1/6/2015 5.-IROs and the PIT

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5. IRQs and the PIT

In this chapter we're going to be learning about interrupt requests (IRQs) and the programmable interval timer (PIT).

5.1. Interrupt requests (theory)

There are several methods for communicating with external devices. Two of the most useful and popular are polling and interrupting.

Polling

Spin in a loop, occasionally checking if the device is ready.

Interrupts

Do lots of useful stuff. When the device is ready it will cause a CPU interrupt, causing your handler to be run.

As can probably be gleaned from my biased descriptions, interrupting is considered better for many situations. Polling has lots of uses - some CPUs may not have an interrupt mechanism, or you may have many devices, or maybe you just need to check so infrequently that it's not worth the hassle of interrupts. Any rate, interrupts are a very useful method of hardware communication. They are used by the keyboard when keys are pressed, and also by the programmable interval timer (PIT).

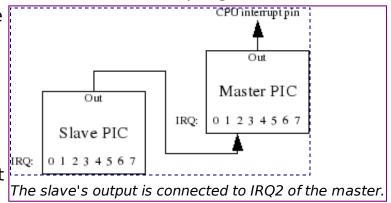
The low-level concepts behind external interrupts are not very complex. All devices that are interrupt-capable have a line connecting them to the PIC (programmable interrupt controller). The PIC is the only device that is directly connected to the CPU's interrupt pin. It is used as a multiplexer, and has the ability to prioritise between interrupting devices. It is

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essentially, a glorified 8-1 multiplexer. At some point, someone somewhere realised that 8 IRQ lines just wasn't enough, and they daisy-chained another 8-1 PIC beside the original. So in all modern PCs, you have 2 PICs, the master and the slave, serving a total of 15 interruptable devices (one line is used to signal the slave PIC).

The other clever thing about the PIC is that you can change the interrupt number it delivers for each IRQ line. This is referred to as *remapping the PIC* and is actually extremely useful. When the computer boots, the default interrupt mappings are:



- IRQ 0..7 INT 0x8..0xF
- IRO 8..15 INT 0x70..0x77

This causes us somewhat of a problem. The master's IRQ mappings (0x8-0xF) conflict with the interrupt numbers used by the CPU to signal exceptions and faults (see last <u>chapter</u>). The normal thing to do is to remap the PICs so that IRQs 0..15 correspond to ISRs 32..47 (31 being the last CPU-used ISR).

5.2. Interrupt requests (practical)

The PICs are communicated with via the I/O bus. Each has a command port and a data port:

- Master command: 0x20, data: 0x21
- Slave command: 0xA0, data: 0xA1

The code for remapping the PICs is the most difficult and obfusticated. To remap them, you have to do a full reinitialisation of them, which is why the code is so long. If you're interested in what's actually happening, there is a nice description <u>here</u>.

```
static void init_idt()
{
    ...
    // Remap the irq table.
    outb(0x20, 0x11);
    outb(0xA0, 0x11);
    outb(0x21, 0x20);
    outb(0xA1, 0x20);
```

```
UULU(UXAI, UXZO);
outb(0x21, 0x04):
outb(0xA1, 0x02):
outb(0x21, 0x01):
outb(0xA1, 0x01);
outb(0x21, 0x0);
outb(0xA1, 0x0);
idt set gate(32, (u32int)irq0, 0x08, 0x8E);
idt set gate(47, (u32int)irg15, 0x08, 0x8E);
```

Notice that now we are also setting IDT gates for numbers 32-47, for our IRQ handlers. We must, therefore, also add stubs for these in interrupt.s. Also, though, we need a new macro in interrupt.s - an IRQ stub will have 2 numbers associated with it - it's IRO number (0-15) and it's interrupt number (32-47):

```
; This macro creates a stub for an IRQ - the first parameter is
; the IRQ number, the second is the ISR number it is remapped to.
%macro IRO 2
 global irg%1
  irg%1:
   cli
   push byte 0
   push byte %2
    imp irg common stub
%endmacro
    Ο,
IR0
           32
IR0
    1.
         33
```

```
IRO 15.
```

We also have a new common stub - irg common stub. This is because IRQs behave subtly differently - before you return from an IRQ handler, you must inform the PIC that you have finished, so it can dispatch the next (if there is one waiting). This is known as an EOI (end of interrupt). There is a slight complication though. If the master PIC sent the IRQ (number 0-7), we must send an EOI to the master (obviously). If the slave sent the IRO (8-15), we must send an EOI to both the master and the slave (because of the daisy-chaining of the two).

