N+1
*/

const int MOD = 1e9 + 7;
ll mul(ll x, ll y) { return (x*y)%MOD; }
ll pot(ll x, ll y) {
   if(y==0) return 1;
   ll ans = pot(x,y/2);
   ans = mul(ans,ans);
```

```
if (y&1)ans=mul(ans,x);
    return ans;
}
11 inv(11 x) { return pot(x, MOD-2); }
// mxN it the double of the max input 'n'
const int mxN = 2e6 + 10;
vl fact(mxN,1);
void init() {
   for (int i =1;i<=mxN;i++) {</pre>
       fact[i] = mul(fact[i-1],i);
    }
}
11 catalan(ll n) {
    if (n<0) return 0;</pre>
    11 up = fact[2*n];
   ll down = mul(fact[n],fact[n+1]);
   return mul(up,inv(down));
```

3 CatalanConvolution

```
/*
Return Catalan Convolution.

Convolution for k=3
((( A ) B ) C ) D

Where A + B + C + D = N, for N + 1
*/

const int MOD = 1e9 + 7;
ll mul(ll x, ll y) { return (x*y)%MOD; }
ll pot(ll x, ll y) {
   if(y==0) return 1;
   ll ans = pot(x,y/2);
   ans = mul(ans,ans);
   if (y&1)ans=mul(ans,x);
   return ans;
}
ll inv(ll x) { return pot(x, MOD-2); }
```

```
// mxN it the double of the max input N, plus max K
const int mxN = 2e6 + 1e6 + 10;
vl fact(mxN,1);
11 cnk(ll n, ll k) {
   if (k < 0 || n < k) return 0;</pre>
   11 nOverK = mul(fact[n],inv(fact[k]));
   return mul(nOverK,inv(fact[n-k]));
void init() {
   for (int i =1;i<=mxN;i++) {</pre>
       fact[i] = mul(fact[i-1],i);
}
// for parethesis example
// number of n+k pairs having k open parethesis at beginning
// (cnk(2n+k,n)*(k+1))/(n+k+1)
ll catalanCov(ll n, ll k) {
   11 up = mul(cnk(2*n+k,n),(k+1)\%MOD);
   ll\ down = (n+k+1)\%MOD;
   return mul(up,inv(down));
}
6
(()
ans: 2
// size, and prefix
ll countParenthesisWithPrefix(ll n, string &p) {
   if (n&1) return 0;
   11 k = 0;
   for (auto c : p) {
       if (c=='(') k++;
       else k--;
       if (k<0) return 0;</pre>
   n=(n-(11)p.size()-k)/2;
   return catalanCov(n,k);
```

4 ClosestPairOfPoints

```
// It seems O(n log n), not sure but it worked for 50000
#define x first
#define y second
long long dist2(pair<int, int> a, pair<int, int> b) {
       return 1LL * (a.x - b.x) * (a.x - b.x) + 1LL * (a.y - b.y) * (a.y)
pair<int, int> closest_pair(vector<pair<int, int>> a) {
       int n = a.size();
       assert(n >= 2);
       vector<pair<int, int>, int>> p(n);
       for (int i = 0; i < n; i++) p[i] = {a[i], i};
       sort(p.begin(), p.end());
       int 1 = 0, r = 2;
       long long ans = dist2(p[0].x, p[1].x);
       pair<int, int> ret = {p[0].y, p[1].y};
       while (r < n) {
              while (1 < r \&\& 1LL * (p[r].x.x - p[1].x.x) * (p[r].x.x -
                   p[1].x.x) >= ans) 1++;
              for (int i = 1; i < r; i++) {</pre>
                      long long nw = dist2(p[i].x, p[r].x);
                      if (nw < ans) {</pre>
                             ans = nw;
                             ret = {p[i].y, p[r].y};
                      }
              }
              r++;
       return ret:
}
// Tested: https://vjudge.net/solution/52922194/ccPUXODAMWTzpzCEvXbV
void test_case() {
   11 n:
   cin >> n;
   vector<pair<int,int>> points(n);
   for (int i = 0;i<n;i++) cin >> points[i].x >> points[i].y;
   auto ans = closest_pair(points);
   cout << fixed << setprecision(6);</pre>
   if (ans.F > ans.S) swap(ans.F,ans.S);
   ld dist = sqrtl(dist2(points[ans.F],points[ans.S]));
```