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First our asm common stub. It is almost identical to *isr_common_stub*.

```
: In isr.c
[EXTERN irg handler]
; This is our common IRQ stub. It saves the processor state, sets
; up for kernel mode segments, calls the C-level fault handler,
; and finally restores the stack frame.
irq common stub:
                            ; Pushes edi, esi, ebp, esp, ebx, edx, ecx, eax
   pusha
                            ; Lower 16-bits of eax = ds.
   mov ax, ds
                            ; save the data segment descriptor
   push eax
   mov ax, 0x10 ; load the kernel data segment descriptor
   mov ds, ax
   mov es, ax
   mov fs, ax
   mov qs, ax
   call irq_handler
                  ; reload the original data segment descriptor
   pop ebx
   mov ds, bx
   mov es, bx
   mov fs, bx
   mov qs, bx
   popa ; Pops edi,esi,ebp...
add esp, 8 ; Cleans up the pushed error code and pushed ISR number
   sti
                  ; pops 5 things at once: CS, EIP, EFLAGS, SS, and ESP
```

Now the C code (goes in isr.c):

```
// This gets called from our ASM interrupt handler stub.
void irq_handler(registers_t regs)
{
    // Send an EOI (end of interrupt) signal to the PICs.
    // If this interrupt involved the slave.
    if (regs.int_no >= 40)
    {
```

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```
// Send reset signal to slave.
   outb(0xA0, 0x20);
}
// Send reset signal to master. (As well as slave, if necessary).
outb(0x20, 0x20);

if (interrupt_handlers[regs.int_no] != 0)
{
   isr_t handler = interrupt_handlers[regs.int_no];
   handler(regs);
}
```

This is fairly straightforward - if the IRQ was > 7 (interrupt number > 40), we send a reset signal to the slave. In either case, we send one to the master also.

You may also notice that I have added a small custom handler mechanism, allowing you to register custom interrupt handlers. This can be very useful as an abstraction technique, and will neaten up our code nicely.

Some other declarations are needed:

5.2.1. isr.h

```
// A few defines to make life a little easier
#define IRQ0 32
...
#define IRQ15 47

// Enables registration of callbacks for interrupts or IRQs.
// For IRQs, to ease confusion, use the #defines above as the
// first parameter.
typedef void (*isr_t)(registers_t);
void register_interrupt_handler(u8int n, isr_t handler);
```

5.2.2. isr.c

```
isr_t interrupt_handlers[256];
void register_interrupt_handler(u8int n, isr_t handler)
```

```
{
    interrupt_handlers[n] = handler;
}
```

And there we go! We can now handle interrupt requests from external devices, and dispatch them to custom handlers. Now all we need is some interrupt requests to handle!

5.3. The PIT (theory)

The programmable interval timer is a chip connected to IRQ0. It can interrupt the CPU at a user-defined rate (between 18.2Hz and 1.1931 MHz). The PIT is the primary method used for implementing a system clock and the only method available for implementing multitasking (switch processes on interrupt).

The PIT has an internal clock which oscillates at approximately 1.1931MHz. This clock signal is fed through a <u>frequency</u> <u>divider</u>, to modulate the final output frequency. It has 3 channels, each with it's own frequency divider.

- Channel 0 is the most useful. It's output is connected to IRQ0.
- Channel 1 is very un-useful and on modern hardware is no longer implemented. It used to control refresh rates for DRAM.
- Channel 2 controls the PC speaker.

Channel 0 is the only one of use to us at the moment.

OK, so we want to set the PIT up so it interrupts us at regular intervals, at frequency f. I generally set f to be about 100Hz (once every 10 milliseconds), but feel free to set it to whatever you like. To do this, we send the PIT a 'divisor'. This is the number that it should divide it's input frequency (1.9131MHz) by. It's dead easy to work out:

```
divisor = 1193180 Hz / frequency (in Hz)
```

Also worthy of note is that the PIT has 4 registers in I/O space - 0x40-0x42 are the data ports for channels 0-2 respectively, and 0x43 is the command port.

5.4. The PIT (practical)

```
// timer.h -- Defines the interface for all PIT-related functions.
// Written for JamesM's kernel development tutorials.
#ifndef TIMER H
#define TIMER H
#include "common.h"
void init timer(u32int frequency);
#endif
And timer.c doesn't have much in either:
// timer.c -- Initialises the PIT, and handles clock updates.
// Written for JamesM's kernel development tutorials.
#include "timer.h"
#include "isr.h"
#include "monitor.h"
u32int tick = 0:
static void timer callback(registers t regs)
   tick++:
   monitor write("Tick: ");
   monitor write dec(tick);
   monitor write("\n");
void init timer(u32int frequency)
   // Firstly, register our timer callback.
   register interrupt handler(IRQO, &timer callback);
   // The value we send to the PIT is the value to divide it's input clock
   // (1193180 Hz) by, to get our required frequency. Important to note is
   // that the divisor must be small enough to fit into 16-bits.
   u32int divisor = 1193180 / frequency;
   // Send the command byte.
   outh(0x43 0x36)
```

```
// Divisor has to be sent byte-wise, so split here into upper/lower bytes.
u8int l = (u8int)(divisor & 0xFF);
u8int h = (u8int)( (divisor>>8) & 0xFF );

// Send the frequency divisor.
outb(0x40, l);
outb(0x40, h);
}
```

OK, lets go through this code. Firstly, we have our *init_timer* function. This tells our interrupt mechanism that we want to handle IRQ0 with the function *timer_callback*. This will be called whenever a timer interrupt is recieved. We then calculate the divisor to be sent to the PIT (see theory above). Then, we send a command byte to the PIT's command port. This byte (0x36) sets the PIT to repeating mode (so that when the divisor counter reaches zero it's automatically refreshed) and tells it we want to set the divisor value.

We then send the divisor value. Note that it must be sent as two seperate bytes, not as one 16-bit value.

When this is done, all we have to do is edit our Makefile, add one line to main.c

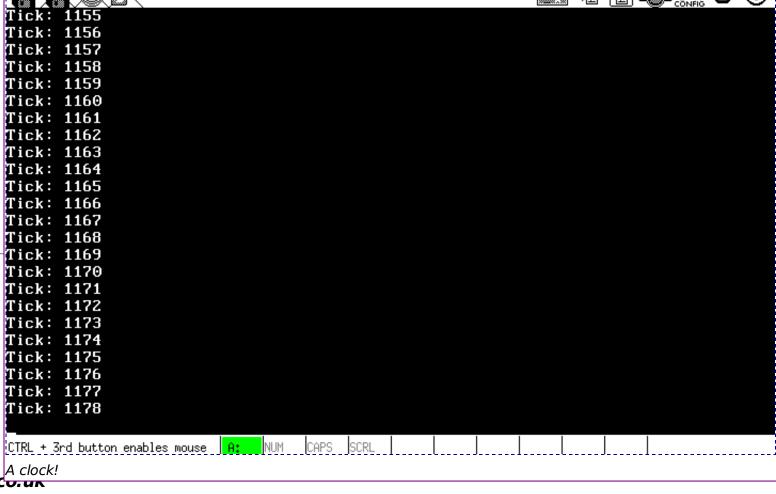
```
init_timer(50); // Initialise timer to 50Hz
```

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compile, and run! You should get output like that on the right. Note however that bochs does not accurately emulate the timer chip, so although your code will run at the correct speed on a real machine, it probably won't in bochs!

Full source code for this tutorial can be found here.

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