```
cout << ans.F << " " << ans.S << " " << dist << endl;
}</pre>
```

5 ConvexHullTrick

```
* Author: Simon Lindholm
 * Date: 2017-04-20
 * License: CCO
 * Source: own work
 * Description: Container where you can add lines of the form kx+m, and
     query maximum values at points x.
 * Useful for dynamic programming (''convex hull trick'').
 * Time: O(\log N)
 * Status: stress-tested
 // For minimum you can multiply by -1 'k' and 'm' when adding, and the
     answer when querying.
 // Tested in https://atcoder.jp/contests/dp/submissions/55836691
#pragma once
struct Line {
       mutable ll k, m, p;
       bool operator<(const Line& o) const { return k < o.k; }</pre>
       bool operator<(ll x) const { return p < x; }</pre>
};
struct LineContainer : multiset<Line, less<>>> {
       // (for doubles, use inf = 1/.0, div(a,b) = a/b)
       static const ll inf = LLONG_MAX;
       ll div(ll a, ll b) { // floored division
              return a / b - ((a ^ b) < 0 && a % b); }
       bool isect(iterator x, iterator y) {
              if (y == end()) return x->p = inf, 0;
              if (x->k == y->k) x->p = x->m > y->m ? inf : -inf;
              else x->p = div(y->m - x->m, x->k - y->k);
              return x->p >= y->p;
       void add(ll k. ll m) {
              auto z = insert(\{k, m, 0\}), y = z++, x = y;
              while (isect(v, z)) z = erase(z);
```

6 DP-Mask-Over-Submasks

```
// DP of submask over submasks
// 0(3^n)
// j&(-j); get a '1' bit of j
// for (int j=i;j;j = (j-1)\&i){...} j is submask of 'i' the mask
11 dp[1<<18]; // answer of mask
ll cst[1<<18]; // cost of use a submask
ll a[18][18]; // elements
ll pos[1<<18]; // trick to get fast the pos
void test_case() {
    11 n;
    cin >> n;
    for (int i =0;i<n;i++) {</pre>
       for (int j =0; j<n; j++) {</pre>
           cin >> a[i][j];
    }
    for (int i =0;i<n;i++) {</pre>
       pos[1 << i] = i;
    for (int i = 0; i < (1 < n); i++) {
       11 j = i;
       vl idxs;
       while (j) {
           11 k = j\&(-j);
           idxs.pb(pos[k]);
           j^=k;
       }
```

```
for (int j = 0; j < idxs.size(); j ++) {
    for (int k = j + 1; k < idxs.size(); k ++) {
        cst[i] += a[idxs[j]][idxs[k]];
    }
    dp[i] = cst[i];
    for (int j = i; j; j = (j - 1)&i) {
        dp[i] = max(dp[i], cst[j] + dp[i^j]);
    }
}
cout << dp[(1 << n) - 1] << "\n";
}</pre>
```

7 Optimized-Polard-Rho

```
#define fore(i, b, e) for(int i = b; i < e; i++)</pre>
11 gcd(ll a, ll b){return a?gcd(b%a,a):b;}
11 mulmod(ll a, ll b, ll m) {
 11 r=a*b-(11)((long double)a*b/m+.5)*m;
 return r<0?r+m:r;</pre>
ll expmod(ll b, ll e, ll m){
 if(!e)return 1;
 11 q=expmod(b,e/2,m);q=mulmod(q,q,m);
 return e&1?mulmod(b,q,m):q;
bool is_prime_prob(ll n, int a){
 if(n==a)return true;
 11 s=0, d=n-1;
 while (d\%2==0)s++, d/=2;
 11 x=expmod(a,d,n);
  if((x==1)||(x+1==n))return true;
  fore(_,0,s-1){}
   x=mulmod(x,x,n);
   if(x==1)return false;
   if(x+1==n)return true;
 return false;
bool rabin(ll n){ // true iff n is prime
 if(n==1)return false;
  int ar[]={2,3,5,7,11,13,17,19,23};
```

```
fore(i,0,9)if(!is_prime_prob(n,ar[i]))return false;
 return true;
}
// optimized version: replace rho and fact with the following:
const int MAXP=1e6+1; // sieve size
int sv[MAXP]: // sieve
11 add(11 a, 11 b, 11 m){return (a+=b)<m?a:a-m;}</pre>
ll rho(ll n){
 static ll s[MAXP];
  while(1){
   11 x=rand()%n,y=x,c=rand()%n;
   ll *px=s,*py=s,v=0,p=1;
   while(1){
     *py++=y=add(mulmod(y,y,n),c,n);
     *py++=y=add(mulmod(y,y,n),c,n);
     if((x=*px++)==y)break;
     ll t=p;
     p=mulmod(p,abs(y-x),n);
     if(!p)return gcd(t,n);
     if(++v==26){
       if((p=gcd(p,n))>1&&p<n)return p;</pre>
       v=0;
     }
   }
   if(v&&(p=gcd(p,n))>1&&p<n)return p;</pre>
 }
}
void init_sv(){
 fore(i,2,MAXP)if(!sv[i])for(ll j=i;j<MAXP;j+=i)sv[j]=i;</pre>
void fact(ll n, map<ll,int>& f){ // call init_sv first!!!
 for(auto&& p:f){
   while (n\%p.F==0) {
     p.S++; n/=p.F;
   }
 }
 if(n<MAXP)while(n>1)f[sv[n]]++,n/=sv[n];
 else if(rabin(n))f[n]++;
 else {ll q=rho(n);fact(q,f);fact(n/q,f);}
```

8 PointInConvexPolygon

```
11 IN = 0:
11 \text{ ON} = 1;
11 \text{ OUT} = 2;
vector<string> ANS = {"IN", "ON", "OUT"};
#define pt pair<11,11>
#define x first
#define v second
pt sub(pt a, pt b) { return {a.x - b.x, a.y - b.y}; }
11 cross(pt a, pt b) { return a.x*b.y - a.y*b.x; } // x = 180 -> sin = 0
11 orient(pt a, pt b, pt c) { return cross(sub(b,a),sub(c,a)); }//
    clockwise = -
// poly is in clock wise order
ll insidePoly(vector<pt> &poly, pt query) {
    11 n = poly.size();
   ll left = 1;
   11 \text{ right} = n - 2;
    11 \text{ ans} = -1;
   if (!(orient(poly[0], poly[1], query) <= 0</pre>
         && orient(poly[0], poly[n-1], query) >= 0)) {
       return OUT;
    while (left <= right) {</pre>
       11 mid = (left + right) / 2;
       if (orient(poly[0], poly[mid], query) <= 0) {</pre>
           left = mid + 1;
           ans = mid;
       } else {
           right = mid - 1;
       }
    left = ans;
    right = ans + 1;
    if (orient(poly[left], query, poly[right]) < 0) {</pre>
       return OUT;
    }
    if (orient(poly[left], poly[right], query) == 0
       || (left == 1 && orient(poly[0], poly[left], query) == 0)
       || (right == n-1 && orient(poly[0], poly[right], query) == 0)) {
       return ON;
    return IN;
```

}

9 PolygonDiameter

```
// Given a set of points, it returns
// the diameter (the biggest distance between 2 points)
// tested: https://open.kattis.com/submissions/13937489
const double eps = 1e-9;
int sign(double x) { return (x > eps) - (x < -eps); }</pre>
struct PT {
   double x, y;
   PT() \{ x = 0, y = 0; \}
   PT(double x, double y) : x(x), y(y) {}
   PT operator - (const PT &a) const { return PT(x - a.x, y - a.y); }
   bool operator < (PT a) const { return sign(a.x - x) == 0 ? y < a.y :</pre>
        x < a.x; 
   bool operator == (PT a) const { return sign(a.x - x) == 0 && sign(a.y
        - v) == 0; }
};
inline double dot(PT a, PT b) { return a.x * b.x + a.y * b.y; }
inline double dist2(PT a, PT b) { return dot(a - b, a - b); }
inline double dist(PT a, PT b) { return sqrt(dot(a - b, a - b)); }
inline double cross(PT a, PT b) { return a.x * b.y - a.y * b.x; }
inline int orientation(PT a, PT b, PT c) { return sign(cross(b - a, c -
    a)); }
double diameter(vector<PT> &p) {
   int n = (int)p.size();
   if (n == 1) return 0;
   if (n == 2) return dist(p[0], p[1]);
   double ans = 0:
   int i = 0, j = 1;
   while (i < n) {</pre>
       while (cross(p[(i + 1) % n] - p[i], p[(j + 1) % n] - p[j]) >= 0) {
              ans = max(ans, dist2(p[i], p[j]));
              j = (j + 1) \% n;
       ans = max(ans, dist2(p[i], p[j]));
       i++:
   }
```

```
return sqrt(ans);
vector<PT> convex_hull(vector<PT> &p) {
       if (p.size() <= 1) return p;</pre>
       vector < PT > v = p;
   sort(v.begin(), v.end());
   vector<PT> up, dn;
   for (auto& p : v) {
       while (up.size() > 1 && orientation(up[up.size() - 2], up.back(),
            \} (0 = < (q
           up.pop_back();
       while (dn.size() > 1 && orientation(dn[dn.size() - 2], dn.back(),
            p) <= 0) {
           dn.pop_back();
       }
       up.push_back(p);
       dn.push_back(p);
   }
   v = dn;
   if (v.size() > 1) v.pop_back();
   reverse(up.begin(), up.end());
   up.pop_back();
   for (auto& p : up) {
       v.push_back(p);
   if (v.size() == 2 && v[0] == v[1]) v.pop_back();
   return v;
void test_case() {
   11 n:
   cin >>n;
   vector<PT> p(n);
   for (int i = 0; i < n; i++) cin >> p[i].x >> p[i].y;
   p = convex_hull(p);
   cout << fixed<<setprecision(10) << diameter(p) << "\n";</pre>
```

10 fast-hadamard-transform

// like polynomial multiplication, but XORing exponents

```
// instead of adding them (also ANDing, ORing)
const int MAXN=1<<18;</pre>
#define fore(i,1,r) for(int i=int(1);i<int(r);++i)</pre>
#define SZ(x) ((int)(x).size())
ll c1[MAXN+9],c2[MAXN+9];//MAXN must be power of 2!
void fht(ll* p, int n, bool inv){
   for(int l=1;2*l<=n;1*=2)for(int i=0;i<n;i+=2*l)fore(j,0,1){</pre>
       11 u=p[i+j],v=p[i+l+j];
       // if(!inv)p[i+j]=u+v,p[i+l+j]=u-v; // XOR
       // else p[i+j]=(u+v)/2, p[i+l+j]=(u-v)/2;
       //if(!inv)p[i+j]=v,p[i+l+j]=u+v; // AND
       //else p[i+j]=-u+v,p[i+l+j]=u;
       if(!inv)p[i+j]=u+v,p[i+l+j]=u; // OR
       else p[i+j]=v,p[i+l+j]=u-v;
   }
}
// like polynomial multiplication, but XORing exponents
// instead of adding them (also ANDing, ORing)
vector<ll> multiply(vector<ll>& p1, vector<ll>& p2){
   int n=1<<(32-__builtin_clz(max(SZ(p1),SZ(p2))-1));</pre>
   fore(i,0,n)c1[i]=0,c2[i]=0;
   fore(i,0,SZ(p1))c1[i]=p1[i];
   fore(i,0,SZ(p2))c2[i]=p2[i];
   fht(c1,n,false);fht(c2,n,false);
   fore(i,0,n)c1[i]*=c2[i];
   fht(c1,n,true);
   return vector<ll>(c1,c1+n);
}
// maxime the OR of a pair of given nums and count
// how many pairs can get that maximum OR
// tested: https://csacademy.com/contest/archive/task/maxor
void test_case() {
   ll n; cin >> n;
   vl a(MAXN),b(MAXN);
   for (int i =0;i<n;i++) {</pre>
       11 x;
       cin >> x;
       a[x]++;
       b[x]++;
   }
   vl c = multiply(a,b);
   pair<11,11> best = {0, c[0]};
```

```
for (int i = 0;i<MAXN;i++) {
    if (c[i]) best = {i,(c[i]-a[i])/2};
}
cout <<best.F << " " << best.S << endl;
}</pre>
```

11 general₁ azy_tree

```
struct Node {
    11 mn:
    11 \text{ size} = 1;
    Node(ll mn):mn(mn) {
}:
struct Func {
   11 a = 0;
};
Node e() { // op(x, e()) = x
    Node a(INT64_MAX);
    return a;
};
Func id() { // \text{ mapping}(x, id()) = x}
    Func 1 = \{0\};
    return 1;
}
Node op(Node &a, Node &b) { // associative property
    Node c = e();
    c.size = a.size + b.size:
    c.mn = min(a.mn, b.mn);
    return c;
Node mapping(Node node, Func &lazy) {
    node.mn += lazy.a;
    return node:
}
```

```
Func composicion(Func &prev, Func &actual) {
   prev.a = prev.a + actual.a;
   return prev;
}
struct lazytree {
   int n;
   vector<Node> nodes;
   vector<Func> lazy;
   void init(int nn) {
       n = nn:
       int size = 1;
       while (size < n) {</pre>
           size *= 2;
       11 m = size *2;
       nodes.assign(m, e());
       lazy.assign(m, id());
   }
   void push(int i, int sl, int sr) {
       nodes[i] = mapping(nodes[i], lazy[i]);
       if (sl != sr) {
           lazy[i * 2 + 1] = composicion(lazy[i*2+1],lazy[i]);
           lazy[i * 2 + 2] = composicion(lazy[i*2+2],lazy[i]);
       }
       lazy[i] = id();
   }
   void apply(int i, int sl, int sr, int l, int r, Func f) {
       push(i, sl, sr);
       if (1 <= sl && sr <= r) {</pre>
           lazy[i] = f;
           push(i,sl,sr);
       } else if (sr < 1 || r < sl) {</pre>
       } else {
           int mid = (sl + sr) >> 1;
           apply(i * 2 + 1, sl, mid, l, r, f);
           apply(i * 2 + 2, mid + 1, sr, 1, r, f);
           nodes[i] = op(nodes[i*2+1],nodes[i*2+2]);
       }
   }
   void apply(int 1, int r, Func f) {
```

```
assert(1 <= r):
       assert(r < n);
       apply(0, 0, n - 1, 1, r, f);
   void update(int i, Node node) {
       assert(i < n);</pre>
       update(0, 0, n-1, i, node);
   }
   void update(int i, int sl, int sr, int pos, Node node) {
       if (sl <= pos && pos <= sr) {</pre>
           push(i,sl,sr);
           if (sl == sr) {
               nodes[i] = node;
           } else {
               int mid = (sl + sr) >> 1;
               update(i * 2 + 1, sl, mid, pos, node);
               update(i * 2 + 2, mid + 1, sr, pos, node);
               nodes[i] = op(nodes[i*2+1], nodes[i*2+2]);
           }
       }
   }
   Node query(int i, int sl, int sr, int l, int r) {
       push(i,sl,sr);
       if (1 <= s1 && sr <= r) {</pre>
           return nodes[i];
       } else if (sr < 1 || r < sl) {</pre>
           return e():
       } else {
           int mid = (sl + sr) >> 1;
           auto a = query(i * 2 + 1, sl, mid, l, r);
           auto b = query(i * 2 + 2, mid + 1, sr, 1, r);
           return op(a,b);
       }
   }
   Node query(int 1, int r) {
       assert(1 <= r);</pre>
       assert(r < n);
       return query(0, 0, n - 1, 1, r);
   }
};
```